Computer-Aided Reasoning for Software

CSE507

Program Synthesis

courses.cs.washington.edu/courses/cse507/18sp/

Emina Torlak
emina@cs.washington.edu

Today

Last lecture

- Solvers as angelic runtime oracle

Today

- Program synthesis: computers programming computers

Reminders

- HW3 is due tonight.

topics

Computers programming computers?

"Information technology has been praised as a labor saver and cursed as a destroyer of obsolete jobs. But the entire edifice of modern computing rests on a fundamental irony: **the software that makes it all possible is, in a very real sense, handmade.** Every miraculous thing computers can accomplish begins with a human programmer entering lines of code by hand, character by character."

Interview with Moshe Vardi

Program synthesis aims to automate (tedious parts of) programming.

The program synthesis problem

$$\exists\ P.\ \forall\ x.\ \varphi(x, P(x))$$

Find a program P that satisfies the specification φ on all inputs.

The program synthesis problem

φ may be a formula, a reference implementation, input/output pairs, traces, demonstrations, etc.

$$\exists P. \forall x. \; \varphi(x, P(x))$$

Find a program P that satisfies the specification φ on all inputs.

The program synthesis problem

φ may be a formula, a reference implementation, input/output pairs, traces, demonstrations, etc.

Synthesis improves

- Productivity (when writing φ is easier than writing P).
- Correctness (when verifying φ is easier than verifying P).

$$\exists\, P.\ \forall\, x.\ \ \varphi(x, P(x))$$

Find a program P that satisfies the specification φ on all inputs.

Two kinds of program synthesis

$\exists P. \forall x. \ \varphi(x, P(x))$

Deductive (classic) synthesis

Inductive (syntax-guided) synthesis

Two kinds of program synthesis

$\exists P. \forall x. \; \varphi(x, P(x))$

Deductive (classic) synthesis

Derive the program P from the constructive proof of the theorem $\forall x. \exists y. \varphi(y, x)$.

Inductive (syntax-guided) synthesis

Two kinds of program synthesis

$\exists P. \forall x. \ \varphi(x, P(x))$

Deductive (classic) synthesis

Derive the program P from the constructive proof of the theorem $\forall x. \exists y. \varphi(y, x)$.

Inductive (syntax-guided) synthesis

Discover the program P by searching a restricted space of candidate programs for one that satisfies φ on all inputs.

Two kinds of program synthesis

$$\exists\, P.\ \forall\, x.\ \ \varphi(x, P(x))$$

SPIRAL

Synthesis as a problem in deductive theorem proving.

FlashFill

Synthesis as a search problem.

Deductive synthesis with axioms and E-graphs

Denali Superoptimizer
[Joshi, Nelson, Randall, PLDI'02]

Deductive synthesis with axioms and E-graphs

Specification φ, given as a reference implementation.

```
reg6 * 4 + 1
```

Denali Superoptimizer [Joshi, Nelson, Randall, PLDI'02]

Deductive synthesis with axioms and E-graphs

Specification φ, given as a reference implementation.

Optimal (lowest cost) program P that is equivalent to φ on all inputs (values of reg6).

`reg6 * 4 + 1` → Denali Superoptimizer [Joshi, Nelson, Randall, PLDI'02] → `s4addl(reg6, 1)`

Deductive synthesis with axioms and E-graphs

Specification φ, given as a reference implementation.

`reg6 * 4 + 1`

Denali Superoptimizer [Joshi, Nelson, Randall, PLDI'02]

`s4addl(reg6, 1)`

Optimal (lowest cost) program P that is equivalent to φ on all inputs (values of reg6).

∀ k, n. 2^n = 2**n

∀ k, n. k*2^n = k << n

∀ k, n. k*4 + n = **s4addl**(k, n)

…

Two kinds of axioms:

- Instruction semantics.
- Algebraic properties of functions and relations used for specifying instruction semantics.

Deductive synthesis with axioms and E-graphs

Specification φ, given as a reference implementation.

1. Construct an E-graph.
2. Use a SAT solver to search the E-graph for a K-cycle program.

Optimal (lowest cost) program P that is equivalent to φ on all inputs (values of reg6).

`reg6 * 4 + 1` → Denali Superoptimizer [Joshi, Nelson, Randall, PLDI'02] → `s4addl(reg6, 1)`

$\forall$ k, n. 2^n = 2**n

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Two kinds of axioms:

- Instruction semantics.
- Algebraic properties of functions and relations used for specifying instruction semantics.

