

"Puzzle"

$$A \subseteq \{0, 1, 2, \dots, p-1\} \quad |A| \geq \frac{p}{2}$$

Show $\exists$ int n s.t. all gaps
of $\{nA \bmod p\} \leq 5\sqrt{p}$

First Moment Method

r.v. X takes some value $\geq E(X)$
 $\leq E(X)$

$X \geq 0$ int valued $\Rightarrow \Pr(X \leq 1) \leq E(X)$
 $\Rightarrow X=0$ w.h.p. if $E(X) \ll 1$

Proving $P(X \geq 1)$ large?

$E(X)$ large insufficient - could be HUGE w/ negligible prob.

2nd Moment Method

pretty much any use of 2nd moment $E(X^2)$

Chebyshev's Inequality

$\forall \lambda > 0$

$$\Pr(|X - \mu| \geq \lambda \sigma) \leq \frac{1}{\lambda^2}$$

X nonnegative integer-valued r.v.

Another version

$$\Pr(X=0) \leq \frac{\text{Var}(X)}{E(X^2)}$$

Pf :

$$\Pr(X=0) \leq \Pr(|X-\mu| \geq \mu) \leq \frac{\sigma^2}{\mu^2} = \frac{\text{Var}(X)}{(E(X))^2}$$

Corollary

If $\text{Var}(X) = o(E(X)^2)$ then $\Pr(X > 0) = 1 - o(1)$

Typically $X = X_1 + X_2 + \dots + X_n$

Claim:
$$\text{Var}(X) = \sum_{i=1}^n \text{Var}(X_i) + \sum_{i \neq j} \text{Cov}(X_i, X_j)$$

$$\text{Var}(X+Y) = E[(X+Y - E(X+Y))^2]$$

$$= E[(X - E(X))^2 + (Y - E(Y))^2 + 2(X - E(X))(Y - E(Y))]$$

$$= \text{Var}(X) + \text{Var}(Y) + 2\text{Cov}(X, Y)$$

$$\text{Cov}(X, Y) = E(XY) - E(X)E(Y)$$

Observe $\text{Cov}(X, Y) = 0$
if X & Y indep

Often X_i 's are indicator r.v.'s

$$\text{Var}(X_i) = p_i(1-p_i) \leq p_i = E(X_i)$$

$$\Rightarrow \text{Var}(X) \leq E(X) + \sum_{i \neq j} \text{Cov}(X_i, X_j)$$

$$X_i = \begin{cases} 1 & p_i \\ 0 & 1-p_i \end{cases}$$

Example 1: Consider $G_{n,p}$ graph on n vertices where each edge present w/ prob p (independently)

For what p does $G_{n,p}$ have a clique with 4 or more vertices?

Claim:

① almost surely doesn't for $p = o(n^{-2/3})$

② almost surely does for $p = \omega(n^{-2/3})$

Proof of ①

$\binom{n}{4}$ possible 4-cliques

$C_1, \dots, C_{\binom{n}{4}}$ enumeration

$$X_i = \begin{cases} 1 & C_i \text{ clique} \\ 0 & \text{o.w.} \end{cases}$$

$$X = \sum_{i=1}^{\binom{n}{4}} X_i \quad \# \text{ of 4-cliques}$$

$$E(X) = \binom{n}{4} p^6 \approx \frac{n^4}{4!} p^6$$

$$\text{for } p = o(n^{-2/3}) \quad E(X) = o(1)$$

$$X \geq 0, \text{ integer} \Rightarrow \Pr(X \geq 1) \leq E(X) = o(1)$$

first moment method

Proof of ②

For $p = w(n^{-2/3})$

$E(X) \xrightarrow{n \rightarrow \infty} \infty$

Insufficient by itself

So we use 2nd moment method, we'll show $\text{Var}(X) = o(E(X)^2)$

need to compute $\text{Var}(X)$

$\text{Cov}(X_i, X_j)$

type 1

$|C_i \cap C_j| = \emptyset$

$\text{Cov}(X_i, X_j) = 0$



type 2

$|C_i \cap C_j| = 1$

$\text{Cov}(X_i, X_j) = 0$



type 3

$|C_i \cap C_j| = 2$



$\text{Cov}(X_i, X_j) \leq E(X_i X_j) = p^2$

type 4

$|C_i \cap C_j| = 3$



$\text{Cov}(X_i, X_j) \leq E(X_i X_j) = p^3$

$$\text{Var}(X) \leq E(X) + \sum_{\substack{1 \leq i, j \leq m \\ i \neq j}} c_{ij}(X_i, X_j)$$

