CSE 521 Design and Analysis of Algorithms I

Overview

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University of Washington

Computer Science & Engineering

CSE 521, Wi '13: Design and Analysis of Algorithms I

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Administrative

Course Email/BBoard

Subscription Options Class List Archive E-mail Course Staff GoPost BBoard

Lecture Notes

1: Overview & Example

2: Analysis

3: Graphs, B/DFS

Lecture: MGH 271 (schematic)

Office Hours Location Phone

Instructor: Larry Ruzzo, ruzzo@cs CSE 554 206-543-6298 tha

http://www.cs.washington.edu/521 TA: Cyrus Rashtchian, cyrash@cs tba CSE xxx

staff Q&A about

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agorithms: recursion, divide and conquer, balancing

Julework, etc.

analysis of algorithms. NP-completeness.

ading: Homework, Final. Overall weights 60%, 40%, roughly.

Extra Credit: Assignments may include "extra credit" sections. These will enrich your understanding of the them for the glory, not the points, and don't start extra credit until the basics are complete.

Collaboration: Homeworks are all individual, not group, exercises. Discussing them with others is fine, ever homework solutions. Follow the "Gilligan's Island Rule": if you discuss the assignment with someone else, discussion, then go watch 30+ minutes of TV (Gilligan's Island reruns especially recommended) before you You may not look at other people's written solutions to these problems, not in your friends' notes, not in the doubt about whether your activities cross allowable boundaries, tell us before, not after, you turn in your ass Misconduct Policy, and the links there.

Textbook: Algorithm Design by Jon Kleinberg and Eva Tardos, Addison Wesley, 2006. (Available from U.)

What you have to do

Weekly Homework (~60% of grade)

Programming?

perhaps some small projects

Written homework assignments

English exposition and pseudo-code

Analysis and argument as well as design

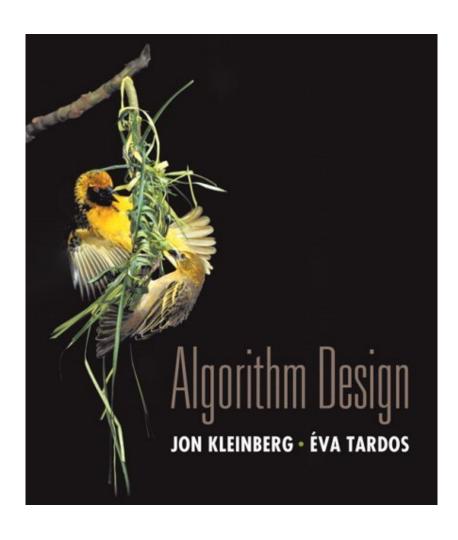
Final Exam

(~40%)

Late Policy:

Papers and/or electronic turnins generally due in class; minus 20% per day thereafter

Textbook



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

What the course is about

Design of Algorithms

design methods
common or important types of problems
analysis of algorithms - efficiency
correctness proofs

What the course is about

Complexity, NP-completeness and intractability solving problems in principle is not enough algorithms must be efficient

some problems have no efficient solution

NP-complete problems

important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

Very Rough Division of Time

Algorithms (6-7 weeks)

Analysis

Techniques: greedy, divide&conquer, dynamic programming,

Toolkit: flows & matchings, linear programming

Applications

Complexity & NP-completeness (2-3 weeks)

Check online schedule page for (evolving) details



University of Washington Computer Science & Engineering

CSE 417, Wi '06: Approximate Schedule

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		Due	Lecture Topic	Reading
Week 1 1/2-1/6	М		Holiday	
	w		Intro, Examples & Complexity	Ch. 1; Ch. 2
	F		Intro, Examples & Complexity	
Week 2	М		Intro, Examples & Complexity	
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Complexity Example