Denali by example

`reg6 * 4 + 1`

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E-graph matching

SAT

s`4addl(reg6, 1)`

Denali by example

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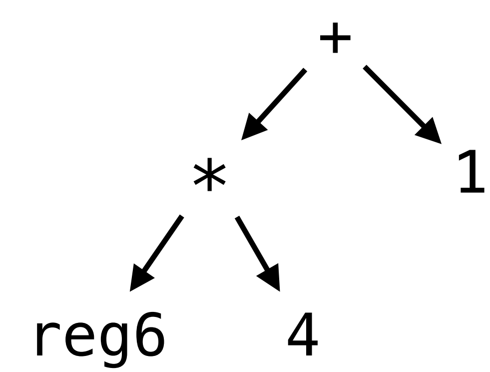
E-graph matching



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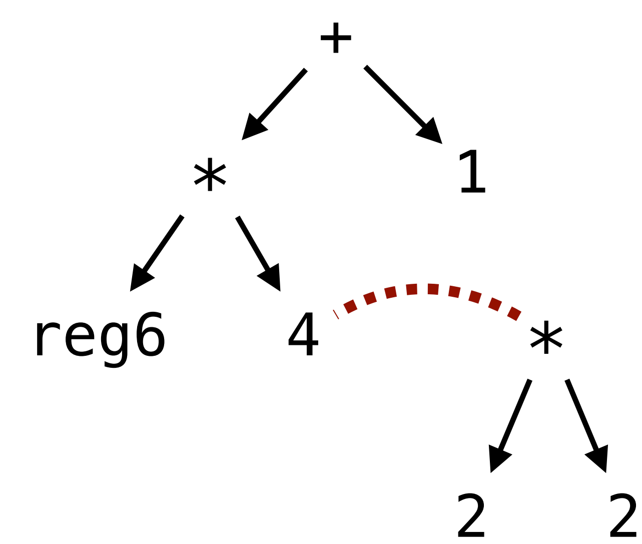

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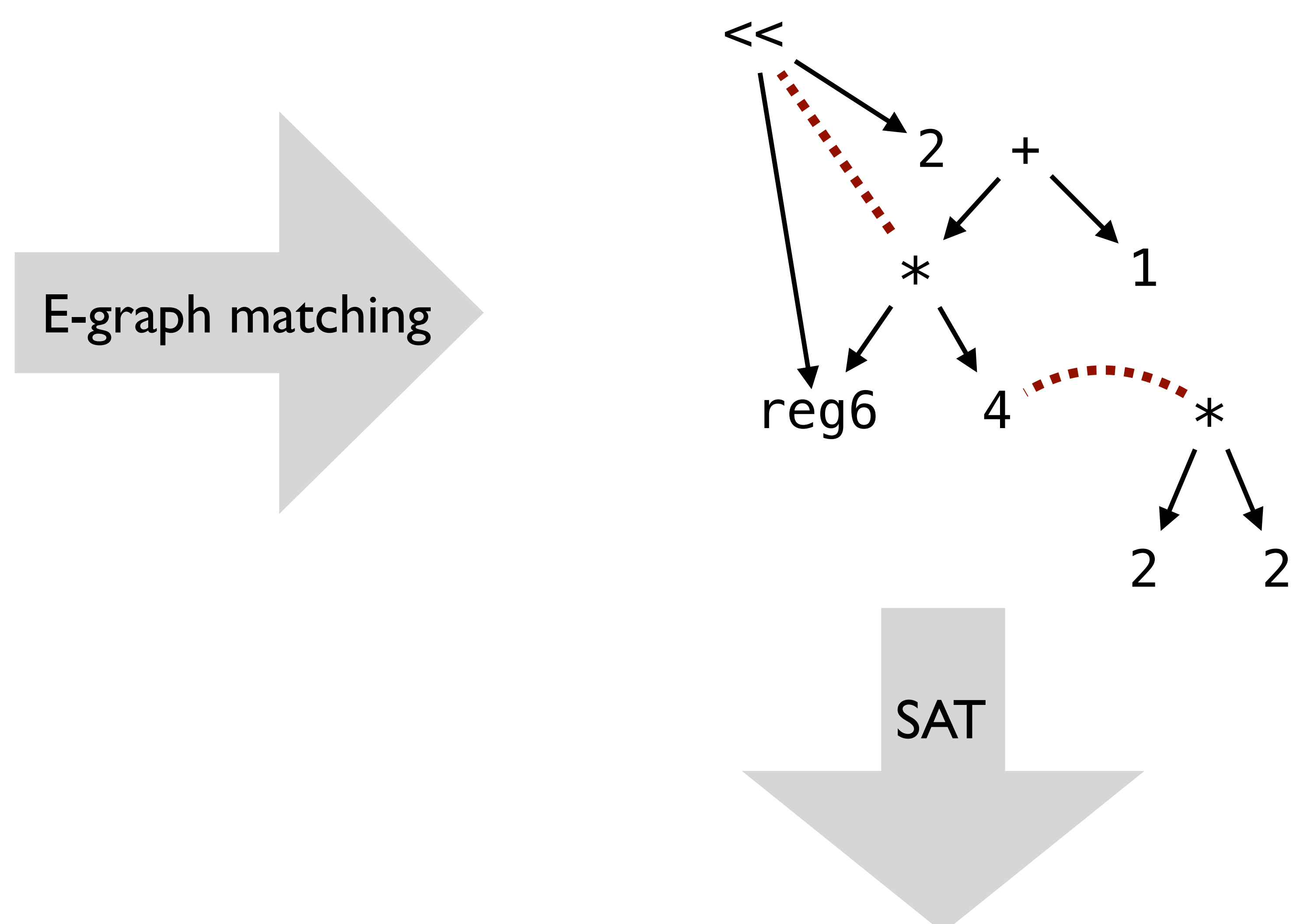

