Computer-Aided Reasoning for Software

CSE507

Practical Applications of SAT

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Today

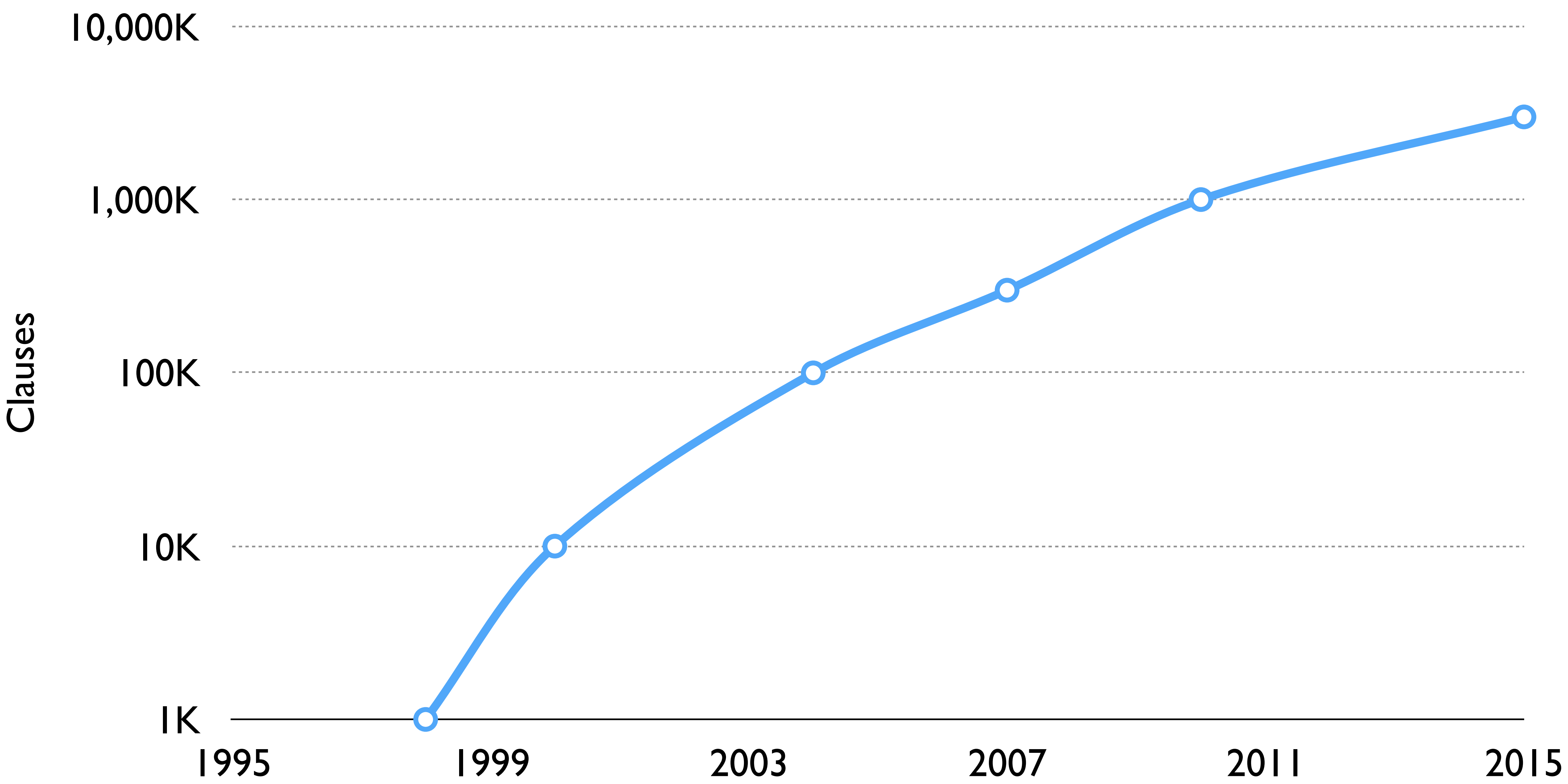
Past 2 lectures

- The theory and mechanics of SAT solving

Today

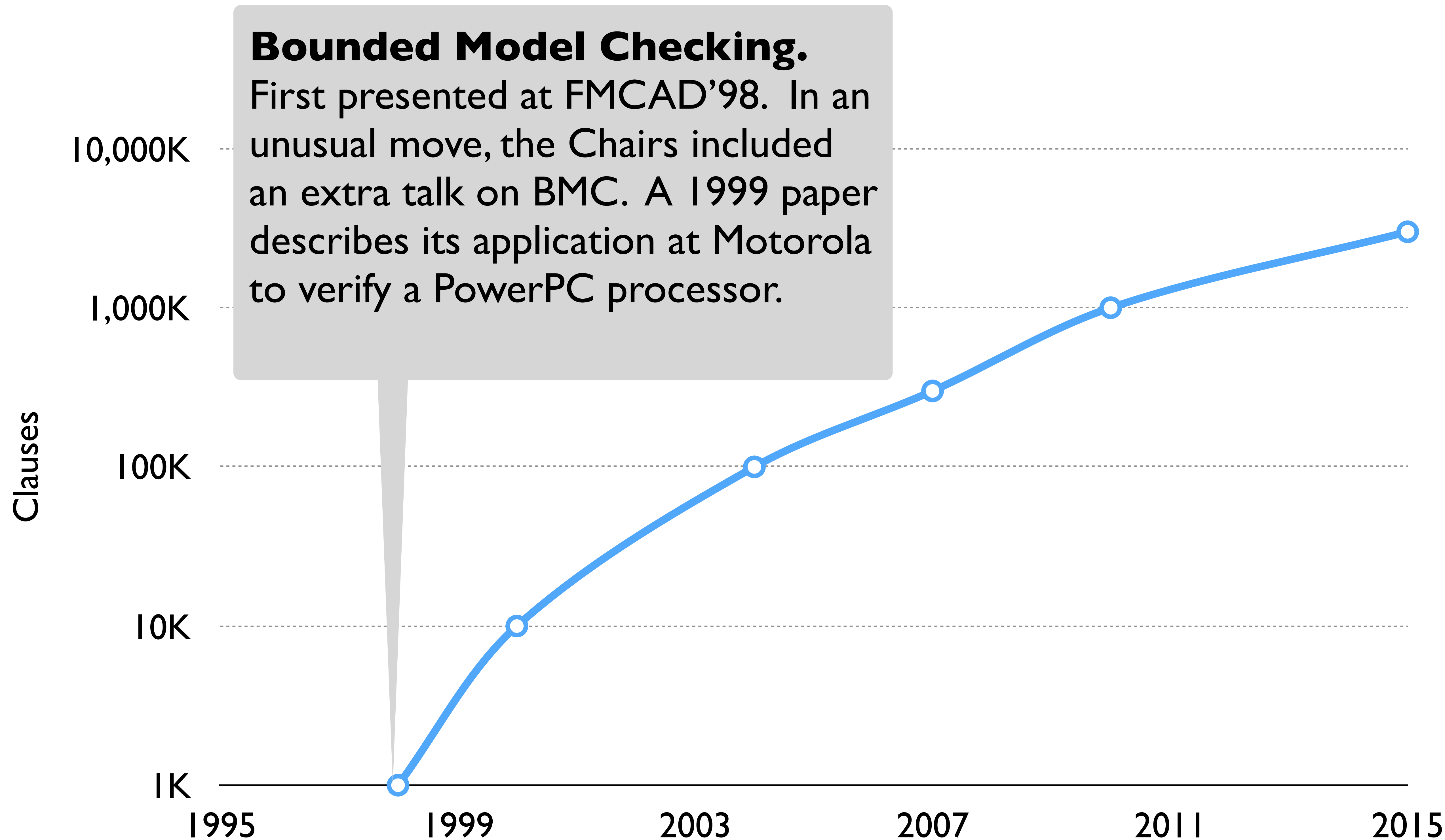
- Practical applications of SAT
- Variants of the SAT problem
- Motivating the next lecture on SMT

A brief history of SAT solving and applications



Based on a slide from Vijay Ganesh

A brief history of SAT solving and applications



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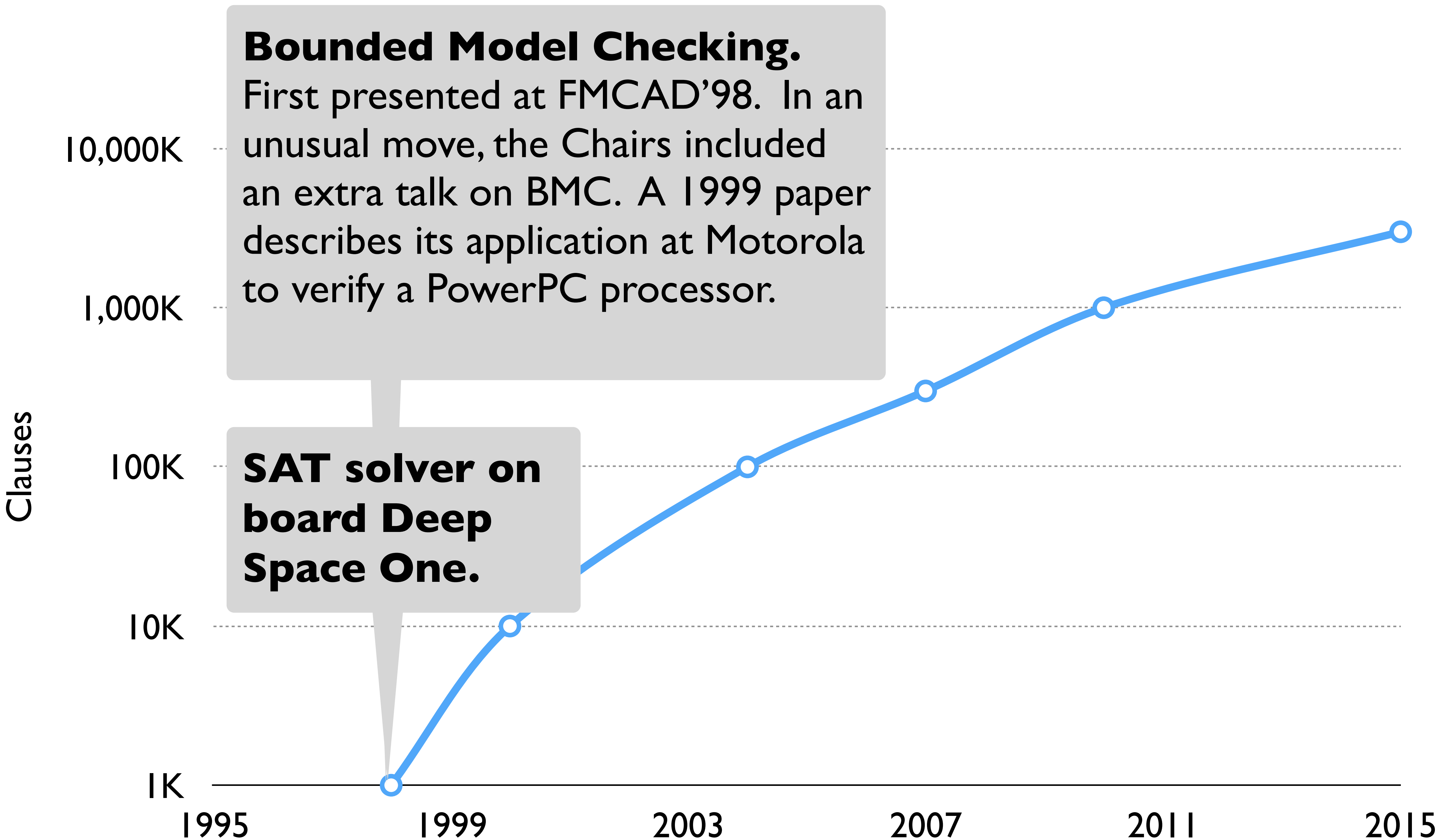

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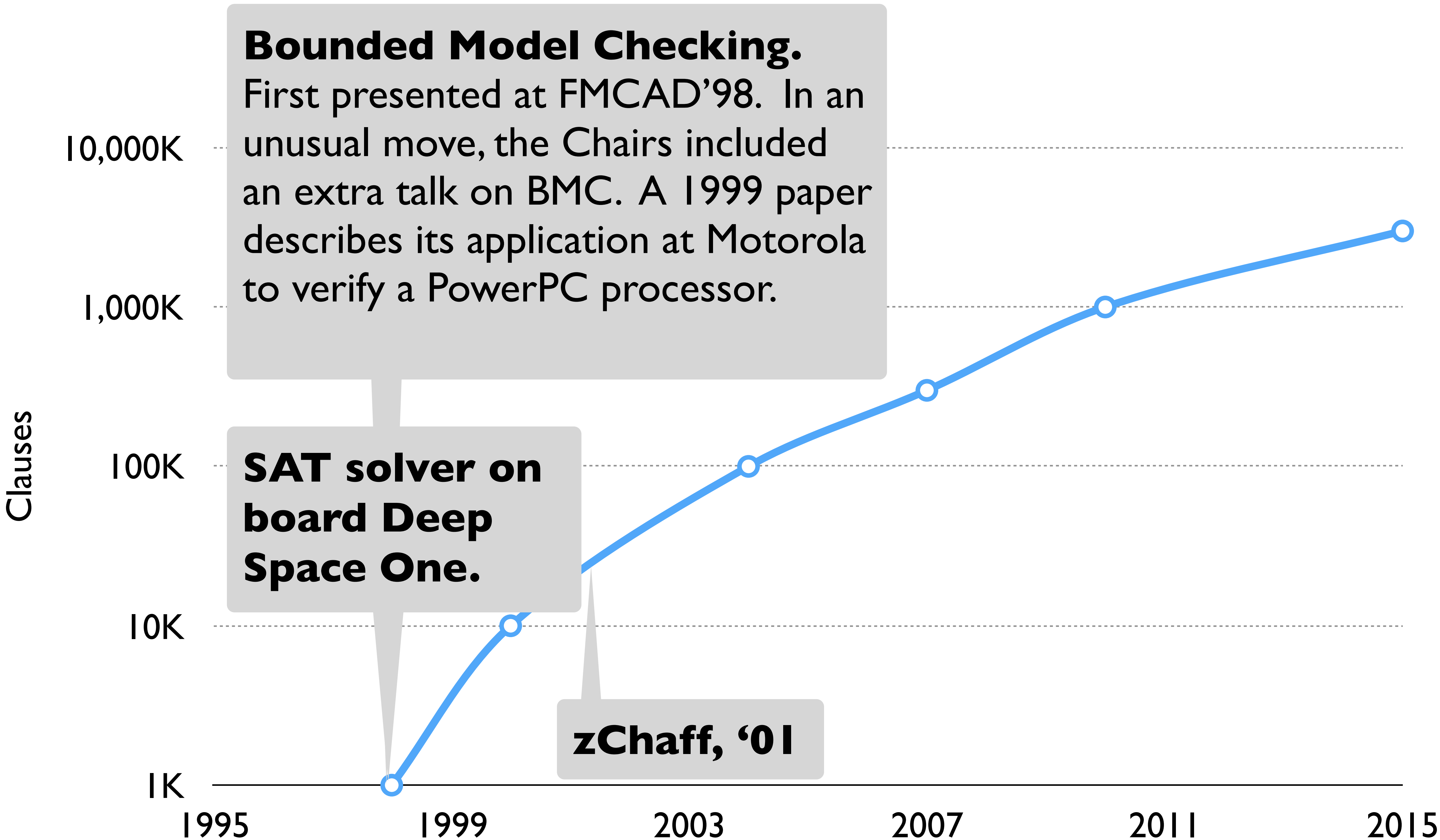

