Computer-Aided Reasoning for Software

CSE507

Reasoning about Programs II

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Overview

topics

Last week

- Reasoning about (partial) correctness with Hoare Logic
- Reasoning about total correctness with Dafny

Today

- Automating Hoare Logic with verification condition generation—how tools like Dafny work

Recap: Imperative Programming Language (IMP)

Expression E

- $Z \mid V \mid E_1 + E_2 \mid E_1 * E_2$

Conditional C

- $true \mid false \mid E_1 = E_2 \mid E_1 \leq E_2$

Statement S

• **skip**	(Skip)
• **abort**	(Abort)
• V := E	(Assignment)
• $S_1; S_2$	(Composition)
• **if** C **then** S_1 **else** S_2	(If)
• **while** C **do** S	(While)

Recap: inference rules for Hoare logic

$$\frac{}{\vdash \{P\}\ \textbf{skip}\ \{P\}}$$

$$\frac{\vdash \{P\}\ S_1\ \{R\} \qquad \vdash \{R\}\ S_2\ \{Q\}}{\vdash \{P\}\ S_1; S_2\ \{Q\}}$$

$$\frac{}{\vdash \{\mathrm{true}\}\ \textbf{abort}\ \{\mathrm{false}\}}$$

$$\frac{\vdash \{P \wedge C\}\ S_1\ \{Q\} \quad \vdash \{P \wedge \neg C\}\ S_2\ \{Q\}}{\vdash \{P\}\ \textbf{if}\ C\ \textbf{then}\ S_1\ \textbf{else}\ S_2\ \{Q\}}$$

$$\frac{}{\vdash \{Q[E/x]\}\ x := E\ \{Q\}}$$

$$\frac{\vdash \{P_1\}\ S\ \{Q_1\} \qquad P \Rightarrow P_1 \quad Q_1 \Rightarrow Q}{\vdash \{P\}\ S\ \{Q\}}$$

$$\frac{\vdash \{P \wedge C\}\ S\ \{P\}}{\vdash \{P\}\ \textbf{while}\ C\ \textbf{do}\ S\ \{P \wedge \neg C\}}$$

loop invariant

Challenge: manual proof construction is tedious!

```
{x ≤ n}
while (x < n) do
  {x ≤ n ∧ x < n}
  {x+1 ≤ n}               // consequence
  x := x + 1
  {x ≤ n}                 // assignment
{x ≤ n ∧ x ≥ n}           // while
{x = n}                   // consequence
```

Hoare Logic proofs are highly manual:

- When to apply the rule of consequence?
- What loop invariants to use?

Challenge: manual proof construction is tedious!

```
{x ≤ n}                    // precondition
while (x < n) do
  {x ≤ n          }        // loop invariant

  x := x + 1


{x = n}                    // postcondition
```

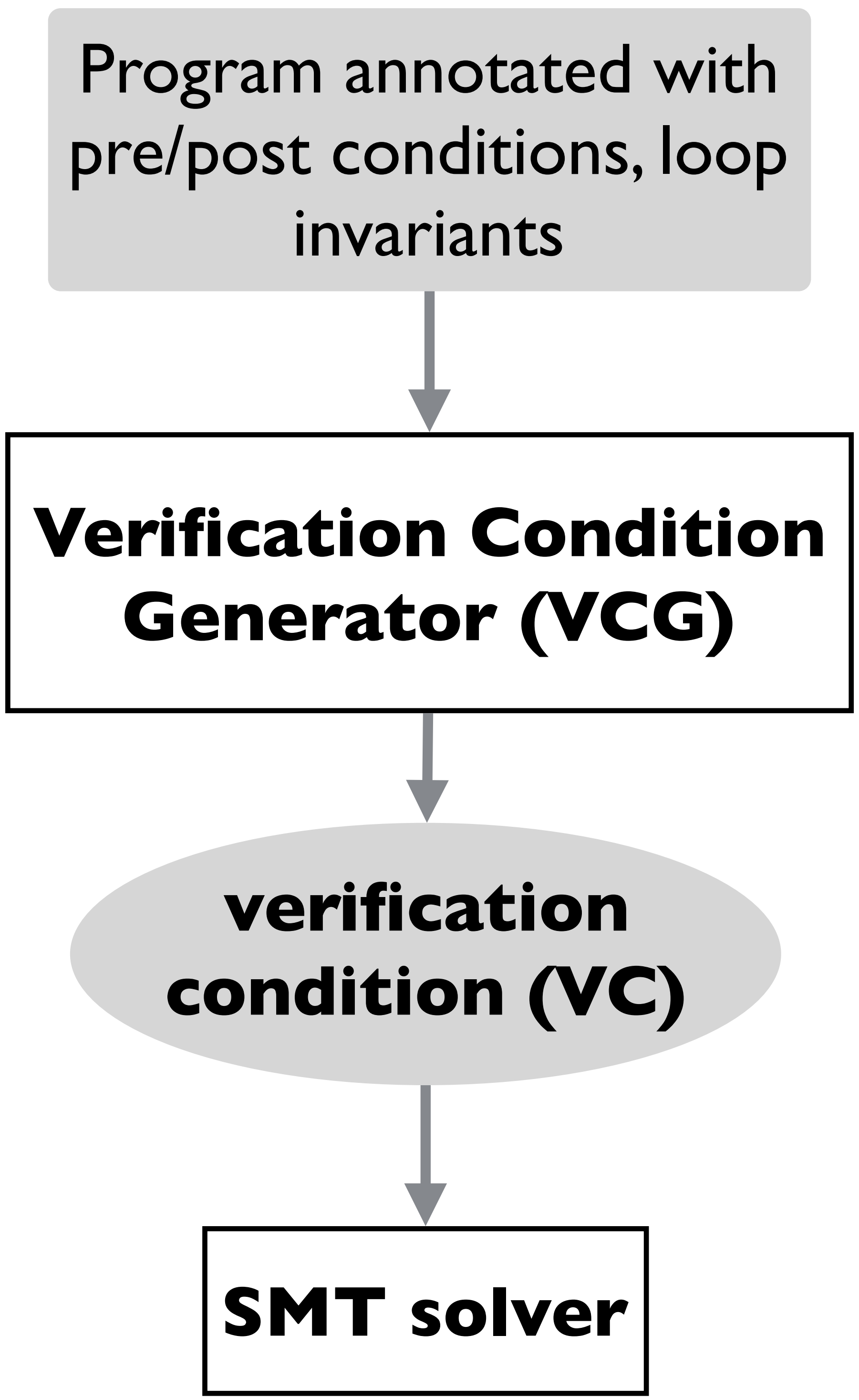
Hoare Logic proofs are highly manual:

- When to apply the rule of consequence?
- What loop invariants to use?

We can automate much of the proof process with verification condition generation!

- But loop invariants still need to be provided …

Automating Hoare logic with VC generation



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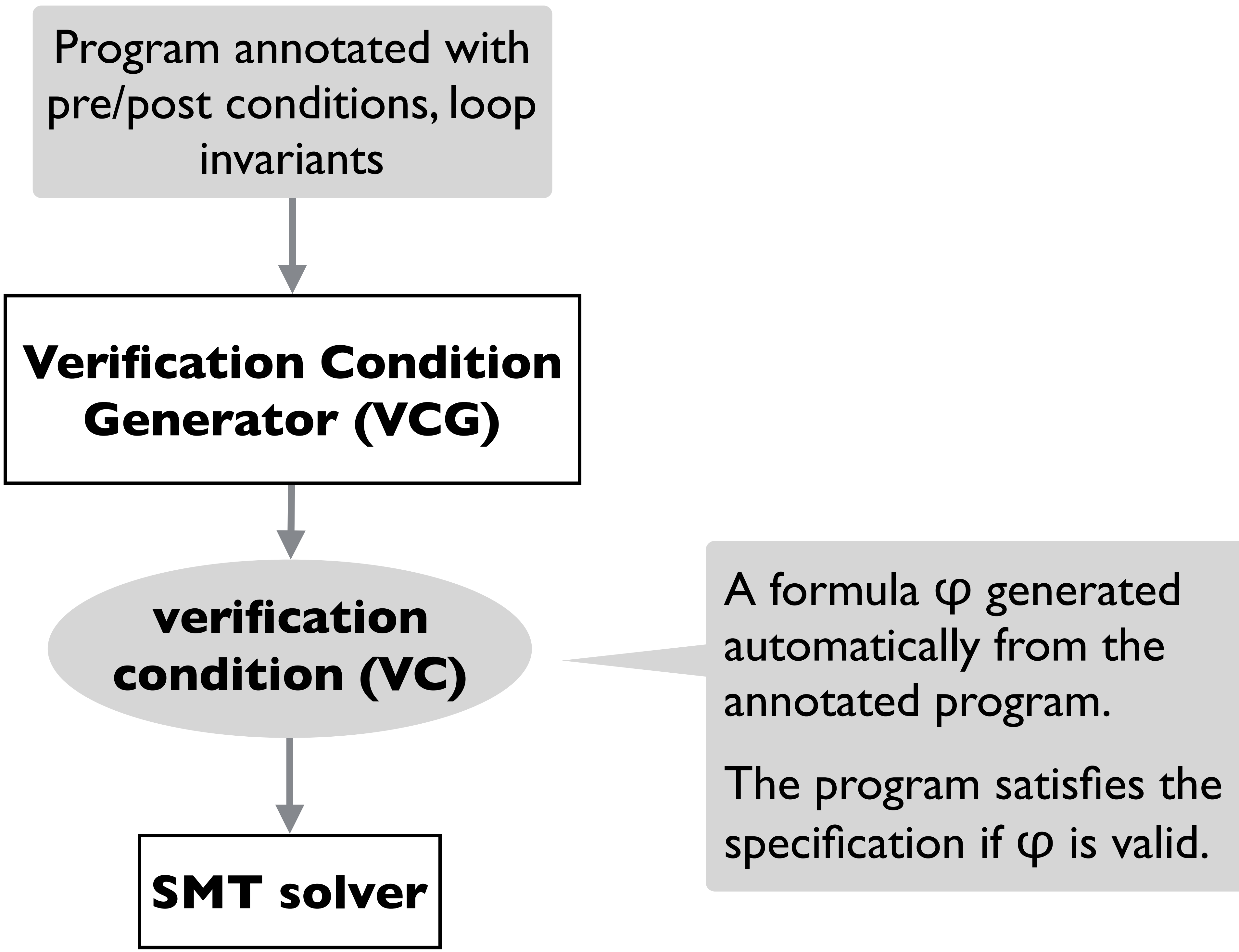