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Denali by example

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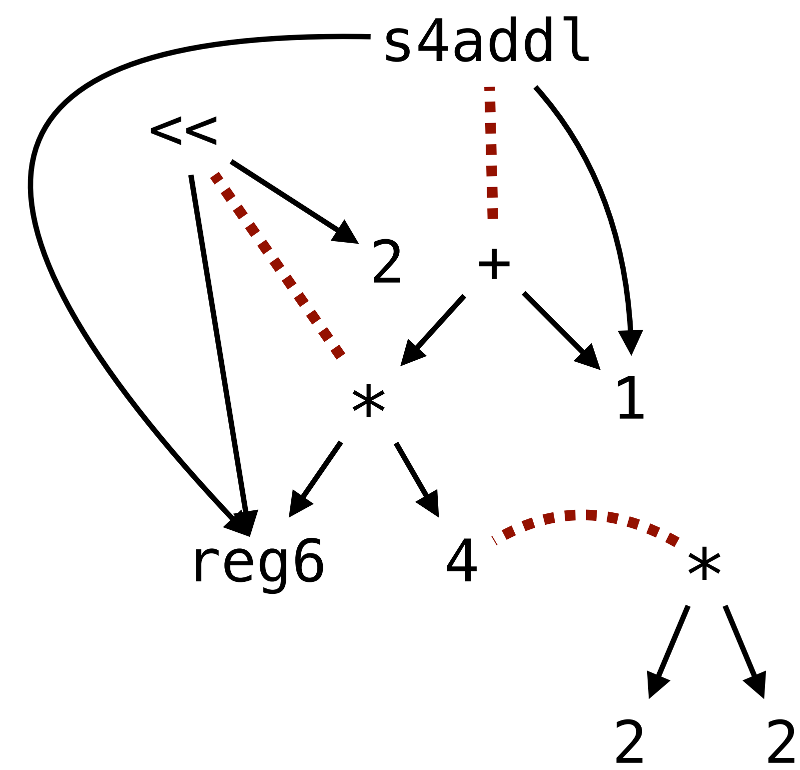
E-graph matching



SAT

`s4addl(reg6, 1)`

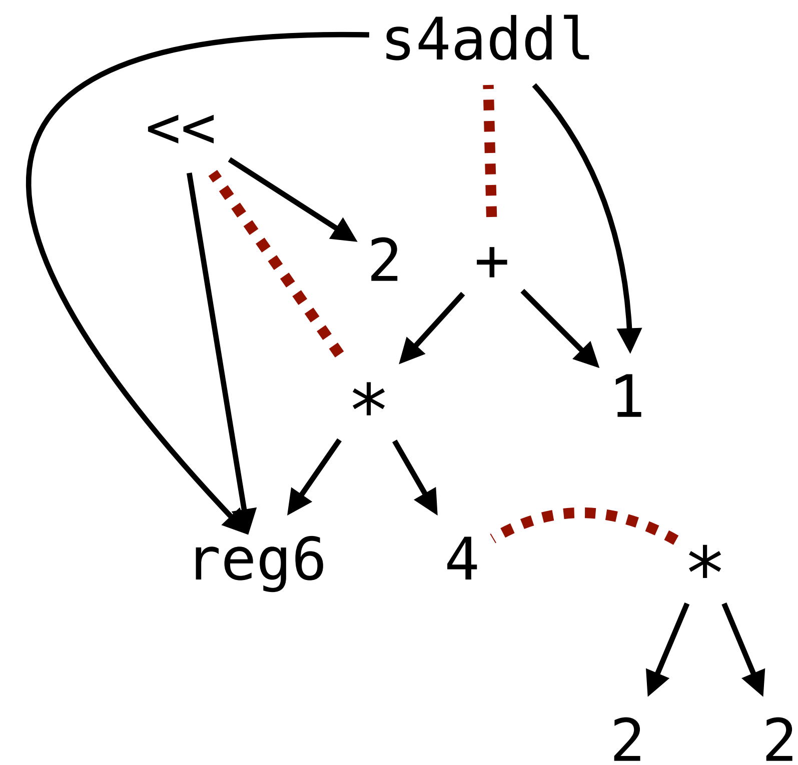
Deductive synthesis versus compilation

Deductive synthesizer

- Non-deterministic.
- Searches *all correct rewrites* for one that is optimal.