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Cryptography (e.g., RSA, SSL in browsers)
    Secret: p,q prime, say 512 bits each
    Public: n which equals p \times q, 1024 bits
In principle
    there is an algorithm that given n will find p and q:
   try all 2^{5/2} > 1.3 \times 10^{154} possible p's: kinda slow...
In practice
    no fast algorithm known for this problem (on non-quantum computers)
    security of RSA depends on this fact
    ("quantum computing": strongly driven by possibility of changing this)
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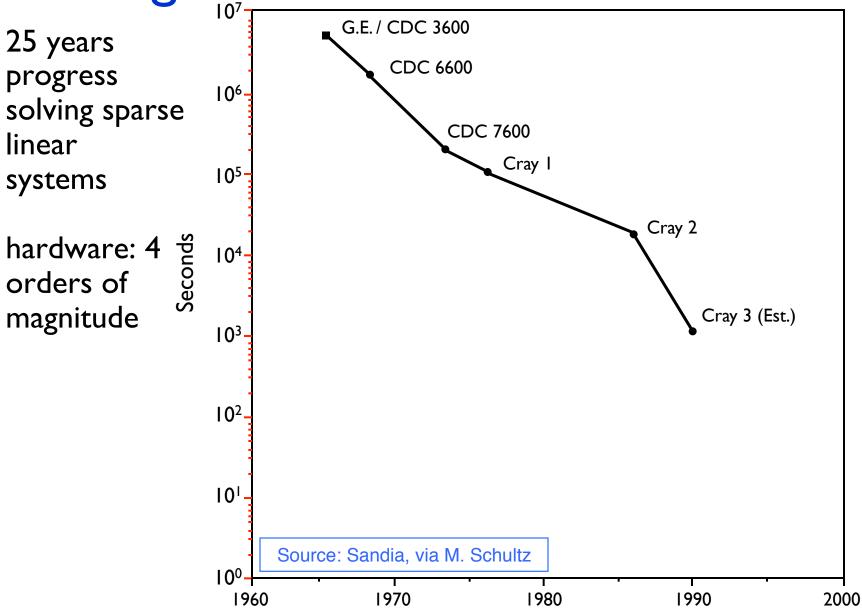
Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

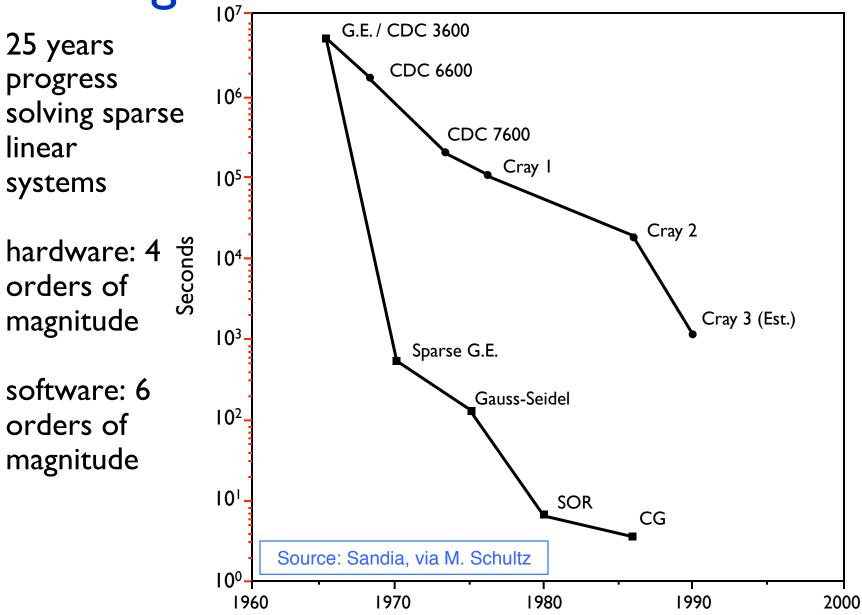
Ex: sparse linear equations over 25 years

10 orders of magnitude improvement!

Algorithms or Hardware?



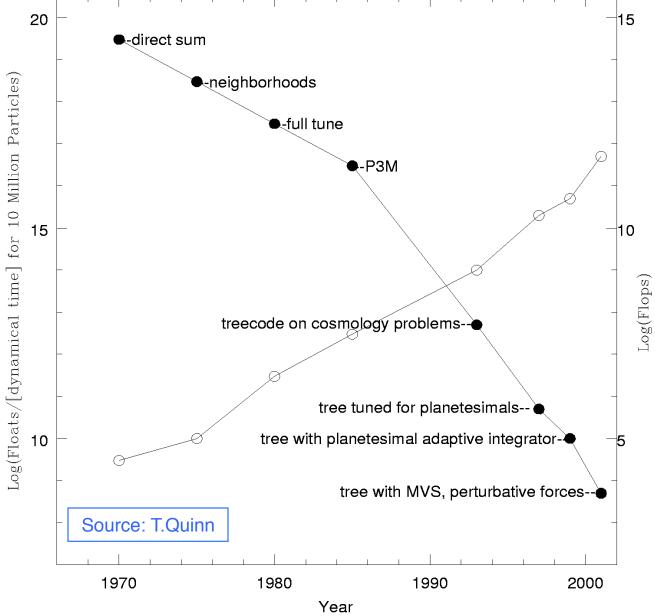
Algorithms or Hardware?



