

Martingales

Today
- more martingales
Azuma-Hoeffding
- start online decision-making

Sequence of r.v.s X_0, X_1, X_2, \dots called a discrete time martingale

- $E(|X_n|) < \infty$
- • $E(X_{n+1} | X_0, X_1, \dots, X_n) = X_n$
- $E(X_{n+1} - X_n | X_0, \dots, X_n) = 0$

A sequence of r.v.s X_0, X_1, \dots is a martingale with respect to the sequence Y_0, Y_1, \dots if $\forall n \geq 0$ the following conditions hold:

- X_n is a fn of Y_0, Y_1, \dots, Y_n
 - $E(|X_n|) < \infty$
 - • $E(X_{n+1} | Y_0, \dots, Y_n) = X_n$
- think of Y_0, \dots, Y_n as information up to time n

Examples

① Sums of indep random variables

$Y_0 = 0$ Y_1, \dots, Y_n iid w/ $E(Y_k) = 0 \forall k$
 Define $X_n = Y_0 + Y_1 + Y_2 + \dots + Y_n$
 $\{X_n\}$ is a martingale wrt. $\{Y_n\}$

② "Doob's" martingale process

Y_1, Y_2, \dots arbitrary seq of random vars

X r.v. with finite expectation

$X_n = E(X | Y_1, \dots, Y_n)$ forms martingale wrt $\{Y_n\}$

$X_0 = E(X)$

Example: Edge exposure martingale

$G(n, p)$ random graph
label $m = \binom{n}{2}$ potential edges e_1, e_2, \dots, e_m

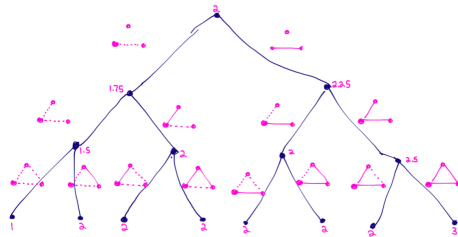
Let $f(G)$ be some function of the graph $f: 2^{\binom{n}{2}} \rightarrow \mathbb{R}$
which edges of graph

$Y_j = \begin{cases} 1 & \text{if edge } e_j \text{ present} \\ 0 & \text{otherwise} \end{cases} \quad \Pr(Y_j=1) = p$

$X_k = E[f(G) | Y_1, \dots, Y_k]$ $X_0 = E[f(G)]$
 $X_m = f(G) = E(f(G) | Y_1, \dots, Y_m)$

Example: $f(G)$: sized max clique

$G(n, \frac{1}{2})$



$X_0 = E(\text{max clique})$

$X_1 = E(\text{max clique} | Y_1)$

$X_2 = E(\text{max clique} | Y_1, Y_2)$

$X_3 = E(\text{max clique} | Y_1, Y_2, Y_3)$

Some useful facts about martingales:

① $E(X_n) = E(X_0)$

by induction

$$E(X_{n+1} | Y_0, \dots, Y_n) = X_n$$

$$E[E(X_{n+1} | Y_0, \dots, Y_n)] = E(X_n)$$

$$= E(X_{n+1})$$

$$E(E(X|Y)) \stackrel{E(X|Y=y) \cdot \Pr(Y=y)}{=} \sum_y E(X|Y=y) \Pr(Y=y)$$

$$= E(X)$$

② Definition

A r.v. T is called a "stopping time" wrt $\{Y_t\}$ if

T takes values in $\{0, 1, 2, \dots\}$

and if $\forall n \geq 0$, the event $\{T \leq n\}$ is determined by Y_0, \dots, Y_n

i.e. can determine if $T \leq n$ or $T > n$ from knowledge of values Y_0, \dots, Y_n

"know it when you see it"

Examples

- first time I win 5 games in row
- first time I win \$100

Non-example:

- last time I win 5 games in a row

Optional Stopping Thm

$\{Z_t\}$ is a martingale wrt $\{X_t\}$

For T a stopping time "know it when you see it"

$$E(Z_T) = E(Z_0)$$

whenever any of the following hold

- Z_t 's bounded ($\exists c$ s.t. $\forall i |Z_i| \leq c$)
- T is bounded
- $E(T) < \infty$ and $\exists c$ s.t. $E(|Z_{t+1} - Z_t| | X_{1..t}, X_t) \leq c$

