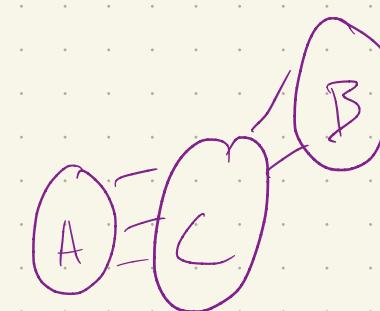


Undirected GMs aka MRFs

Graph factorizes wrt G as $P(X) = \frac{1}{Z} \prod f_c(x_c)$ (GE)

$$Z \triangleq \sum_{X \in X^n} \left(\prod f_c(x_c) \right)$$

Graph separation (GS) A & C are separated by B
"A - B - C" 

Global Markov (GM) wrt G

$P(X)$ satisfies GM wrt G iff $\nexists A - B - C$

$$X_A \perp\!\!\!\perp X_C \mid X_B$$

[Factorization according to G \Rightarrow GM wrt G]

Local Markov Property (L)

$P(X)$ satisfies L wrt G if

$$X_i \perp\!\!\!\perp X_{V \setminus (i \cup S(i))} \mid S(i)$$

Neighbors
of i



Pair Wise Property (P)

$P(x)$ satisfies (P) wrt G if

$$X_i \perp\!\!\! \perp X_j \mid X \cup \{S_{i,j}\} \quad \forall (i,j) \in E$$

$$\begin{array}{c} G \Rightarrow L \Rightarrow P \\ \textcircled{1} \qquad \textcircled{2} \end{array}$$

Lemme [Data Processing \neq]

$X - z - Y - h(Y)$

(1) $X \perp\!\!\! \perp Y \mid z \Rightarrow X \perp\!\!\! \perp h(Y) \mid z$

(2) $X \perp\!\!\! \perp Y \mid z, h(Y)$

Proof that $G \Rightarrow L$

We know $\forall A, B, C \ s.t. A - B - C \text{ in } G$,

$$X_A \perp\!\!\! \perp X_C \mid X_B$$

In particular $A = \{i\}$ $B = \{S(i)\}$ $C = V \setminus (A \cup B)$

then $A - B - C$ since $(i, j) \in E \Rightarrow j \in B$

$$\Rightarrow a - b_1 - \dots - c \Rightarrow X_i \perp\!\!\! \perp X_{V_i} \mid X_B$$

Proof that $L \Rightarrow P$

We know $x_i \perp\!\!\!\perp X_{V \setminus i \cup \{j\}} | X_{S_i} \forall i$.

$\nexists \forall j: (i, j) \notin E \Rightarrow j \notin S_i \cup \{i\}$

(DP2)(a) $\Rightarrow x_i \perp\!\!\!\perp X_{V \setminus (i \cup \{j\})} | X_{V \setminus \{i, j\}}$

(DP1)(b) $\Rightarrow x_i \perp\!\!\!\perp \underbrace{x_j}_{\text{---}} | X_{V \setminus \{i, j\}}$

which is (P)

□

Assume our PMF P has $P(X) = P(x_1, \dots, x_n) > 0$

$\forall x = [x_1, \dots, x_n] \in \mathcal{X}^n$

Claim: If $P(x) > 0 \ \forall x \in \mathbb{X}^n$ then $(P) \Rightarrow (G)$
pairwise global

$(G \Rightarrow L) \Rightarrow (P)$

[if $P(x) > 0$]

Lemma: Intersection Lemma

Fix $P(x) > 0$.

suppose $\cdot X_A \perp\!\!\!\perp X_B \mid (x_C, x_D)$

and

$\cdot X_A \perp\!\!\!\perp X_C \mid (x_B, x_D)$

then

$X_A \perp\!\!\!\perp X_B, X_C \mid X_D$

Proof of Claim $[P(x) > 0 \ \forall x \text{ implies } (P) \Rightarrow (G)]$

$(P) \star \forall (i, j) \notin E \rightarrow X_i \perp\!\!\!\perp X_j \mid X_{V \setminus \{i, j\}}$

Want to show $\star \Rightarrow X_A \perp\!\!\!\perp X_C \mid X_B \ \forall A - B - C$

Will show by induction on $|B| = n-2, n-3, \dots, 1$

(1) $|B| = n-2$, $A = \{i, j\}$, $C = \{j\}$. Since $(i, j) \notin E$,

$$\{i, j\} - \{(i, j)\} = \{j\}$$

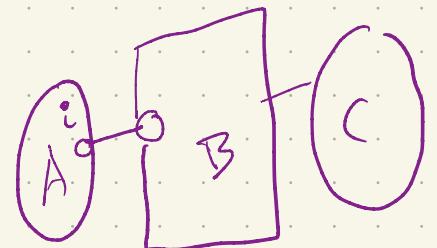
and so $x_A \perp\!\!\! \perp x_C | x_B$

(2) Assume (6) holds $\wedge |B| \geq s$.