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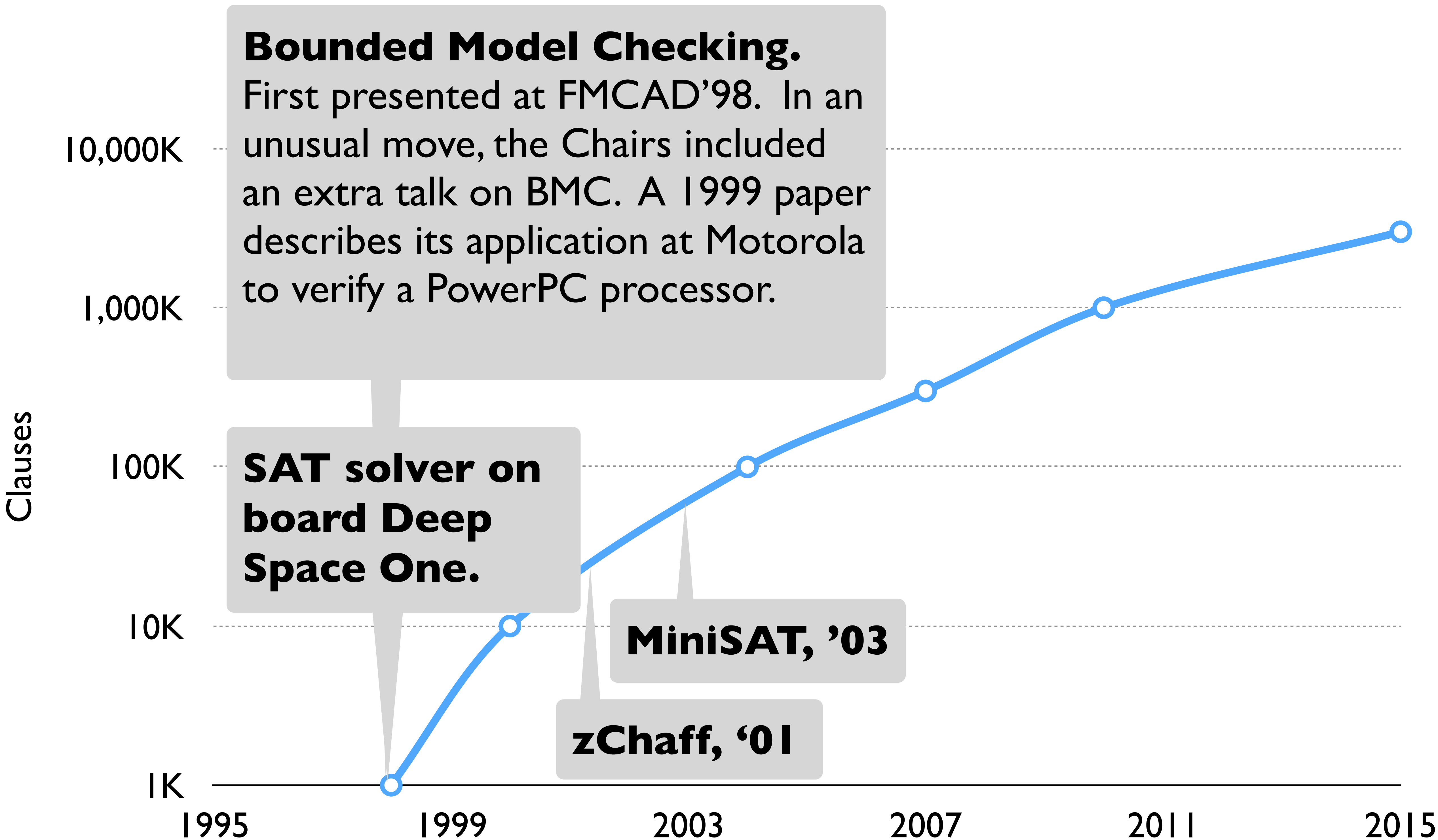

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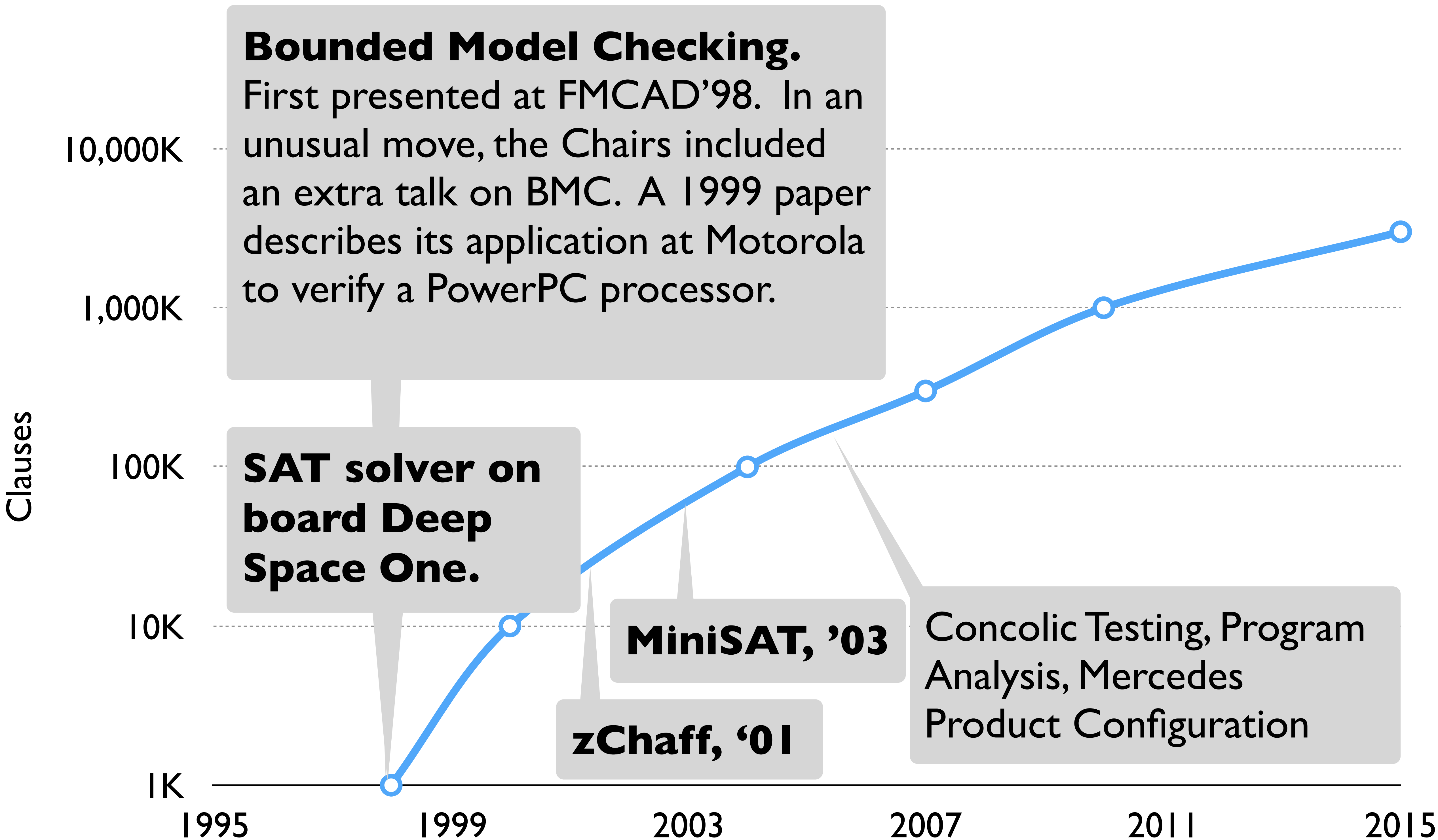

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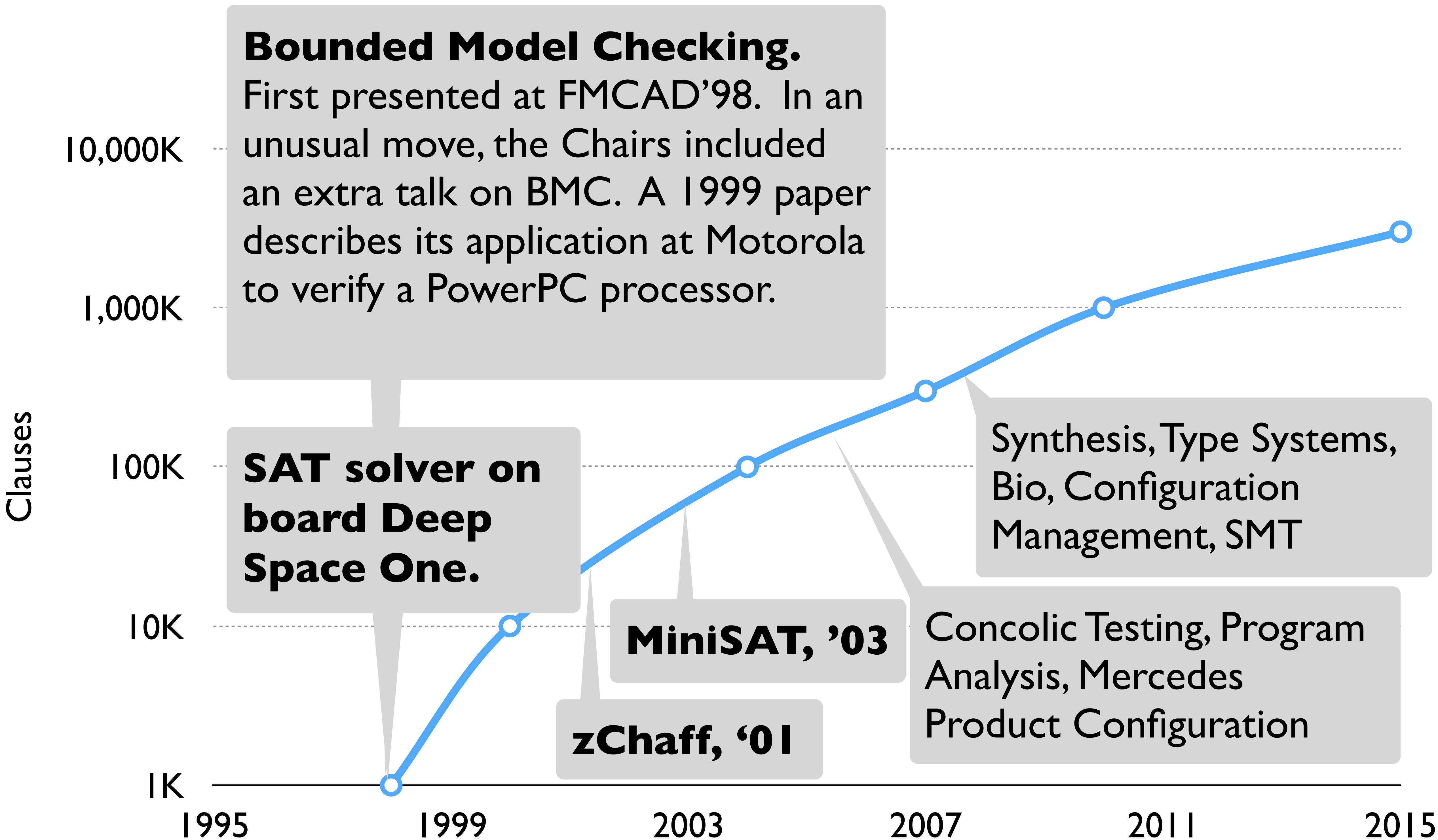
A brief history of SAT solving and applications



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today

Bounded Model Checking (BMC) & Configuration Management

Bounded Model Checking (in general)

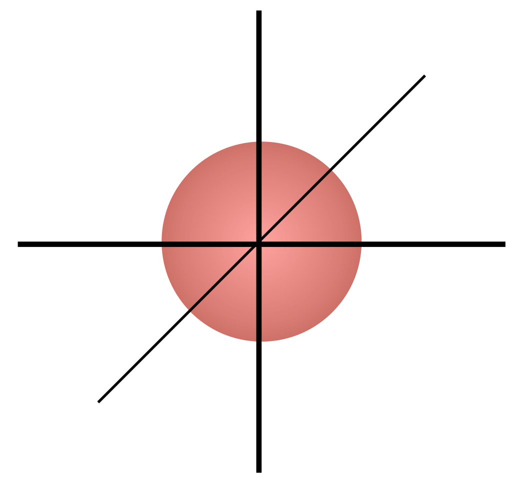
Given a system and a property, BMC checks if the property is satisfied by all executions of the system with ≤k steps, on all inputs of size ≤n.

Bounded Model Checking (in general)

Given a system and a property, BMC checks if the property is satisfied by all executions of the system with ≤k steps, on all inputs of size ≤n.