Compiler

- Deterministic.
- Lowers a source program into a target program using a *fixed sequence of rewrite steps*.




```
reg6 * 4 + 1
     ↓
reg6 << 2 + 1
```

Deductive synthesis versus inductive synthesis

$$\exists\, P.\ \forall\, x.\ \ \varphi(x, P(x))$$

Deductive synthesis

- Efficient and provably correct: thanks to the semantics-preserving rules, only correct programs are explored.
- Requires *sufficient axiomatization* of the domain.
- Requires *complete* specifications to seed the derivation.

Deductive synthesis versus inductive synthesis

$\exists$ P. $\forall$ x. φ(x, P(x))

Deductive synthesis

- Efficient and provably correct: thanks to the semantics-preserving rules, only correct programs are explored.
- Requires *sufficient axiomatization* of the domain.
- Requires *complete* specifications to seed the derivation.

Inductive synthesis

- Works with *multi-modal and partial* specifications.
- Requires *no axioms*.
- But often at the cost of *lower efficiency* and *weaker (bounded) guarantees* on the correctness/ optimality of synthesized code.

Inductive syntax-guided synthesis

CEGIS:
Counterexample-Guided
Inductive Synthesis
[**Solar-Lezama et al,
ASPLOS'06**]

Inductive syntax-guided synthesis

A partial or multimodal specification φ of the desired program (e.g., assertions, i/o pairs).

```
reg6 * 4 + 1
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CEGIS: Counterexample-Guided Inductive Synthesis [Solar-Lezama et al, ASPLOS'06]

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CEGIS:
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[Solar-Lezama et al, ASPLOS'06]

```
expr :=
  const | reg6 |
  s4addl(expr, expr) |
  …
```

A syntactic *sketch* (e.g., a grammar) describing the shape of the desired program P.

This defines the space of candidate programs to search. Can be fine-tuned for better performance.

Inductive syntax-guided synthesis

A partial or multimodal specification φ of the desired program (e.g., assertions, i/o pairs).

```
reg6 * 4 + 1
```

CEGIS: Counterexample-Guided Inductive Synthesis [Solar-Lezama et al, ASPLOS'06]

A program P from the given space of candidates that satisfies φ on all (usually bounded) inputs.

```
s4addl(reg6, 1)
```

```
expr :=
 const | reg6 |
 s4addl(expr, expr) |
 …
```

A syntactic *sketch* (e.g., a grammar) describing the shape of the desired program P.

This defines the space of candidate programs to search. Can be fine-tuned for better performance.

Inductive syntax-guided synthesis

A partial or multimodal specification φ of the desired program (e.g., assertions, i/o pairs).

Guess a program that works on a finite set of inputs, verify it, and learn from bad guesses.

A program P from the given space of candidates that satisfies φ on all (usually bounded) inputs.

