

Computer-Aided Reasoning for Software

CSLE507

The DPLL(T) Framework

Today

Last lecture

- Deciding conjunctions of $(T_1 \cup T_2)$ -constraints with Nelson-Oppen

Today

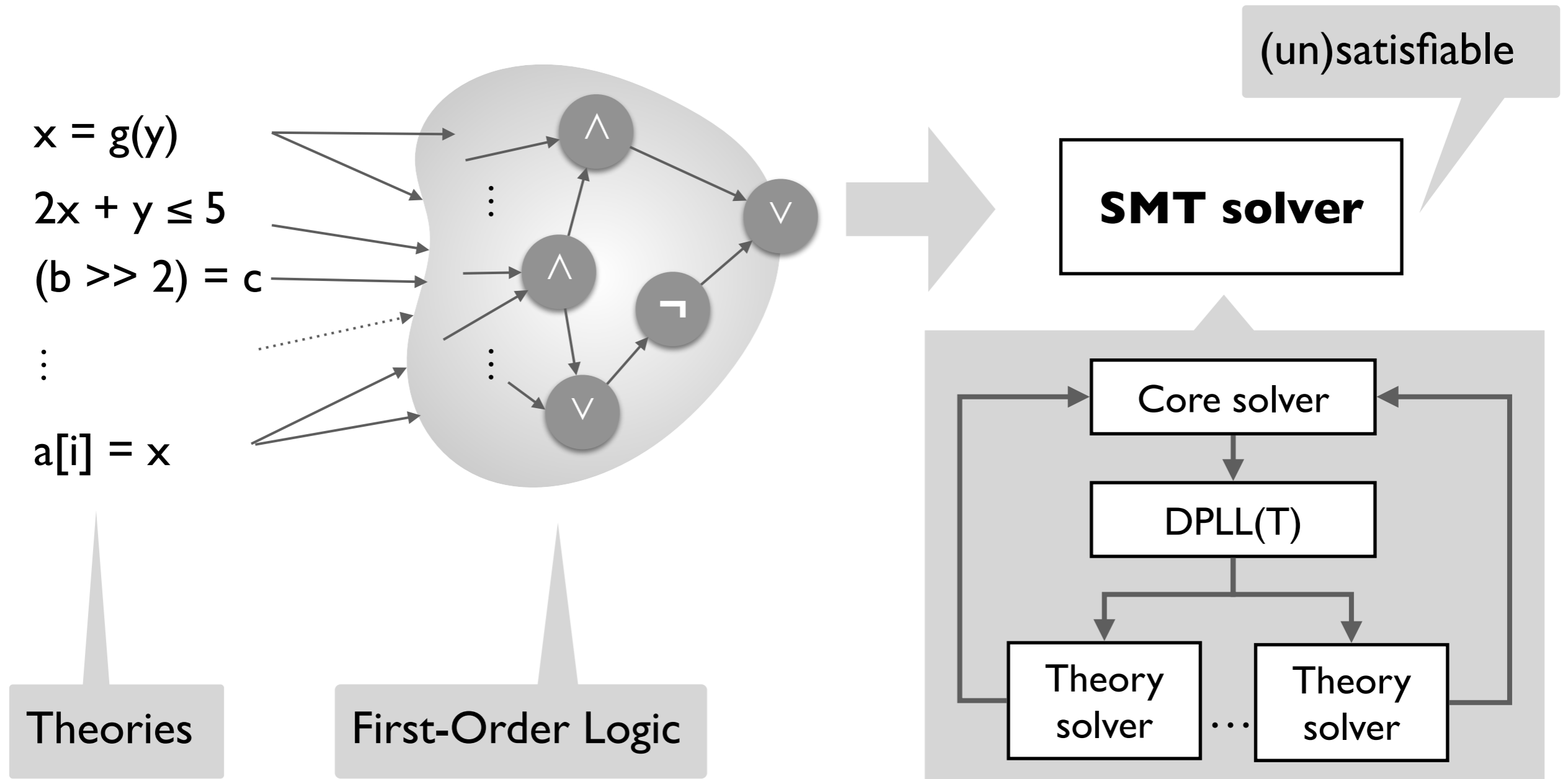
- Deciding arbitrary boolean combinations of theory constraints