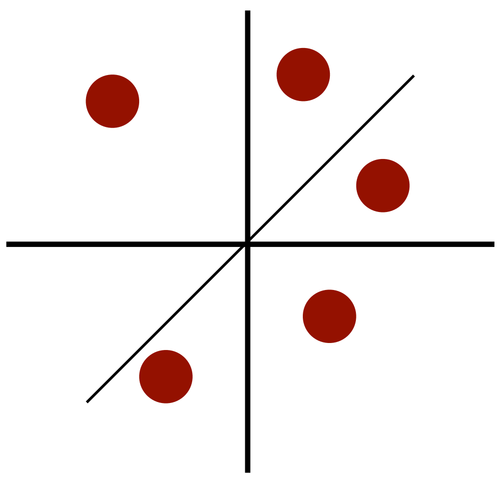
We will focus on **safety properties** (i.e., making sure a bad state, such as an assertion violation, is not reached).

Bounded Model Checking (in general)

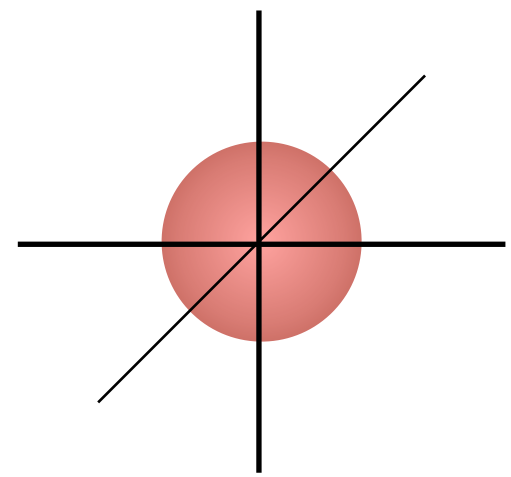


BMC: checks all executions of size ≤k

Bounded Model Checking (in general)



Testing: checks a few executions of arbitrary size

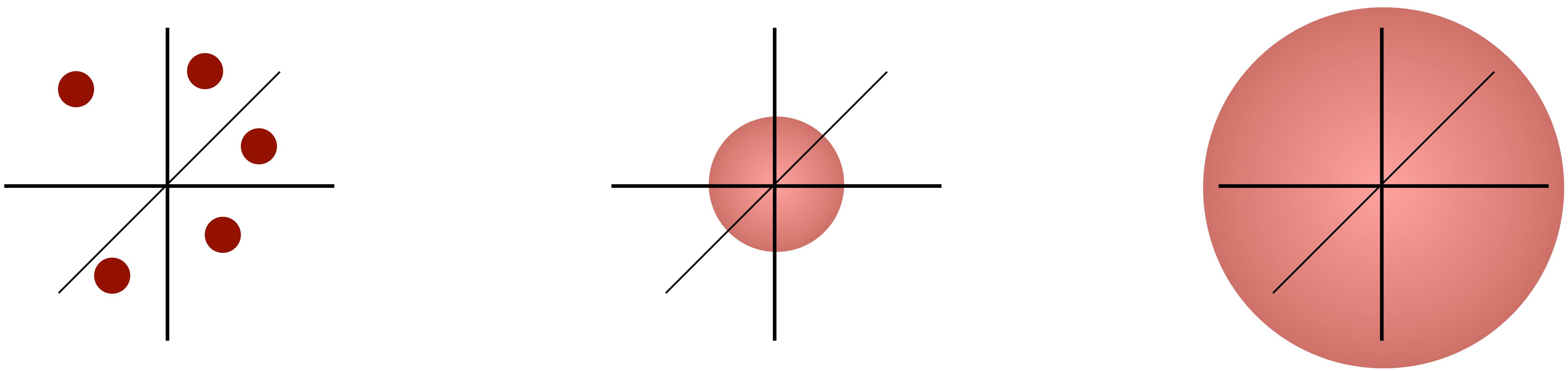


BMC: checks all executions of size $\leq k$

low confidence — high confidence

low human labor — high human labor

Bounded Model Checking (in general)



Testing: checks a few executions of arbitrary size

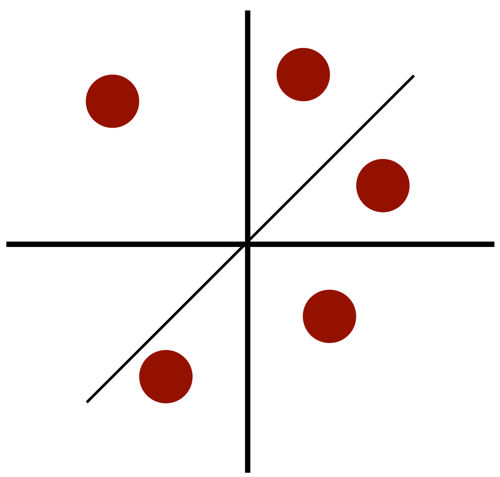
BMC: checks all executions of size ≤k

Verification: checks all executions of every size

low confidence — high confidence

low human labor — high human labor

Bounded Model Checking (in general)



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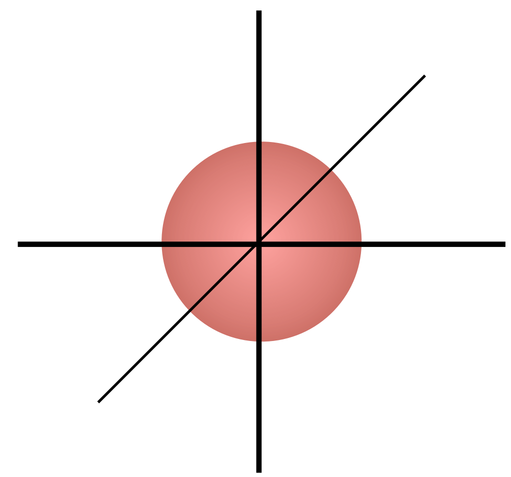
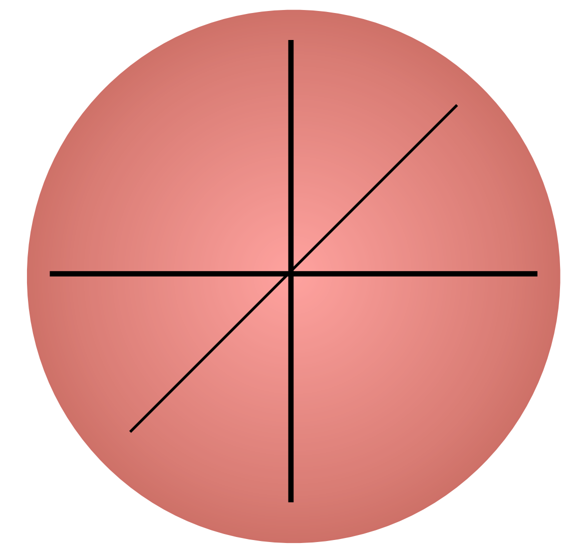


BMC: checks all executions of size ≤k



Verification: checks all executions of every size

low confidence

low human labor

The **small scope hypothesis**: most bugs can be triggered with small inputs and executions.

high confidence

high human labor

BMC by example

BMC by example

```
int daysToYear(int days) {
  int year = 1980;
  while (days > 365) {

    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }

  }
  return year;
}
```

The Zune Bug: on December 31, 2008, all first generation Zune players from Microsoft became unresponsive because of this code. What's wrong?

BMC by example

```
int daysToYear(int days) {
  int year = 1980;
  while (days > 365) {

    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
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      year += 1;
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  }
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}
```

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Infinite loop triggered on the last day of every leap year.

BMC by example

```
int daysToYear(int days) {
  int year = 1980;
  while (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
  }
  return year;
}
```

The Zune Bug: on December 31, 2008, all first generation Zune players from Microsoft became unresponsive because of this code. What's wrong?

Infinite loop triggered on the last day of every leap year.

A desired safety property: the value of the days variable decreases in every loop iteration.

BMC step 1 of 4: finitize loops

```
int daysToYear(int days) {
  int year = 1980;
  while (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
  }
  return year;
}
```

BMC step 1 of 4: finitize loops

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

- Unwind all loops k times (e.g., k=1), and add an **unwinding assertion** at the end.

BMC step 1 of 4: finitize loops

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

- Unwind all loops k times (e.g., k=1), and add an **unwinding assertion** at the end.
- If a CEX violates a program assertion, we have found a buggy behavior of length ≤k.

BMC step 1 of 4: finitize loops

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

- Unwind all loops k times (e.g., k=1), and add an **unwinding assertion** at the end.
- If a CEX violates a program assertion, we have found a buggy behavior of length ≤k.
- If a CEX violates an unwinding assertion, the program has no buggy behavior of length ≤k, but it may have a longer one.

BMC step 1 of 4: finitize loops

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

- Unwind all loops k times (e.g., k=1), and add an **unwinding assertion** at the end.
- If a CEX violates a program assertion, we have found a buggy behavior of length ≤k.
- If a CEX violates an unwinding assertion, the program has no buggy behavior of length ≤k, but it may have a longer one.
- If there is no CEX, the program is correct for all k!

BMC step 1 of 4: finitize loops & inline calls

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

Assume call to isLeapYear is inlined (replaced with the procedure body). We'll keep it for readability.

BMC step 2 of 4: eliminate side effects

```
int daysToYear(int days) {
  int year = 1980;
  if (days > 365) {
    int oldDays = days;
    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }
    assert days < oldDays;
    assert days <= 365;
  }
  return year;
}
```

BMC step 2 of 4: eliminate side effects

```
int days;
int year = 1980;
if (days > 365) {
  int oldDays = days;
  if (isLeapYear(year)) {
    if (days > 366) {
      days = days − 366;
      year = year + 1;
    }
  } else {
    days = days − 365;
    year = year + 1;
  }
  assert days < oldDays;
  assert days <= 365;
}
return year;
```

Convert to **Static Single Assignment** (SSA) form:

BMC step 2 of 4: eliminate side effects

```
int days0;
int year0 = 1980;
if (days0 > 365) {
  int oldDays0 = days0;
  if (isLeapYear(year0)) {
    if (days0 > 366) {
      days1 = days0 - 366;
      year1 = year0 + 1;
    }
  } else {
    days3 = days0 - 365;
    year3 = year0 + 1;
  }
  assert days4 < oldDays0;
  assert days4 <= 365;
}
return year5;
```

Convert to **Static Single Assignment** (SSA) form:

- Replace each assignment to a variable v with a definition of a fresh variable v_i.
- Change uses of variables so that they refer to the correct definition (version).

BMC step 2 of 4: eliminate side effects

```
int days₀;
int year₀ = 1980;
boolean g₀ = (days₀ > 365);
int oldDays₀ = days₀;
boolean g₁ = isLeapYear(year₀);
boolean g₂ = days₀ > 366;
days₁ = days₀ - 366;
year₁ = year₀ + 1;
days₂ = φ(g₁ && g₂, days₁, days₀);
year₂ = φ(g₁ && g₂, year₁, year₀);
days₃ = days₀ - 365;
year₃ = year₀ + 1;
days₄ = φ(g₁, days₂, days₃);
year₄ = φ(g₁, year₂, year₃);
assert days₄ < oldDays₀;
assert days₄ <= 365;
year₅ = φ(g₀, year₄, year₀);
return year₅;
```

Convert to **Static Single Assignment** (SSA) form:

- Replace each assignment to a variable v with a definition of a fresh variable v_i.
- Change uses of variables so that they refer to the correct definition (version).
- Make conditional dependences explicit with gated φ nodes.

BMC step 2 of 4: eliminate side effects

```
int days_0;
int year_0 = 1980;
boolean g_0 = (days_0 > 365);
int oldDays_0 = days_0;
boolean g_1 = isLeapYear(year_0);
boolean g_2 = days_0 > 366;
days_1 = days_0 - 366;
year_1 = year_0 + 1;
days_2 = φ(g_1 && g_2, days_1, days_0);
year_2 = φ(g_1 && g_2, year_1, year_0);
days_3 = days_0 - 365;
year_3 = year_0 + 1;
days_4 = φ(g_1, days_2, days_3);
year_4 = φ(g_1, year_2, year_3);
assert days_4 < oldDays_0;
assert days_4 <= 365;
year_5 = φ(g_0, year_4, year_0);
return year_5;
```

```
int days_0;
int year_0 = 1980;
if (days_0 > 365) {
  int oldDays_0 = days_0;
  if (isLeapYear(year_0)) {
    if (days_0 > 366) {
      days_1 = days_0 - 366;
      year_1 = year_0 + 1;
    }
  } else {
    days_3 = days_0 - 365;
    year_3 = year_0 + 1;
  }
  assert days_4 < oldDays_0;
  assert days_4 <= 365;
}
return year_5;
```

BMC step 3 of 4: convert into equations

```
int days0;
int year0 = 1980;
boolean g0 = (days0 > 365);
int oldDays0 = days0;
boolean g1 = isLeapYear(year0);
boolean g2 = days0 > 366;
days1 = days0 - 366;
year1 = year0 + 1;
days2 = φ(g1 && g2, days1, days0);
year2 = φ(g1 && g2, year1, year0);
days3 = days0 - 365;
year3 = year0 + 1;
days4 = φ(g1, days2, days3);
year4 = φ(g1, year2, year3);
assert days4 < oldDays0;
assert days4 <= 365;
year5 = φ(g0, year4, year0);
return year5;
```

We can now read off the equations that encode the program semantics, and the assertions to be checked.

