Lecture ⁴

·

Oct ⁸ , ²⁰²⁴

$$
\begin{array}{lll}\n\mathcal{D}_{\vert\rho} & \text{in} & \text{app} & \mathcal{U} & \text{then each state updates to} \\
\mathcal{U}|\mathcal{Y}_{c}\rangle. & \text{So} & \\
& & \mathcal{P}_{c} & \sum_{a} p_{a} \mathcal{U}_{a} | \mathcal{Y}_{a} \times \mathcal{Y}_{a} | \mathcal{U}^{+} \\
& & \mathcal{U} \left(\sum_{a} p_{a} | \mathcal{Y}_{a} \times \mathcal{Y}_{a} | \mathcal{U}^{+} \right) \mathcal{U}^{+} \\
& & \mathcal{U} \left(\sum_{a} p_{a} | \mathcal{Y}_{a} \times \mathcal{Y}_{a} | \mathcal{U}^{+} \right) \mathcal{U}^{+} \\
& & \mathcal{U} \left(\mathcal{U}_{c} & \mathcal{U}_{c} \right) & \\
& & \mathcal{U} \left(\mathcal{U}_{c} & \mathcal{U}_{c} \right) & \\
& & \mathcal{U} \left(\mathcal{U}_{c} & \mathcal{U}_{c} \right) & \\
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& & \mathcal{U} \left(\mathcal{U}_{c} & \mathcal{U}_{c} \right) & \\
& & \mathcal{U} \left(\mathcal{U}_{c} & \mathcal{U}_{c} \right) & \\
& & \mathcal{U} \
$$

$$
p_{a}^{b} - n_{tunneled.m.} = (10 \times 010 \text{ L}) |4a \times 4. | (10 \times 010 \text{ L})
$$
\n
$$
p'_{a} = (10 \times 010 \text{ L}) |4a \times 4. | (10 \times 010 \text{ L})
$$
\n
$$
p_{t} + (10 \times 010 \text{ L}) |4 \times 4. |
$$
\n
$$
p_{t} + 100 \text{ L}
$$
\n
$$
p'_{t} + p_{t} + 5 \text{ mod}
$$
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$$
p'_{t} + p_{t} + 5 \text{ mod}
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\n
$$
p'_{t} + p_{t} + 5 \text{ mod}
$$
\n
$$
p'_{a} = 5 \text{ pr}(a | 7 = 0) \cdot p'_{a}
$$
\n
$$
= 5 \text{ pr}(a | 7 = 0) \cdot (10 \times 010 \text{ L}) |4a \times 4a| (10 \times 010 \text{ L})
$$
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$$
= 5 \text{ pr}(7 = 0) \cdot (10 \times 010 \text{ L}) |4a \times 4a| (10 \times 010 \text{ L})
$$
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= \text{ [10000 L)} \cdot (10 \times 010 \text{ L}) \cdot (10 \times 010 \text{ L})
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= \text{ [10000 L)} \cdot (10 \times 010 \text{ L}) \cdot (10 \times 010 \text{ L})
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= 5 \text{ pr}(7 = 010) \cdot p_{a}
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= 5 \text{ pr}(7 = 010) \cdot p_{a}
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= 5 \text{ pr}(7 = 010) \cdot p_{a}
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= 5 \text{ pr}(7 = 010) \cdot p_{a}
$$
\n
$$
= 5 \text{ pr}(10 \times 010 \text{ L}) \cdot (10 \times 010 \text{ L})
$$

$$
= \frac{(10\times01041) |4a\times4a| (10\times01041)}{4r((10\times0104) p)}
$$

Exercise Show that
$$
tr(\rho^2) = 1
$$
 iff ρ is pure (i.e.)
 $\rho = 145(41)$.

$$
\mathbb{P}f. \quad If \quad \rho = |\psi\rangle\langle\psi|, \text{ then } \rho^2 = |\psi\rangle\langle\psi|\psi\rangle\langle\psi| = \rho.
$$
\n
$$
\mathcal{S}_{\nu} + r(\rho^2) = tr(\rho) = 1, \text{ powo} \Leftarrow \text{obiction}.
$$