③ Tail inequalities

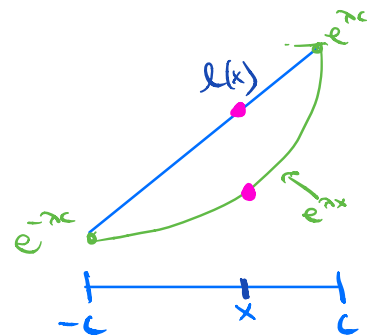
$E(X_n) = E(X_0)$ how far can it be from its expectation

Azuma-Hoeffding Inequality

X_0, \dots, X_n martingale s.t. $\forall k \quad |X_k - X_{k-1}| \leq c_k$
 c_k may depend on k

Then $\forall t \geq 0$, any $R > 0$

$$\Pr(|X_t - X_0| > R) \leq 2 e^{-\left[\frac{R^2}{2 \sum_{k=1}^t c_k^2}\right]}$$



Fact: X r.v. s.t. $|X| \leq c$
 and $E(X) = 0$

Then $E(e^{\lambda X}) \leq e^{\frac{(\lambda c)^2}{2}}$

Proof: By convexity of $f(x) = e^{\lambda x}$
 for any $x \in [-c, c]$, we have

$$e^{\lambda x} \leq \frac{(1 - \frac{x}{c})e^{-\lambda c} + (1 + \frac{x}{c})e^{\lambda c}}{2} = l(x)$$

If X with $E(X) = 0$ and $|X| \leq c$

then $E(e^{\lambda X}) \leq E[l(X)] = \frac{e^{-\lambda c} + e^{\lambda c}}{2} = \sum_{k=0}^{\infty} \frac{(\lambda c)^{2k}}{(2k)!}$

$$\frac{1}{2} \left[\begin{array}{l} 1 - \lambda c + \frac{(\lambda c)^2}{2!} - \frac{(\lambda c)^3}{3!} + \dots \\ 1 + \lambda c + \frac{(\lambda c)^2}{2!} + \frac{(\lambda c)^3}{3!} + \dots \end{array} \right]$$

$$\leq \sum_{k=0}^{\infty} \frac{(\lambda c)^{2k}}{2^{2k} k!} = e^{\frac{(\lambda c)^2}{2}}$$

Corollary: $E(X_{t+1} - X_t | H_t) = 0$ $|X_{t+1} - X_t| \leq c \Rightarrow E[e^{\lambda(X_{t+1} - X_t)} | H_t] \leq e^{\frac{(\lambda c)^2}{2}}$

$$\begin{aligned}
 E[e^{\lambda X_{t+1}} | H_t] &= E[e^{\lambda(X_{t+1}-X_t)} e^{\lambda X_t} | H_t] \\
 &= e^{\lambda X_t} E[e^{\lambda(X_{t+1}-X_t)} | H_t] \\
 &\leq e^{\lambda X_t} e^{\frac{\lambda^2 c_t^2}{2}}
 \end{aligned}$$

\uparrow
 X_0, \dots, X_t

n.v. $E[e^{\lambda(X_{t+1}-X_t)} | X_0, \dots, X_t]$
 \uparrow
 X_0, \dots, X_t

\uparrow
 X_0, \dots, X_t

\uparrow
 $\Pr(X_{t+1}=x_t)$

$$\begin{aligned}
 &E(X|Y=y) \cdot p \cdot \Pr(Y=y) \\
 &E(E(X|Y)) \\
 &= \sum_y E(X|Y=y) \Pr(Y=y) \\
 &= E(X)
 \end{aligned}$$

\Rightarrow taking expectations on both sides

$$E[e^{\lambda X_{t+1}}] \leq E[e^{\lambda X_t}] e^{\frac{\lambda^2 c_t^2}{2}}$$

so by induction $E[e^{\lambda X_{t+1}}] \leq E[e^{\lambda X_0}] e^{\lambda^2 \sum_{i=0}^t \frac{c_i^2}{2}}$

$$\begin{aligned}
 \text{Finally, } \Pr(X_t - X_0 \geq R) &= \Pr(e^{\lambda(X_t - X_0)} \geq e^{\lambda R}) \leq e^{-\lambda R} E[e^{\lambda(X_t - X_0)}] \\
 &\leq e^{-\lambda R} e^{\lambda^2 \sum_{i=0}^{t-1} \frac{c_i^2}{2}}
 \end{aligned}$$