BMC step 3 of 4: convert into equations

```
int year0 = 1980;
boolean g0 = (days0 > 365);
int oldDays0 = days0;
boolean g1 = isLeapYear(year0);
boolean g2 = days0 > 366;
days1 = days0 - 366;
year1 = year0 + 1;
days2 = φ(g1 && g2, days1, days0);
year2 = φ(g1 && g2, year1, year0);
days3 = days0 - 365;
year3 = year0 + 1;
days4 = φ(g1, days2, days3);
year4 = φ(g1, year2, year3);
assert days4 < oldDays0;
assert days4 <= 365;
```

We can now read off the equations that encode the program semantics …

BMC step 3 of 4: convert into equations

```
year₀ = 1980;
g₀ = (days₀ > 365);
oldDays₀ = days₀;
g₁ = isLeapYear(year₀);
g₂ = days₀ > 366;
days₁ = days₀ − 366;
year₁ = year₀ + 1;
days₂ = φ(g₁ && g₂, days₁, days₀);
year₂ = φ(g₁ && g₂, year₁, year₀);
days₃ = days₀ − 365;
year₃ = year₀ + 1;
days₄ = φ(g₁, days₂, days₃);
year₄ = φ(g₁, year₂, year₃);
assert days₄ < oldDays₀;
assert days₄ <= 365;
```

We can now read off the equations that encode the program semantics …

BMC step 3 of 4: convert into equations

```
year₀ = 1980 ∧
g₀ = (days₀ > 365)∧
oldDays₀ = days₀ ∧
g₁ = isLeapYear(year₀)∧
g₂ = days₀ > 366 ∧
days₁ = days₀ − 366 ∧
year₁ = year₀ + 1 ∧
days₂ = φ(g₁ ∧ g₂, days₁, days₀) ∧
year₂ = φ(g₁ ∧ g₂, year₁, year₀) ∧
days₃ = days₀ − 365 ∧
year₃ = year₀ + 1 ∧
days₄ = φ(g₁, days₂, days₃) ∧
year₄ = φ(g₁, year₂, year₃) ∧
assert days₄ < oldDays₀;
assert days₄ <= 365;
```

We can now read off the equations that encode the program semantics …

BMC step 3 of 4: convert into equations

We can now read off the equations that encode the program semantics …

$$
\begin{aligned}
&\mathtt{year}_0 = 1980 \land \\
&\mathtt{g}_0 = (\mathtt{days}_0 > 365) \land \\
&\mathtt{oldDays}_0 = \mathtt{days}_0 \land \\
&\mathtt{g}_1 = \mathtt{isLeapYear}(\mathtt{year}_0) \land \\
&\mathtt{g}_2 = \mathtt{days}_0 > 366 \land \\
&\mathtt{days}_1 = \mathtt{days}_0 - 366 \land \\
&\mathtt{year}_1 = \mathtt{year}_0 + 1 \land \\
&\mathtt{days}_2 = \mathtt{ite}(\mathtt{g}_1 \land \mathtt{g}_2, \mathtt{days}_1, \mathtt{days}_0) \land \\
&\mathtt{year}_2 = \mathtt{ite}(\mathtt{g}_1 \land \mathtt{g}_2, \mathtt{year}_1, \mathtt{year}_0) \land \\
&\mathtt{days}_3 = \mathtt{days}_0 - 365 \land \\
&\mathtt{year}_3 = \mathtt{year}_0 + 1 \land \\
&\mathtt{days}_4 = \mathtt{ite}(\mathtt{g}_1, \mathtt{days}_2, \mathtt{days}_3) \land \\
&\mathtt{year}_4 = \mathtt{ite}(\mathtt{g}_1, \mathtt{year}_2, \mathtt{year}_3) \land \\
&\mathbf{assert}\ \mathtt{days}_4 < \mathtt{oldDays}_0; \\
&\mathbf{assert}\ \mathtt{days}_4 <= 365;
\end{aligned}
$$

BMC step 3 of 4: convert into equations

$$
\begin{aligned}
&year_0 = 1980 \wedge \\
&g_0 = (days_0 > 365) \wedge \\
&oldDays_0 = days_0 \wedge \\
&g_1 = isLeapYear(year_0) \wedge \\
&g_2 = days_0 > 366 \wedge \\
&days_1 = days_0 - 366 \wedge \\
&year_1 = year_0 + 1 \wedge \\
&days_2 = ite(g_1 \wedge g_2, days_1, days_0) \wedge \\
&year_2 = ite(g_1 \wedge g_2, year_1, year_0) \wedge \\
&days_3 = days_0 - 365 \wedge \\
&year_3 = year_0 + 1 \wedge \\
&days_4 = ite(g_1, days_2, days_3) \wedge \\
&year_4 = ite(g_1, year_2, year_3) \wedge \\
&(\neg(days_4 < oldDays_0) \vee \\
&\ \neg(days_4 <= 365))
\end{aligned}
$$

We can now read off the equations that encode the program semantics, and the assertions to be checked.

BMC step 3 of 4: convert into equations

$$
\begin{aligned}
&year_0 = 1980 \land \\
&g_0 = (days_0 > 365) \land \\
&oldDays_0 = days_0 \land \\
&g_1 = isLeapYear(year_0) \land \\
&g_2 = days_0 > 366 \land \\
&days_1 = days_0 - 366 \land \\
&year_1 = year_0 + 1 \land \\
&days_2 = ite(g_1 \land g_2, days_1, days_0) \land \\
&year_2 = ite(g_1 \land g_2, year_1, year_0) \land \\
&days_3 = days_0 - 365 \land \\
&year_3 = year_0 + 1 \land \\
&days_4 = ite(g_1, days_2, days_3) \land \\
&year_4 = ite(g_1, year_2, year_3) \land \\
&(\neg(days_4 < oldDays_0) \lor \\
&\ \neg(days_4 <= 365))
\end{aligned}
$$

We can now read off the equations that encode the program semantics, and the assertions to be checked.

A solution to this formula is a sound *counterexample*: an interpretation for all logical variables that satisfies the program semantics (for up to k unwindings) but violates at least one of the assertions.

BMC step 4 of 4: convert into CNF

$$year_1 = year_0 + 1$$

BMC step 4 of 4: convert into CNF

```
year1 = year0 + 1
```

$year_0 = \underset{31}{0}\,\underset{30}{0}\,\underset{29}{0} \ldots \underset{2}{0}\,\underset{1}{0}\,\underset{0}{0}$

Represent numbers as arrays of bits.

BMC step 4 of 4: convert into CNF

$year_1 = year_0 + 1$

$year_0 = 000 \ldots 000$

31 30 29 … 2 1 0

Represent numbers as arrays of bits.
Use one boolean variable per bit for each number.

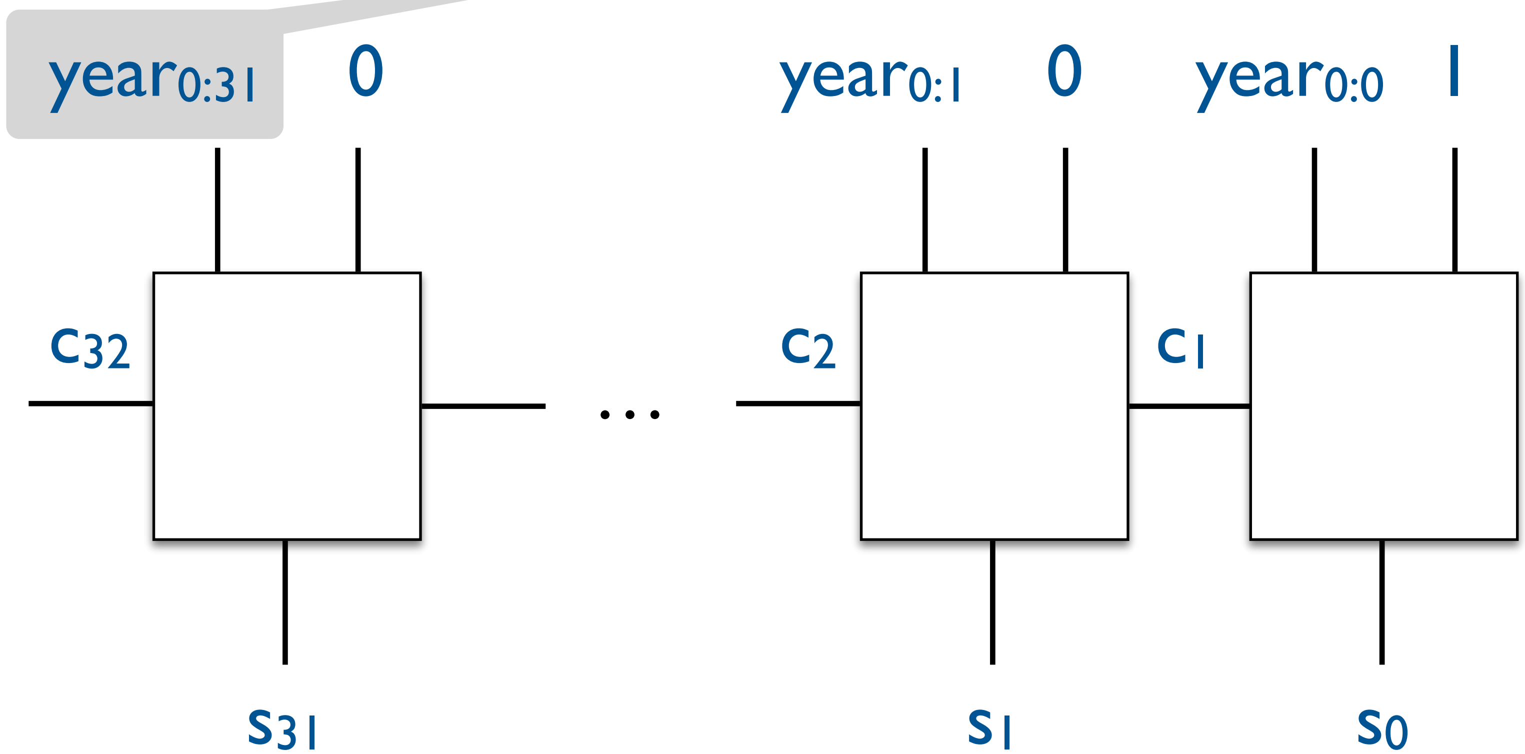
$year_{0:31}$

BMC step 4 of 4: convert into CNF

$year_1 = year_0 + 1$

$year_0 = 0_{31}0_{30}0_{29} \ldots 0_{2}0_{1}0_{0}$

Represent numbers as arrays of bits.
Use one boolean variable per bit for each number.



Construct an adder circuit for $year_0 + 1$.

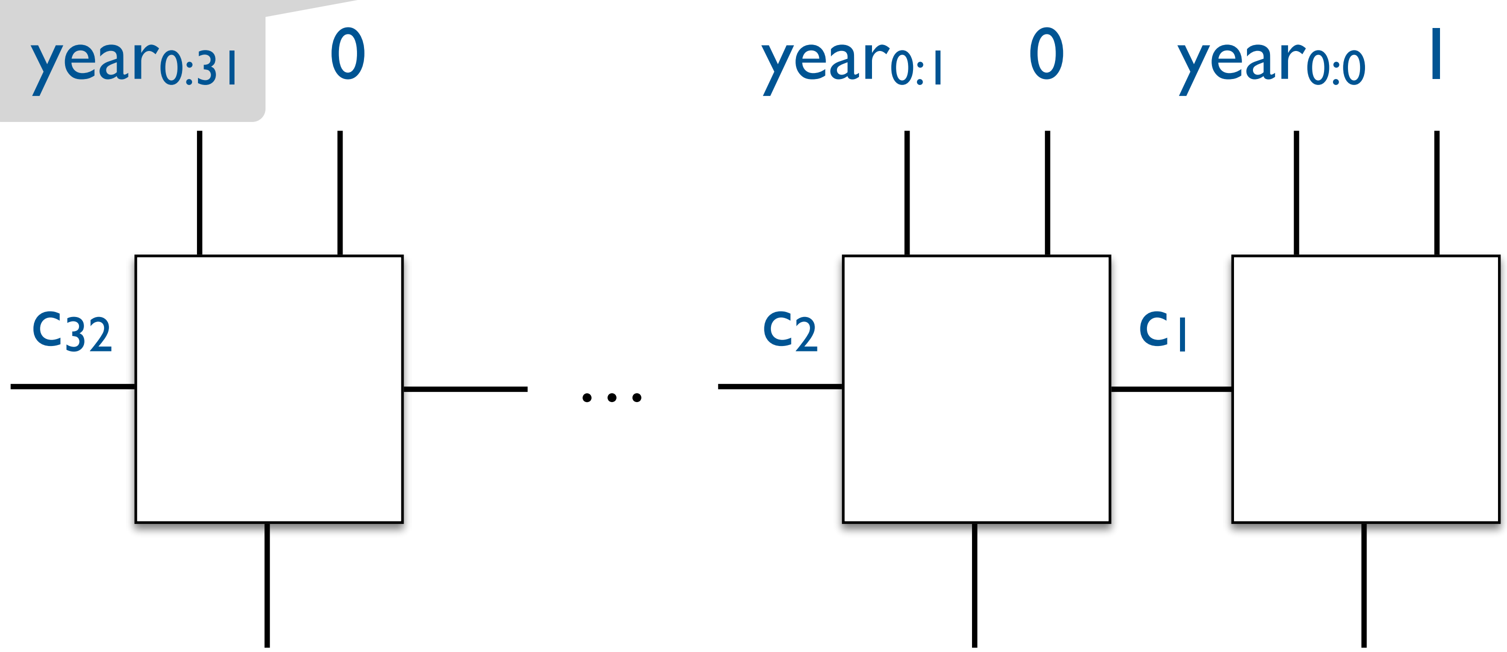
BMC step 4 of 4: convert into CNF

$year_1 = year_0 + 1$

$year_0 = 000 \ldots 000$

31 30 29 … 2 1 0

Represent numbers as arrays of bits.
Use one boolean variable per bit for each number.



Construct an adder circuit for $year_0 + 1$.

$year_{1:31} \leftrightarrow s_{31} \wedge \ldots \wedge year_{1:1} \leftrightarrow s_1 \wedge s_0 \leftrightarrow year_{1:0}$

Introduce new clauses to constrain bits in $year_1$ to match bits in the sum.

BMC counterexample for k=1

```
int daysToYear(int days) {
  int year = 1980;
  while (days > 365) {

    if (isLeapYear(year)) {
      if (days > 366) {
        days -= 366;
        year += 1;
      }
    } else {
      days -= 365;
      year += 1;
    }

  }
  return year;
}
```

days = 366

today

Bounded Model Checking (BMC) & Configuration Management

Configuration Management

Given a configuration, consisting of a set of components, their dependencies, and conflicts:

- Decide if a new component can be added to the configuration.
- Add the component while optimizing some linear function.
- If the component cannot be added, find a way to add it by removing as few conflicting components from the current configuration as possible.

Configuration Management

Given a configuration, consisting of a set of components, their dependencies, and conflicts:

- Decide if a new component can be added to the configuration.
- Add the component while optimizing some linear function.
- If the component cannot be added, find a way to add it by removing as few conflicting components from the current configuration as possible.