For
$$
\Rightarrow
$$
 direction, write ρ in eigenbasis so $\rho = \sum \lambda_i |\psi_i \times \psi_i|$
with $\{|\psi_i\rangle\}$ orthonormal and $\sum \lambda_i = 1$ since $+ \sigma(\rho) = 1$.
 $+ \sigma(\rho^*) = + \sigma(\sum_{ij} \lambda_i \lambda_j \psi_i \times \psi_i |\psi_j \times \psi_i|)$
 $= + \sigma(\sum_i \lambda_i^2 |\psi_i \times \psi_i|) = \sum \lambda_i^2 \le \sum \lambda_i = 1$

with equality only if
$$
\lambda_i^2 = \lambda_i
$$
, $\forall i$, so $\lambda_i = 1$
\n $\lambda_{2...} = 0$.
\nso $\rho \cdot |\Psi_i \times \Psi_i|$ pure.
\nAside: $\rho \cdot \Sigma \lambda_i |\Psi_i \times \Psi_i|$ means $|\Psi_i \rangle$ we $p_i \lambda_i$.
\nAside: For $\neg \Sigma \lambda_i |\Psi_i \times \Psi_i|$ means $|\Psi_i \rangle$ we $p_i \lambda_i$.
\nMultiply do one need to understand. Obviously, making
\n \neg Denoting the original state of
\na system which is entangled.
\nFor instance, when Alice ϕ Bob hold an EPR
\npair, what is the state of Alice system by
\nitself?
\nNotation: $|\varphi|$ ρ_{AB} is the state of the
\nsystem and $\mathbb{C}_{AB}^{d_B}$
\nThis is the state of the
\nSystom now $\mathbb{C}_{AB}^{d_B}$ $\mathbb{C}_{AB}^{d_B}$
\nAiles's gubb Bals, gubib,

How call
$$
\rho_A
$$
 the stato on Alia's system.

\n ρ_B the s.taton Bob's system.

\nZasy case: $\rho_{AB} = \rho_A \otimes \rho_B$.

\nThen ρ_A is obvious.

\nHow about $\rho_{AB} = \sum P_i \rho_A^i \otimes \rho_B^i$ { ρ_{B}^i } ρ_{B} .

\nThus we expect $\rho_A = \sum P_i \rho_A^i$

\n $\rho_B = \sum P_i \rho_B^i$.

It becomes harder when the states are entangled and not just prob. mixtures .

What are the properties that
$$
\rho_A
$$
 should have?
\n① The statistics of measuring ρ_A should
\nbe identified to that of meaning only

Alice's qubits of
$$
\rho_{AB}
$$
.

\n(2) Let $\sigma_{AB} = (U_A \otimes \underline{U}) \rho_{AB} (U_A^+ \otimes \underline{U})$

\nThen $\sigma_A = U_A \rho_A U_A^+$. Meaning, unitary transform by Alice can be calculated from just ρ_A .

\n(3) Any action on Bub's qubits control change the stata of Alice's qubits.

$$
\begin{array}{ll}\n\text{Def } & \text{Partial} + \text{Rec} \\
\text{for } \text{Bib's qubits/system,} & \text{thus for any matrix } \text{PAB} \\
\text{tr}_{B}(\rho_{AB}) & \text{if } \sum_{i=1}^{d_{B}} (\mathbb{1}_{A} \otimes \langle b_{i}|) \rho_{AB} (\mathbb{1}_{A} \otimes |b_{i} \rangle) \\
\end{array}
$$

If
$$
Ph_B
$$
 is a density matrix, what is the operational
meaning of $H_B(P_{AB})$?