\uparrow
 Markov's Ineq

Optimizing, we choose $\lambda = \frac{R}{\sum_{i=1}^t c_i^2}$

$$\Rightarrow \Pr(X_t - X_0 \geq R) \leq e^{-\frac{R^2}{2 \sum_{i=1}^t c_i^2}}$$

$$-\left(\lambda R - \frac{\lambda^2 \sum c_i^2}{2}\right) = -\left(\frac{R^2}{\sum c_i^2} - \frac{R^2}{2 \sum c_i^2}\right)$$

Factor of 2 comes from $\Pr(X_t - X_0 < -R)$

Applications:

① n Balls in n bins # empty bins

B_1, B_2, \dots, B_n where B_i is bin that ball i goes into.

X : # empty bins.

$$X_k = E(X | B_1, \dots, B_k)$$

$$|X_k - X_{k-1}| \leq 1$$

$$\Pr(|X_n - X_0| > c\sqrt{n}) \leq 2e^{-\frac{cn}{2n}} = 2e^{-\frac{c}{2}}$$

\uparrow X \uparrow $E(X)$

take $R = c\sqrt{n}$.

$$E(X) = n \left(1 - \frac{1}{n}\right)^n \approx \frac{n}{e}$$

$$X = \frac{n}{e} \pm O(\sqrt{n})$$

Azuma-Hoeffding Inequality

X_0, \dots, X_n martingale st. $\forall k \quad |X_k - X_{k-1}| \leq c_k$
 c_k may depend on k

Then $\forall t \geq 0$, any $R > 0$
 $\Pr(|X_t - X_0| > R) \leq 2e^{-\frac{R^2}{\sum_{k=1}^t c_k^2}}$

$Z_i = 1$ bin i empty.

$$X = \sum_{i=1}^n Z_i$$

② Chromatic # in random graph $G(n, \frac{1}{2})$

Vertex exposure martingale

$$X_k = E[\chi(G) | N(v_1), N(v_2), \dots, N(v_k)]$$

$N(v_i)$ = edges from v_i to v_1, \dots, v_{i-1}

$$X_0 = E[\chi(G)]$$

$$X_n = \chi(G)$$

$$|X_k - X_{k-1}| \leq 1$$



apply A-H

$$\Pr(|\chi(G) - E(\chi(G))| > R) \leq 2e^{-\frac{R^2}{2n}}$$

$$R = \omega(\sqrt{n}) \rightarrow 0$$

② Finding "interesting" patterns (e.g. in DNA seqs)

Let $X = (X_1, \dots, X_n)$ be sequence of characters chosen independently & u.a.r. from Σ $|\Sigma| = s$ Pr(A)Pr(G)Pr(T)Pr(C) = 1

e.g. $\Sigma = \{A, T, C, G\}$

Let $B = (b_1, \dots, b_k)$ fixed string of characters AATAATGAC

F r.v. = # occurrences of B in string X

$F_i = E(F | X_1, \dots, X_i)$ Doob martingale

$F_0 = E(F)$ $F_n = F$

$$E(F) = (n-k+1) \frac{1}{s^k}$$

$|\Sigma| = s$
each char equally likely

$$|F_i - F_{i+1}| \leq k \quad \text{each char in at most } k \text{ matches}$$

$$nk^2 = \sum_{i=1}^n c_i^2$$

\Rightarrow By Azuma-Hoeffding

$$\Pr\left(\left| \frac{F_n}{F_0} - E\left(\frac{F_n}{F_0}\right) \right| \geq \lambda\right) \leq 2e^{-\frac{\lambda^2}{2nk^2}}$$

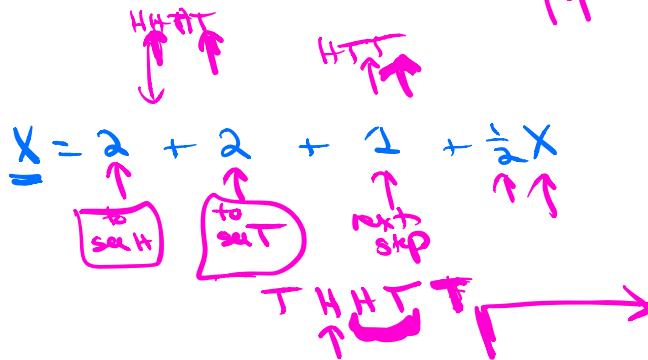
\Rightarrow for $\lambda = ck\sqrt{n}$

$$\Pr\left(\left| \frac{F_n}{F_0} - E\left(\frac{F_n}{F_0}\right) \right| \geq ck\sqrt{n}\right) \leq 2e^{-\frac{c^2}{2}}$$

One more cute (and useful) application of O.S.T.