Reminders

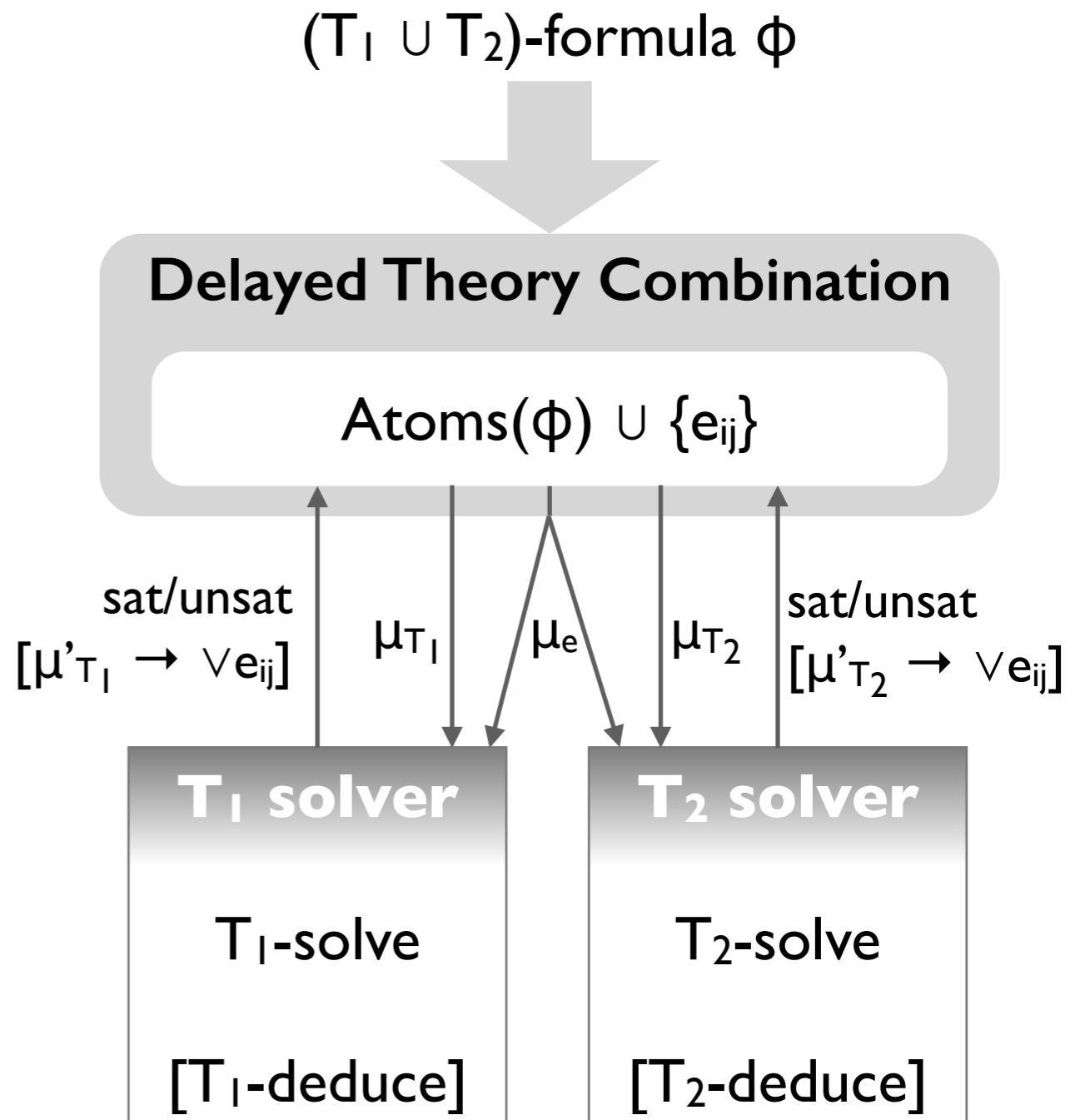
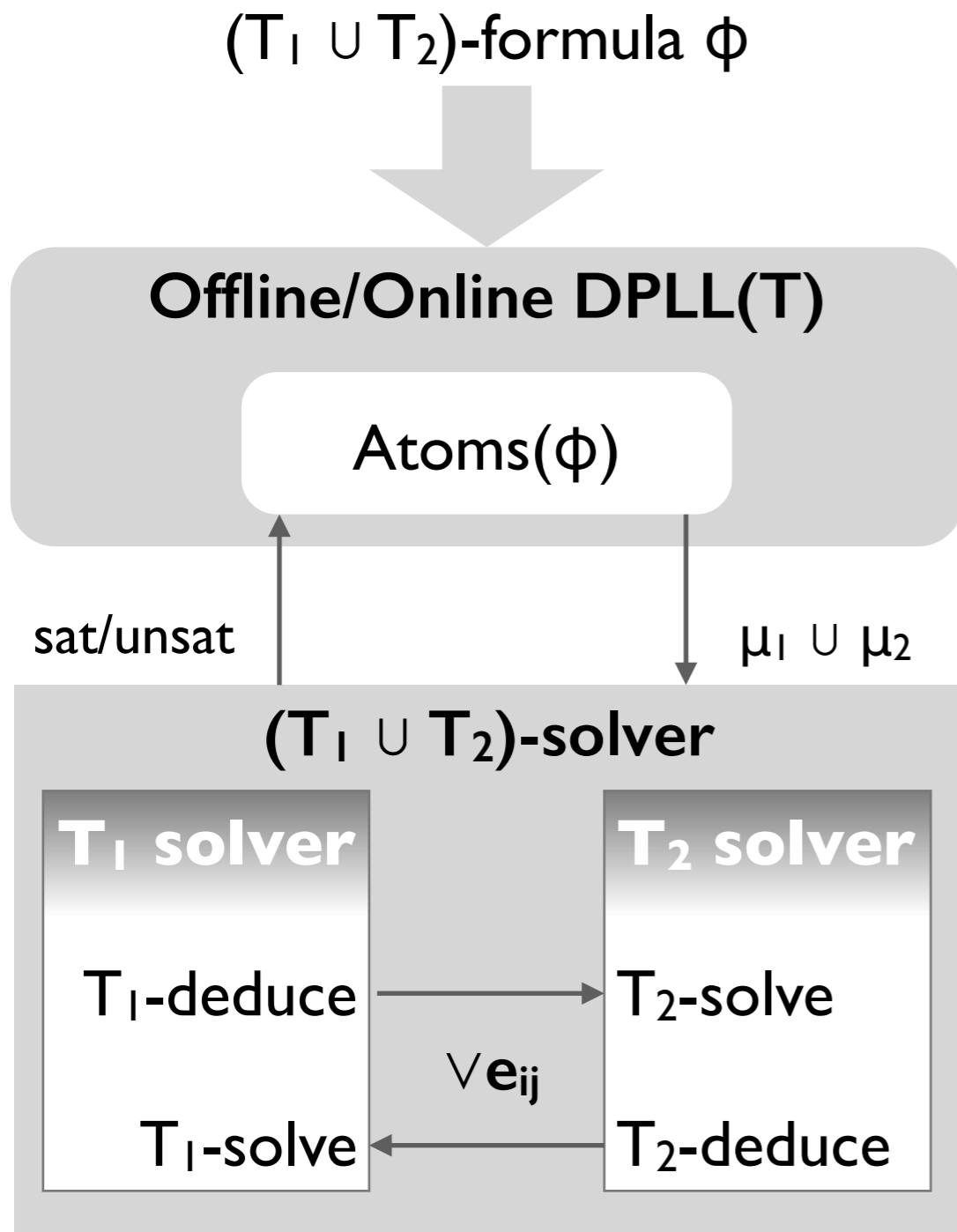
- Project proposal due tonight.



Recall: Satisfiability Modulo Theories (SMT)



The DPLL(T) Framework



Offline DPLL(T)

Assume ϕ is in CNF.

```
Offline-DPLLT(T-formula  $\phi$ )  
 $\phi^P \leftarrow$  T2B( $\phi$ )  
while (TRUE) do  
   $\mu^P, \text{res} \leftarrow$  CDCL( $\phi^P$ )  
  if res = UNSAT then return UNSAT  
  else  
    T-res  $\leftarrow$  T-solve(B2T( $\mu^P$ ))  
    if T-res = SAT then return SAT  
    else  $\phi^P \leftarrow \phi^P \wedge \neg \mu^P$ 
```

T2B computes the *boolean abstraction* (aka *boolean skeleton*) of ϕ by replacing every atom in ϕ with a fresh boolean variable.

B2T computes the *boolean refinement* of the current propositional assignment μ^P .

Theory conflict clause.

Boolean abstraction (T2B) and refinement (B2T)

T2B(ϕ)

- T2B(a_i) = b_i , if a_i is a theory atom and b_i is a fresh boolean variable
- T2B(b_j) = b_j , if b_j is a boolean variable
- T2B($\phi_1 \wedge \phi_2$) = T2B(ϕ_1) \wedge T2B(ϕ_2)
- T2B($\phi_1 \vee \phi_2$) = T2B(ϕ_1) \vee T2B(ϕ_2)
- T2B($\neg\phi_1$) = \neg T2B(ϕ_1)

$$\phi: (x = 1) \wedge ((x = 2) \vee (x = 3))$$

$$\text{T2B}(\phi): b_1 \wedge (b_2 \vee b_3)$$

$$\text{B2T}(b_1 \wedge b_3): (x = 1) \wedge (x = 3)$$

$$\text{B2T}(\phi^p) = \text{T2B}^{-1}(\phi^p)$$

Soundness and termination of offline DPLL(T)

```
Offline-DPLLT(T-formula  $\phi$ )  
 $\phi^P \leftarrow \mathbf{T2B}(\phi)$   
while (TRUE) do  
   $\mu^P, \text{res} \leftarrow \mathbf{CDCL}(\phi^P)$   
  if res = UNSAT then return UNSAT  
  else  
    T-res  $\leftarrow$  T-solve( $\mathbf{B2T}(\mu^P)$ )  
    if T-res = SAT then return SAT  
    else  $\phi^P \leftarrow \phi^P \wedge \neg \mu^P$ 
```

The algorithm is sound because ϕ^P overapproximates the satisfiability of ϕ at the beginning of each loop iteration (so it never returns UNSAT for SAT formulas), and it only returns SAT when μ^P both satisfies ϕ^P and is consistent with the theory axioms.