SAT

Configuration Management

Given a configuration, consisting of a set of components, their dependencies, and conflicts:

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SAT

Pseudo-Boolean Constraints

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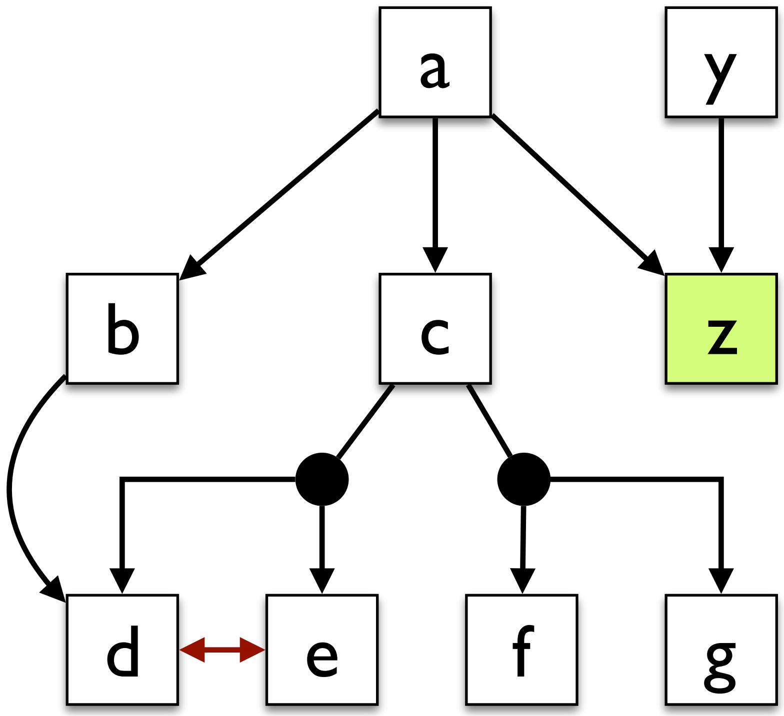
SAT

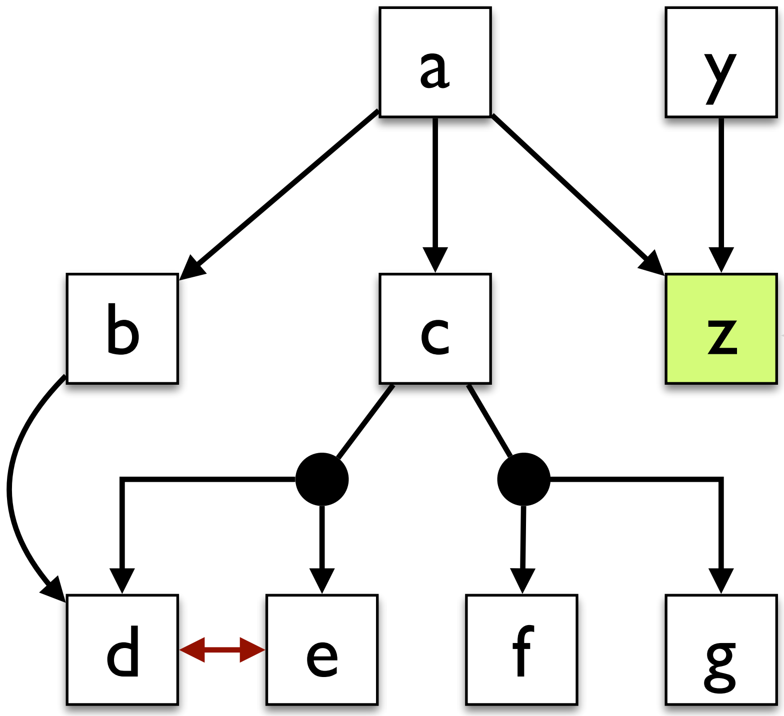
Pseudo-Boolean Constraints

Partial (Weighted) MaxSAT

Deciding if a component can be installed

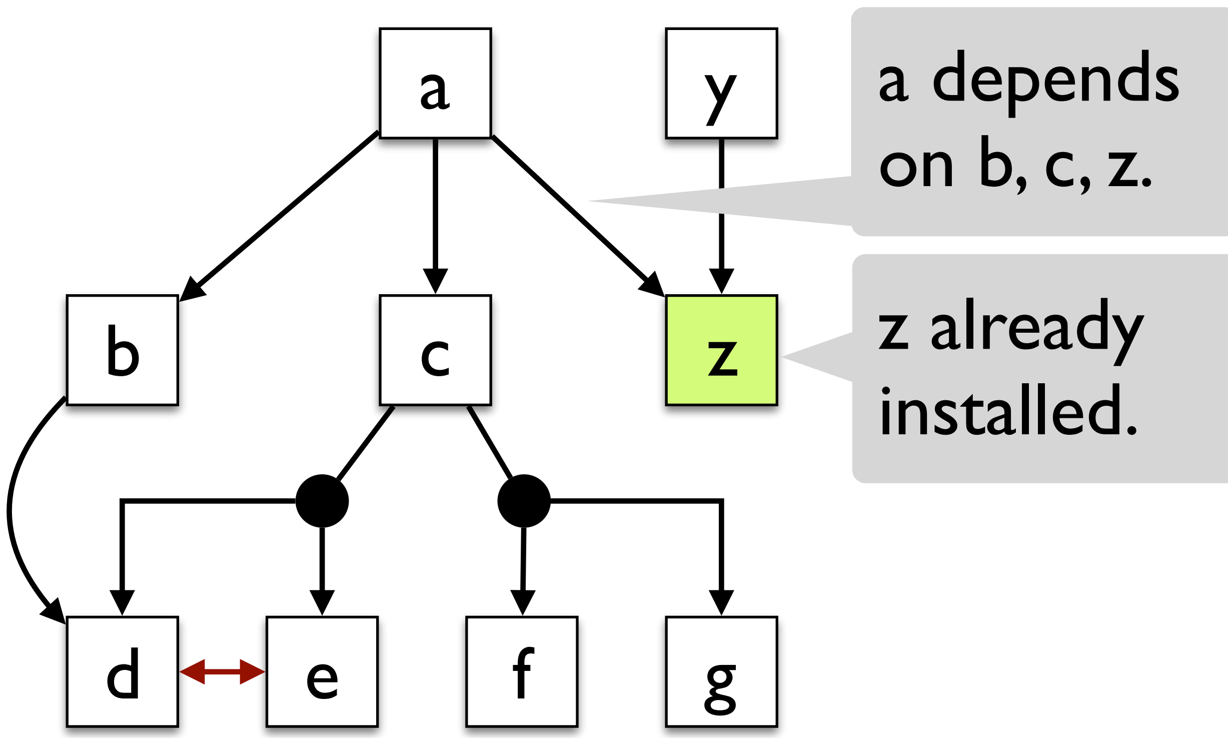


Deciding if a component can be installed

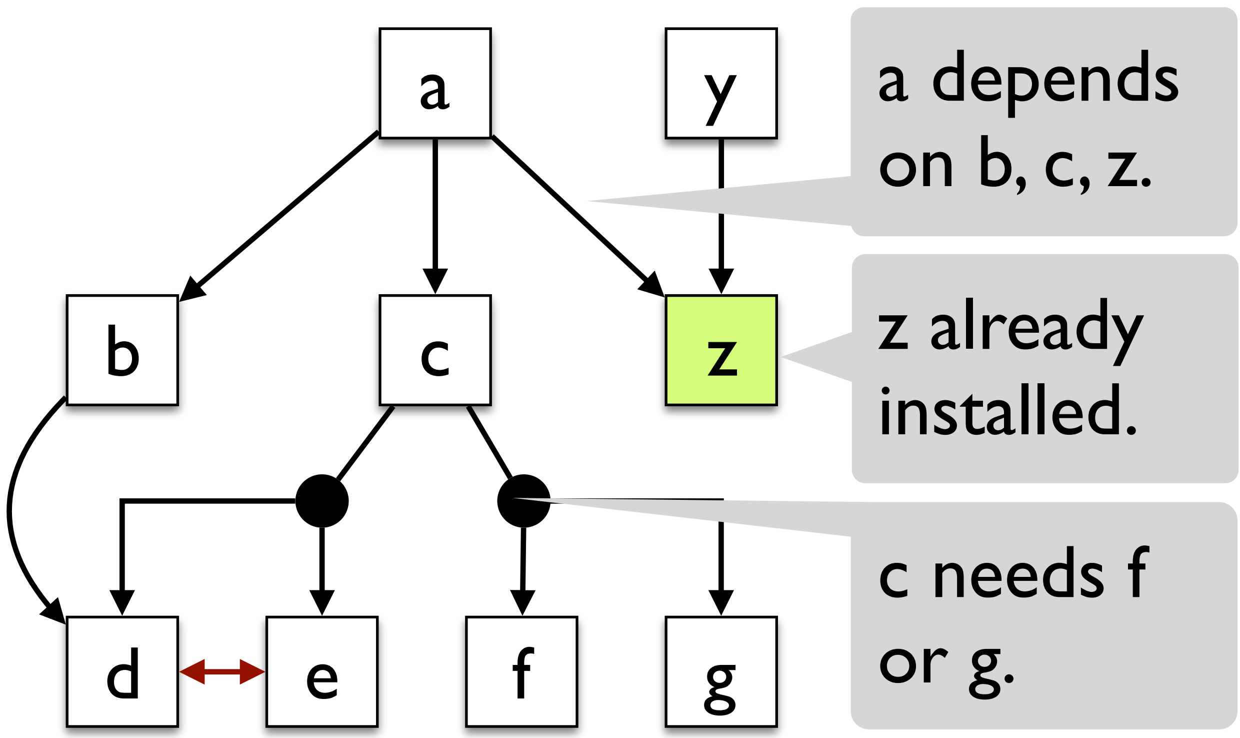


z already installed.

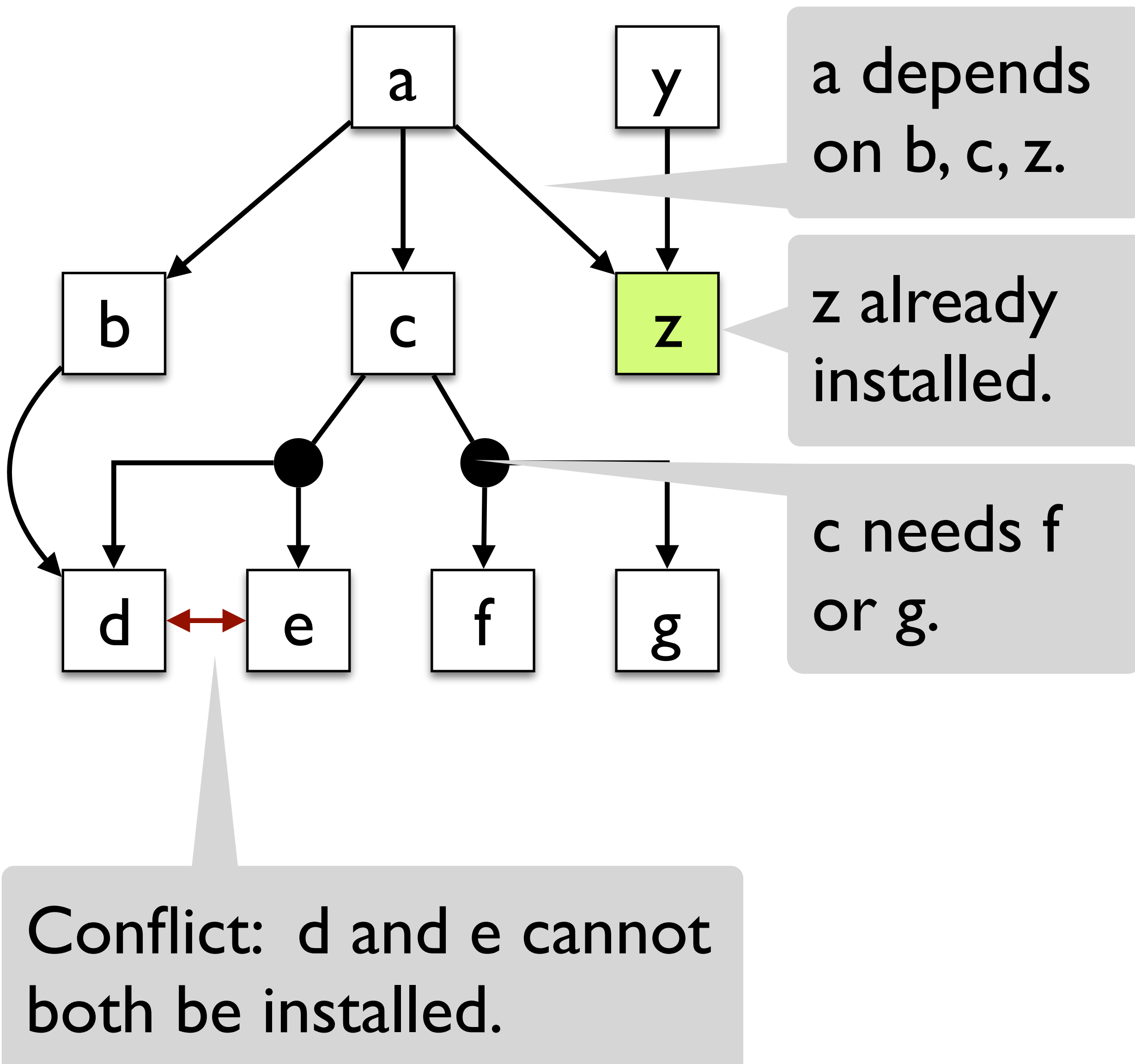
Deciding if a component can be installed