$$\leq \underbrace{\binom{n}{4} p^6}_{\text{type 3}} + \underbrace{\binom{n}{6} \binom{6}{2,2,2} p^{11}}_{\text{type 4}} + \binom{n}{5} \binom{5}{3,1,1} p^9$$

$$\leq c(n^4 p^6 + n^6 p^{11} + n^5 p^9) = o([E(X)]^2)$$

since

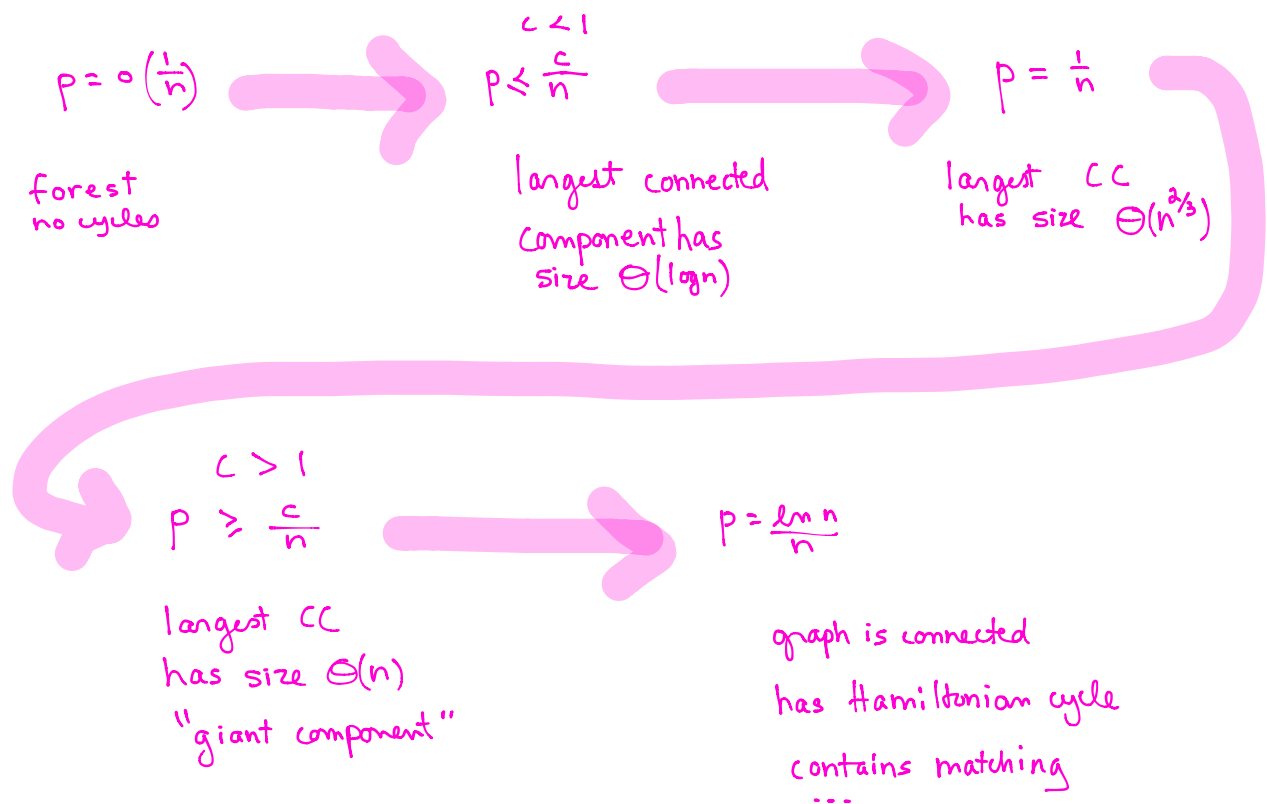
$$[E(X)]^2 = \left(\binom{n}{4} p^6\right)^2 = \Theta(n^8 p^{12})$$

$$p = \omega(n^{-2/3})$$

$$\Rightarrow \Pr(X=0) \leq \frac{\text{Var}(X)}{[E(X)]^2} = o(1)$$

Evolution of Random graphs

The following statements hold almost surely $\equiv$ with prob $\rightarrow 1$ as $n \rightarrow \infty$



i

Graph properties expressible in "First-Order theory" of graphs

language: vars $(x, y, z, \dots)$ represent vertices

equality $x = y$; adjacency $x \sim y$

$\wedge, \vee, \neg$ $\forall, \exists$

Examples:

G contains triangle: $\exists x \exists y \exists z \ x \sim y \wedge y \sim z \wedge z \sim x$

no isolated pt $\forall x \exists y (x \sim y)$

diameter 2 $\forall x \forall y \exists z ((x=y) \vee (x \sim y) \vee ((x \sim z) \wedge (z \sim y)))$

(G Hamiltonian, G connected) not FO properties

For A a FO property, let

$$\Pr(G(n, p) \models A) \equiv \Pr(\text{random graph from } G(n, p) \text{ has property } A)$$

Thm [Fagin] [Glebskii, Kogan, Liogonkii, Talanov]

$\forall$ fixed p , $0 < p < 1$, and any FO graph property

$$\lim_{n \rightarrow \infty} \Pr(G(n, p) \models A) = 0 \text{ or } 1$$

Thm [Shelah, Spencer]

$\forall$ irrational α , $0 < \alpha < 1$, setting $p = p(n) = n^{-\alpha}$ and for any FO property A

$$\lim_{n \rightarrow \infty} \Pr(G(n, p) \models A) = 0 \text{ or } 1$$

this is about threshold for

Ex $p = n^{-2/3}$ threshold for K_4

$p \ll n^{-2/3}$
no K_4

$$p = n^{-2/3} \quad \Pr(K_4) = 1 - e^{-1/4}$$

$p \gg n^{-2/3}$
 K_4

What happens at $n^{-2/3}$?
Nothing!!

Another useful 2nd moment inequality

$$\Pr(X > 0) \geq \frac{(E(X))^2}{E(X^2)}$$

Follows from Cauchy-Schwartz

$$[E(XY)]^2 \leq E(X^2) E(Y^2)$$

Set $Y = 1_{X>0}$

$$[E(X)]^2 \leq E(X^2) \underbrace{E([1_{X>0}]^2)}_{\Pr(X>0)}$$

Proof:

wlog $E(X^2) > 0$
 $E(Y^2) > 0$

$$\text{Let } U = \frac{X}{\sqrt{E(X^2)}} \quad V = \frac{Y}{\sqrt{E(Y^2)}}$$

$$2|UV| \leq U^2 + V^2$$

$$\Rightarrow 2|E(UV)| \leq 2E(|UV|)$$

$$\leq E(U^2) + E(V^2) = 2$$

$$\Rightarrow [E(UV)]^2 \leq 1$$

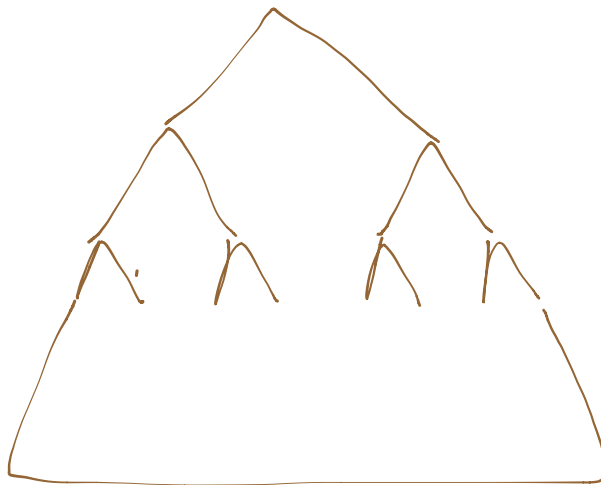
$$= [E(XY)]^2 \leq E(X^2) E(Y^2)$$

Percolation on Tree

Complete binary tree of depth n

each edge present w/ prob p

Is there a surviving
path from root to a leaf?