The algorithm terminates because there are finitely many satisfying assignments to the boolean abstraction, and we get a different one every time.

Offline DPLL(T) example

```
Offline-DPLLT(T-formula  $\phi$ )  
 $\phi^P \leftarrow \mathbf{T2B}(\phi)$   
while (TRUE) do  
   $\mu^P, \text{res} \leftarrow \mathbf{CDCL}(\phi^P)$   
  if res = UNSAT then return UNSAT  
  else  
    T-res  $\leftarrow$  T-solve( $\mathbf{B2T}(\mu^P)$ )  
    if T-res = SAT then return SAT  
    else  $\phi^P \leftarrow \phi^P \wedge \neg \mu^P$ 
```

- $\phi \leftarrow (x = 1) \wedge ((x = 2) \vee (x = 3))$
- $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3)$
- $\mu^P \leftarrow b_1 \wedge b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3)$
- $\mu^P \leftarrow b_1 \wedge \neg b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3) \wedge (\neg b_1 \vee b_2 \vee \neg b_3)$
- $\mu^P \leftarrow b_1 \wedge b_2 \wedge \neg b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3) \wedge (\neg b_1 \vee b_2 \vee \neg b_3) \wedge (\neg b_1 \vee \neg b_2 \vee b_3)$
- UNSAT

Offline DPLL(T): improvements

```
Offline-DPLLT(T-formula  $\phi$ )  
 $\phi^P \leftarrow \mathbf{T2B}(\phi)$   
while (TRUE) do  
   $\mu^P, \text{res} \leftarrow \mathbf{CDCL}(\phi^P)$   
  if res = UNSAT then return UNSAT  
  else  
    T-res  $\leftarrow$  T-solve( $\mathbf{B2T}(\mu^P)$ )  
    if T-res = SAT then return SAT  
    else  $\phi^P \leftarrow \phi^P \wedge \neg\mu^P$ 
```

The *theory conflict clause* $\neg\mu^P$ is too weak; it blocks one assignment at a time. What is a better clause?

$\neg\mathbf{T2B}(\mathbf{MINIMALUNSATCORE}(\mathbf{B2T}(\mu^P)))$

- $\phi \leftarrow (x = 1) \wedge ((x = 2) \vee (x = 3))$
- $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3)$
- $\mu^P \leftarrow b_1 \wedge b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3)$
- $\mu^P \leftarrow b_1 \wedge \neg b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3) \wedge (\neg b_1 \vee b_2 \vee \neg b_3)$
- $\mu^P \leftarrow b_1 \wedge b_2 \wedge \neg b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2 \vee \neg b_3) \wedge (\neg b_1 \vee b_2 \vee \neg b_3) \wedge (\neg b_1 \vee \neg b_2 \vee b_3)$
- UNSAT

Offline DPLL(T): improvements

```
Offline-DPLLT(T-formula  $\phi$ )  
 $\phi^P \leftarrow \mathbf{T2B}(\phi)$   
while (TRUE) do  
   $\mu^P, \text{res} \leftarrow \mathbf{CDCL}(\phi^P)$   
  if res = UNSAT then return UNSAT  
  else  
    T-res  $\leftarrow$  T-solve( $\mathbf{B2T}(\mu^P)$ )  
    if T-res = SAT then return SAT  
    else  
       $t \leftarrow \mathbf{T2B}(\mathbf{MINUNSATCORE}(\mathbf{B2T}(\mu^P)))$   
       $\phi^P \leftarrow \phi^P \wedge \neg t$ 
```

- $\phi \leftarrow (x = 1) \wedge ((x = 2) \vee (x = 3))$
- $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3)$
- $\mu^P \leftarrow b_1 \wedge b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2)$
- $\mu^P \leftarrow b_1 \wedge \neg b_2 \wedge b_3$
 - $\phi^P \leftarrow b_1 \wedge (b_2 \vee b_3) \wedge (\neg b_1 \vee \neg b_2) \wedge (\neg b_1 \vee \neg b_3)$
- UNSAT

Better but still need a *full assignment* to the boolean abstraction in order to generate a conflict clause.

Online DPLL(T) address this issue.

Online DPLL(\mathcal{T})

```
Online-DPLL $\mathcal{T}$ ( $\mathcal{T}$ -formula  $\phi$ ,  $\mathcal{T}$ -assignment  $\mu$ )  
if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
  return UNSAT  
 $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
while (TRUE) do  
  T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
  while (TRUE) do  
     $res \leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
    if  $res =$  SAT then return SAT  
    else if  $res =$  CONFLICT  
       $blevel \leftarrow$  T-ANALZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
      if ( $blevel < 0$ ) then return UNSAT  
      else T-BACKTRACK( $blevel$ ,  $\phi^P$ ,  $\mu^P$ )  
    else break
```

Everything passed by reference.

All procedures have access to T2B and B2T.

Online DPLL(T): T -PREPROCESS

```
Online-DPLL $_T$ ( $T$ -formula  $\phi$ ,  $T$ -assignment  $\mu$ )  
  if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
    return UNSAT  
   $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
  while (TRUE) do  
    T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
    while (TRUE) do  
       $res \leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
      if  $res = SAT$  then return SAT  
      else if  $res = CONFLICT$   
         $blevel \leftarrow$  T-ANALZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
        if ( $blevel < 0$ ) then return UNSAT  
        else T-BACKTRACK( $blevel$ ,  $\phi^P$ ,  $\mu^P$ )  
      else break
```

Simplifies ϕ and updates μ , if needed, so that equisatisfiability is preserved.

Common simplifications:

- Drop dual operators
- Exploit associativity
- Sort arguments
- Exploit theory-specific properties

Online DPLL(\mathcal{T}): \mathcal{T} -DECIDE

```
Online-DPLL $\mathcal{T}$ ( $\mathcal{T}$ -formula  $\phi$ ,  $\mathcal{T}$ -assignment  $\mu$ )  
  if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
    return UNSAT  
   $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
  while (TRUE) do  
    T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
    while (TRUE) do  
      res  $\leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
      if res = SAT then return SAT  
      else if res = CONFLICT  
        blevel  $\leftarrow$  T-ANALZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
        if (blevel < 0) then return UNSAT  
        else T-BACKTRACK(blevel,  $\phi^P$ ,  $\mu^P$ )  
      else break
```

Analogous to DECIDE in CDCL:

- Selects an unassigned l^P literal and adds it to μ^P .
- May consider the semantics of literals in \mathcal{T} .

