# Entropy and Shannon information

For a random variable *X* with distribution *p*(*x*), the **entropy** is

 $H[X] = -\Sigma_x p(x) \log_2 p(x)$ 

#### **Information** is defined as

 $I[X] = -\log_2 p(x)$ 

#### Typically, "information" = *mutual information*:

how much knowing the value of one random variable *r* (the response) reduces uncertainty about another random variable *s* (the stimulus).

Variability in response is due both to different **stimuli** and to **noise**. How much response variability is "useful", i.e. can represent different messages, depends on the noise. Noise can be specific to a given stimulus.

#### Information quantifies how *independent* r and s are:

 $I(S;R) = D_{KL} [P(R,S), P(R)P(S)]$ 

Alternatively:

 $I(S;R) = H[R] - \Sigma_s P(s) H[R|s].$ 

Mutual information is the difference between the total response entropy and the mean noise entropy:

 $I(S;R) = H[R] - \Sigma_s P(s) H[R|s)].$ 

→ Need to know the conditional distribution P(s|r) or P(r|s).

Take a particular stimulus  $s=s_0$  and repeat many times to obtain P(r|s<sub>0</sub>). Compute variability due to noise: *noise entropy*  Information is symmetric in r and s

Extremes:

1. response is unrelated to stimulus: p[r|s] = ?, MI = ?

2. response is perfectly predicted by stimulus: p[r|s] = ?

r<sub>+</sub> encodes stimulus +, r<sub>-</sub> encodes stimulus -

but with a probability of error:  $P(r_+|+) = 1 - p$  $P(r_-|-) = 1 - p$ 

What is the response entropy H[r]?

What is the noise entropy?

## **Entropy and Shannon information**



 $H[r] = -p_{+} \log p_{+} - (1-p_{+}) \log(1-p_{+})$ 

When  $p_{+} = \frac{1}{2}$ , H[r|s] = -p log p - (1-p)log(1-p)

### Noise limits information

#### A communication channel $S \rightarrow R$ is defined by P(R|S)

$$I(S;R) = \sum_{s,r} P(s) P(r|s) \log[P(r|s)/P(r)]$$

The **channel capacity** gives an upper bound on transmission through the channel:

 $C(R|S) = \sup I(S;R)$ 

#### Perfect decodability through the channel:



If the entropy of T is less than the channel capacity, then T' can be perfectly decoded to recover T.

#### Transform S by some function F(S):

$$R \xrightarrow{encode} S \xrightarrow{transmit} F(S)$$

The transformed variable F(S) cannot contain more information about R than S.

How can one compute the entropy and information of spike trains?

**Entropy:** 

Discretize the spike train into binary words w with letter size  $\Delta t$ , length T. This takes into account correlations between spikes on timescales T $\Delta t$ .

Compute  $p_i = p(w_i)$ , then the naïve entropy is

$$S_{\text{naive}}(T, \Delta \tau; \text{size}) = -\sum_{i} \tilde{p}_i \log_2 \tilde{p}_i;$$

Strong et al., 1997; Panzeri et al.



Many information calculations are limited by sampling: hard to determine P(w) and P(w|s)

Systematic bias from undersampling.

Correction for finite size effects:

$$S_{\text{naive}}(T, \Delta \tau; \text{size}) = S(T, \Delta \tau) + \frac{S_1(T, \Delta \tau)}{\text{size}} + \frac{S_2(T, \Delta \tau)}{\text{size}^2}.$$



*Information* : difference between the variability driven by stimuli and that due to noise.

Take a stimulus sequence *s* and repeat many times.

For each time in the repeated stimulus, get a set of words P(w|s(t)).

Average over  $s \rightarrow$  average over time:

$$H_{noise} = \langle H[P(w|s_i)] \rangle_i.$$

Choose length of repeated sequence long enough to sample the noise entropy adequately.

Finally, do as a function of word length T and extrapolate to infinite T.







### Calculating information in the LGN

Another example: temporal coding in the LGN (Reinagel and Reid '00)



# Calculating information in the LGN



Apply the same procedure: collect word distributions for a random, then repeated stimulus.



# Information in the LGN

Use this to quantify how precise the code is, and over what timescales correlations are important.