Toss fair coin

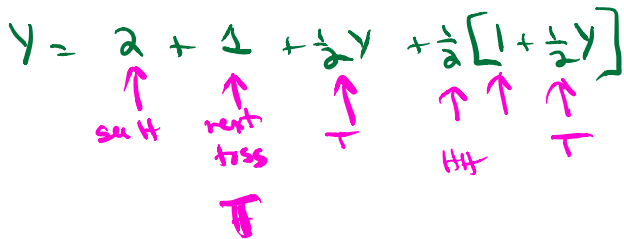
$X = \text{Exp \# steps to see } \underline{HTH?}$



$$X = 2 + 2 + 1 + \frac{1}{2}X$$

$$\frac{X}{2} = 5 \Rightarrow X = 10$$

$Y = \text{Exp \# steps to see } \underline{HHH?}$



$$Y = 2 + 1 + \frac{1}{2}Y + \frac{1}{2} \left[1 + \frac{1}{2}Y \right]$$

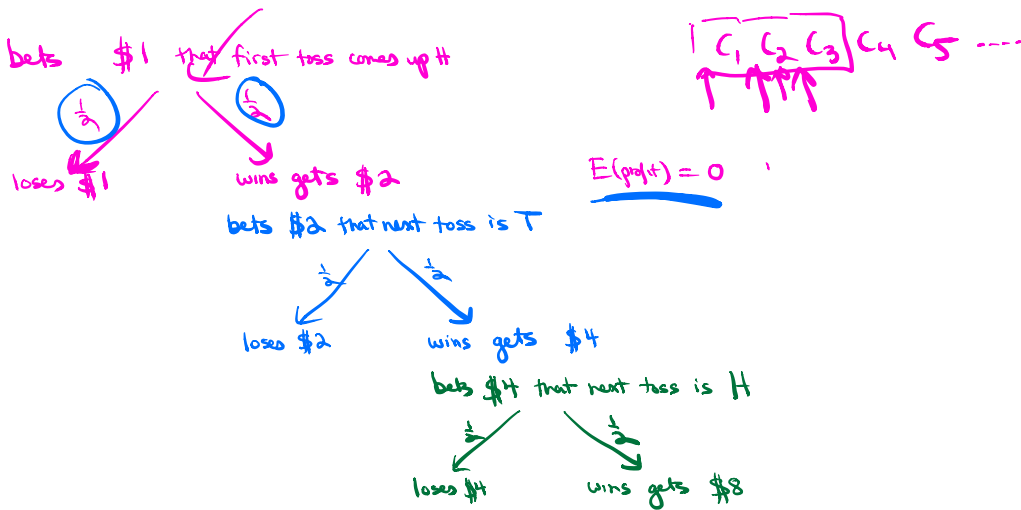
$$Y = 3\frac{1}{2} + \frac{3}{4}Y \Rightarrow Y = 14$$

A martingale approach (for any pattern σ)

e.g. $\sigma = HTH$

Sequence of indep coin tosses C_1, C_2, C_3, \dots $C_i = \begin{cases} H & \text{up } \frac{1}{2} \\ T & \text{up } \frac{1}{2} \end{cases}$

At each time step t , a new gambler arrives
- makes a series of double or nothing bets on σ



Let X_t exp profit of all gamblers up to step t $X_0 = 0$

$\{X_t\}$ is a martingale w.r.t $\{C_t\}$

$$E(X_{t+1} | C_1, \dots, C_t) = X_t$$

$$X_{t+1} = \sum_{j=1}^{t+1} \text{profit of gambler that arrived at beginning of step } j \text{ upto end of step } t+1$$