Online DPLL(T): T-DEDUCE

```
Online-DPLLT(T-formula  $\phi$ , T-assignment  $\mu$ )  
  if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
    return UNSAT  
   $\phi^P, \mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
  while (TRUE) do  
    T-DECIDE( $\phi^P, \mu^P$ )  
    while (TRUE) do  
      res  $\leftarrow$  T-DEDUCE( $\phi^P, \mu^P$ )  
      if res = SAT then return SAT  
      else if res = CONFLICT  
        blevel  $\leftarrow$  T-ANALZECONFLICT( $\phi^P, \mu^P$ )  
        if (blevel < 0) then return UNSAT  
        else T-BACKTRACK(blevel,  $\phi^P, \mu^P$ )  
      else break
```

Applies BCP to ϕ^P and μ^P until

- μ^P propositionally violates ϕ^P : returns CONFLICT.
- μ^P propositionally satisfies ϕ^P : invokes T-solver on $B2T(\mu^P)$ and returns SAT if T-solver does. Otherwise returns CONFLICT.
- no more literals can be deduced: invokes T-solver on partial assignment $B2T(\mu^P)$ and returns CONFLICT if T-solver returns UNSAT. This is *early pruning*. May also do *theory propagation*.

Online DPLL(T): T-ANALYZECONFLICT

```
Online-DPLLT(T-formula  $\phi$ , T-assignment  $\mu$ )  
  if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
    return UNSAT  
   $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
  while (TRUE) do  
    T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
    while (TRUE) do  
       $res \leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
      if  $res =$  SAT then return SAT  
      else if  $res =$  CONFLICT  
         $blevel \leftarrow$  T-ANALYZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
        if ( $blevel < 0$ ) then return UNSAT  
        else T-BACKTRACK( $blevel$ ,  $\phi^P$ ,  $\mu^P$ )  
      else break
```

Extends ANALYZECONFLICT from CDCL:

- if the conflict is caused by a boolean (BCP) failure, returns the same level and conflict clause as ANALYZECONFLICT
- if the conflict is caused by a theory failure, returns a mixed boolean+theory conflict clause

Online DPLL(**T**): **T**-BACKTRACK

```
Online-DPLLT( $T$ -formula  $\phi$ ,  $T$ -assignment  $\mu$ )  
  if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
    return UNSAT  
   $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
  while (TRUE) do  
    T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
    while (TRUE) do  
       $res \leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
      if  $res =$  SAT then return SAT  
      else if  $res =$  CONFLICT  
         $blevel \leftarrow$  T-ANALZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
        if ( $blevel < 0$ ) then return UNSAT  
        else T-BACKTRACK( $blevel$ ,  $\phi^P$ ,  $\mu^P$ )  
        else break
```

Analogous to BACKTRACK in CDCL:

- adds learned clause to ϕ^P (*T-learning*).
- backtracks to $blevel$ by undoing all the assignments $> blevel$ (*T-backjumping*).

Online DPLL(T) example

T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee$
 $\neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1,$
 $b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$

Online DPLL(T) example: T-Decide

T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$

$\neg b_5@1$

$\neg b_1@4$

$b_6@3$

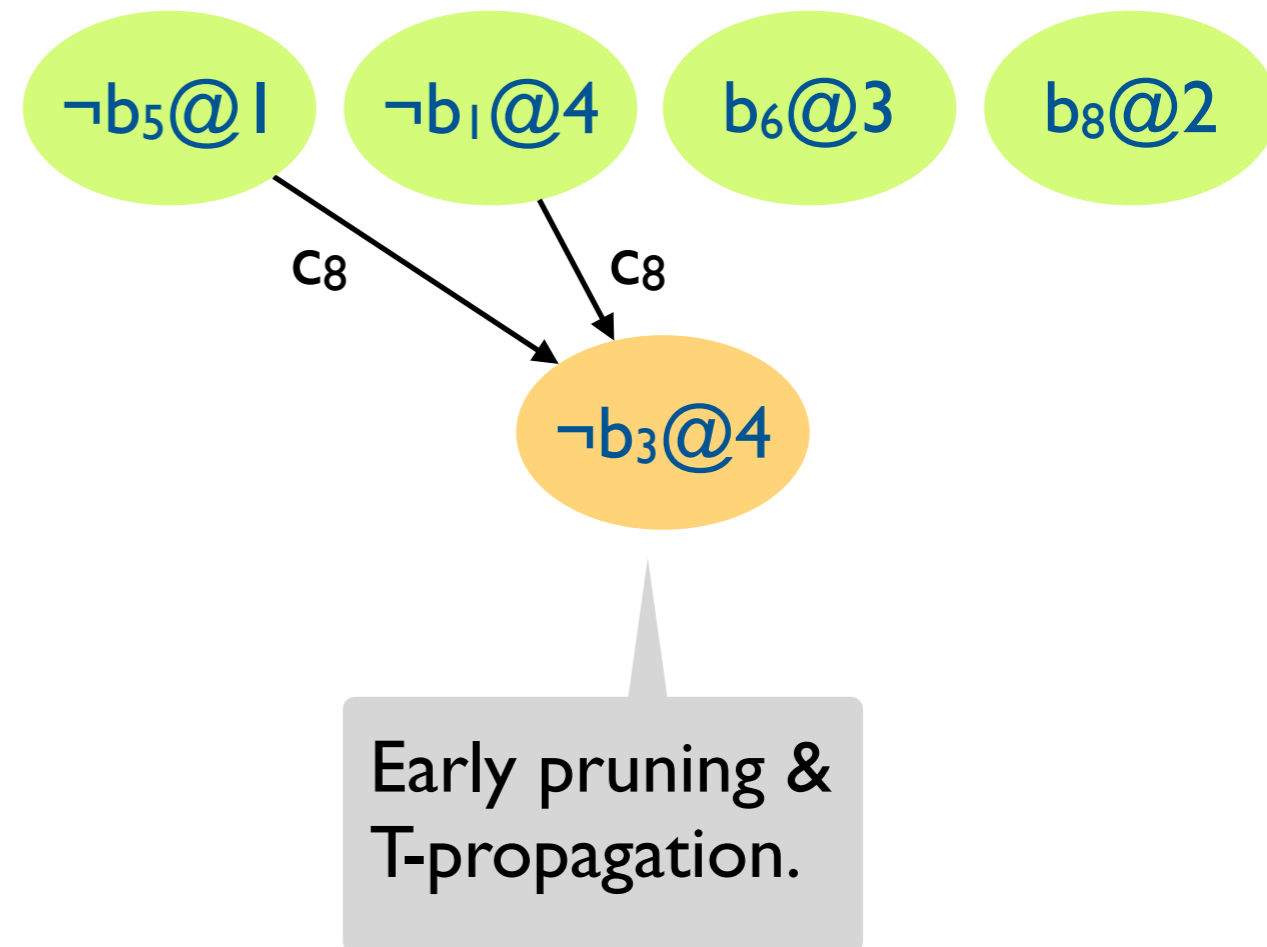
$b_8@2$

T-DECIDE makes 4 decisions.