$$|B| = s-1 < n-2$$

Assume wlog $|A| \geq 2$ & $A - B - C$ are separated in G

Take $i \in A$ adjacent to B . Then,



$$A \setminus \{i\} - B \cup \{i\} - C \quad \text{also separated}$$

$$\tilde{B}, |B| = s \Rightarrow^{IH} (1) x_C \perp\!\!\! \perp x_{A \setminus \{i\}} | x_B \cup \{i\}}$$

$$\{i\} - B \cup (A \setminus \{i\}) - C \Rightarrow^{IH} (2) x_C \perp\!\!\! \perp x_i | x_B, x_{A \setminus \{i\}}$$

Intersection
 \Rightarrow Lemma(1, 2)

$$x_C \perp\!\!\! \perp x_A | x_B \wedge \begin{matrix} A - B - C \text{ sep} \\ + | B | = s \end{matrix}.$$

The intersection lemma uses $P(x) \geq 0$, and we actually need it or \exists counterexample $G \not\models P$ [HS]

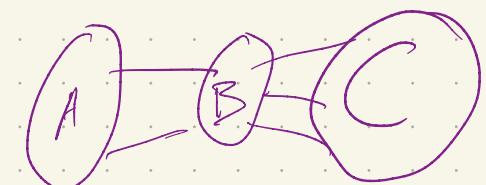
Def Factorization (F): We say $P(x)$ factorizes wrt G if

$$P(x) = \frac{1}{Z} \sum_{c \in C} f_c(x_c)$$

Maximal clique set

Claim: $(F) \Rightarrow (G)$

Proof: We have $\stackrel{\textcircled{1}}{P(x)} \geq \prod_{c \in C} f_c(x_c)$ & $\stackrel{\textcircled{2}}{A - B - C}$



want to show that $x_A \perp\!\!\!\perp x_c \mid x_B$

$$\propto F_{AB}(x_A, x_B) F_{BC}(x_B, x_C)$$

(HWI)

$x_A \perp\!\!\!\perp x_c \mid x_B$

No factor has both $a \in A, c \in C$ \square

Claim $(G) \Rightarrow (F)$

if $P(x) > 0$

Hammer Seley - Clifford thm

So, if $P(x) > 0$

$$(F) \Leftrightarrow (G) \Leftrightarrow L \Leftrightarrow (P)$$

Not true $G \not\models F$ when $P(x) > 0$

Proof $G \Rightarrow F$
 Fix (P, G) s.t. $P(x) > 0$, P satisfies (G) local Markov wrt G . ($(x_A \perp\!\!\!\perp x_C | x_B \wedge A - B - C \in G)$)

Want to find factorization s.t. $P(x) \propto \prod_{c \in \mathcal{C}} f_c(x_c)$

Def "Pseudo-factors"

$$\tilde{f}_S(x_S) \triangleq \prod_{U \subseteq S} P(x_U, O_{V \setminus U})$$

$$\hookrightarrow x_{S \setminus U} = \vec{0}$$

$$O \in \mathcal{X}$$

Example: the singleton pseudofactors:

$$\binom{n}{1} \text{ singletons } \tilde{f}_i(x_i) = \frac{P(x_i, O^{n-1})}{P(x_i=0, x_{-i}=O^{n-1})}$$

$$\text{Pairwise } f_{ij}(x_{ij}) = \frac{P(x_i, x_j, O^{n-2}) \cdot P(x_i=x_j=0, O^{n-2})}{P(x_i, x_j=0, O^{n-2}) P(x_i=0, x_j, O^{n-2})}$$

Recall
[Bayes Nets / Directed GMs]
 $I(G) = \text{The set of all } \perp\!\!\!\perp \text{ implied by } G = (V, E)$
 $I(P) = \text{the set of all } \perp\!\!\!\perp \text{ implied by}$
 factorization of P

- $\exists P(x)$ where no $\not\models G$ with $I(G) = I(P)$

$$x_1 = x_2$$

$$x_3 = x_1 + z = x_2 + z$$

$$P(x_1, x_2, x_3)$$

$$= \begin{cases} 0 & \text{if } x_1 \neq x_2 \\ P(x_1)P(z=x_3-x_1 | x_1) & \end{cases}$$

$$\underline{P(x_1, x_2) P(x_3 | x_1, x_2)}$$

- $\forall G, \exists P$ s.t. $I(G) = I(P)$
- $\exists G_1 \neq G_2$ s.t. $I(G_1) = I(G_2)$

Fix \rightarrow \leftarrow
Find example $x_1 \rightarrow x_2 | x_1$, $x_3 | x_2$, $x_1 | x_2 | x_1$, $x_3 | x_1$

Def G is a perfect map for P if $I(G) = I(P)$

G is an I -map for P if $I(G) \subseteq I(P)$

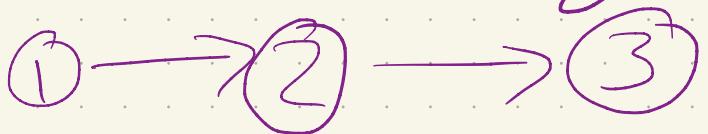
G is a minimal I -map for P if

• It is an I -map for P

• $G' = (V, E \setminus \{(i, j)\})$ isn't an imap
 $\forall (i, j) \in E$

④ Minimal IMap isn't unique

- Ordering of variables
- Even fixing ordering



$$x_1 = x_2$$

$$x_3 = x_1 + z$$

⑤ $I(G) = \cap I(P)$ where P factorizes wrt G