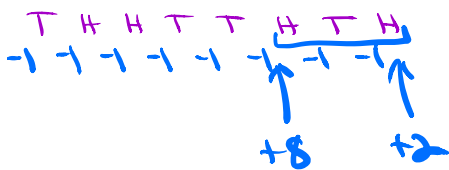


T be first time see $\sigma = HTH$ $E(X_T) = E(X_0) = 0$

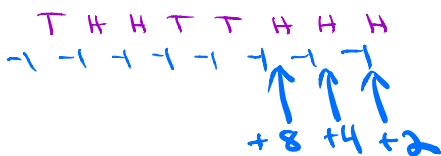
$$X_T = -T + 8 + 2$$

$$0 = E(X_T) = -E(T) + 10$$

$$E(T) = 10$$



$\sigma = HTH$



$$E(T) = 8 + 4 + 2 = 14$$

$$X_T = -T + 8 + 4 + 2$$

$$\underline{E(T) = 14}$$

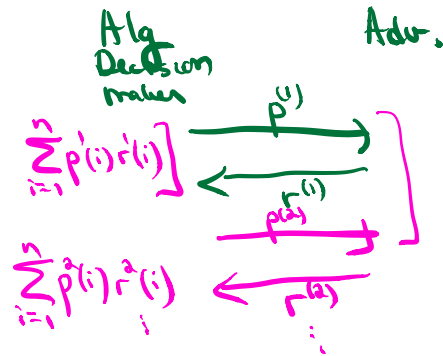
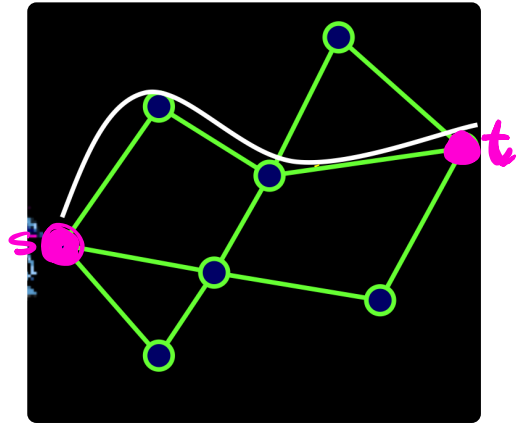
Repeated Online Decision Making and the Multiplicative Weights Algorithm

A set of possible actions $|A|=n$
 T a time horizon $A=\{1,2,\dots,n\}$

Setup:

- At each time step $t=1..T$
 - a decision maker picks an action $a_t \in A$ where $p^t(i) = \Pr(a_t = i)$
 - $\vec{p}^t = (p^t(1), p^t(2), \dots, p^t(n))$
 - an adversary picks reward vector $\vec{r}^t = (r^t(1), r^t(2), \dots, r^t(n))$ where $r^t(i) =$ reward to alg if picked action i
 - decision maker learns r^t

Goal of alg
 maximize $\sum_{t=1}^T \sum_{i=1}^n p^t(i) r^t(i)$
 $\underbrace{\hspace{10em}}_{p^{(t)}, r^{(t)}}$



Examples:

① Choosing a route

② Choosing stocks to buy



Best possible world $\sum_{t=1}^T \max_i r^+(i)$ (*) $|r^+(i)| \leq 1$
 $\forall i$

This benchmark way too strong.

Ex: $A = \{1, 2\}$
 $p^+(1) + p^+(2) = 1$

$p^+(1) \geq \frac{1}{2} \Rightarrow$
 $r^+(1) = -1$
 $r^+(2) = 1$

o.w. \Rightarrow
 $r^+(1) = 1$
 $r^+(2) = -1$

$E(\text{reward of alg}) \leq 0$

(*) = T

$$\text{Regret}(\vec{p}^1, \dots, \vec{p}^T) = \frac{1}{T} \left[\underbrace{\max_{a \in A} \sum_{t=1}^T r^+(a)}_{\substack{\text{best reward} \\ \text{possible if} \\ \text{use same} \\ \text{action every} \\ \text{day}}} - \underbrace{\sum_{t=1}^T p^t \cdot r^t}_{\substack{\text{alg total} \\ \text{exp reward.}}} \right]$$

Goal: get $\text{Reg} \rightarrow 0$ as T gets large.

Obvious ^{det} alg: "Follow the leader"

upto $t-1$, let a be action with max

$$\sum_{\sigma=1}^{t-1} r^+(\sigma)$$

use that on step t

all actions unif.