


Undirected GM / Markov Random Fields

A distribution $P(X)$ factorizes wrt $G = (V, E)$

$$\text{if } P(X) \propto \prod_{c \in \mathcal{C}} f_c(X_c)$$

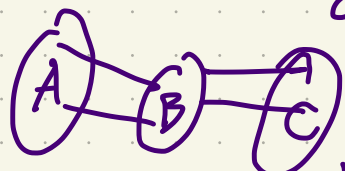
is proportional to $\rightarrow$ set of all cliques
or, equivalently, all maximal cliques

A normalization term can convert this into a probability distribution.

Graph Separation: We say A and C are separated by B

if $\forall A - B - C$

any path from any $a \in A$ to any $c \in C$ must pass through $b \in B$



[Global Markov Property (G)]

A distribution $P(X)$ is consistent w G if $\forall$ $A - B - C$ sets

$$X_A \perp\!\!\!\perp X_C \mid X_B$$

$$A, B, C \subseteq V$$

Local Markov Property (L)

$P(x)$ satisfies (L) wrt. $G = (V, E)$ if

$$x_i \perp\!\!\!\perp x_{V \setminus \{i\} \cup S_i} \mid x_{S_i} \quad S_i = \{j \mid (i,j) \in E\}$$

$$\forall i \in V$$

Pairwise Property (P)

$P(x)$ satisfies (P) wrt graph G if

$$\forall (i,j) \notin E \quad x_i \perp\!\!\!\perp x_j \mid x_{V \setminus \{i,j\}}$$

Claim $(G) \xRightarrow{(1)} (L) \xRightarrow{(2)} (P)$

Why $(G) \Rightarrow (L)$

(G) $\exists$ sets A, B, C s.t. $A - B - C$, we have that $x_A \perp\!\!\!\perp x_C \mid x_B$
for $P(x)$, Graph G $\rightarrow$ assume

$\Rightarrow (L)$: $x_i \perp\!\!\!\perp x_{V \setminus \{i\} \cup S_i} \mid x_{S_i}$ Want to show $\forall i$

Proof why? $A = \{i\}$ $B = \delta_i$ $C = V \setminus (\{i\} \cup \delta_i)$
 $A - B - C$ (b/y def of the sets)

$$X_A \perp\!\!\!\perp X_C \mid X_B$$

$$X_i \perp\!\!\!\perp X_{V \setminus \{i\} \cup \delta_i} \mid X_{\delta_i} (L)!$$

Claim $(L) \Rightarrow (P)$

$P(x)$ satisfies (P) if $\forall (i, j) \notin E \quad X_i \perp\!\!\!\perp X_j \mid X_{V \setminus \{i, j\}}$

Lemma [Data Processing inequalities]

suppose we have a graph s.t. $X - Z - Y - h(Y)$

Then we know $X \perp\!\!\!\perp Y \mid Z$, and

(DP1) $\bullet X \perp\!\!\!\perp h(Y) \mid Z \quad \forall \text{ functions } h(\cdot)$

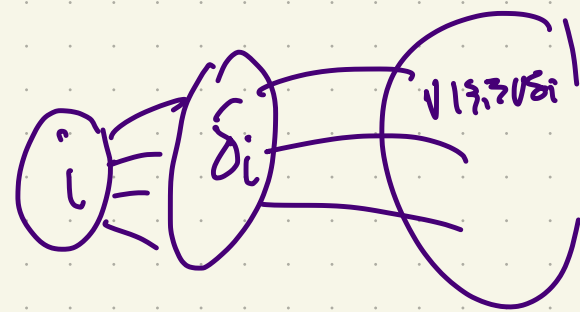
(Postprocessing doesn't break independences)

(DP2) $\bullet X \perp\!\!\!\perp Y \mid Z, h(Y)$

Take arbitrary $(i, j) \notin E$. Want to show

$$X_{\{i\}} \perp\!\!\!\perp X_{\{j\}} \mid X_{V \setminus \{i, j\}}$$

We know: $X_i \perp\!\!\!\perp X_{V \setminus (\{i\} \cup \delta_i)} \mid X_{\delta_i}$



(that $(i, j) \notin E \Rightarrow j \notin \delta_i$)

DP2 $\Rightarrow X_i \perp\!\!\!\perp X_{V \setminus (\{i\} \cup \delta_i)} \mid X_{\delta_i} \cup X_{V \setminus \{i, j\}}$

DP1 $\Rightarrow X_i \perp\!\!\!\perp X_j \mid X_{V \setminus \{i, j\}}$

□

Claim $(P) \Rightarrow (G)$ when $P(X) > 0$
 $\forall X \in \mathcal{X}^n$.

If $P(X) > 0 \quad \forall X \in \mathcal{X}$

(A)

(B)

(C)

(D)

Lemma

And

Intersection Lemma

(1) $X_A \perp\!\!\!\perp X_B \mid X_C, X_D$

(2) $X_A \perp\!\!\!\perp X_C \mid X_B, X_D$

$\Rightarrow X_A \perp\!\!\!\perp X_B, X_C \mid X_D$

Proof of Claim 1

We know $(P) \forall (i,j) \notin E \quad X_i \perp\!\!\!\perp X_j \mid X_{V \setminus \{i,j\}}$

Want to show $\forall A-B-C \quad X_A \perp\!\!\!\perp X_C \mid X_B$

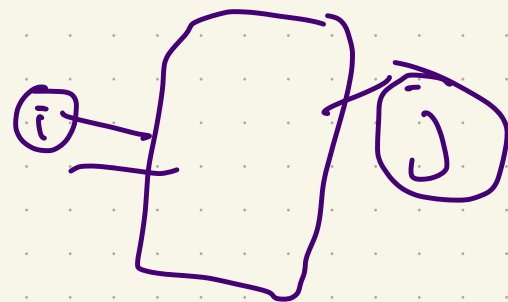
By induction on $|B| = n-2, n-3, \dots, 2$

(1) Start w $|B| = n-2$

$$A = \{i\}$$

$$C = \{j\}$$

$$X_i \perp\!\!\!\perp X_j$$

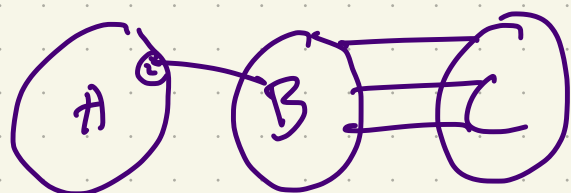


$$\hookrightarrow X_A \perp\!\!\!\perp X_C \mid X_{V \setminus \{i,j\}}$$

(2) Want to show (G) for $|B| = s-1 < n-2$

Assume (G) holds $\forall |B| \geq s$

Fix $|A| \geq 2$ (WLOG), assumption $A-B-C$



$$A \setminus \{i\} - B \cup \{i\} - C$$

$$\Rightarrow X_{A \setminus \{i\}} \perp\!\!\!\perp X_C \mid X_{B \cup \{i\}}$$

$\{i\} - \text{BUA} \setminus \{i\} - C$

[Why?]

$(IH) \Rightarrow X_i \perp\!\!\!\perp X_C \mid X_{\text{BUA} \setminus \{i\}} \quad \square$

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