Online DPLL(T) example: T-Deduce

T_R -formula ϕ and ϕ^P :

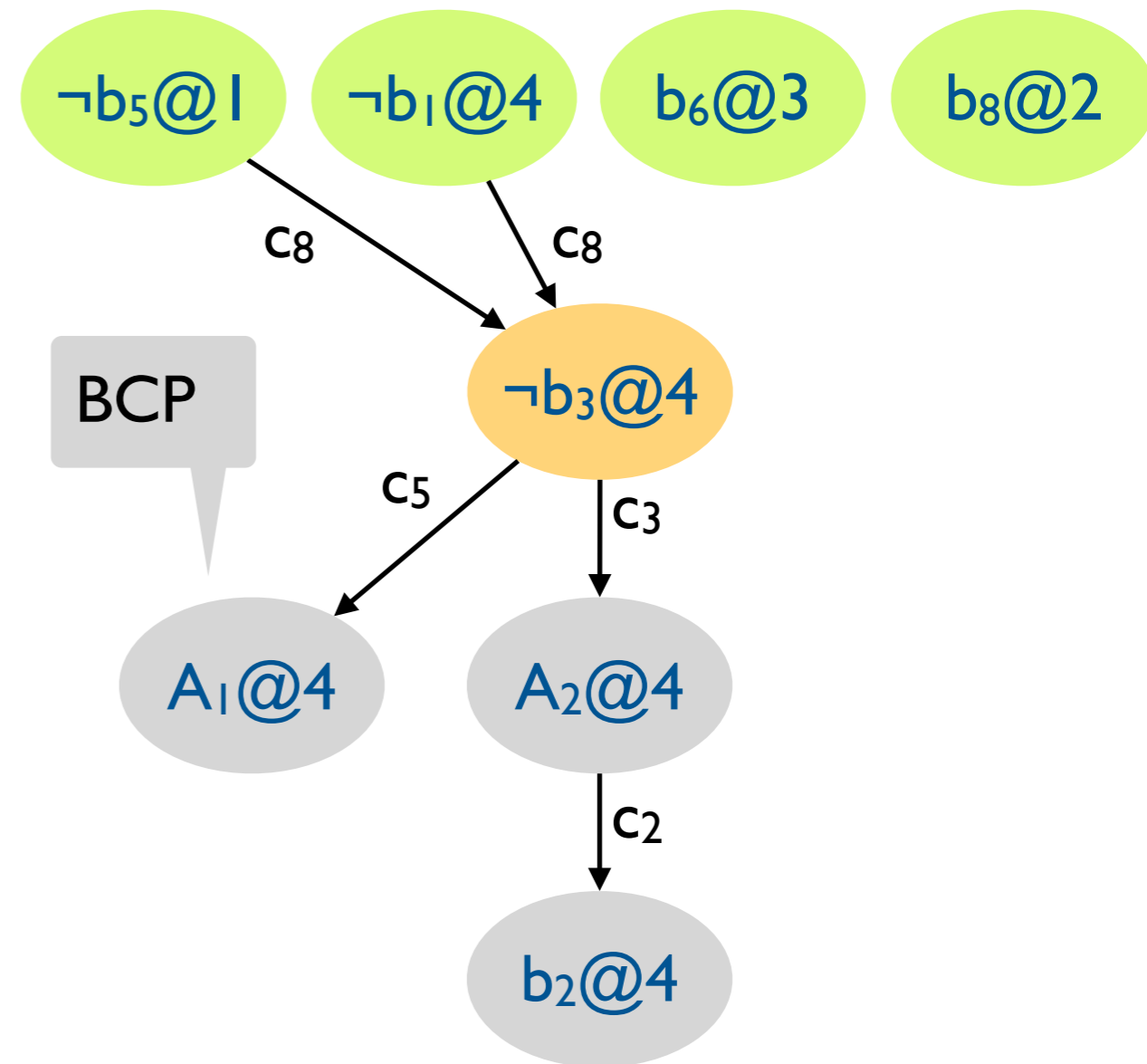
1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$
8. $b_5 \vee b_1 \vee \neg b_3$



Online DPLL(T) example: T-Deduce

T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$
8. $b_5 \vee b_1 \vee \neg b_3$