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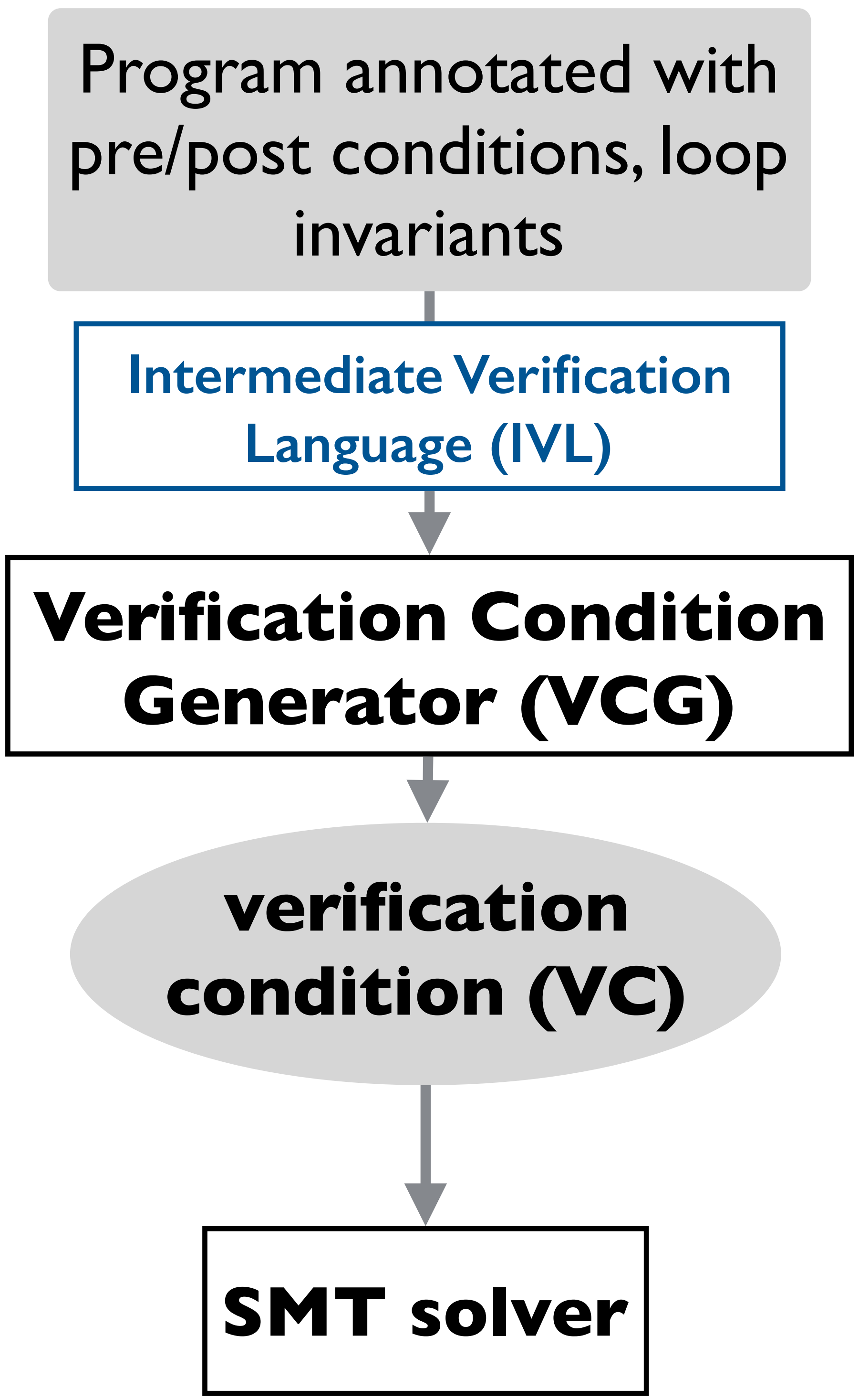

