

general random walk on graph

(reversible)

Associated with each edge $e=(x,y)$ is

$$c_{xy} = c_e \quad \text{"conductance of edge"}$$

$$P_{xy} = \frac{c_{xy}}{c_x}$$

$$c_x = \sum_{y | (x,y) \in E} c_{xy}$$

simple random walk $c_{xy}=1 \quad \forall (x,y) \in E$

$$\pi_x = \frac{c_x}{c_G}$$

$$c_G = \sum_{x \in V} c_x$$

Connections with Electrical Networks

a fn $h: V \rightarrow \mathbb{R}$ is harmonic for P at v if

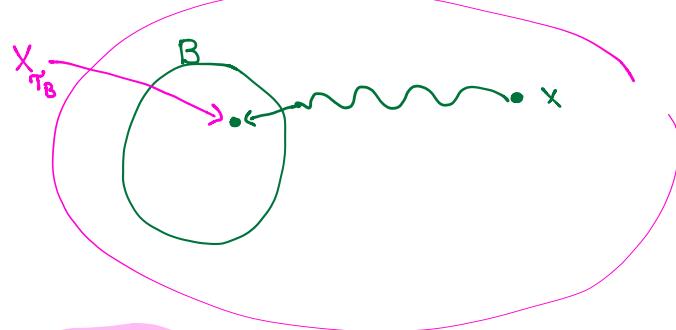
$$h(i) = \sum_j P_{ij} h(j) \quad h \text{ is avg of neighbors}$$

Proposition: (X_+) irreducible MC

$B \subset S$ $h_B: B \rightarrow \mathbb{R}$ fn defined on B

$$h(x) = E_x h_B(X_{T_B})$$

T_B : time to reach B



is unique harmonic extension of h_B

$$\textcircled{1} \quad h(x) = h_B(x) \quad \forall x \in B$$

\textcircled{2} h harmonic on $x \notin B$

Proof

\textcircled{1} ✓

$$\begin{aligned}
 ② h(x) &= E_x h_B(X_{T_B}) = \sum_y P_{xy} E_x \underbrace{[h_B(X_{T_B}) | X_i=y]}_{E_y [h_B(X_{T_B})]} \\
 &= \sum_y P_{xy} h(y) \quad \checkmark
 \end{aligned}$$

Uniqueness:

Suppose not. h & g both harmonic extensions

$$\text{then } h(x) - g(x) = 0 \text{ on } x \in B$$

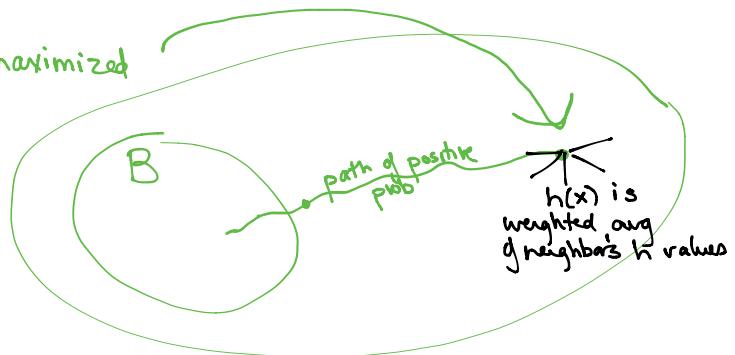
is harmonic on $x \notin B$

Consider x where $h(x) - g(x)$ maximized

$$\Rightarrow \max \leq 0$$

$$\text{similarly } \min \geq 0$$

$$\Rightarrow h(x) = g(x) \quad \forall x$$



Electrical Networks, Voltages & Current Flows

$$G = (V, E)$$

Fix 2 nodes s & t

\uparrow
source \uparrow
 sink

A fn ϕ on V which is harmonic on $V \setminus \{s, t\}$

called a **voltage**

(completely determined
by $\phi(s)$ & $\phi(t)$)

$$e = (x, y)$$

An (electrical) flow is a fn on oriented edges $\vec{e} = x \rightarrow y$
 $\vec{e} = x \leftarrow y$

is

- anti-symmetric $f(\vec{e}) = -f(\vec{e}) \quad \forall e$

- $\sum_{y | (x, y) \in E} f(x \rightarrow y) = 0 \quad \forall x \notin \{s, t\}$

- $\sum_{y | (s, y) \in E} f(s \rightarrow y) \geq 0$
flow out of $s \geq 0$

strength of flow
or value

Kirchhoff's Law =
Flow in = Flow out

$$\sum_x \sum_{y | (x, y) \in E} f(x \rightarrow y) = 0$$

To construct electrical network. If edge e , put resistor

$$\text{where } r_e = \frac{1}{c_e}.$$

resistance r_e conductance c_e

Given voltage ϕ , the current flow I on each edge $\vec{e} = x \rightarrow y$ satisfies

