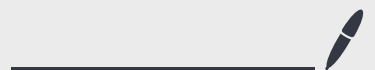


Lecture 15

Nov 14, 2024



Proving local Hamiltonian is QMA-hard:

Given input circuit  $C = g_T \dots g_1$  acting on  $n$  qubit input,  $m$  ancilla.

Construct a local Ham  $H$  and quantities  $(a, b)$ ,

①  $\exists |\psi\rangle$  s.t.  $C(|\psi\rangle)$  accepts w pr  $\geq 1 - \epsilon$ ,

then  $\exists |\psi\rangle$  s.t.  $\langle \psi | H | \psi \rangle \leq a$

②  $\forall |\psi\rangle$ ,  $C(|\psi\rangle)$  accepts w pr  $\leq \epsilon$ ,

then  $\forall |\psi\rangle$ ,  $\langle \psi | H | \psi \rangle \geq b$ .

In class, we will only prove that  $\underbrace{O(\log n)}_{=k}$ -LH is QMA-hard.

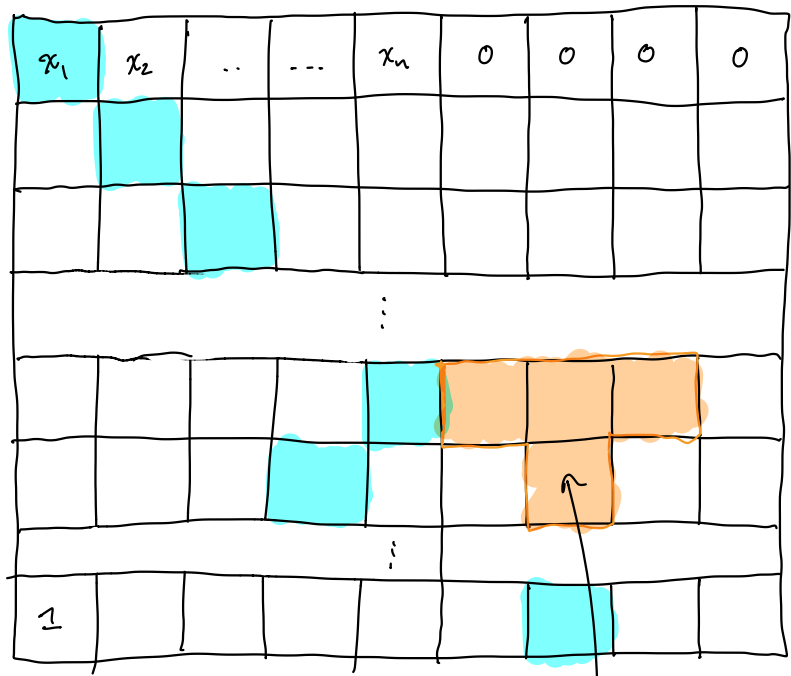
HW will include problem for proving  $S$ -local hardness.

Quantum analog of the Cook-Levin theorem proving that 3-SAT is NP-complete.

Cook-Levin Tableau Review:

given classical computation of  $T$  time steps using total space  $S$ ,  
write a  $T \cdot S$  tableau with the states of the machine at  
time  $t$  in row  $t$ .

checks for ancilla  
↓



 = where  
TM  
head is.

↑  
check for accept

check that evolution of  
TM was appropriately listed.

Can we construct the same for quantum comp.?

Consider a circuit  $C = g_T \dots g_1$  can we create a table with rows

$$|\Psi_0\rangle = |\Psi_{\text{witness}}\rangle \otimes |0^m\rangle$$

$$|\Psi_1\rangle = g_1 |\Psi_0\rangle$$

$$|\Psi_2\rangle = g_2 g_1 |\Psi_0\rangle = g_2 |\Psi_1\rangle$$

⋮

$$|\Psi_T\rangle = C |\Psi_0\rangle.$$

Why does this not work? Local checks cannot verify the evolutions.