Ans: Bob measures his system according to basis $\{1b_i>\}$

but does not tell Alive the outcome. Note It remains to show that this dep does not depend on the choice of basis. Ex . Par ⁼ (E) (0-1)(2014) * = (Inil) Pa (11: 7) -(10 <01) (101 - 102) ((011 - (101) (10 10) + (10())(10k - 1ro))((01) - (01)(10(1)] : (((1))(- (1) ⁺ (10) ((x)] : Laim Pa ⁼ tris(Pab).

Why is this the carrect definition?

First let us prove that the def of partial true down
clipard on choice of
$$
\{16:>3\}
$$
.
Ex. If $\langle \Psi | \rho | \psi \rangle = \langle \Psi | \rho' | \Psi \rangle$ for all unit vectors $|\Psi \rangle$
then $\rho = \rho'$.
For any $|\Psi \rangle$, let $\rho_B^{\Psi} := (\langle \Psi_A | \Phi \mathbb{1}_B) \rho_{AB} (|\Psi_A \rangle \otimes \mathbb{1}_B)$.
Now,

$$
\langle\psi| tr_{B}(\rho_{AB}) |\psi\rangle
$$

\n= $\langle\psi_{A}|\left(\sum_{i=1}^{d_{B}}(\mathbb{1}_{A}\otimes\langle b_{i}|)\rho_{AB}(\mathbb{1}_{A}\otimes b_{i})\right)|\psi_{A}\rangle$
\n= $\sum_{i=1}^{d_{B}}\langle b_{i}| \rho_{B}^{\psi}|b_{i}\rangle = \sum_{i=1}^{d_{B}}tr(\rho_{B}^{\psi})$

which is independent of $\{|b_i\rangle\}$ as trace is basis indep.

Out V₈ be any unitary. Let
$$
1/3
$$
... $1d_8$ be basis for B.

\n
$$
+r_{B}((1_{A} \otimes V_{B}) \rho_{AB} (1_{A} \otimes V_{B}^{+}))
$$
\n
$$
= \sum_{i=1}^{d_8} (1_{A} \otimes \langle i | V_{B} \rangle) \rho_{AB} (1_{B} \otimes V_{B}^{+} | i \rangle)
$$
\n
$$
= \sum_{i=1}^{d_9} (1_{A} \otimes \langle i | V_{B} \rangle) \rho_{AB} (1_{B} \otimes V_{B}^{+} | i \rangle)
$$
\n
$$
= tr_{B}(\rho_{AB}) \leftarrow
$$
\n
$$
- 1 + tr_{B}(\rho_{AB}) \leftarrow
$$
\n
$$
= \sum_{i=1}^{d_8} (1_{A} \otimes \langle i | \rangle) \rho_{AB} (1_{A} \otimes 1_{B})
$$
\n
$$
= \sum_{i=1}^{d_8} (1_{A} \otimes \langle i | \rangle) \rho_{AB} (1_{A} \otimes 1_{B}^{+})
$$
\n
$$
= \sum_{i=1}^{d_8} (1_{A} \otimes \langle i | \rangle) \rho_{AB} (1_{A} \otimes 1_{B}^{+})
$$
\n
$$
= tr_{B}(\eta_{AB} \eta_{AB} + \eta_{AB} \eta_{AB}
$$

Def. The maximally mixed state in
$$
d
$$
 - dimensions
\nis $\frac{1}{d}$ 11_d. For qubit, $\frac{1}{2^n}$ 11_{2^n}.
\nDef. A state ρ_{AB} between Alice and Bob event of n
\ngubits is maximally entangled if $\rho_A = \frac{1}{2^n}$ 11_{2^n}.

is
$$
\frac{1}{d} \mathcal{A}_{d}
$$
. For qubit, $\frac{1}{2^n} \mathcal{A}_{2^n}$
Def. A state β_{AB} between Alice and Bob each of n
qubits is maximally entangled if $\beta_A = \frac{1}{2^n} \mathcal{A}_{2^n}$.
In class exercise: Every maximally entangled static
is $\sum_{i=1}^{2^n} P_i (\mathcal{A}_A \otimes \mathcal{U}_B^i) \mathbb{P} \times \mathbb{E} \left(\mathcal{A}_A \otimes \mathcal{U}_B^{i^+} \right)$
where $|\mathcal{D}| = \sum_{i=1}^{2^n} \frac{1}{\sqrt{2^n}} |i \rangle |i \rangle$.

