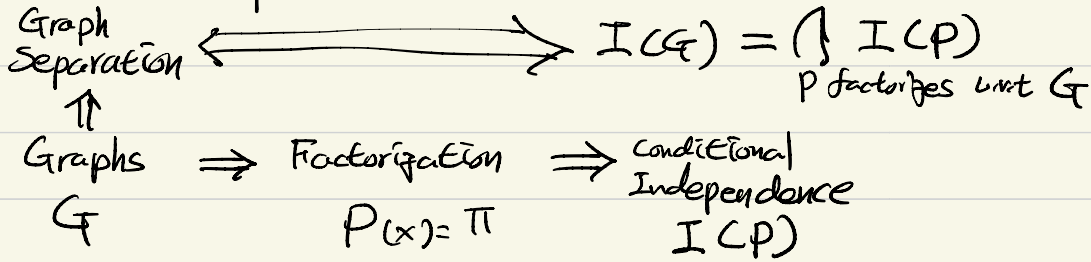


Probabilistic Graphical Models



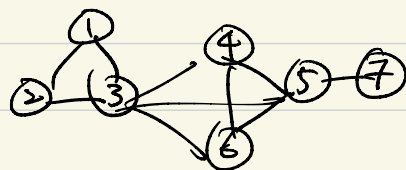
* Undirected Graphical Models

Def. Undirected Graph $\mathcal{G} = (V, E)$

$(i, j) = (j, i)$

- a set of nodes $C = \{i_1, i_2, \dots, i_k\}$ is a clique in \mathcal{G} if all pairs are connected in \mathcal{G} .

- a clique C is maximal clique if you cannot add any more nodes without breaking clique.



$\{1, 2, 3\}$

$\{3, 4, 5, 6\}$

$\{5, 7\}$

- \mathcal{C} to denote set of all maximal cliques.

Def. Undirected Graphical Model on $\mathcal{G} = (V, E)$ is

a family of distributions on $X = [X_1 \dots X_n]$ that factorizes as

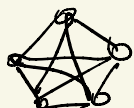
Markov
Random
Fields (MRF)

$$P(x) = \frac{1}{Z} \prod_{C \in \mathcal{C}} f_C(x_C)$$

$f_C: \mathcal{X}^{|C|} \rightarrow \mathbb{R}_+$ factors or compatibility functions

$Z \triangleq \sum_{x \in \mathcal{X}^n} \left\{ \prod_{C \in \mathcal{C}} f_C(x_C) \right\}$ called a Partition Function.

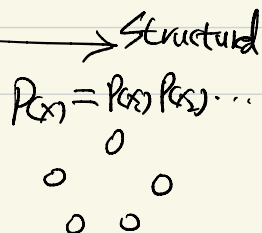
Unstructured



$P(x)$

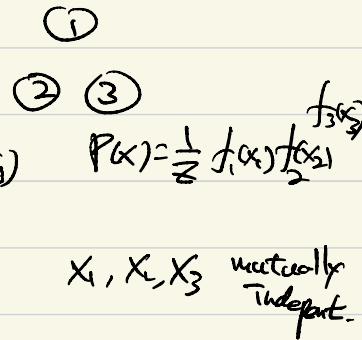
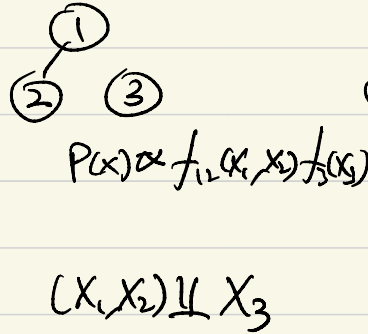
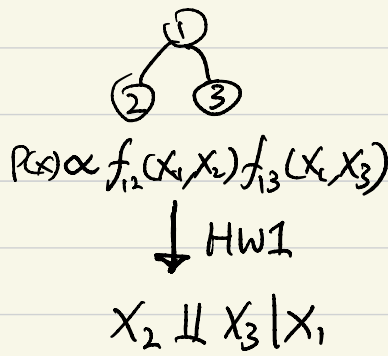
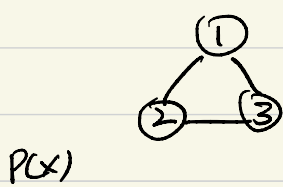
distribution: $|I(P)|$ is small \longleftrightarrow $|I(P)|$ large

graph: dense \longleftrightarrow sparse



Warm up examples >

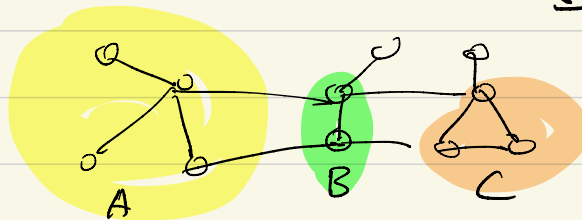
G:



Def. [Graph Separation]

$B \subset V$ separates $A \subset V$ if every path from any $a \in A$ to $c \in C$ passes through B .