```
reg6 * 4 + 1
```

CEGIS: Counterexample-Guided Inductive Synthesis [**Solar-Lezama et al, ASPLOS'06**]

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A syntactic *sketch* (e.g., a grammar) describing the shape of the desired program P.

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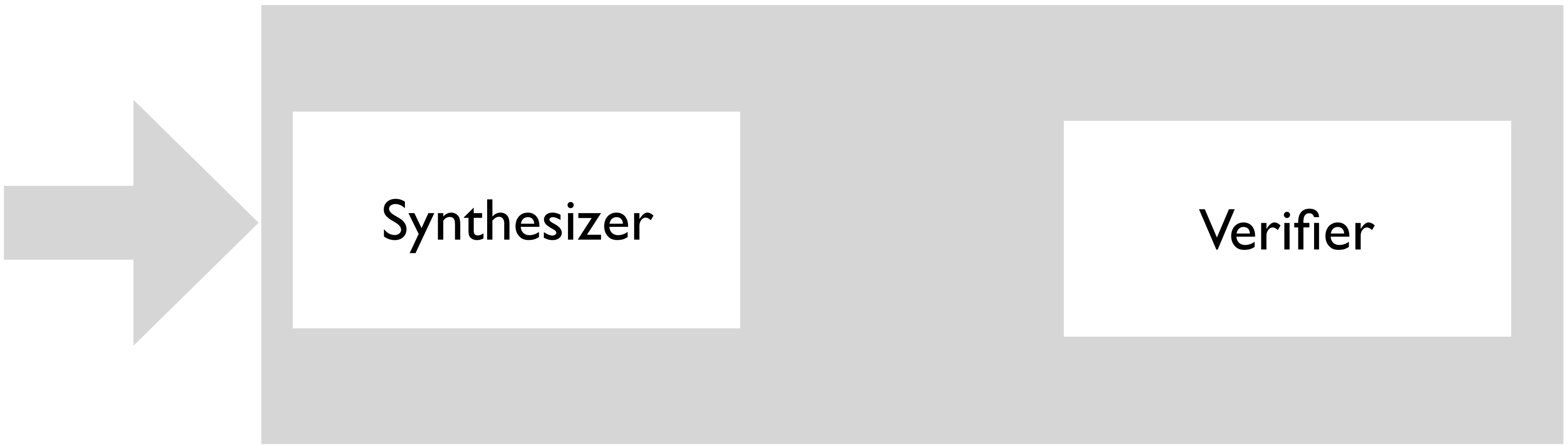
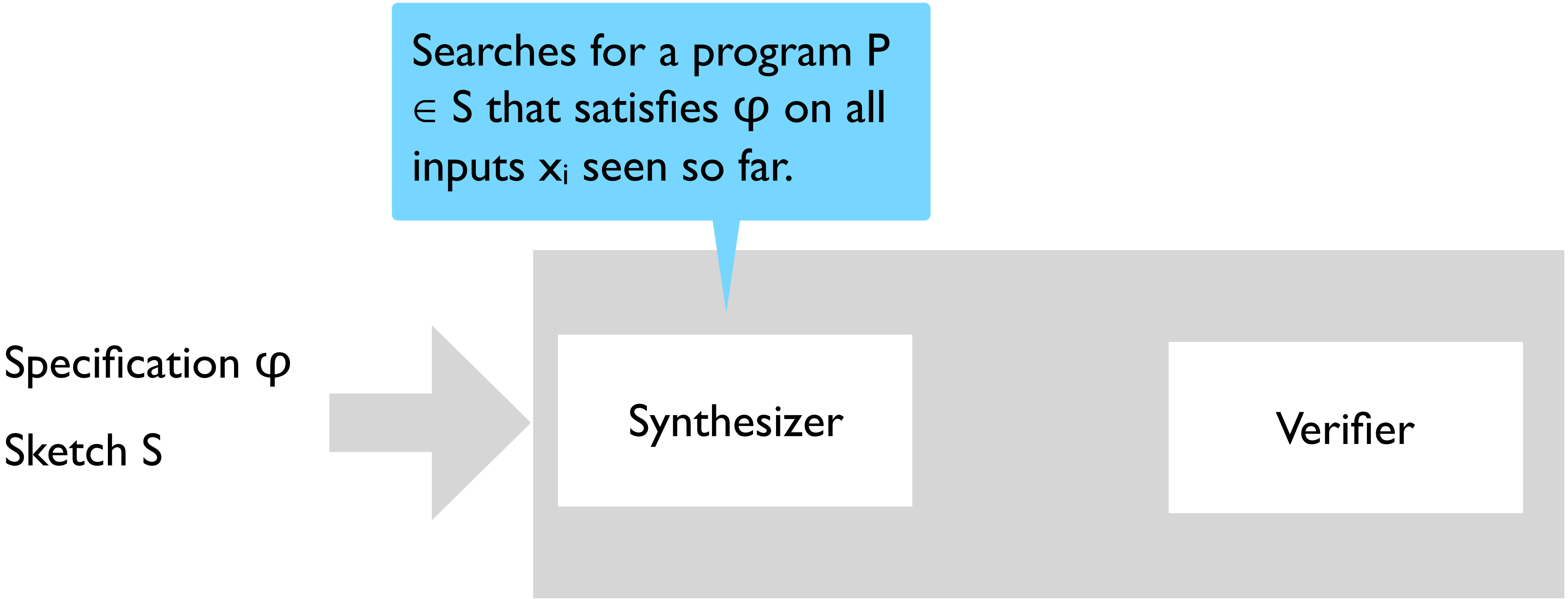
Overview of CEGIS

Specification φ

Sketch S



Overview of CEGIS