Ohm's Law

$$I(x \rightarrow y) = \frac{\phi(x) - \phi(y)}{r_{xy}} = c_{xy} [\phi(x) - \phi(y)]$$

resistance on edge

Observations:

$$\textcircled{1} \quad \sum_{y \mid (x,y) \in E} I(x \rightarrow y) = \sum_{y \mid (x,y) \in E} \phi(x) - \phi(y) = d_x \phi(x) - \sum_{y \mid (x,y) \in E} \phi(y)$$

conservation of current

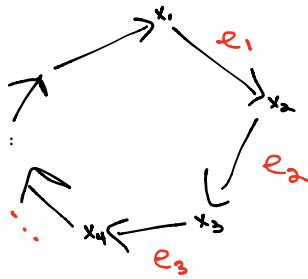
$$= d_x \left[\phi(x) - \sum_{y \mid (x,y) \in E} \frac{1}{d_x} \phi(y) \right]$$

= 0 since ϕ harmonic

$$\sum_{y \mid (x,y) \in E} I(x \rightarrow y) = \sum_{y \mid (x,y) \in E} c_{xy} (\phi(x) - \phi(y)) = \dots = 0$$

② Cycle law

$$\sum_i r_{e_i} I(x_i \rightarrow x_{i+1}) = \sum_i \phi(x_i) - \phi(x_{i+1}) \\ = 0$$



③ $\phi'(x) = \phi(x) + c$ $\forall x$ has no effect on current or harmonicity

\Rightarrow wlog $\phi(t) = 0$ & ϕ determined by $\phi(s)$

If f is a flow satisfying the cycle law for every cycle, and

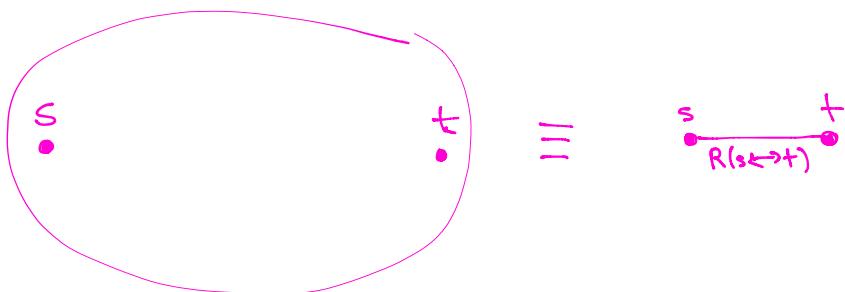
value of flow = value of current flow, then $f = I$

given resistances on edges

ϕ voltage \iff I current flow

Definition The effective resistance between
 $s \& t$

$$R(s \leftrightarrow t) = \frac{\phi(s) - \phi(t)}{II} \quad \leftarrow \text{indep of } \phi$$



Lemma 1

$$\Pr(T_{st} < T_{ss}^+) = \frac{1}{d_s R(s \leftrightarrow t)}$$

Proof

$$g: x \mapsto \Pr_x(T_{xt} < T_{xs})$$

is a harmonic fn on $V \setminus \{s, t\}$

$$\text{s.t. } g(s) = 0 \quad g(t) = 1$$

$$h(x) = \frac{\phi(s) - \phi(x)}{\phi(s) - \phi(t)}$$

also harmonic with same
boundary values

$$\Rightarrow \Pr_x(T_{xt} < T_{xs}) = \frac{\phi(s) - \phi(x)}{\phi(s) - \phi(t)}$$

$$\begin{aligned} \Pr_s(T_{st} < T_{ss}^+) &= \sum_x P_{sx} \Pr_x(T_{xt} < T_{xs}) \\ &= \sum_x \frac{c_{sx}}{c_s} \frac{\phi(s) - \phi(x)}{\phi(s) - \phi(t)} \\ &= \frac{\sum_x I(s \rightarrow x)}{c_s (\phi(s) - \phi(t))} = \frac{\|I\|}{c_s (\phi(s) - \phi(t))} \\ &= \frac{1}{c_s R(s \leftrightarrow t)} \end{aligned}$$

Lemma 2

$$E_s(\# \text{q visits to } s \text{ before hit } t) = c_s R(s \leftrightarrow t)$$

Proof: Lemma 1 \Rightarrow geometric distn

Lemma 3

Let T be a

stopping time

in finite irreducible MC st. $\Pr_s(X_T = s) = 1$

Then $\frac{E_s(\# \text{q visits to } x \text{ before } T)}{E_s(T)} = \pi_x \quad \forall x$

Commute Time Identity

Consider r.w. on G (c_e conductance on edge e)