#### How much information does a single spike convey about the stimulus?

Key idea: the information that a spike gives about the stimulus is the reduction in entropy between the distribution of spike times not knowing the stimulus, and the distribution of times knowing the stimulus.

The response to an (arbitrary) stimulus sequence  $\mathbf{s}$  is r(t).

Without knowing that the stimulus was  $\mathbf{s}$ , the probability of observing a spike in a given bin is proportional to  $\bar{r}$ , the mean rate, and the size of the bin.

Consider a bin  $\Delta t$  small enough that it can only contain a single spike. Then in the bin at time t,

$$P(r = 1) = \bar{r}\Delta t,$$
  

$$P(r = 0) = 1 - \bar{r}\Delta t,$$
  

$$P(r = 1|s) = r(t)\Delta t,$$
  

$$P(r = 0|s) = 1 - r(t)\Delta t.$$

Now compute the entropy difference:  $p = \bar{r}\Delta t$ ,  $p(t) = r(t)\Delta t$ .  $I(r,s) = -p \log p - (1-p) \log(1-p) + \leftarrow \text{prior}$   $+ \frac{1}{T} \int_0^T dt \ [p(t) \log p(t) + (1-p(t)) \log(1-p(t))]$ .  $\leftarrow \text{ conditional}$ Note substitution of a time average for an average over the *r* ensemble.

Assuming  $p \ll 1$ ,  $\log(1-p) \sim p$  and using  $\frac{1}{T} \int_0^T dt \, p(t) \to p$ 

$$I(r,s) = \frac{1}{T} \int_0^T dt \,\Delta t \, r(t) \log \frac{r(t)}{\bar{r}} + Var(p(t))/2ln2 + O(p^3).$$

In terms of information per spike (divide by  $ar{r}\Delta t$  ):

$$I(r,s) = \frac{1}{T} \int_0^T dt \, \frac{r(t)}{\bar{r}} \log \frac{r(t)}{\bar{r}}$$

Given 
$$I(r,s) = \frac{1}{T} \int_0^T dt \, \frac{r(t)}{\bar{r}} \log \frac{r(t)}{\bar{r}}$$

note that:

- It doesn't depend explicitly on the stimulus
- The rate *r* does not have to mean rate of spikes; rate of any event.
- Information is limited by spike precision, which blurs r(t), and the mean spike rate.

Compute as a function of  $\Delta t$ :

Undersampled for small bins



## Adaptation and coding efficiency



















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# Efficient coding

In order to encode stimuli effectively, an encoder should match its outputs to the statistical distribution of the inputs



Fly visual system



Laughlin, '81

# Variation in time

Contrast varies hugely in time.

Should a neural system optimize over evolutionary time or locally?





### Time-varying stimulus representation





For fly neuron H1, determine the input/output relations throughout the stimulus presentation



A. Fairhall, G. Lewen, R. R. de Ruyter and W. Bialek (2001)

## Barrel cortex



Extracellular *in vivo* recordings of responses to whisker motion in rat S1 barrel cortex in the anesthetized rat

M. Maravall et al., (2007)



# Single cortical neurons



R. Mease, A. Fairhall and W. Moody, J. Neurosci.

## Using information to evaluate coding



# Adaptive representation of information



As one changes the characteristics of *s*(t), changes can occur both in the *feature* and in the *decision function* 

Barlow '50s, Laughlin '81, Shapley et al, '70s, Atick '91, Brenner '00

### Feature adaptation



Barlow '50s, Laughlin '81, Shapley et al, '70s, Atick '91, Brenner '00

The information in any given event can be computed as:

$$I(E;s) = \left\langle \left(\frac{r_E(t)}{\bar{r}_E}\right) \log_2\left(\frac{r_E(t)}{\bar{r}_E}\right) \right\rangle_s,$$

Define the synergy, the information gained from the joint symbol:

$$Syn[E_1, E_2; s] = I[E_1, E_2; s] - (I[E_1; s] + I[E_2; s]).$$

or equivalently,

 $Syn[E_1, E_2; s] = I[E_1; E_2|s] - I[E_1; E_2].$ 

Negative synergy is called *redundancy*.

In the identified neuron H1, compute information in a spike pair, separated by an interval dt:



 $Syn[E_1, E_2; s] = I[E_1, E_2; s] - (I[E_1; s] + I[E_2; s]).$ 

Brenner et al., '00.