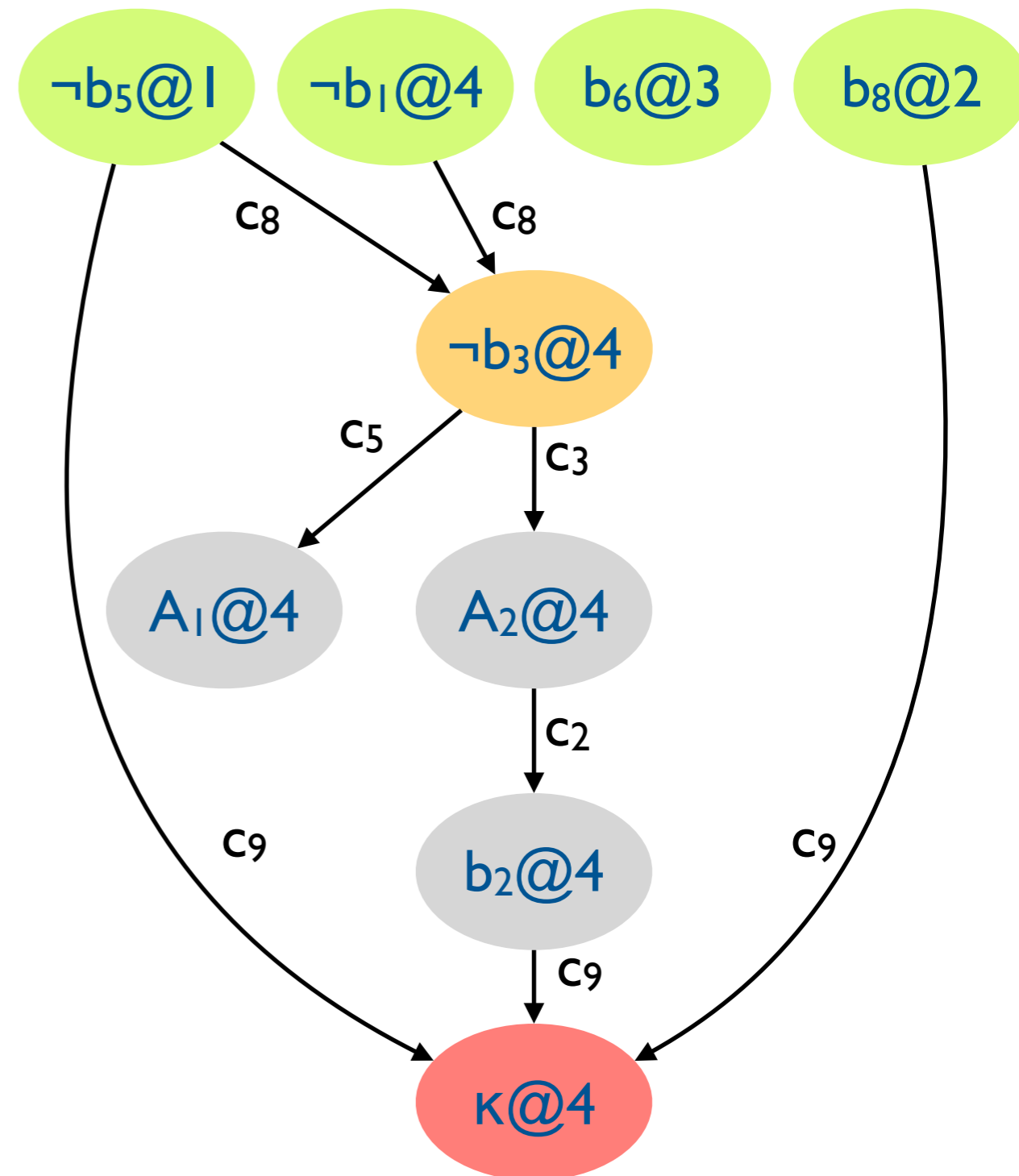
Online DPLL(T) example: T-Deduce

Current partial assignment is T-UNSAT.

T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$
8. $b_5 \vee b_1 \vee \neg b_3$
9. $b_5 \vee \neg b_8 \vee \neg b_2$

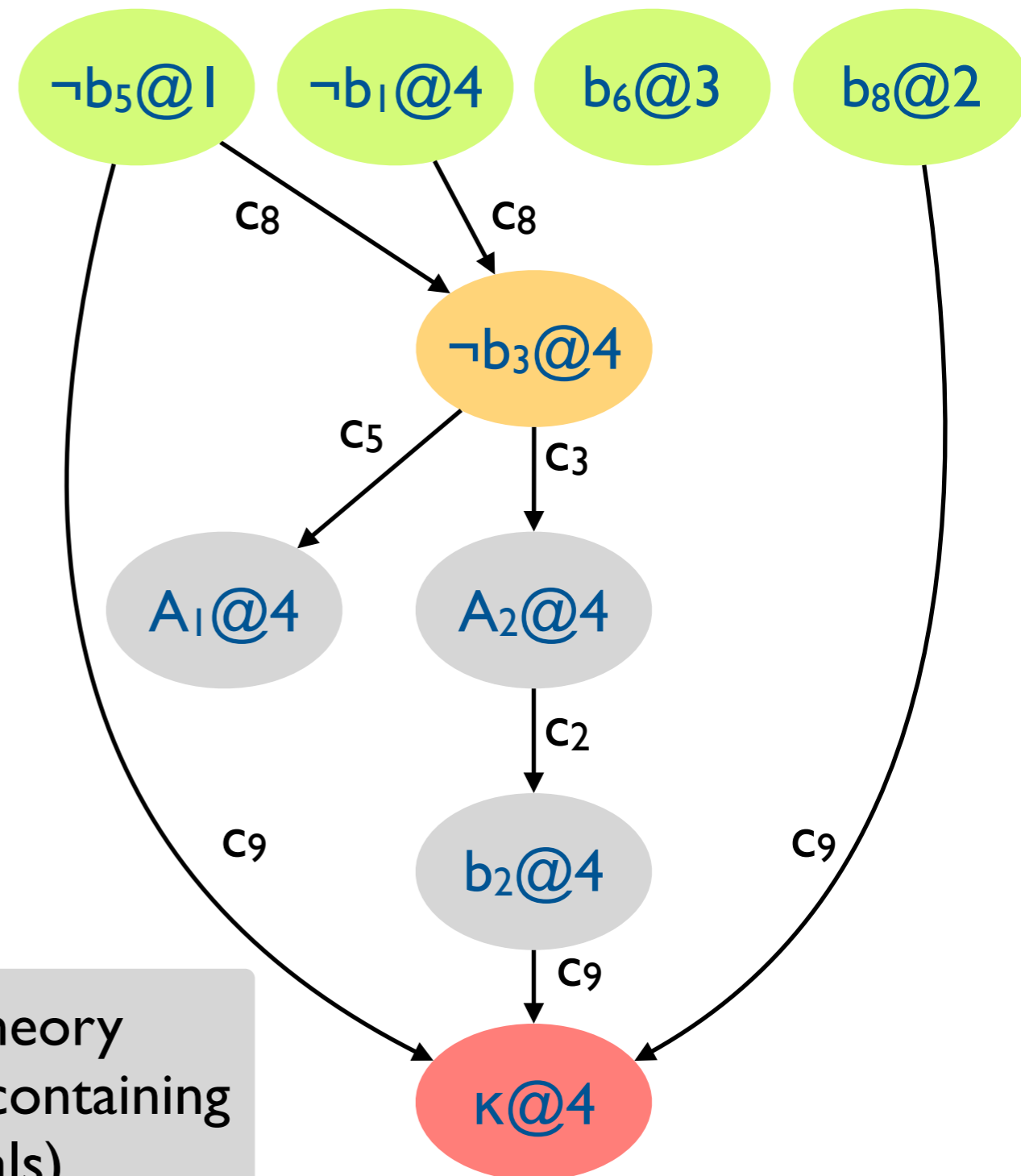
Theory conflict.