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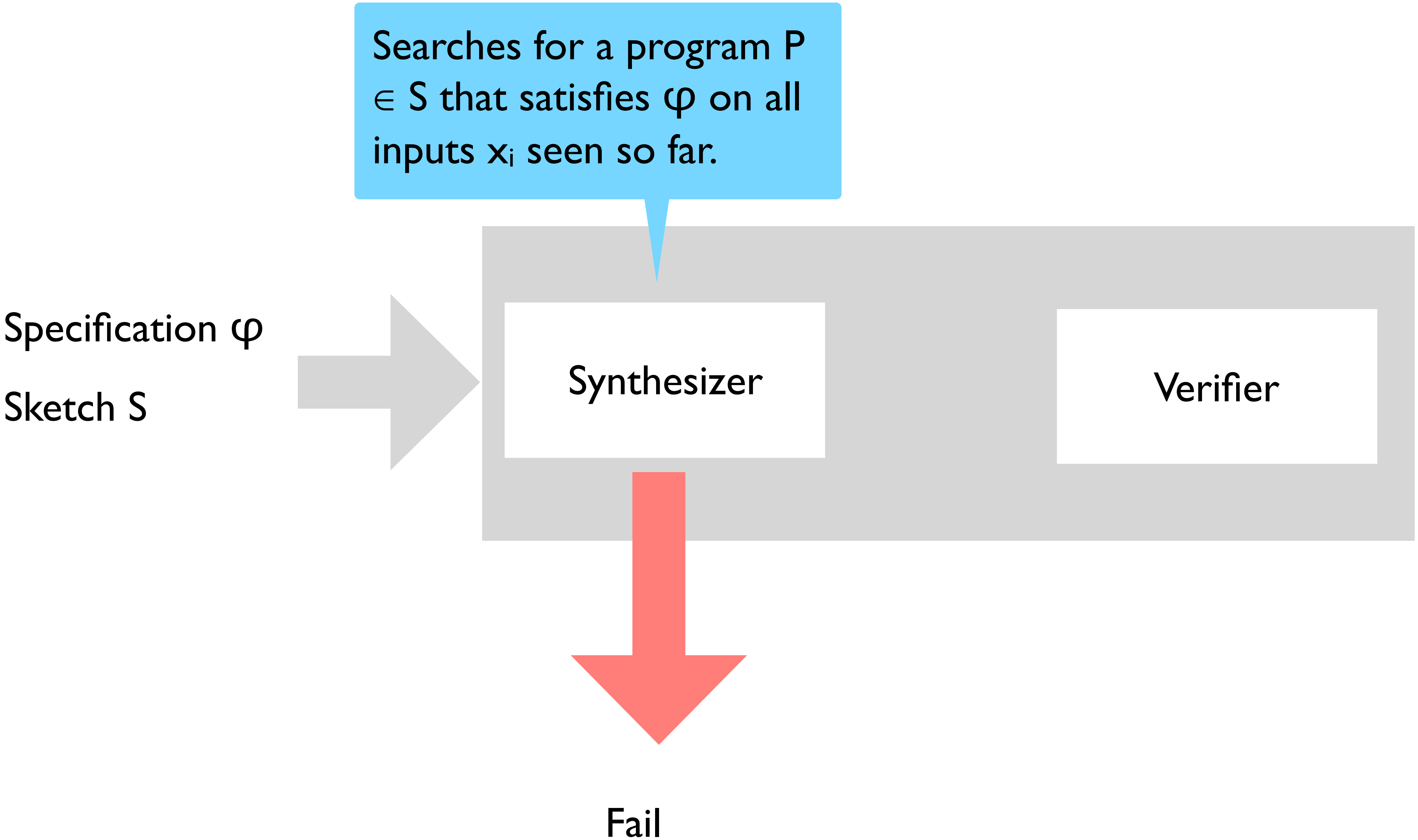

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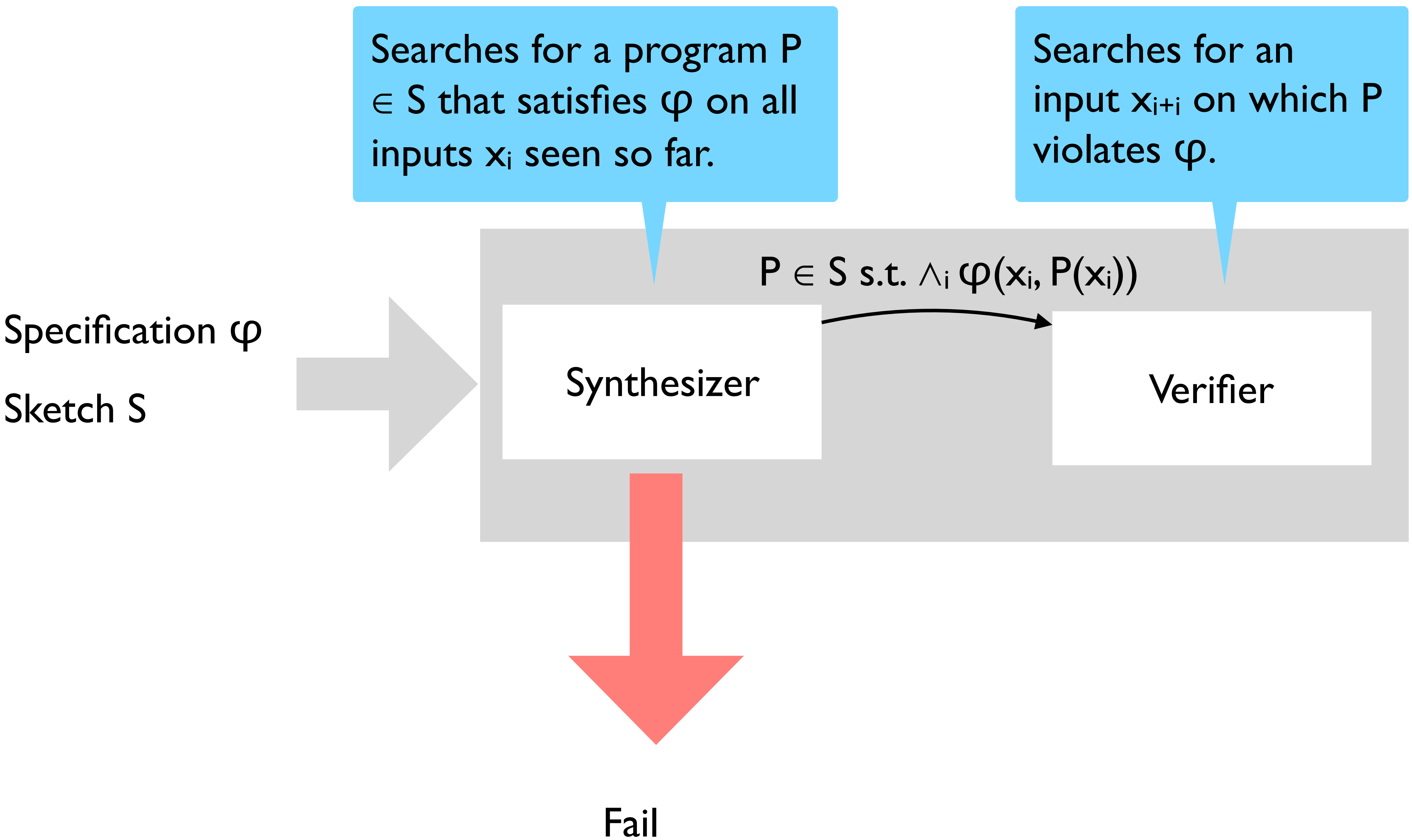