\leftarrow I will casually write
If $A-B-C$

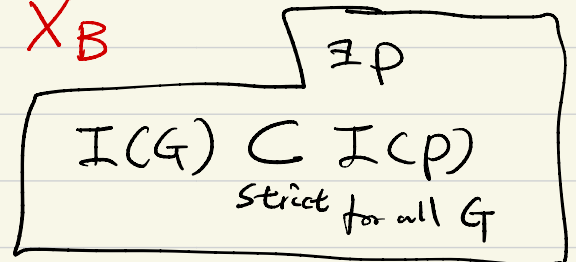


Def. [Global Markov Property (G)]

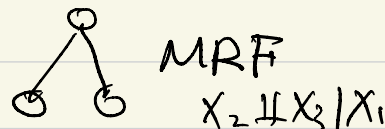
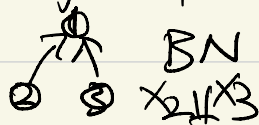
A distribution $P(x)$ satisfies Global Markov Property (G) if for any $A-B-C$ we have

$$X_A \perp\!\!\!\perp X_C \mid X_B$$

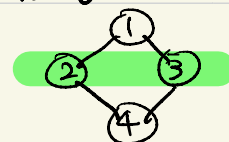
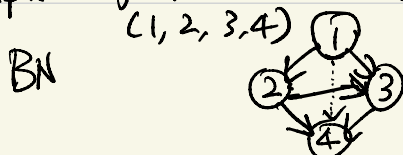
punch line: factorization \Rightarrow (G)



Q. give an example of Bayesian Network that cannot be represented by a MRF.



Q. give an example of MRF that cannot be represented by a BN.



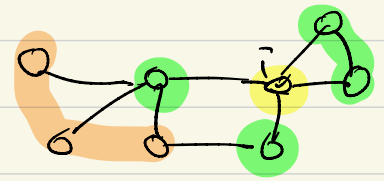
Def [Local Markov Property (L)]

$P(x)$ satisfies (L) wrt G if

$$X_i \perp\!\!\!\perp X_{rest} \mid X_{\partial_i}$$

$\partial_i = \text{neighborhood of node } i \cong \{j \in V \mid (i,j) \in E\}$

$$rest = V \setminus (\{i\} \cup \partial_i)$$

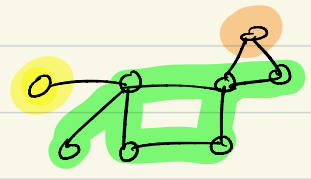


Def [Pairwise Markov Property (P)]

$P(x)$ satisfies (P) wrt G if

$$X_i \perp\!\!\!\perp X_j \mid X_{rest} \text{ for all } i, j \text{ not connected}$$

$$rest = V \setminus \{i, j\}$$



claim: $(G) \Rightarrow (L) \Rightarrow (P)$

proof 1: $(G) \Rightarrow (L)$

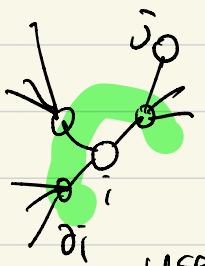
follows from $i - \partial_i - rest$
are separated in Graph.

$$(G): X_A \perp\!\!\!\perp X_C \mid X_B$$

$$(L): X_i \perp\!\!\!\perp X_{rest} \mid X_{\partial_i}$$

$$(P): X_i \perp\!\!\!\perp X_j \mid X_{rest}$$

proof 2: $(L) \Rightarrow (P)$



$$X_i \perp\!\!\!\perp X_{V \setminus (\{i\} \cup \partial_i)} \mid X_{\partial_i}$$

$$\xrightarrow{(a)} X_i \perp\!\!\!\perp X_j \mid X_{V \setminus (\{i\} \cup \partial_i)} \text{ any } j \text{ not in } \partial_i$$

$$\xrightarrow{(b)} X_i \perp\!\!\!\perp X_j \mid X_{V \setminus \{i, j\}} \iff (P)$$