Algorithms or Hardware?

The N-Body Problem:

in 30 years 10⁷ hardware 10¹⁰ software



Goals

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Correctness
often subtle
Analysis
often subtle
Generality, Simplicity, 'Elegance'
Efficiency
time, memory, network bandwidth, ...
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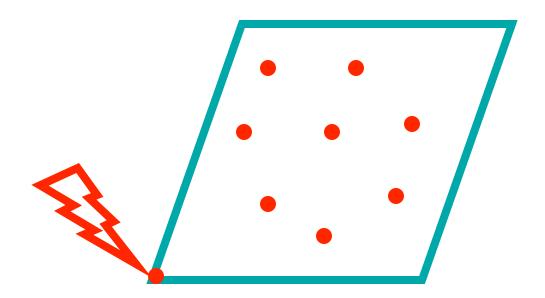
Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

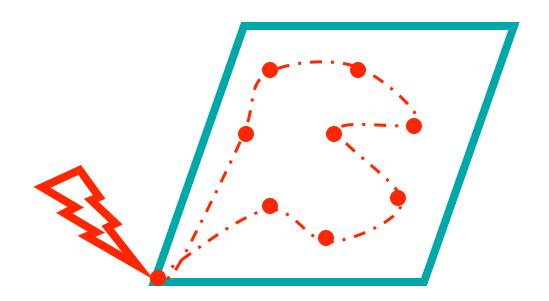
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

Printed Circuit Board



Printed Circuit Board



A Well-defined Problem

Input: Given a set S of n points in the plane Output: The shortest cycle tour that visits each point in the set S.

Better known as "TSP"

How might you solve it?

Nearest Neighbor Heuristic

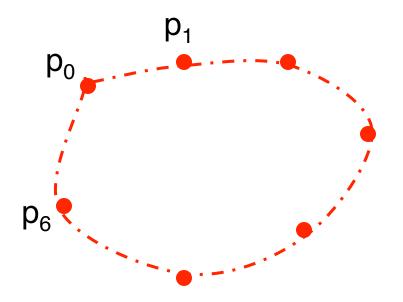
Start at some point p_0 Walk first to its nearest neighbor p_1

heuristic:

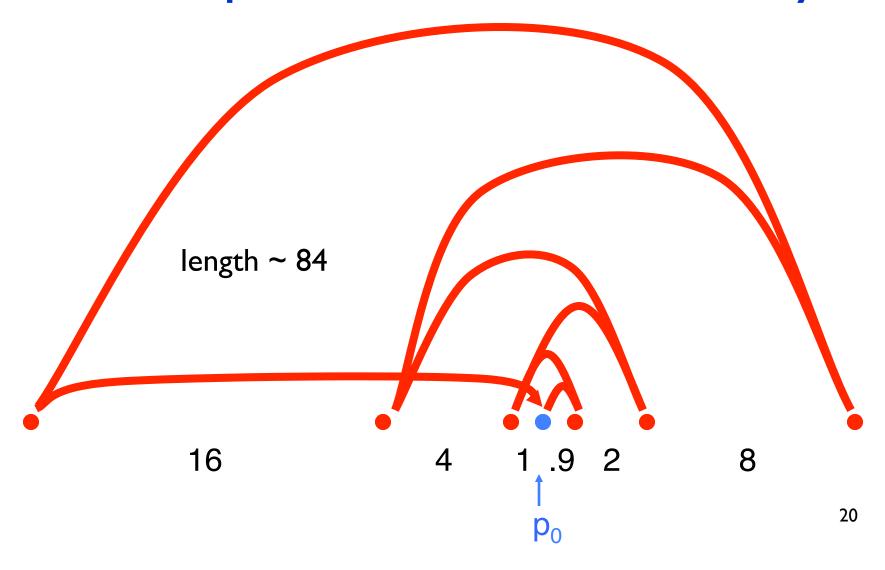
A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually *not* guaranteed to give the best or fastest solution.