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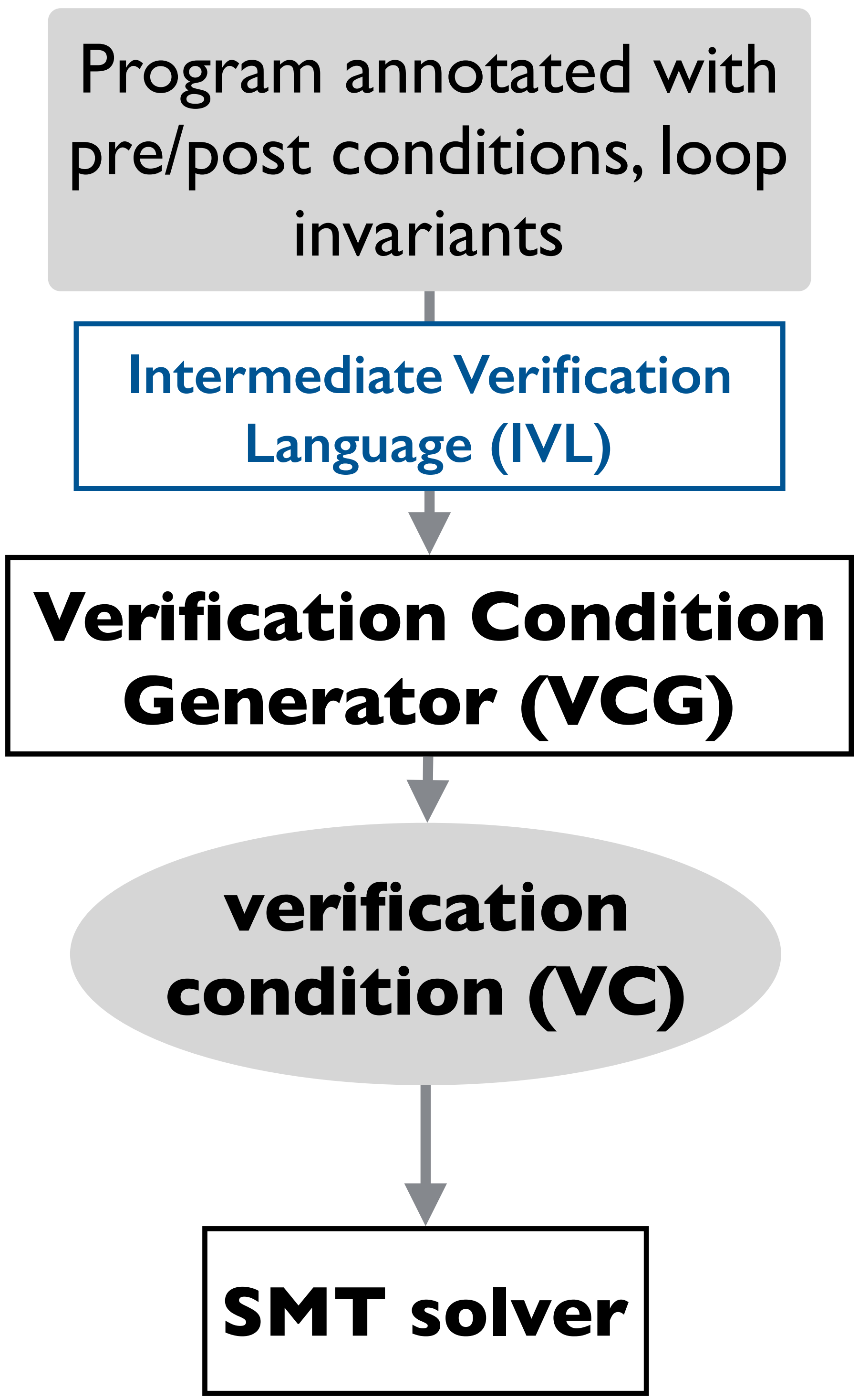

Forwards computation:

- Starting from the precondition, generate formulas to prove the postcondition.
- Based on computing ***strongest postconditions (sp).***

Backwards computation:

- Starting from the postcondition, generate formulas to prove the precondition.
- Based on computing ***weakest liberal preconditions (wp).***

VC generation with WP and SP

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sp(S, P)

- The strongest predicate that holds for states produced by executing S on a state satisfying P.

Symbolic execution, covered in next lecture, computes SPs for finite programs (no unbounded loops).

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wp(S, Q)

- The weakest predicate that guarantees Q will hold for states produced by executing S on a state satisfying that predicate.

Today, we'll see how to compute weakest liberal preconditions (WP) for IMP.

VC generation with WP and SP

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- The strongest predicate that holds for states produced by executing S on a state satisfying P.

wp(S, Q)

- The weakest predicate that guarantees Q will hold for states produced by executing S on a state satisfying that predicate.