(a) follows from $X \perp\!\!\!\perp Y \mid Z \Rightarrow X \perp\!\!\!\perp Y \mid (Z, h(Y))$ for any $h(Y)$

(b) follows from $X \perp\!\!\!\perp Y \mid Z \Rightarrow X \perp\!\!\!\perp h(Y) \mid Z$ for any $h(Y)$

↑ Data Processing Inequality. $X - [Z] - Y - h(Y)$

$$(G) \Rightarrow (L) \Rightarrow (P) \quad \text{PMF: } P(x) > 0 \quad \forall x = [x_1, \dots, x_n] \in \mathcal{X}^n$$

Claim: If $P(x) > 0 \quad \forall x$, then $(P) \Rightarrow (G)$

Lemma [intersection lemma]

For $P(x) > 0$, if $X_A \perp\!\!\!\perp X_B \mid (X_C, X_D)$
 $X_A \perp\!\!\!\perp X_C \mid (X_B, X_D)$ $\Rightarrow X_A \perp\!\!\!\perp (X_B, X_C) \mid X_D$

proof of claim: $X_i \perp\!\!\!\perp X_j \mid X_{\text{rest}} \Rightarrow X_A \perp\!\!\!\perp X_C \mid X_B$

show by induction, on all B of size $|B| = n-2, n-3, \dots$

① initially, $|B| = n-2$, $A = \{i\}, C = \{j\}$

true by (P)

② suppose by induction that (G) holds for all $|B| \geq s$.

we will show for B of size $|B| = s-1 < n-2$

suppose w.l.o.g. $|A| \geq 2$,

$A-B-C$ separated in Graph.

$\Rightarrow A \setminus \{i\} - B \cup \{i\} - C$ also separated.

\tilde{A} \uparrow size $\leq s-1$
 \tilde{B}

$X_C \perp\!\!\!\perp X_{A \setminus \{i\}} \mid X_{B \cup \{i\}}$ (1)

$\Rightarrow \{i\} - B \cup A \setminus \{i\} - C$ also separated

\tilde{A} \uparrow size $> s$
 \tilde{B}

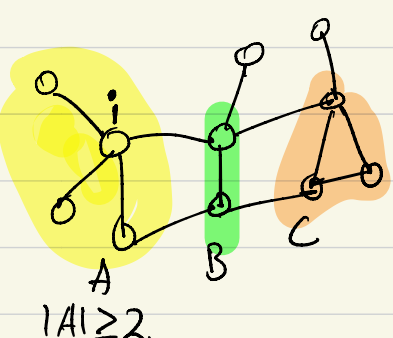
$X_C \perp\!\!\!\perp X_i \mid (X_B, X_{A \setminus \{i\}})$ (2)

Apply intersection lemma: (1) & (2)

$X_C \perp\!\!\!\perp X_A \mid X_B$.

$\forall A, B, C$ s.t. $A-B-C$ separated & $|B| = s$. \square

* If $P(x) \not> 0$, then we can construct a counterexample st.
 $(G) \not\Rightarrow (P)$. [HW1].



Def. [Factorization (F)]

we say $P(x)$ factorizes w.r.t G if

$$P(x) = \frac{1}{Z} \prod_{c \in \mathcal{C}} f_c(x_c)$$

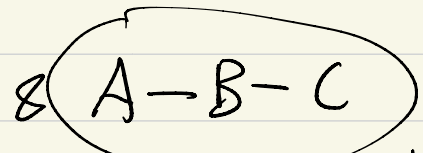
"maximal clique set"

Claim: (F) \Rightarrow (G)

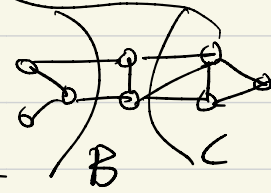
proof: $P(x) \propto \prod_{c \in \mathcal{C}} f_c(x_c)$

$$\propto F_{AB}(x_A, x_B) \cdot F_{BC}(x_B, x_C)$$

HV I $\longrightarrow x_A \perp x_C \mid x_B$



no factor has both $a \in A, c \in C$



Claim: (G) \Rightarrow (F) if $P(x) > 0$
 [Hammersley-Clifford theorem]

implication: for $P(x) > 0$, (F) \Leftrightarrow (G) \Leftrightarrow (L) \Leftrightarrow (P)

Counterexample: If $P(x) \neq 0$, then (G) $\not\Rightarrow$ (F)