Repeatedly walk to the nearest unvisited neighbor p_2 , then p_3 ,... until all points have been visited. Then walk back to p_0

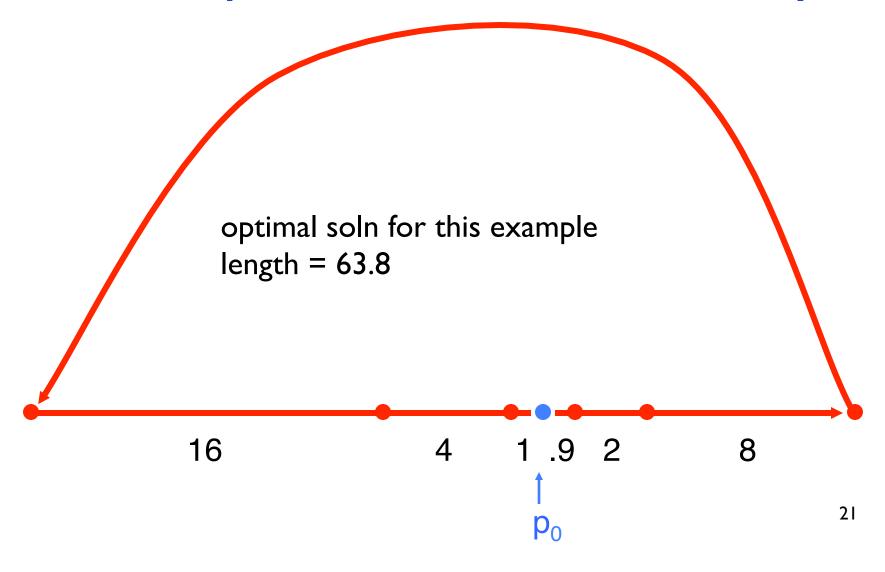
Nearest Neighbor Heuristic



An input where it works badly



An input where it works badly

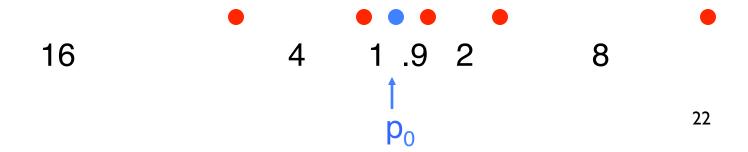


Revised idea - Closest pairs first

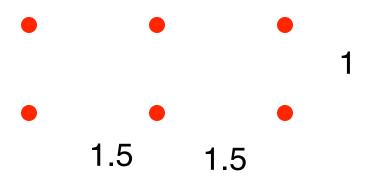
Repeatedly join the closest pair of points

(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)

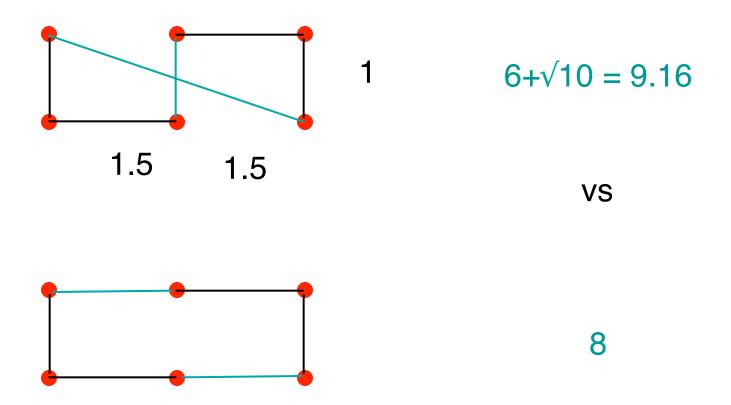




Another bad example



Another bad example



Something that works

"Brute Force Search":

For each of the n! = n(n-1)(n-2)...I orderings of the points, check the length of the cycle you get Keep the best one

Two Notes

The two incorrect algorithms were greedy

Often very natural & tempting ideas

They make choices that look great "locally" (and never reconsider them)

When greed works, the algorithms are typically efficient

BUT: often does not work - you get boxed in

Our correct alg avoids this, but is incredibly slow

20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)

And growing: n! ~ $\sqrt{2 \pi n} \cdot (n/e)^n \sim 2^{O(n \log n)}$

The Morals of the Story

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Algorithms are important
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Many performance gains outstrip Moore's law

Simple problems can be hard

Factoring, TSP

Simple ideas don't always work

Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

Brute-force factoring, TSP

And: for some problems, even the best algorithms are slow