Ex.  $|\Psi_t\rangle = \frac{|0^n\rangle + |1^n\rangle}{\sqrt{2}}$  and  $g_{t+1} = Z_1.$

Then  $|\Psi_{t+1}\rangle = \frac{|0^n\rangle - |1^n\rangle}{\sqrt{2}}.$

Whereas if  $g_{t+1} = \mathbb{1}$ , then  $|\Psi_{t+1}\rangle = \frac{|0^n\rangle + |1^n\rangle}{\sqrt{2}}.$

Claim Any  $k \leq (n-1)$  reduced density matrix of  $\frac{|0^n\rangle \pm |1^n\rangle}{\sqrt{2}}$  is  $\frac{1}{2} (|0^k\rangle\langle 0^k| + |1^k\rangle\langle 1^k|).$

So a check differentiating if  $g_{t+1} = Z_1$  vs.  $\mathbb{1}$  must be non-local if it can distinguish

$$|\Psi_t\rangle \otimes |\Psi_{t+1}\rangle \quad \text{from} \quad |\Psi_t\rangle \otimes |\Psi_t\rangle$$

$$g_t = Z_1. \quad \quad \quad g_t = \mathbb{1}$$

We need a better solution that can detect global changes that occur from local gates.

Let's create a  $O(\log n)$ -local Hamiltonian whose ground states

are of the form  $\frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle \otimes |\Psi_t\rangle$  ← state register

↑ time register

$t$  expressed in  $\lceil \log(T+1) \rceil$  bits

and  $|\Psi_t\rangle = g_t \dots g_1 |\Psi_0\rangle$  (all related).

In order to write out the Ham. let's first understand

$$h = \frac{|1\rangle\langle 1| \otimes \mathbb{1} - |1\rangle\langle 0| \otimes U - |0\rangle\langle 1| \otimes U^\dagger + |0\rangle\langle 0| \otimes \mathbb{1}}{2}$$

$$h = \frac{1}{2} \begin{pmatrix} \mathbb{1} & -u^\dagger \\ -u & \mathbb{1} \end{pmatrix} = \frac{1}{2} \begin{pmatrix} -u^\dagger \\ \mathbb{1} \end{pmatrix} \begin{pmatrix} -u & \mathbb{1} \end{pmatrix}$$

Then  $\langle \psi | h | \psi \rangle$  for  $h = |0\rangle\langle\psi_0| + |1\rangle\langle\psi_1| \leftarrow$  unnormalized eqns.

$$\frac{1}{2} \left( \langle\psi_0| \quad \langle\psi_1| \right) \begin{pmatrix} \mathbb{1} & -u^\dagger \\ -u & \mathbb{1} \end{pmatrix} \begin{pmatrix} |\psi_0\rangle \\ |\psi_1\rangle \end{pmatrix}$$

$$= \frac{1}{2} \left( \langle\psi_0| \quad \langle\psi_1| \right) \begin{pmatrix} |\psi_0\rangle - u^\dagger |\psi_1\rangle \\ -u |\psi_0\rangle + |\psi_1\rangle \end{pmatrix}$$

$$= \frac{1}{2} \left( \langle\psi_0|\psi_0\rangle - \langle\psi_0|u^\dagger|\psi_1\rangle - \langle\psi_1|u|\psi_0\rangle + \langle\psi_1|\psi_1\rangle \right)$$

$$= \frac{1}{2} \left\| |\psi_1\rangle - u |\psi_0\rangle \right\|^2.$$

$h$  measures distance from states of the form

$$\frac{1}{\sqrt{2}} |0\rangle|\psi_0\rangle + \frac{1}{\sqrt{2}} |1\rangle u |\psi_0\rangle.$$

where  $|\psi_0\rangle$  is of norm 1.

Alternate proof of fact:

$$\text{Let } V = |1\rangle\langle 1| \otimes U + |0\rangle\langle 0| \otimes \mathbb{1}$$

$$V^\dagger = |1\rangle\langle 1| \otimes U^\dagger + |0\rangle\langle 0| \otimes \mathbb{1}$$

$$\begin{aligned} V^\dagger h V &= \frac{|1\rangle\langle 1| \otimes \mathbb{1} - |1\rangle\langle 0| \otimes \mathbb{1} - |0\rangle\langle 1| \otimes \mathbb{1} + |0\rangle\langle 0| \otimes \mathbb{1}}{2} \\ &= |-\rangle\langle -| \otimes \mathbb{1}. \end{aligned}$$