{P} S {Q} is valid if

- $P \Rightarrow wp(S, Q)$ or
- $sp(S, P) \Rightarrow Q$

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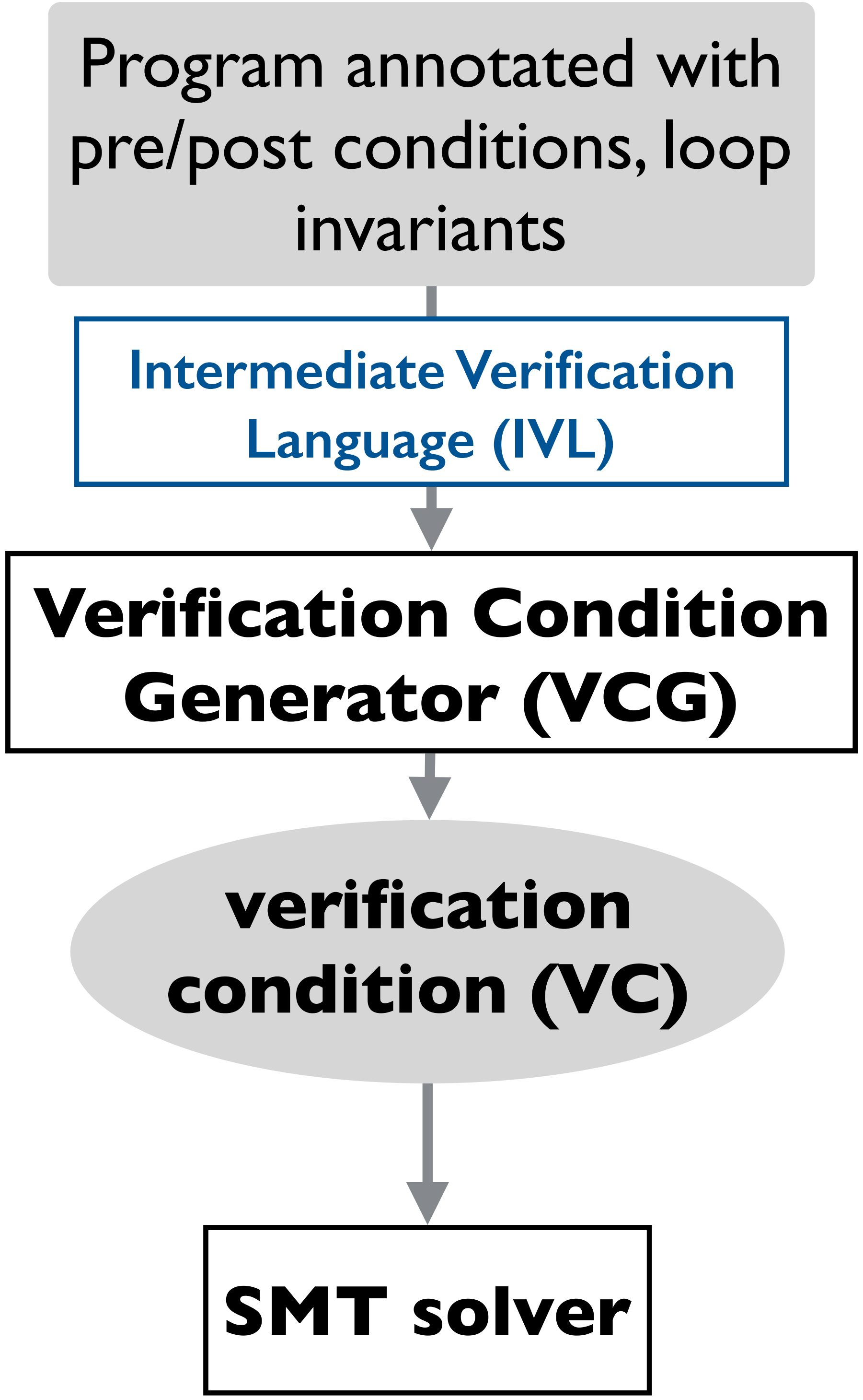
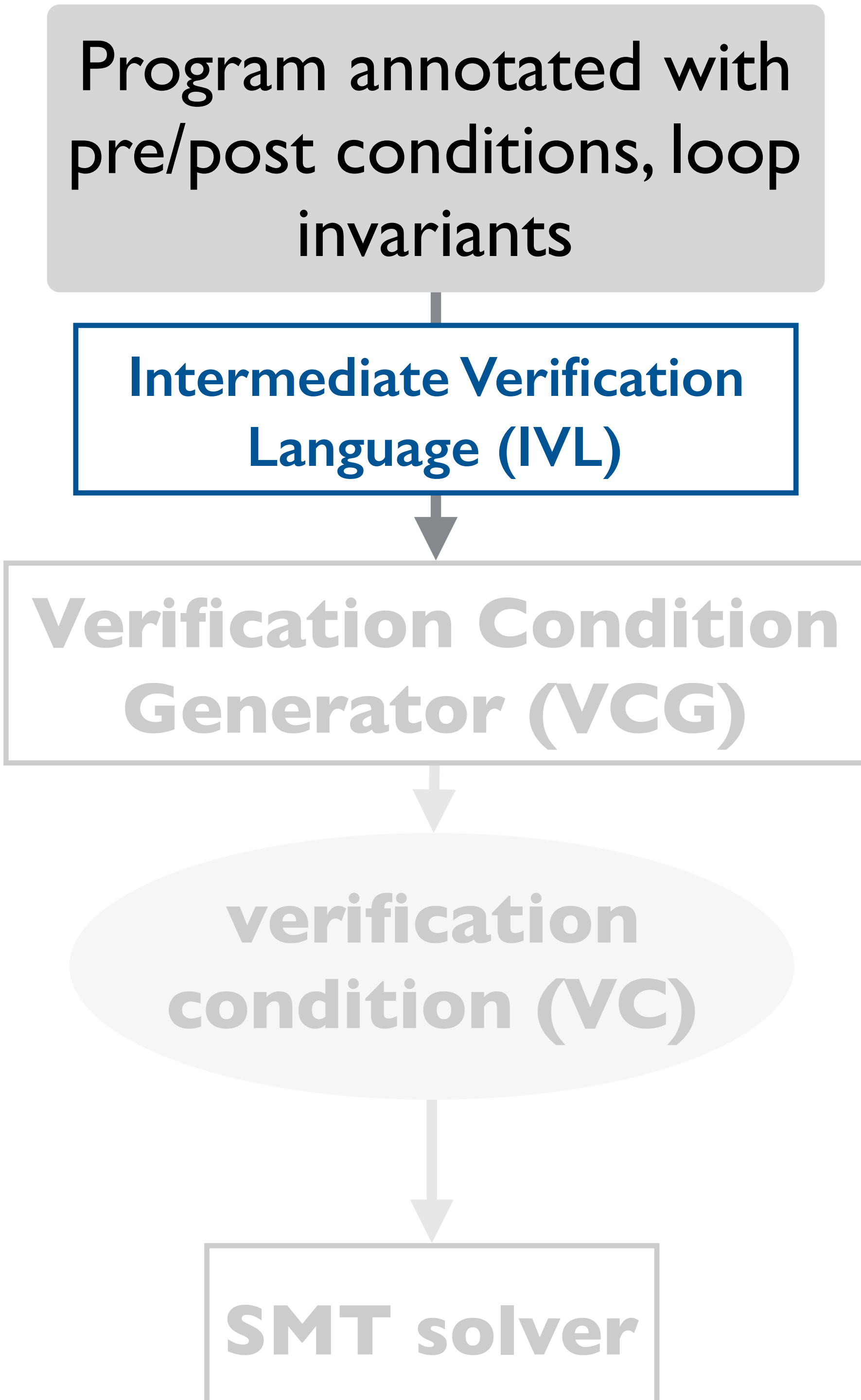
Today, we'll see how to compute weakest liberal preconditions (WP) for IMP.

