

HW3

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Instructions

- You should think about each problem by yourself for at least an hour before choosing to collaborate with others.
- You are allowed to collaborate with fellow students taking the class in solving the problems. But you **must** write your solution on your own.
- You are not allowed to search for answers or hints on the web. You are encouraged to contact the instructor or the TAs for a possible hint.
- You cannot collaborate on Extra credit problems
- Solutions typeset in LATEX are preferred.
- Feel free to use the Discussion Board or email the instructor or the TAs if you have any questions or would like any clarifications about the problems.
- Please upload your solutions to Gradescope.

In solving these assignments, feel free to use these approximations:

$$1 - x \approx e^{-x}, \quad \sqrt{1-x} \approx 1 - x/2, \quad n! \approx (n/e)^n, \quad \left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \left(\frac{en}{k}\right)^k$$

The problems of this HW are hard. To get full points it is enough to solve one of the following problems, although you are encouraged to try both.

- 1) Let $G = (V, E)$ be an undirected graph and suppose each $v \in V$ is associated with a set $S(v)$ of $8r$ colors, where $r \geq 1$. Suppose further that for each $v \in V$ and $c \in S(v)$ there are at most r neighbors u of v such that c lies in $S(u)$.
 - Prove that there exists a proper coloring of G assigning to each vertex v a color from its class $S(v)$ such that, for any edge $u \sim v$, the colors assigned to u and v are different.
 - Design a randomized polynomial time algorithms. Feel free to assume $|S(v)| \geq c \cdot r$ (for all v) for a constant $c \gg 8$.
- 2) Let $\mathcal{A}_1, \dots, \mathcal{A}_n$ be a set of bad events defined on independent variables Z_1, \dots, Z_n with corresponding dependency graph G (ignore conditions of the LLL for this problem). Note that $\mathbb{P}[\neg \mathcal{A}_i]$ is uniquely defined as a function of the probabilities of Z_i 's. Prove or disprove: In the Moser-Tardos algorithm, for a fixed $t > 1$ and $i \in [n]$, the probability that the t -th resampled event is \mathcal{A}_i is at most $\mathbb{P}[\mathcal{A}_i]$. Prove this or find a counterexample.