$\forall s, t$

$$P_{xy} = \frac{c_{xy}}{\sum_{y \in \{(x,y) \in E\}} c_{xy}} = \frac{c_{xy}}{c_x}$$

$$E(T_{st}) + E(T_{ts}) = c_G R(s \leftrightarrow t)$$

$$c_G = \sum_{x \in V} c_x$$

Pf

Use Lemma 3

$$\text{Let } T = T_{st} + T_{ts}$$

$$\frac{E_s(\# q \text{ visits to } s \text{ before } T)}{E_s(T)} = \pi_s = \frac{c_s}{c_G}$$

$$\Rightarrow E_s(\# q \text{ visits to } s \text{ before reaching } t)$$

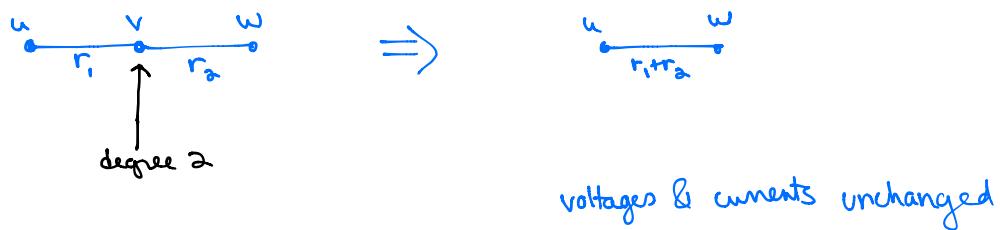
$$\text{Lemma 2} = c_s R(s \leftrightarrow t)$$

Commute effective resistance

→ Simplifying networks

Sometimes can simplify network w/o changing quantities such as effective resistance

① Series Law: resistance in series add



② Parallel Law: conductances in parallel add



Some applications of commute time identity

① Random walk on line

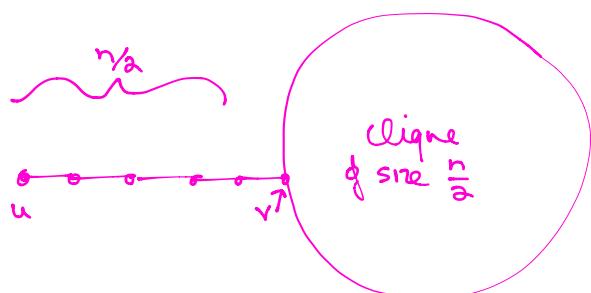


$$R_{0 \leftrightarrow n} = n \quad \text{by series law}$$

$$c_G = n$$

$$\Rightarrow h_{0n} + h_{n0} = 2n^2$$

② lollipop



$$c_G = m = n^2$$

$$R_{u \leftrightarrow v} = \frac{n}{2}$$

$$\Rightarrow h_{uv} + h_{vu} = \Theta(n^3)$$

Cover Time Bounds

Thm:

$$C(G) = \mathcal{O}(mR)$$

$$C(G) = O(mR \log n)$$

where $R = \max_{u,v} R_{u \leftrightarrow v}$

Lower bound immediate.

for upper bound, consider $\ln n$ epochs of length $\frac{2mR}{\alpha}$
 ↑
 const to be determined

$$\forall v \quad \forall u \quad E_u(T_v) \leq 2mR$$

$$\Rightarrow \Pr(v \text{ not visited in particular epoch}) \leq \frac{1}{\alpha} \quad \text{by Markov Inequality}$$

$$\Rightarrow \Pr(v \text{ not visited in any of } \ln n \text{ epochs}) \leq \left(\frac{1}{\alpha}\right)^{\ln n} = n^{-\ln(\alpha)}$$

$$\Rightarrow \Pr(\exists \text{ vertex not visited}) \leq n^{1-\ln(\alpha)}$$

↑
union bound

$$C(G) \leq 2amRlmn + n^{\lceil \log_2 3 \rceil}$$

worst case bound

with $a = e^4$

$$C(G) = O(mRlmn)$$

Example: complete graph

R?

$$H_{uv} = n-1 \quad (\Pr(\text{hit } v \text{ each step}) = \frac{1}{n-1})$$

$$H_{uv} + H_{vu} = 2m R_{uv \leftrightarrow v}$$

$$\frac{2}{n-1} = 2m R_{uv \leftrightarrow v} \quad \Rightarrow \quad R_{uv \leftrightarrow v} = \frac{2}{n}$$

$$\begin{aligned} C(G) &= O(mRlmn) = O\left(\binom{n}{2} \frac{2}{n} lmn\right) \\ &= O(nlmn) \end{aligned}$$

$$\mathcal{E}(f) = \text{energy of flow } f = \sum_e f(e)^2 r_e$$

Thomson's Principle

For any finite connected graph

$$R_{s \leftrightarrow t} = \inf \{ \mathcal{E}(f) \mid f \text{ is unit flow from } s \rightarrow t \}$$



Unique minimizer is unit current flow

Corollary $\{r_e\}_{e \in E}$ & $\{r'_e\}_{e \in E}$ sets of resistances on same G

with $r_e \leq r'_e \forall e$

$$\text{Then } R_{s \leftrightarrow t} \leq R'_{s \leftrightarrow t}$$

Proof $\inf_f \sum_e r_e f(e)^2 \leq \inf_f \sum_e r'_e f(e)^2$

+ Thomson's principle

Cordlang: Adding an edge doesn't \uparrow eff resistance