This lets us verify partial correctness properties.

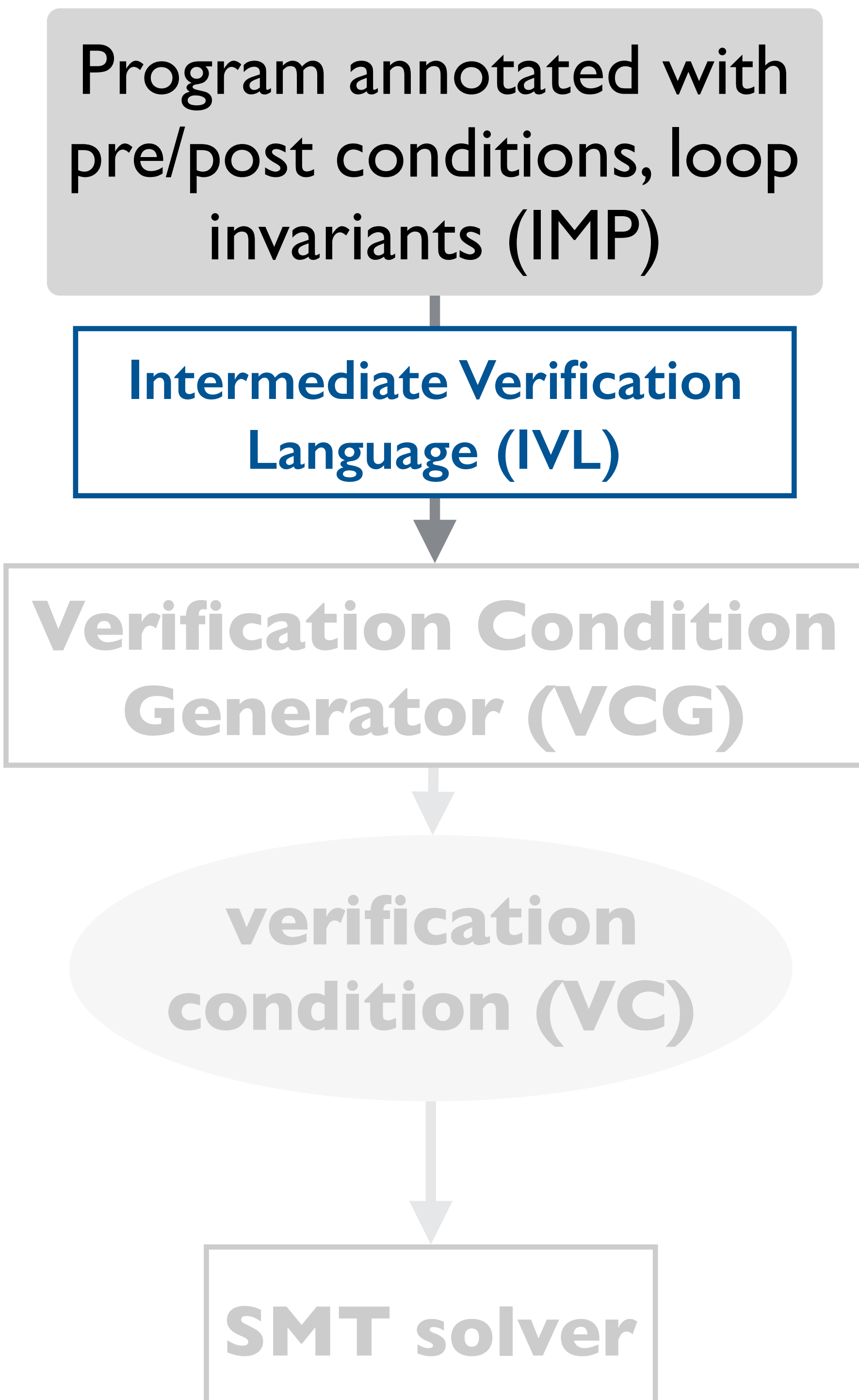
VC generation with WP



VC generation with WP



VC generation with WP: from IMP to IVL