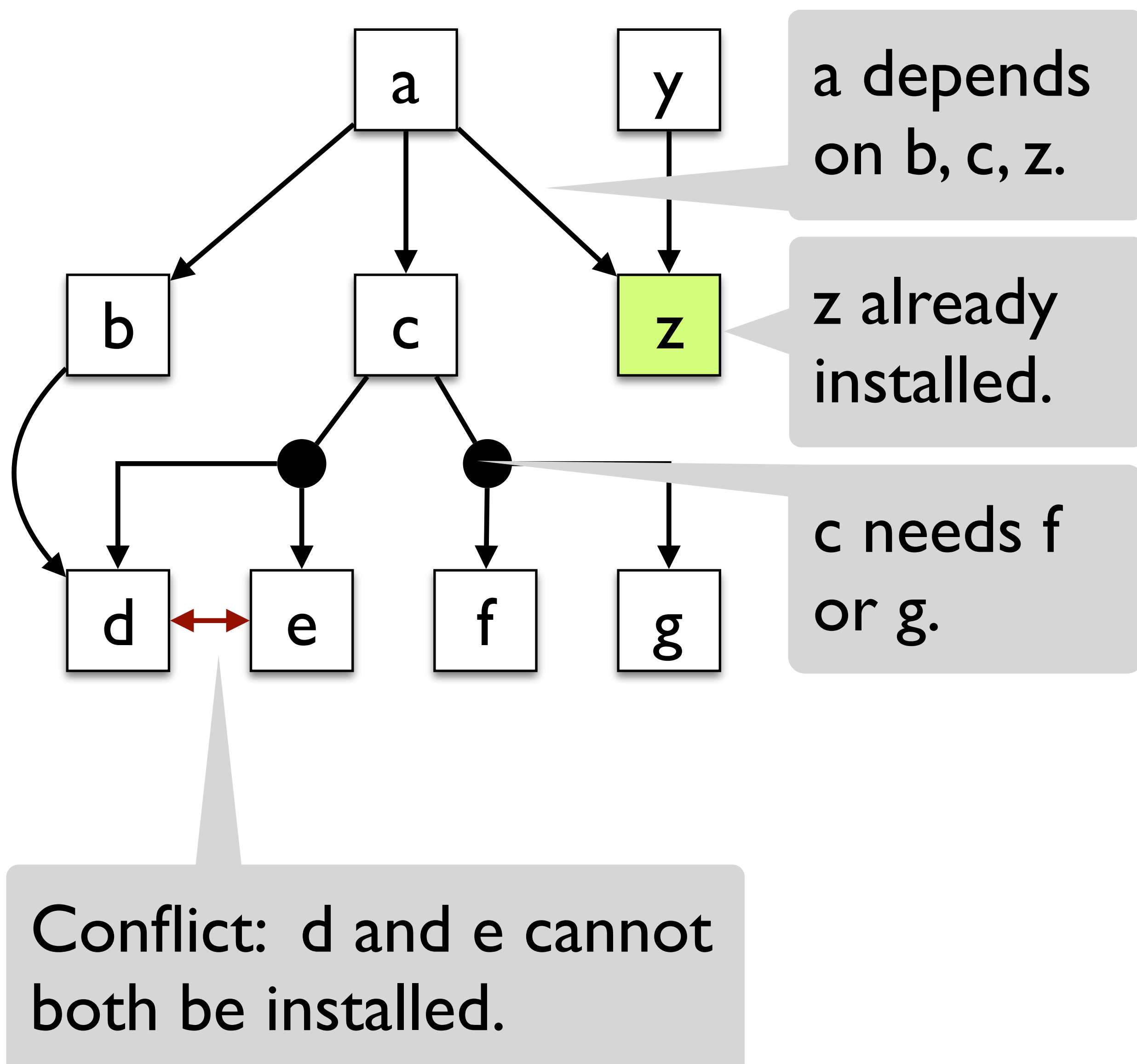
Deciding if a component can be installed



Deciding if a component can be installed

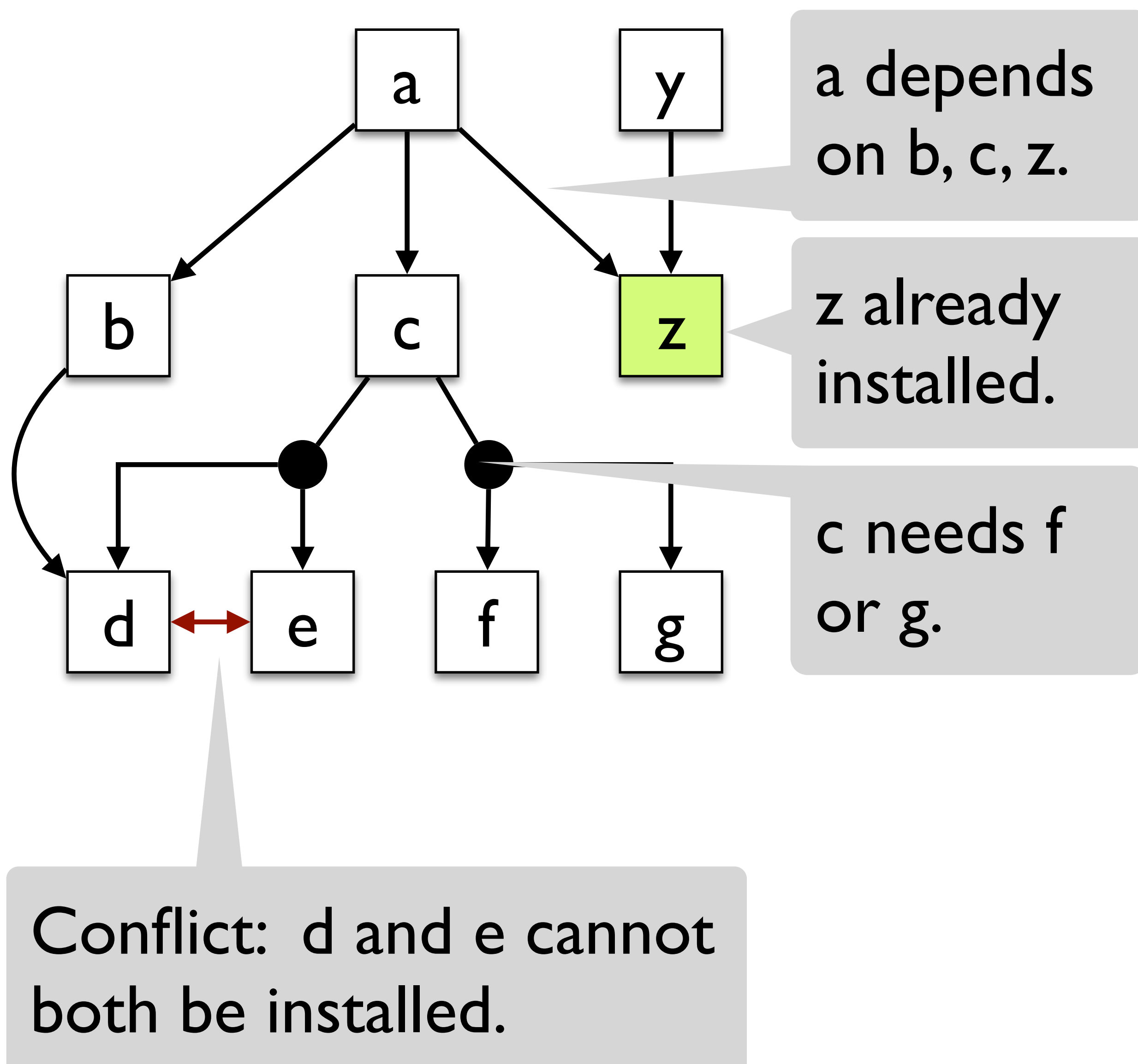


Deciding if a component can be installed



To install a, CNF constraints are:

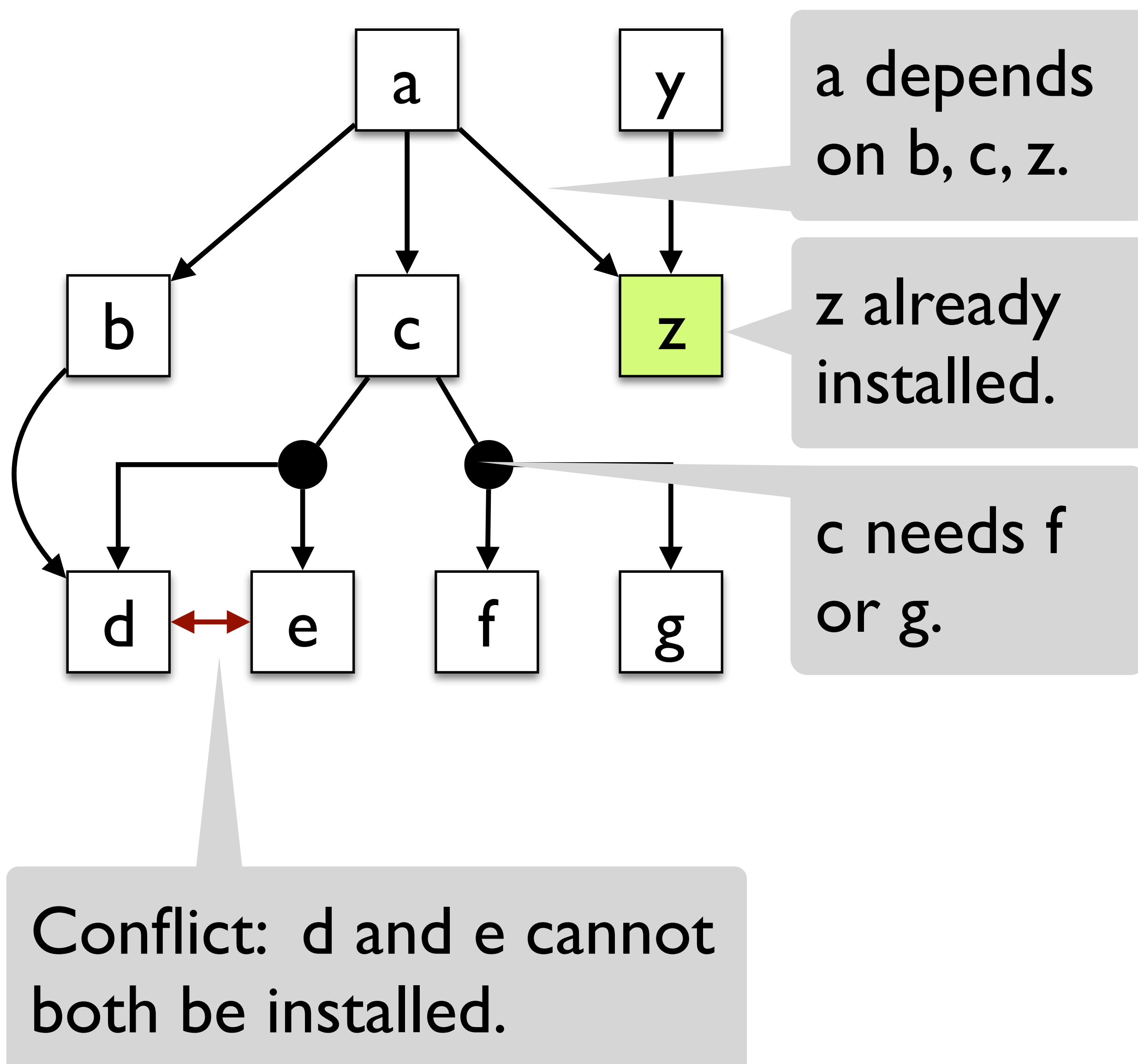
Deciding if a component can be installed



To install a, CNF constraints are:

$(\neg a \lor b) \land (\neg a \lor c) \land (\neg a \lor z) \land$

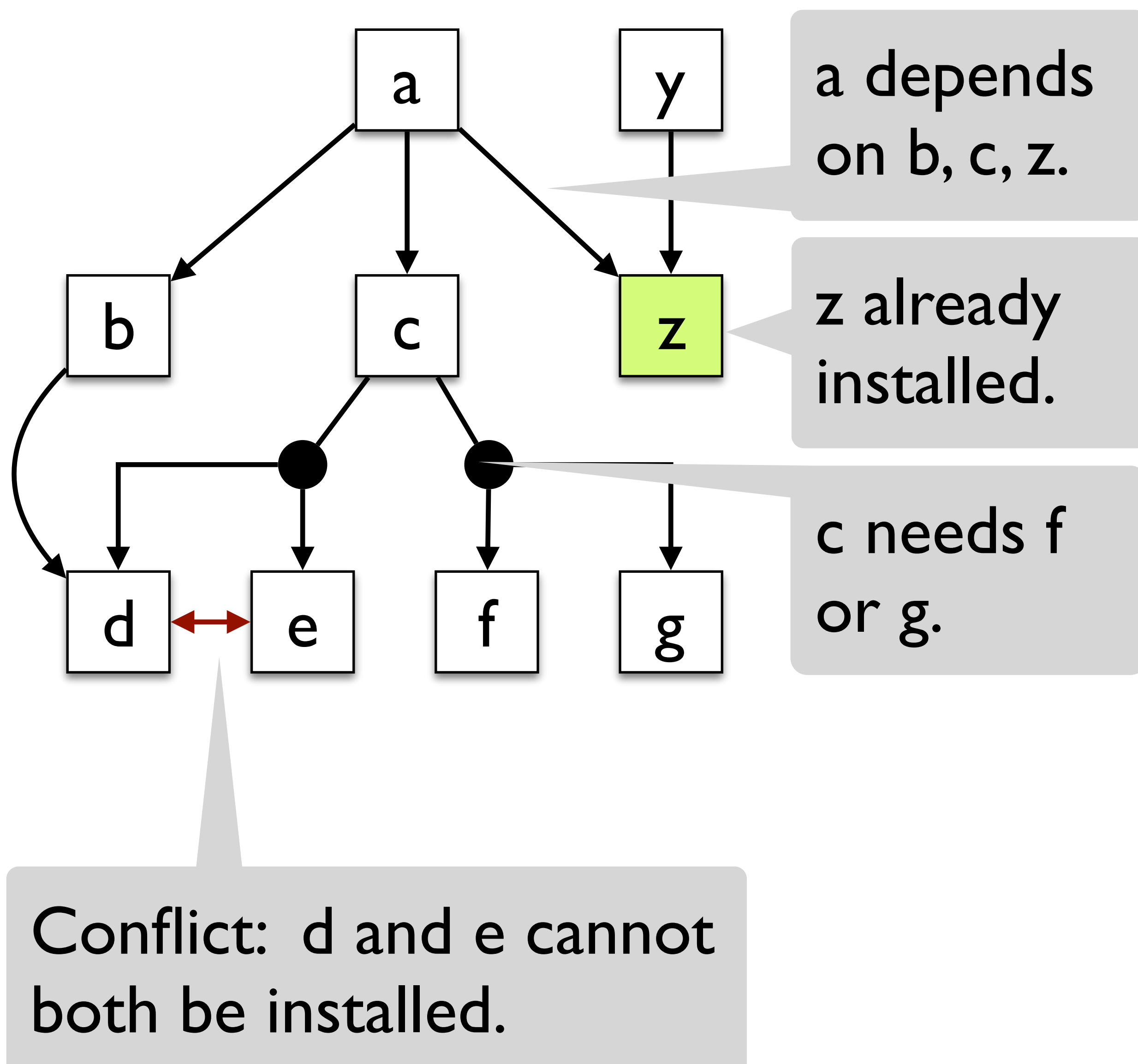
Deciding if a component can be installed



To install a, CNF constraints are:

$(\neg a \vee b) \wedge (\neg a \vee c) \wedge (\neg a \vee z) \wedge$
$(\neg b \vee d) \wedge$

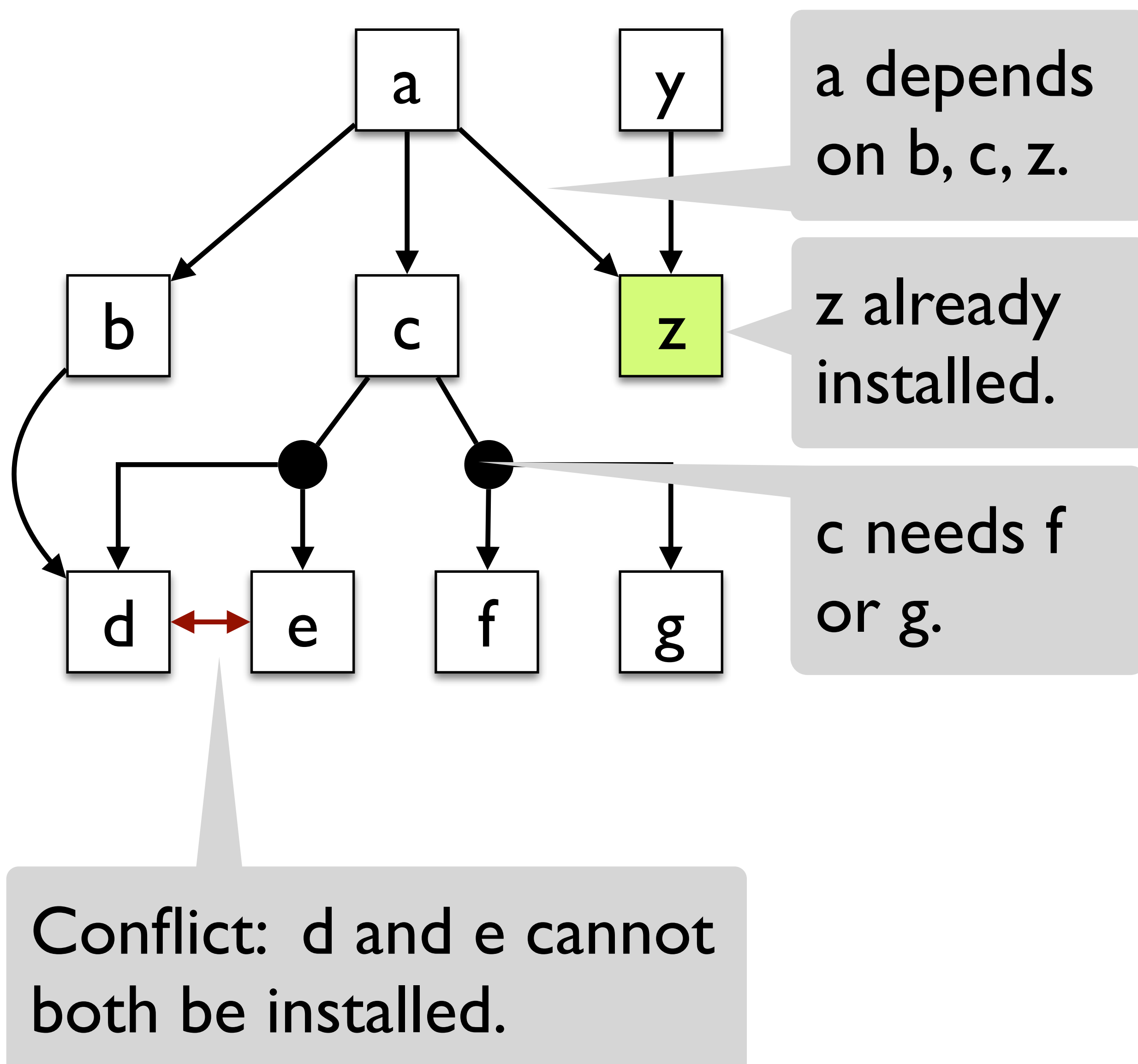
Deciding if a component can be installed



To install a, CNF constraints are:

$(\neg a \vee b) \wedge (\neg a \vee c) \wedge (\neg a \vee z) \wedge$
$(\neg b \vee d) \wedge$
$(\neg c \vee d \vee e) \wedge (\neg c \vee f \vee g) \wedge$

Deciding if a component can be installed



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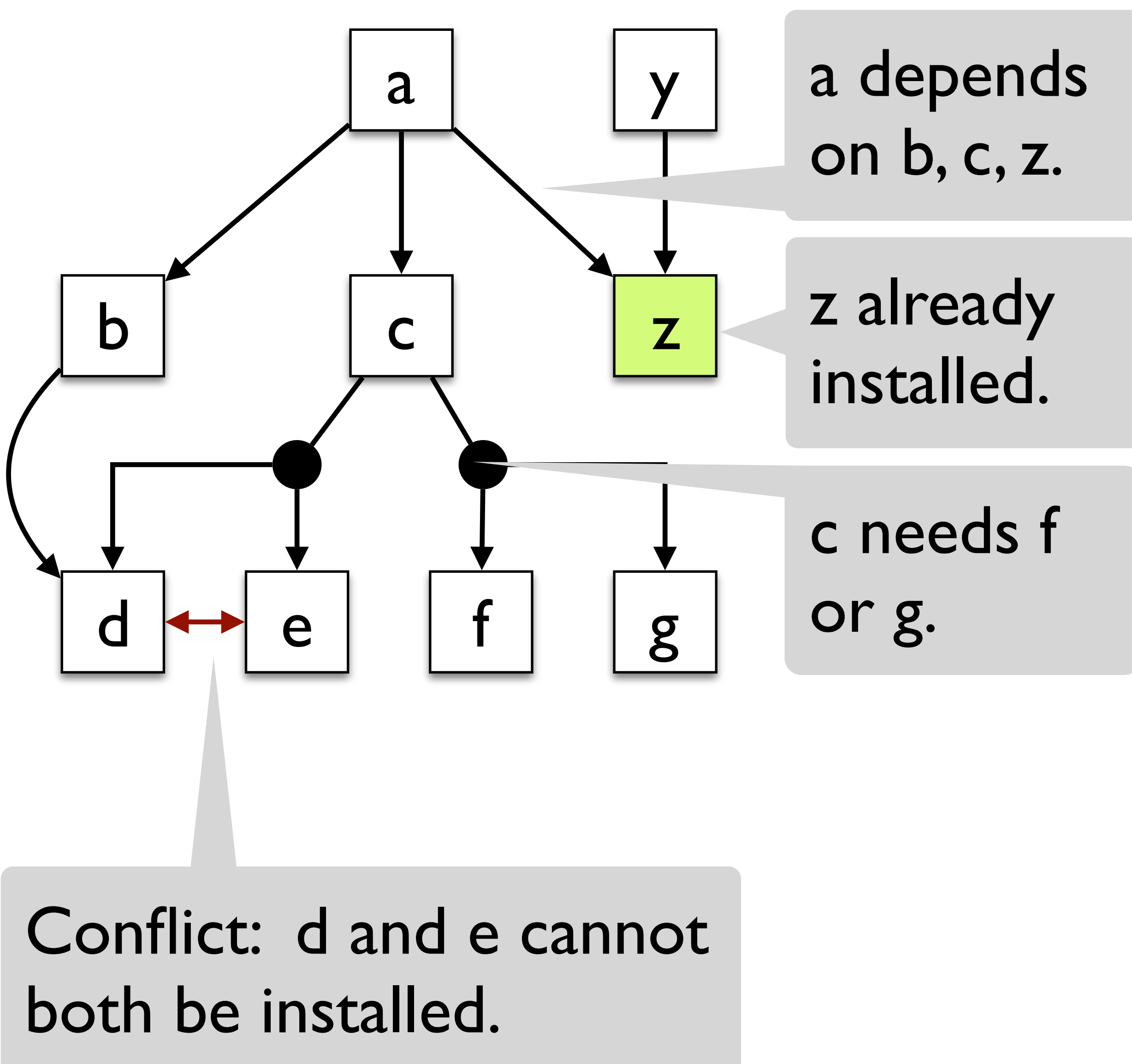
$(\neg b \lor d) \land$
$(\neg c \lor d \lor e) \land (\neg c \lor f \lor g) \land$
$(\neg d \lor \neg e) \land$

Deciding if a component can be installed



To install a, CNF constraints are:

$(\neg a \vee b) \wedge (\neg a \vee c) \wedge (\neg a \vee z) \wedge$

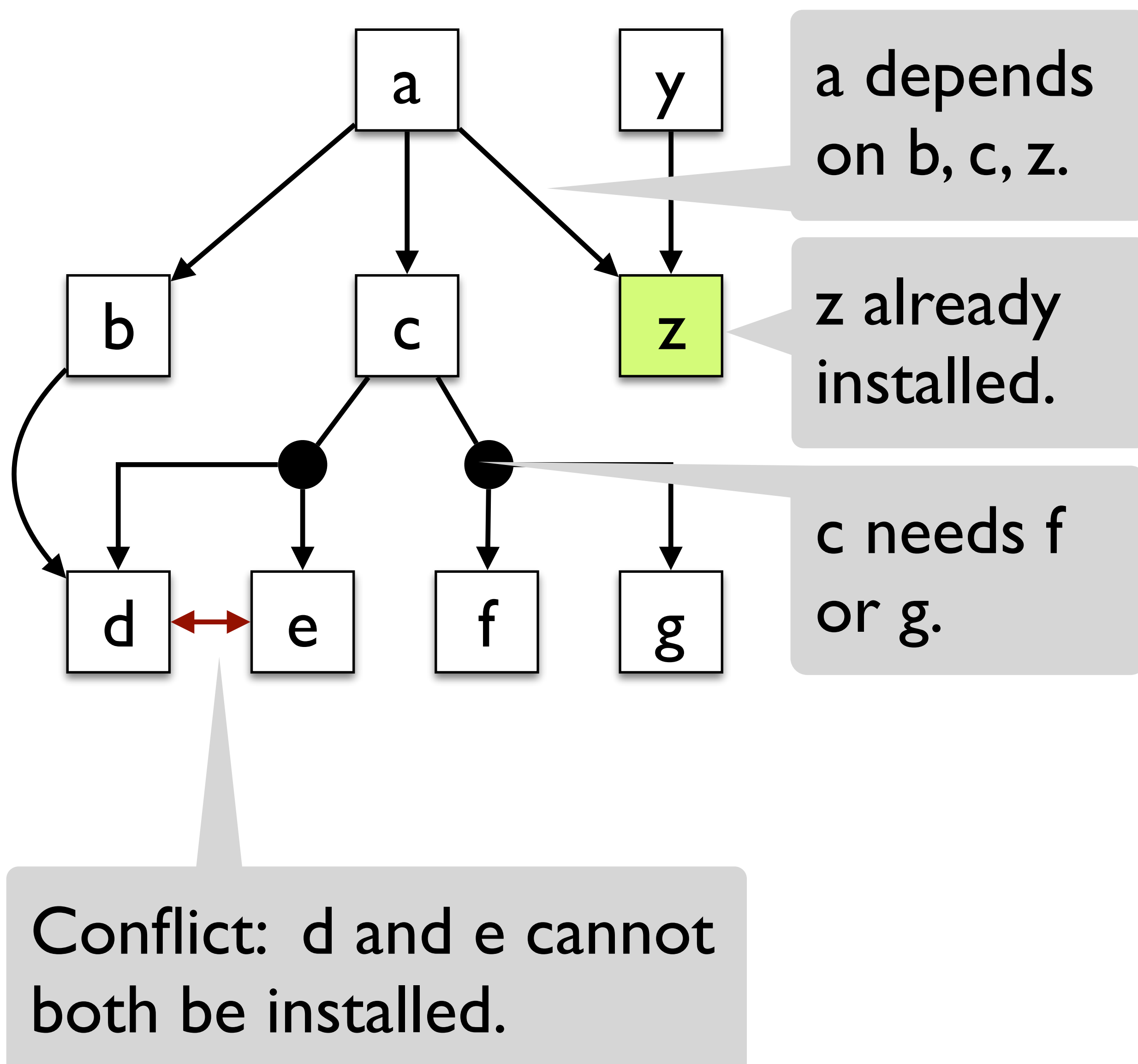
$(\neg b \vee d) \wedge$

$(\neg c \vee d \vee e) \wedge (\neg c \vee f \vee g) \wedge$

$(\neg d \vee \neg e) \wedge$

$(\neg y \vee z) \wedge$

Deciding if a component can be installed



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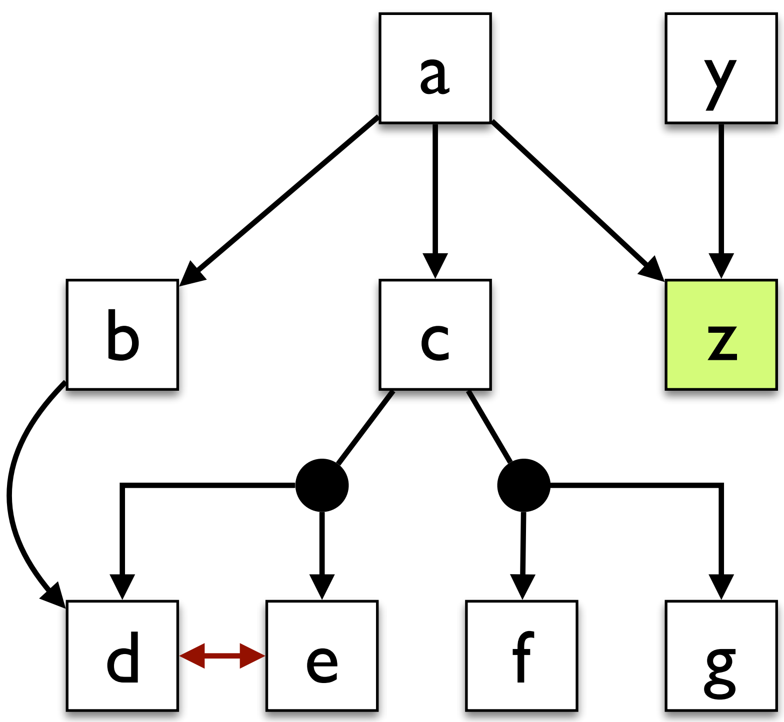
$(\neg a \lor b) \land (\neg a \lor c) \land (\neg a \lor z) \land$
$(\neg b \lor d) \land$
$(\neg c \lor d \lor e) \land (\neg c \lor f \lor g) \land$
$(\neg d \lor \neg e) \land$
$(\neg y \lor z) \land$
$a \land z$