Today : The CHSH game and "spooky-action" at ^a distance . Consider the following game : Fightenlights , ^g Alice Ref Bob

Ref samples
$$
x, y \in \{0, 1\}
$$
 and sends x do Alice
and y to Bob. He asks for answers $a, b \in \{0, 1\}$ from
Alia and Bob, respectively.

Alice and Bob mot answer in 2+6 lightnes.

$$
T_{\text{hy}} \text{ win if } \text{ a} \bullet \text{ b} = \alpha \cdot \gamma.
$$

What is their opt. success prob.?

Space-separation is just to fore no communication.

If Alice and Bob employ classical strategies, the best
stackey only this w pr = 3/4.
But if they we quantum strategies, they can min with
probabilities up to
$$
cos^2 \pi/8 \sim 0.878
$$
.
Next two lectures: understand the quantum strategy,
prove it is optimal, understood how it choosing,
prove it is optimal, understood how it characteristics the
state of Alia and Bob's q.c.'s.
Classical Best alternative strategy a= 0, b= 0
independent of input.
What if Aliu & Bob share classical randomness?
Let $A_{r_1}B_{r_1}: [0,1]^3 \rightarrow [0,1^3]$ be the strategies for mod. C.
The
From
Pr[win] = P_r [A_r(x) ∅ B_r(y) = xy]

 $= \sum_{r} Pr[r] \cdot Pr\left[A_{r}(x) \otimes B_{r}(y) = xy\right]$

$$
\frac{2}{5} \text{ max } P_r \left[A_r(x) \oplus B_r(y) = xy \right]
$$
\n
$$
\frac{2}{5} \text{ max } P_r \left[A_r(x) \oplus B_r(y) = xy \right]
$$
\n
$$
\frac{1}{2} \frac{2}{3} \frac{1}{4} \frac{1}{4
$$

Easiest to draw a picture.

$$
ar\vartheta
$$
 $\langle \theta | \theta + \gamma \rangle = \omega_{s} \gamma$.

(2) For any basis
$$
|w\rangle
$$
, $|w\rangle$ of \mathbb{C}^2 ,
\n $|EPR\rangle = \frac{1}{\sqrt{2}} (|v\rangle|v^*\rangle + |w\rangle|w^*\rangle) \quad (m \;hw2)$

Using these ² facts and the angles between recting we can analyze the success of this game.

Using three 2 facts and the angles between
vechts we can analyze the success of this game.

$$
Pr[\omega in] = \sum_{x,y} Pr[x,y] \cdot \sum_{a,b} Pr[a,b|x,y] \cdot \underline{d} \cdot \sum_{a,b} a \cdot b = xy \}
$$

$$
= \frac{1}{4} \left[\left| \left\langle \alpha_{0}^{\circ}, \beta_{i}^{\circ} | \text{EPR} \right\rangle \right|^{2} + \left| \left\langle \alpha_{0}^{\circ}, \beta_{i}^{\circ} | \text{EPR} \right\rangle \right|^{2} \quad (x,y) = (0,0)
$$

+ $\left| \left\langle \alpha_{0}^{\circ}, \beta_{1}^{\circ} | \text{EPR} \right\rangle \right|^{2} + \left| \left\langle \alpha_{0}^{\circ}, \beta_{1}^{\circ} | \text{EPR} \right\rangle \right|^{2} \quad (x,y) = (0,1)$
+ $\left| \left\langle \alpha_{1}^{\circ}, \beta_{i}^{\circ} | \text{EPR} \right\rangle \right|^{2} + \left| \left\langle \alpha_{1}^{\circ}, \beta_{i}^{\circ} | \text{EPR} \right\rangle \right|^{2} \quad (x,y) = (0,1)$
+ $\left| \left\langle \alpha_{1}^{\circ}, \beta_{1}^{\circ} | \text{EPR} \right\rangle \right|^{2} + \left| \left\langle \alpha_{1}^{\circ}, \beta_{1}^{\circ} | \text{EPR} \right\rangle \right|^{2} \quad (x,y) = (1,0)$