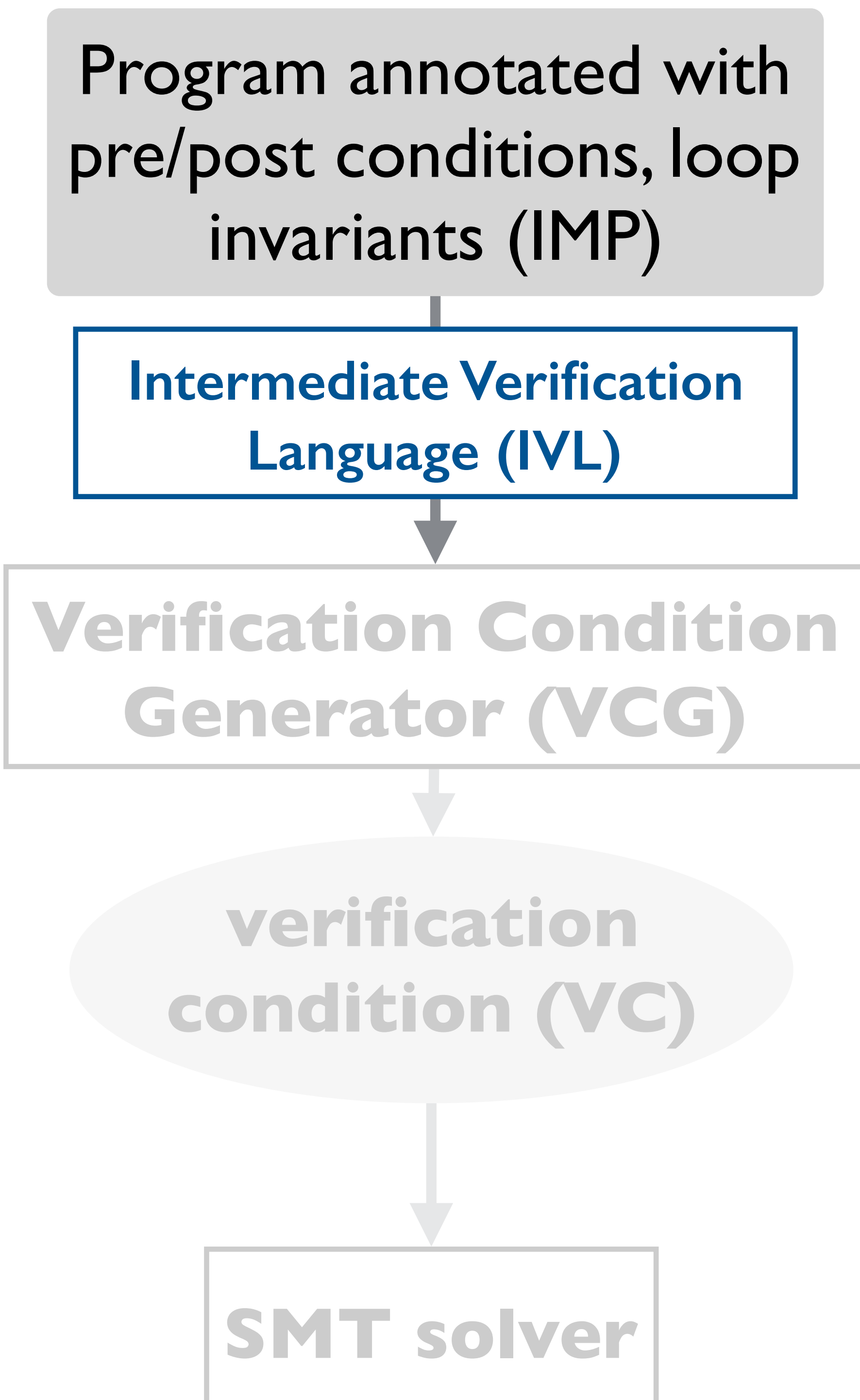
{P} S {Q}

E ::= Z | V | E + E | E * E

C ::= true | false | E = E | E ≤ E

S ::= **skip** | **abort** | V := E | S; S |
if C **then** S **else** S |
while C {I} **do** S