Online DPLL(T) example: T-AnalyzeConflict

T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$
8. $b_5 \vee b_1 \vee \neg b_3$
9. $b_5 \vee \neg b_8 \vee \neg b_2$
10. $b_5 \vee \neg b_8 \vee b_1$

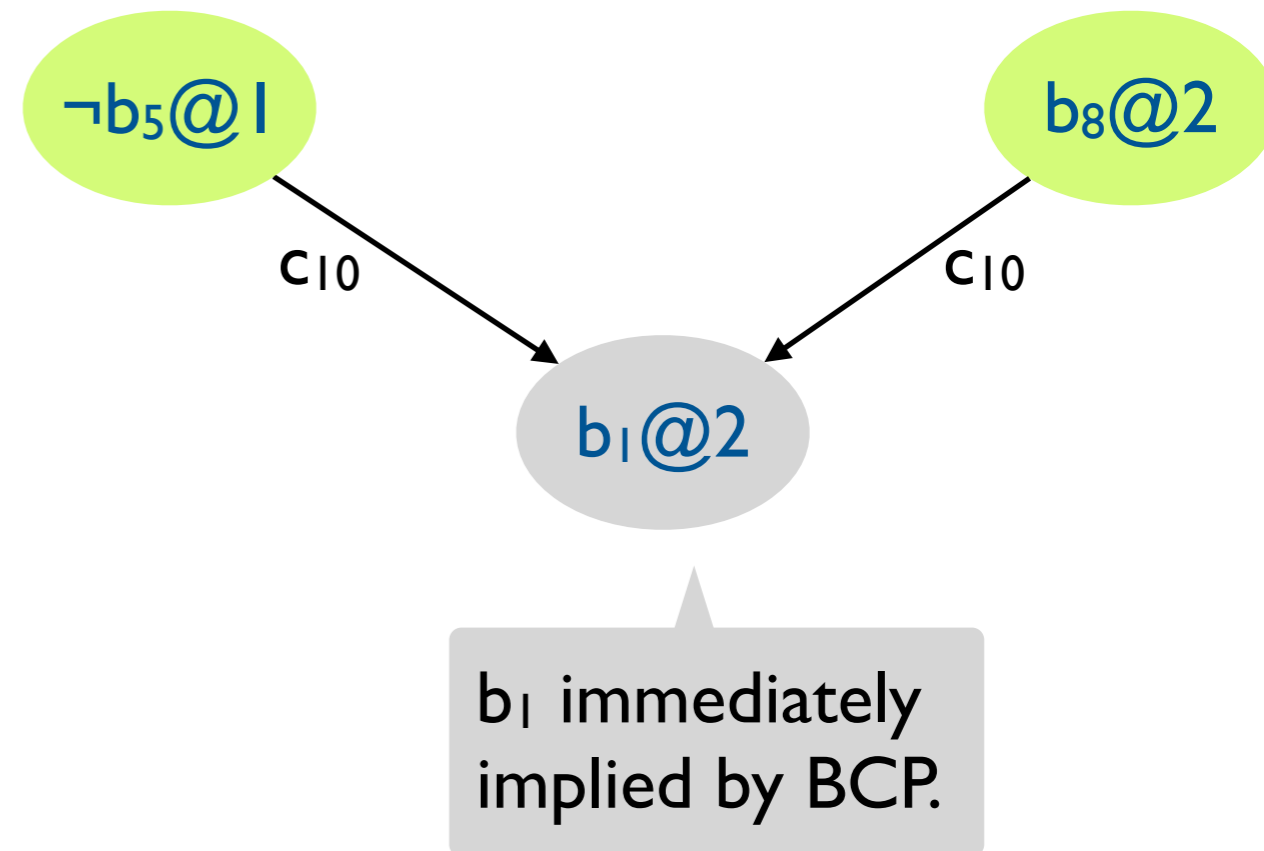


Mixed boolean+theory conflict (last UIP, containing just decision literals).

Online DPLL(T) example: T-Backtrack

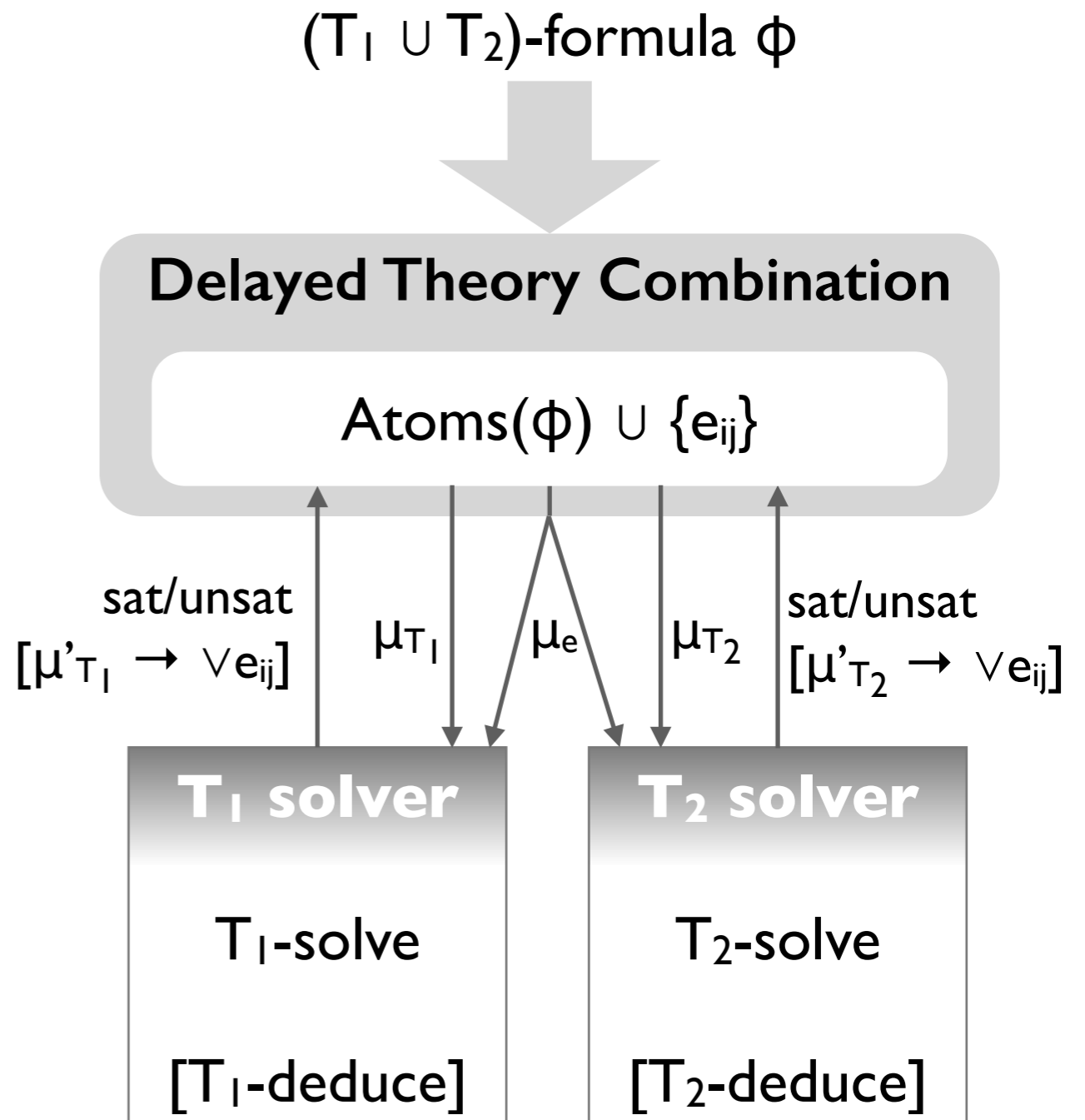
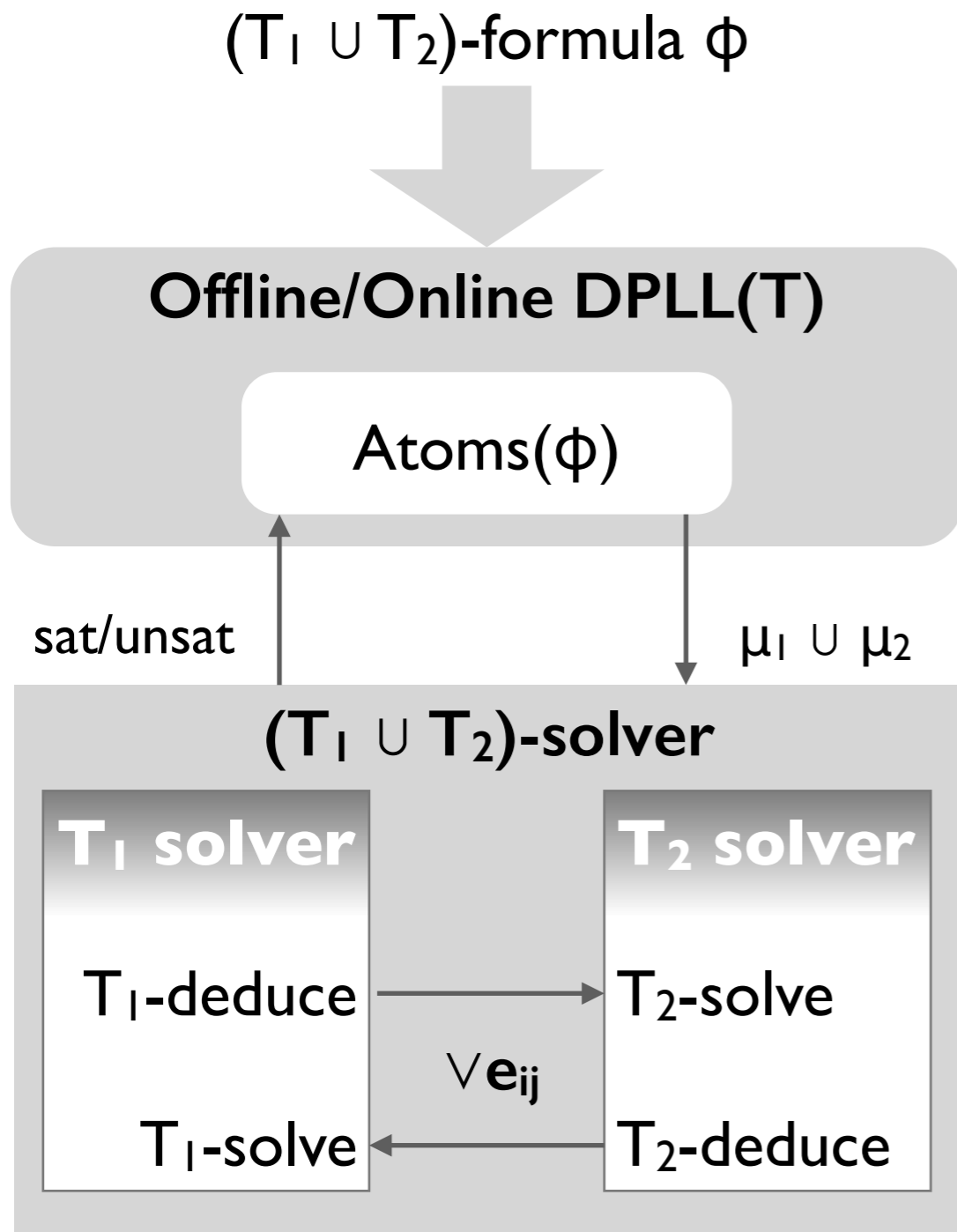
T_R -formula ϕ and ϕ^P :