Recognize

$$
Recegnive
$$
\n
$$
\langle \alpha_{x}^{c}, \beta_{y}^{b} | EPR \rangle = \langle \alpha_{x}^{a}, \beta_{y}^{b} | \cdot \frac{|\beta_{y}^{b}\rangle|\beta_{y}^{c}\rangle + |\beta_{y}^{b}\rangle|\beta_{y}^{c}\rangle
$$
\n
$$
= \frac{1}{\sqrt{2}} \langle \alpha_{x}^{a} | \beta_{y}^{b} \rangle
$$
\n
$$
= \frac{1}{\sqrt{2}} \langle \alpha_{x}^{a} | \beta_{y}^{b} \rangle
$$
\n
$$
\langle \alpha_{x}^{a} | \beta_{y}^{b} \rangle
$$
\n $$

 $= 0.854$

So Alice and Bob can win $\sum_{i=1}^{n}$ with quantum strategies with highes prob than classical.

"Claser, Haven, Shimony , Holt" (CMSH) experiment.

aka Ball inequality experiment.

aka a non-local y exp
game

We know that q mechanics says that n-qubit
states can be described by a
$$
2^n
$$
 -dim vector. Meaning,
that if nature is doing a Tot of computation
"behind the seems"

But value is that
$$
inf_{0}^{1}
$$
 should?

CHSH game powers that
$$
\pi
$$
 picture is incorrect.
\nThe information about the static of the system cannot
\nbe hold locally assuming a spad limit on information.
\nThis is because local info could not produce the
\nabatisfies needed to win the game width prob = 3π .
\nTo describe the static of the system, we must do
\nhade $\rho = IER \times EPR \Big|_{\rho}$ from $\rho_A = \rho_B * \frac{1}{2} \pm \frac{1}{2}$.
\nThe key is that the coin are comebated!
\nCan we exponentially verify GISH?

Can actually conduct this and hevs been done.
\nDoes it have that quantum mechanics is correct?
\n
$$
\frac{N}{m}
$$
. Since the 9: advantage is only probabilistic
\n $\frac{N}{m}$. Since the 9: advantage is only probabilistic
\n $\frac{1}{1 + 15}$ evidence that an world isn't classical.
\nSuppne a ref observes Alice & Bob winning the
\ngame 802 of the time. How many game would
\nthey have to play before he is coming with 1-10°?
\nconfidence That is varied isn't classical?

$$
Clernoff \; bounds: \quad X = \sum_{i}^{N} X_{i} \leftarrow per gene success
$$

$$
\Pr[X \geq (1+\Sigma)\mu] \leq \exp\left(-\frac{S^2\mu}{2\cdot \delta}\right)
$$

$$
\mu = 0.75n
$$
 $\delta = \frac{0.8}{0.75} < 1.07$

$$
P\left\{ X \geq 0.8 \, n \right\} \leq \exp\left(-0.27 \, n\right)
$$

$$
w_{\text{out}} = \exp(-0.27n) \le 10^{-9}
$$

0.27n ≥ 9 ln 10
n ≥ 77 .

Does it prove he world is getur? No For all we know , dure is a stronger thory consistent with quantum mechanics. Ex . Noisy telepatory ! Alice and Bob are telepathic with ^a success rate of ²⁰ %. Herfect telepaly and they win with probability 1 .

Is
$$
cos^2 \pi / 8
$$
, the optimal question strategy?
Can ne du buffer if Alice and Bob shore multiple
EPR pairs?