VC generation with WP: from IMP to IVL



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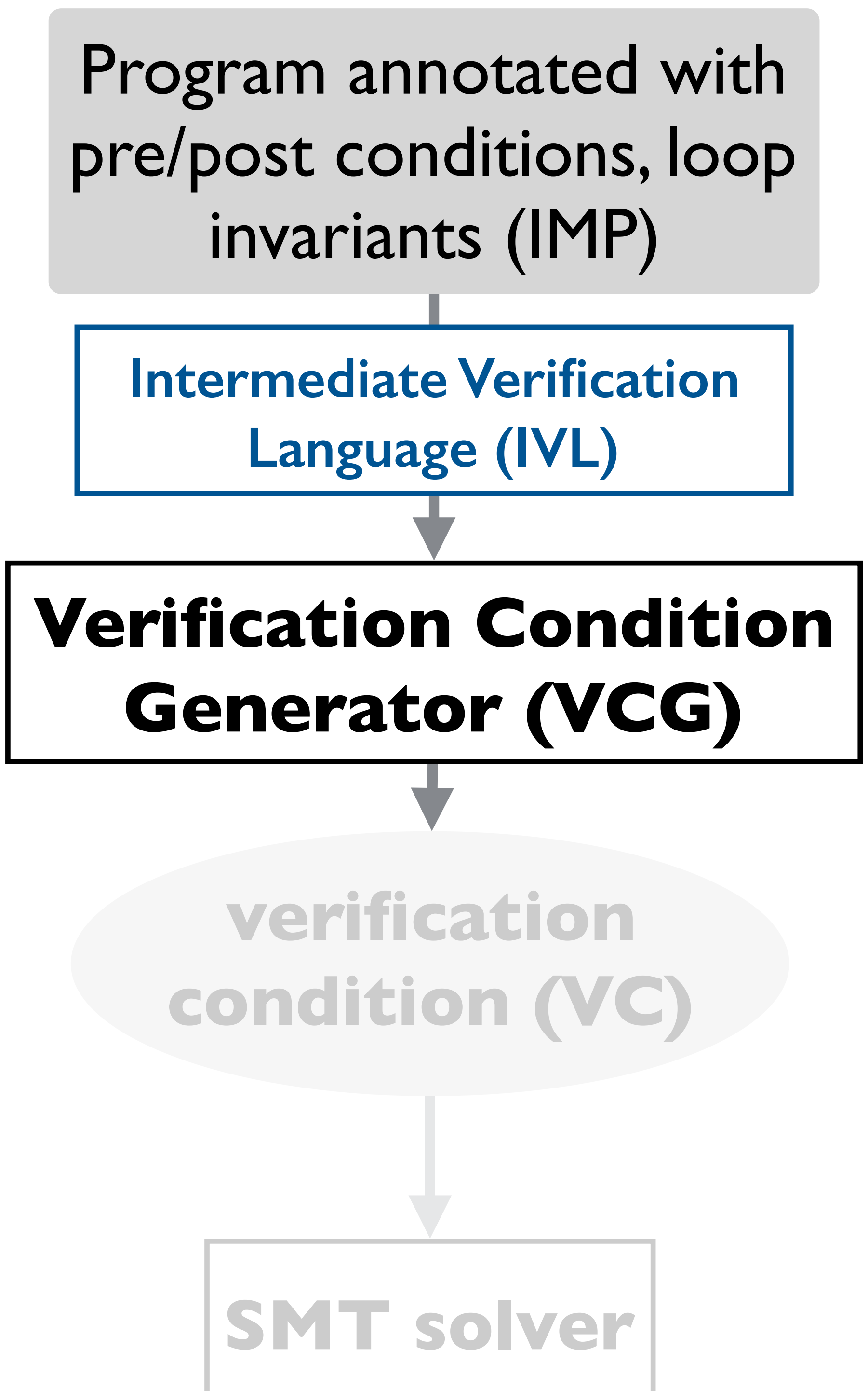
if C **then** S **else** S |
assert C | **assume** C | **havoc** V

VC generation with WP: loop-free code



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VC generation with WP: loop-free code

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- wp(x := E, Q) = Q[E / x]

VC generation with WP: loop-free code

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- wp(x := E, Q) = Q[E / x]
- wp(S_1; S_2, Q) = wp(S_1, wp(S_2, Q))

VC generation with WP: loop-free code

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- wp(x := E, Q) = Q[E / x]
- wp(S_1; S_2, Q) = wp(S_1, wp(S_2, Q))
- wp(**if** C **then** S_1 **else** S_2, Q) = $(C \rightarrow wp(S_1, Q)) \wedge (\neg C \rightarrow wp(S_2, Q))$

VC generation with WP: what about loops?

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- wp(**abort**, Q) = true
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- wp(**if** C **then** S_1 **else** S_2, Q) = (C $\rightarrow$ wp(S_1, Q)) $\wedge$ ($\neg$C $\rightarrow$ wp(S_2, Q))
- wp(**while** C {I} **do** S, Q) = ?

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A fixpoint! In general, cannot be expressed as a syntactic construction in terms of the postcondition.

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Use loop invariants to approximate loop behavior. Then check each invariant is correct and strong enough.

VC generation with WP: what about loops?

while C {I} **do** S

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VC generation with WP: what about loops?

while C {I} **do** S

Cut the loop.

assert I;
havoc x; … *// for each loop target x*
assume I;
if C **then** S; **assert** I; **assume** false;
    else skip;

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wp(S, Q):

- wp(**assert** C, Q) = C ∧ Q
- wp(**assume** C, Q) = C → Q
- wp(**havoc** x, Q) = ∀ x . Q

VC generation with WP: putting it all together

1. Translate IMP to IVL by cutting loops.
2. Compute WP for IVL.

wp(S, Q):

- wp(**skip**, Q) = Q
- wp(**abort**, Q) = true
- wp(x := E, Q) = Q[E / x]
- wp(S_1; S_2, Q) = wp(S_1, wp(S_2, Q))
- wp(**if** C **then** S_1 **else** S_2, Q) =
  (C $\rightarrow$ wp(S_1, Q)) $\wedge$ ($\neg$C $\rightarrow$ wp(S_2, Q))
- wp(**assert** C, Q) = C $\wedge$ Q
- wp(**assume** C, Q) = C $\rightarrow$ Q
- wp(**havoc** x, Q) = $\forall$ x . Q

Verifying a Hoare triple

Theorem: {P} S {Q} is valid if the following formula is valid

$$P \rightarrow wp(S_{IVL}, Q)$$

Verifying a Hoare triple

Theorem: {P} S {Q} is valid if the following formula is valid

$$P \rightarrow wp(S_{IVL}, Q)$$

The other direction doesn't hold because loop invariants may not be strong enough or they may be incorrect. Might get false alarms.

Summary

Today

- Automating Hoare Logic with VCG based on WPs

Next lecture

- Symbolic execution