1. $\neg(2x_2 - x_3 > 2) \vee A_1, \neg b_1 \vee A_1$
2. $\neg A_2 \vee (x_1 - x_5 \leq 1), \neg A_2 \vee b_2$
3. $(3x_1 - 2x_2 \leq 3) \vee A_2, b_3 \vee A_2$
4. $\neg(2x_3 + x_4 \geq 5) \vee \neg(3x_1 - x_3 \leq 6) \vee \neg A_1, \neg b_4 \vee \neg b_5 \vee \neg A_1$
5. $A_1 \vee (3x_1 - 2x_2 \leq 3), A_1 \vee b_3$
6. $(x_2 - x_4 \leq 6) \vee (x_5 = 5 - 3x_4) \vee \neg A_1, b_6 \vee b_7 \vee \neg A_1$
7. $A_1 \vee (x_3 = 3x_5 + 4) \vee A_2, A_1 \vee b_8 \vee A_2$
8. $b_5 \vee b_1 \vee \neg b_3$
9. $b_5 \vee \neg b_8 \vee \neg b_2$
10. $b_5 \vee \neg b_8 \vee b_1$



Backtrack to the second highest decision level (2).

The DPLL(T) Framework



Delayed Theory Combination (DTC)

```
Online-DPLLT(T-formula  $\phi$ , T-assignment  $\mu$ )  
if T-PREPROCESS( $\phi$ ,  $\mu$ ) = CONFLICT then  
  return UNSAT  
 $\phi^P$ ,  $\mu^P \leftarrow$  T2B( $\phi$ ), T2B( $\mu$ )  
while (TRUE) do  
  T-DECIDE( $\phi^P$ ,  $\mu^P$ )  
  while (TRUE) do  
     $res \leftarrow$  T-DEDUCE( $\phi^P$ ,  $\mu^P$ )  
    if  $res = SAT$  then return SAT  
    else if  $res = CONFLICT$   
       $blevel \leftarrow$  T-ANALYZECONFLICT( $\phi^P$ ,  $\mu^P$ )  
      if ( $blevel < 0$ ) then return UNSAT  
      else T-BACKTRACK( $blevel$ ,  $\phi^P$ ,  $\mu^P$ )  
    else break
```

To get DTC, modify Online-DPLL_T so that

- Truth values assigned to both atoms in ϕ and the *interface equalities* e_{ij} not in ϕ .
- T-DECIDE branches on interface equalities e_{ij} after μ propositionally satisfies ϕ .
- T-DEDUCE passes $\mu_i \cup \mu_e$ to each T_i -solver and returns SAT if both return SAT. Otherwise returns CONFLICT.
- T-ANALYZECONFLICT and T-BACKTRACK use the conflict clause (possibly containing interface equalities) from one of the T_i -solvers.
- Early pruning and T-PROPAGATION are performed.

Summary

Today

- The DPLL(T) framework for deciding SMT formulas

Next lecture

- Finite model finding: reasoning about quantified formulas over finite domains
- Last lecture on **computer-aided reasoning**
- It's all **for software** afterwards!