X # leaves reachable from root

$$E(X) = ?$$

$$X = X_1 + X_2 + \dots + X_{2^n}$$

X_i indicates if i^{th}
leaf is reachable

$$E(X) = 2^n \cdot p^n = \begin{cases} \rightarrow \infty & p > \frac{1}{2} \\ = 1 & p = \frac{1}{2} \\ \rightarrow 0 & p < \frac{1}{2} \end{cases}$$

What can we say about $\Pr(X > 0)$?

$$p > \frac{1}{2}$$

$$\Pr(X > 0) \geq \frac{[E(X)]^2}{E(X^2)}$$

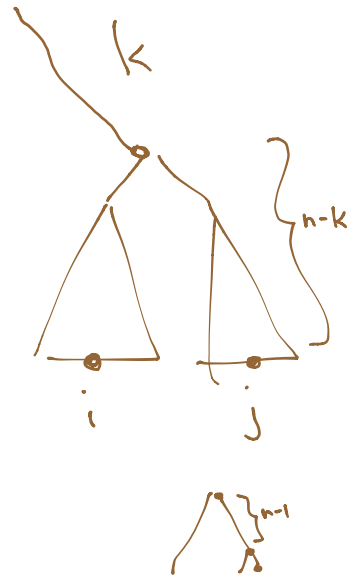
$$X^2 = (\sum X_i)^2 = \sum_i X_i^2 + \sum_i \sum_{j \neq i} X_i X_j$$

$$E(X_i X_j) = p^{2n-k}$$

$$\sum_{i \neq j} E(X_i X_j) = \sum_{k=0}^{n-1} 2^k 2^{n-k} 2^{n-k-1} p^{2n-k}$$

$$= \frac{1}{2} \sum_{k=0}^n 2^{2n} \frac{1}{2^k} p^{2n} \frac{1}{p^k} = \frac{1}{2} (2p)^{2n} \underbrace{\sum_{k=0}^{n-1} \frac{1}{(2p)^k}}$$

$$\leq \frac{1}{2} (2p)^{2n} \frac{2p}{2p-1} = \frac{p}{2p-1} (2p)^{2n}$$



$$\frac{p > \frac{1}{2}}{2p > 1}$$

$$E(X^2) \leq \left(\frac{p}{2p-1}\right) (2p)^n [1+o(1)]$$

$$\Pr(X > 0) \geq \frac{[E(X)]^2}{E(X^2)} \geq \frac{2^{p-1} (1+o(1))}{p} > c \quad \text{for } p > \frac{1}{2}$$

$$p = \frac{1}{2} \quad \sum_{i,j \neq i} E(X_i X_j) = \sum_{k=0}^{n-1} 2^k 2^{n-k} 2^{n-k-1} p^{2n-k} \\ = \frac{n}{2}$$

$$E(X^2) = 1 + \frac{n}{2} \quad [E(X)]^2 = 1$$

$$\Pr(X > 0) \geq \frac{[E(X)]^2}{E(X^2)} \geq \frac{2}{n+2}$$

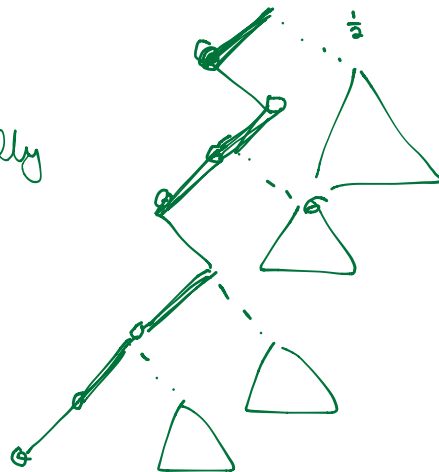
Other direction: $\perp$ event that a vertex with at least $\frac{n}{2}$ left turns survives

$$E(X|L) \geq \frac{n}{4}$$

toss coins to produce lexicographically

leftmost path

$\frac{n}{2}$ left turns for each right child exp



$$\Pr(L) \leq \frac{E(X)}{E(X|L)} \leq \frac{4}{n}$$

$$\Pr(X > 0) = \Pr(L) + \Pr(R) \leq \frac{8}{n}$$