So groundstates of  $V^\dagger h V$  are states of the form

$$|+\rangle \otimes |\psi_0\rangle.$$

and  $V^\dagger h V$  is a projector. So groundstates of  $h$

are states of the form  $V|+\rangle \otimes |\psi_0\rangle$

$$= \frac{1}{\sqrt{2}} \left( |0\rangle|\psi_0\rangle + |1\rangle U|\psi_0\rangle \right).$$

Apply this intuition to general Hamiltonian circuit

$$h_t = \frac{1}{2} \left( |+\rangle\langle +| \otimes \mathbb{1} - |t\rangle\langle t-1| \otimes g_t - |t-1\rangle\langle t| \otimes g_t^\dagger + |t-1\rangle\langle t-1| \otimes \mathbb{1} \right)$$

Each term  $h_t$  acts on time register + 2 extra qubits def. by  $g_t$ .

Some calculation will tell us that for  $|\psi\rangle = \sum_t |t\rangle |\psi_t\rangle$

that

$$\langle \psi | h_t | \psi \rangle = \left\| |\psi_t\rangle - g_t |\psi_{t+1}\rangle \right\|^2.$$

But how do all the pieces act together? Analyze a situation.

$$V = \sum_{t=0}^T |t\rangle \langle t+1| \otimes g_t \dots g_1$$

$$V^\dagger = \sum_{t=0}^T |t+1\rangle \langle t| \otimes g_1^\dagger \dots g_t^\dagger.$$

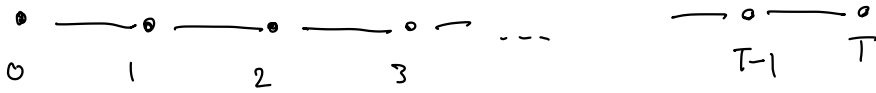
$$\begin{aligned} V^\dagger h_t V &= \frac{1}{2} \left( |t\rangle \langle t+1| \otimes \mathbb{1} - |t+1\rangle \langle t+1| \otimes \mathbb{1} - |t\rangle \langle t-1| \otimes \mathbb{1} + |t-1\rangle \langle t-1| \otimes \mathbb{1} \right) \\ &= \frac{1}{\sqrt{2}} \left( |t\rangle - |t-1\rangle \right) \frac{1}{\sqrt{2}} \left( \langle t+1| - \langle t-1| \right) \otimes \mathbb{1}. \end{aligned}$$

$$\text{So, } V^\dagger \left( \sum_{t=1}^T h_t \right) V =$$

$$\frac{1}{2} \left( \begin{array}{cccc} 1 & -1 & & \\ -1 & 2 & -1 & \\ & \ddots & \ddots & \ddots \\ & & & -1 & 2 & -1 \\ & & & & \ddots & \ddots \\ & & & & & -1 & 1 \end{array} \right) \otimes \mathbb{1}.$$



This is a special matrix. It is the Laplacian of a linear graph and also a circulant matrix. (also tri-diagonal)



$$\lambda_0 = 0, \text{ eigenvector } |\Psi_0\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle \otimes |\Psi_0\rangle \quad \leftarrow \text{for any } |\Psi_0\rangle$$

$$\lambda_1 \geq \frac{c}{T^2} \leftarrow \text{exercise/intuition from graph mixing time.}$$

So, removing rotation by  $V$ :

$$H_{\text{prop}} = \sum_{t=0}^T h_t \text{ is a Hamiltonian with}$$

$$\text{ground-energy} = 0, \text{ groundstates of the form } \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle \otimes g_t \dots g_1 |\Psi_0\rangle$$

for any states  $|\Psi_0\rangle$

known as a history state

highly degenerate groundspace.

and first non-zero energy of  $\geq c/T^2$ .

$$\text{Therefore, } H_{\text{prop}} \geq \frac{c}{T^2} (\mathbb{1} - \Pi_{\text{prop}})$$

projector onto span of history states

Next: Adding checks for ancilla and output and computing overall eigenvalues to make sure cheating provers are detected.

History states just check evolution of the computation

To check output of computation, we want a term that checks that when the time register is set to  $T$ , the first qubit is in the  $|1\rangle$  state.

$$H_{out} = |T\rangle\langle T|_{time} \otimes |0\rangle\langle 0|_1$$

↗ Penalizes being time  $T$  and first qubit = 0.