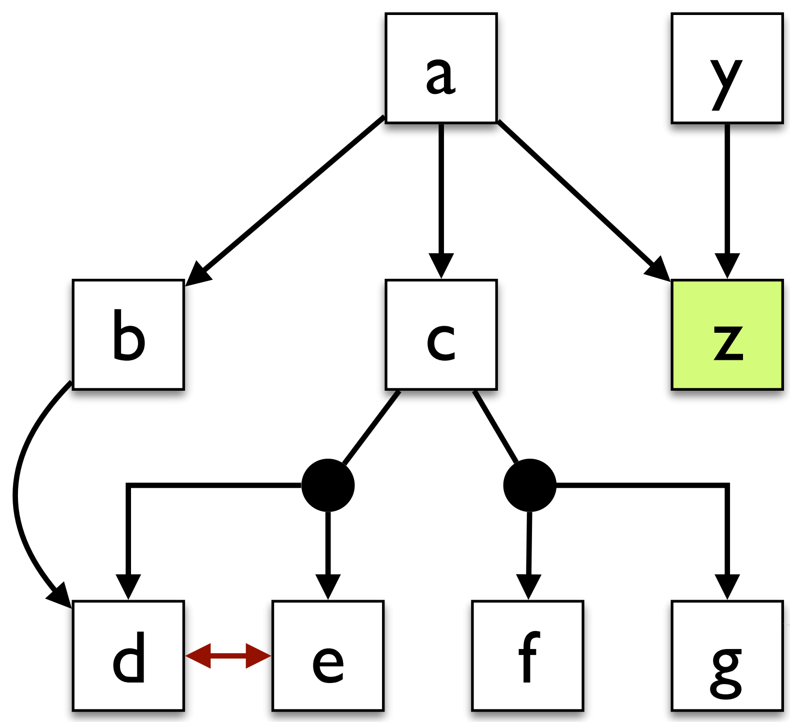
Optimal installation



Assume f and g are 5MB and 2MB each, and all other components are 1MB. How to install a, while minimizing total size?

$(\neg a \vee b) \wedge (\neg a \vee c) \wedge (\neg a \vee z) \wedge$
$(\neg b \vee d) \wedge$
$(\neg c \vee d \vee e) \wedge (\neg c \vee f \vee g) \wedge$
$(\neg d \vee \neg e) \wedge$
$(\neg y \vee z) \wedge$
$a \wedge z$

Optimal installation



Pseudo-boolean solvers accept a linear function to minimize, in addition to a (weighted) CNF.

$$\textbf{min } c_1x_1 + \ldots + c_nx_n$$
$$a_{11}x_1 + \ldots + a_{1n}x_n \geq b_1 \land \ldots \land$$
$$a_{k1}x_1 + \ldots + a_{kn}x_n \geq b_k$$

Assume f and g are 5MB and 2MB each, and all other components are 1MB. How to install a, while minimizing total size?

$(\neg a \lor b) \land (\neg a \lor c) \land (\neg a \lor z) \land$
$(\neg b \lor d) \land$
$(\neg c \lor d \lor e) \land (\neg c \lor f \lor g) \land$
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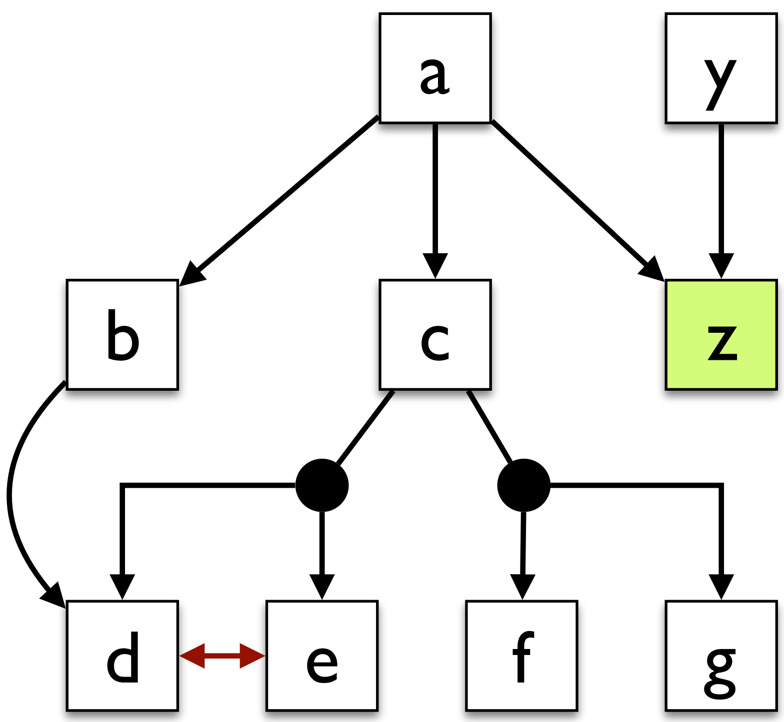
Optimal installation



Assume f and g are 5MB and 2MB each, and all other components are 1MB. How to install a, while minimizing total size?

Pseudo-boolean solvers accept a linear function to minimize, in addition to a (weighted) CNF.

$\textbf{min}\ a + b + c + d + e + 5f + 2g + y + 0z$

$(-a + b \geq 0) \wedge (-a + c \geq 0) \wedge (-a + z \geq 0) \wedge$

$(-b + d \geq 0) \wedge$

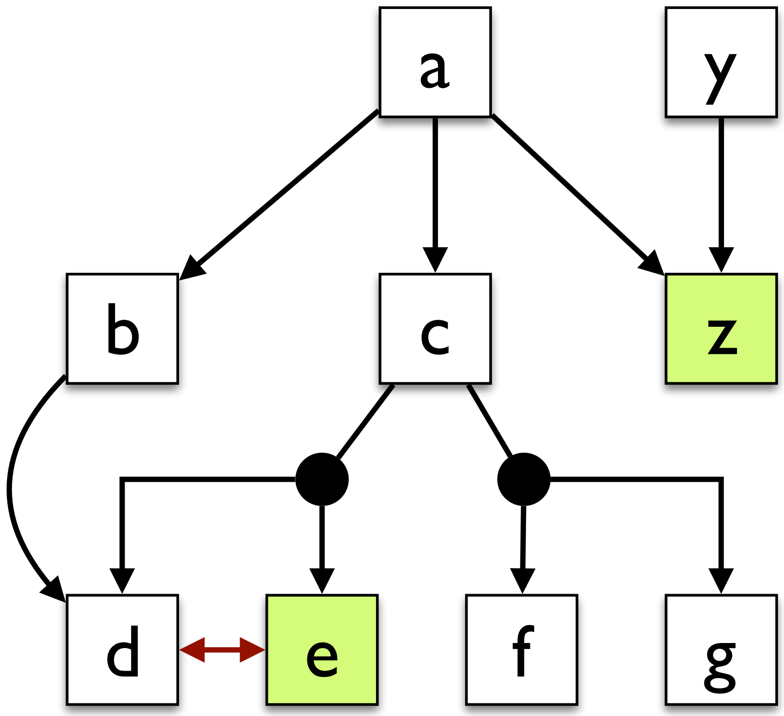
$(-c + d + e \geq 0) \wedge (-c + f + g \geq 0) \wedge$

$(-d + -e \geq -1) \wedge$

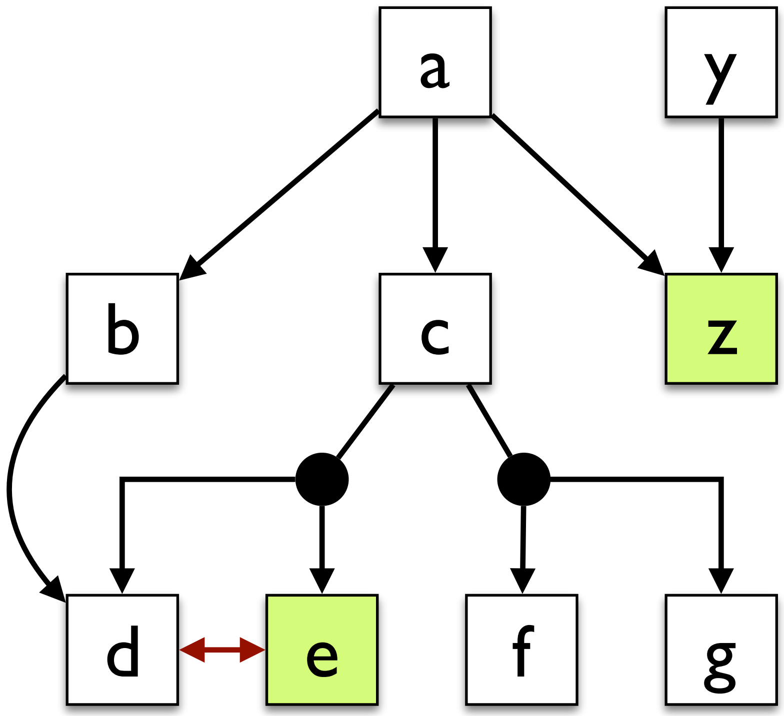
$(-y + z \geq 0) \wedge$

$(a \geq 1) \wedge (z \geq 1)$

Installation in the presence of conflicts

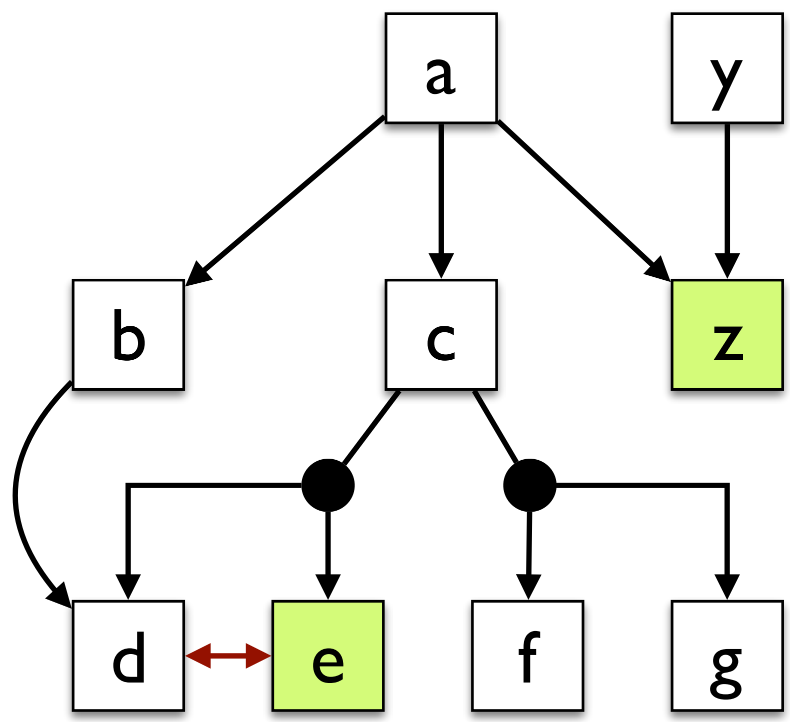


Installation in the presence of conflicts



a cannot be installed because it requires b, which requires d, which conflicts with e.

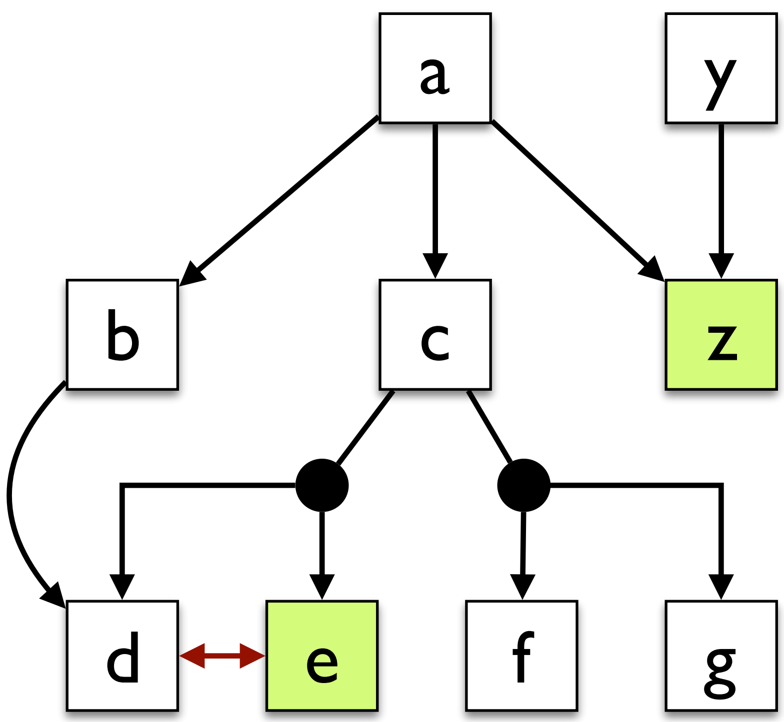
Installation in the presence of conflicts



a cannot be installed because it requires b, which requires d, which conflicts with e.

Partial MaxSAT solver takes as input a set of **hard** clauses and a set of **soft** clauses, and it produces an assignment that satisfies all hard clauses and the greatest number of soft clauses.

Installation in the presence of conflicts



a cannot be installed because it requires b, which requires d, which conflicts with e.

Partial MaxSAT solver takes as input a set of **hard** clauses and a set of **soft** clauses, and it produces an assignment that satisfies all hard clauses and the greatest number of soft clauses.

To install a, while minimizing the number of removed components, Partial MaxSAT constraints are:

hard: $(\neg a \vee b) \wedge (\neg a \vee c) \wedge (\neg a \vee z) \wedge$
$(\neg b \vee d) \wedge$
$(\neg c \vee d \vee e) \wedge (\neg c \vee f \vee g) \wedge$
$(\neg d \vee \neg e) \wedge (\neg y \vee z) \wedge a$

soft: $e \wedge z$

Summary

Today

- SAT solvers have been used successfully in many applications & domains
- But reducing problems to SAT is a lot like programming in assembly …
- We need higher-level logics!

Next lecture

- On to richer logics: introduction to Satisfiability Modulo Theories (SMT)