To check input, we need to make sure that all the ancilla registers of the computation were set to  $|0\rangle$ .

$$H_{in} = \sum_{j=1}^m |0\rangle\langle 0|_{time} \otimes |1\rangle\langle 1|_{ancilla(j)}$$

$H_{out} + H_{in}$  is a commuting Hamiltonian meaning every pair of local terms commute.

Each local term is a projector. So  $H_{out} + H_{in}$  has an integer spectrum.

$$H_{out} + H_{in} \geq (\mathbb{1} - \Pi_{inout})$$

projector onto passing all in & out checks.

Claim  $H = H_{prop} + H_{out} + H_{in}$  suffices as a Ham.

Intuition: Let  $|\Psi\rangle = \sum_{t=0}^T |t\rangle |\Psi_t\rangle$  ← unnormalized.

$$\text{so } \|\Psi\rangle\| = 1 \Rightarrow \sum_{t=0}^T \|\Psi_t\rangle\|^2$$

$h_t$  gives energy  $\|\Psi_t\rangle - g_t |\Psi_{t-1}\rangle\|^2$

$h_{in}$  gives energy  $\sum_{j=1}^m \|\langle 1 | \text{ancilla}_{(j)} | \Psi_0 \rangle\|^2$

$H_{out}$  gives energy  $\|\langle 0 | \Psi_T \rangle\|^2$

Recall  $V = \sum_{t=0}^{\infty} |t\rangle\langle t| \otimes g_t \dots g_1.$

Let  $|\mu\rangle = V^\dagger |\psi\rangle = \sum_t |t\rangle |\mu_t\rangle$

where  $|\mu_t\rangle = g_1^\dagger \dots g_t^\dagger |\psi_t\rangle.$

Then

$H_f$  gives energy  $\| |\mu_t\rangle - |\mu_{t-1}\rangle \|^2$

$H_{in}$  gives energy  $\sum_{j=1}^m \| \langle 1 |_{\text{ancilla}(j)} |\mu_0\rangle \|^2$

$H_{out}$  gives energy  $\| \langle 0 |_C |\mu_T\rangle \|^2$

If satisfiable then  $\exists$  good choice  $|\mu\rangle$  for all  $|\mu_0\rangle \dots |\mu_T\rangle.$

When unsatisfiable, in order for  $H_{in}$  and  $H_{out}$

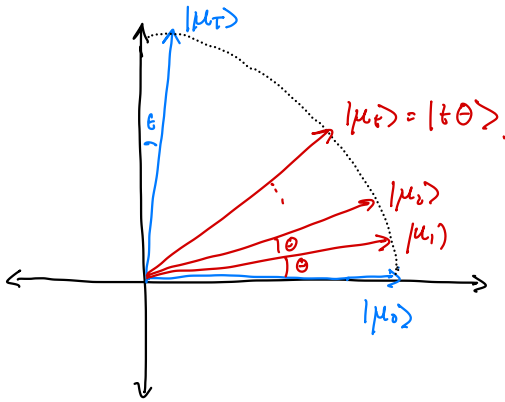
energy to be small...  $|\mu_0\rangle$  and  $|\mu_T\rangle$  are orthogonal.

For intuition only: The following is not the lowest energy state.

Cartoon:  $|\mu_0\rangle = |0\rangle$  and  $|\mu_1\rangle \approx |1\rangle$  in some 2D space

(can be formal due to Jordan's lemma)

Then ideal  $|\mu_t\rangle$  would be  $|\pm\theta\rangle$  for  $\theta = \frac{\pi}{2T}$  to minimize  $H_{\text{prop}}$  energy.



$$\begin{aligned}
 \text{Then } \langle \psi | H | \psi \rangle &= \langle \mu | V^\dagger H V | \mu \rangle \\
 &\leq \langle \mu | V^\dagger h_t V | \mu \rangle \\
 &= T \cdot \sin^2 \theta \\
 &= O\left(\frac{1}{T}\right).
 \end{aligned}$$

We will prove a lower bound of  $\Omega\left(\frac{1}{T^3}\right)$  in the next class.