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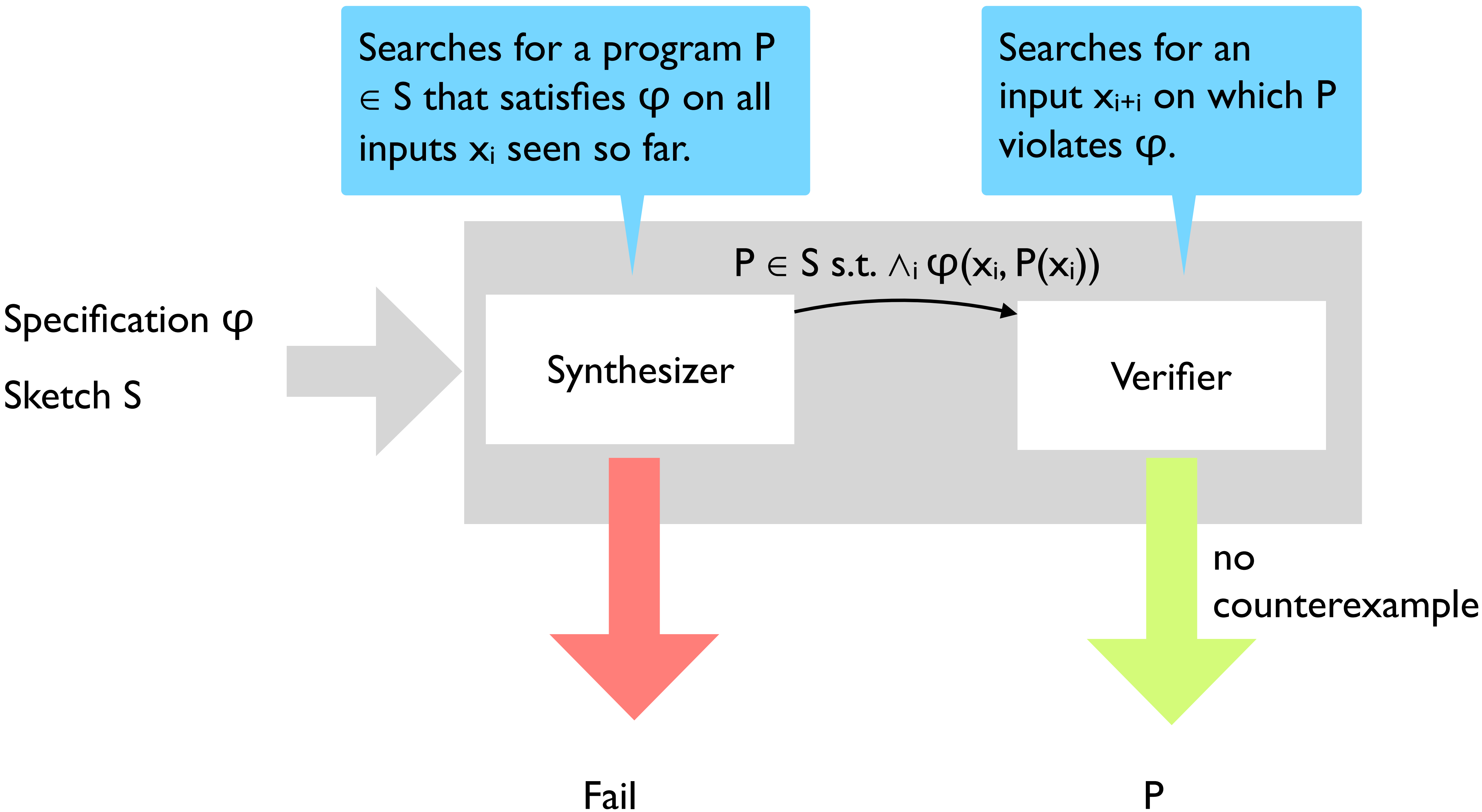

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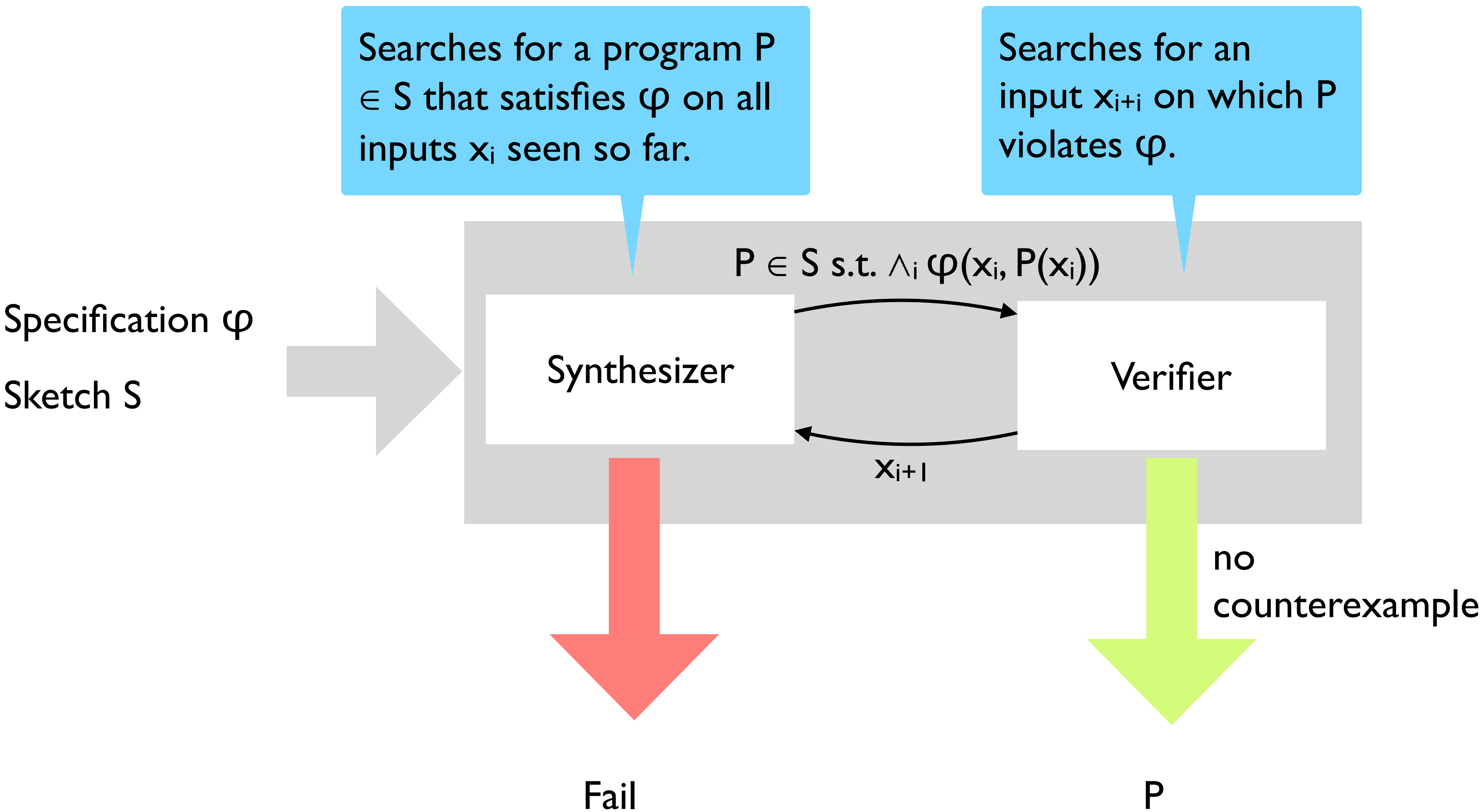

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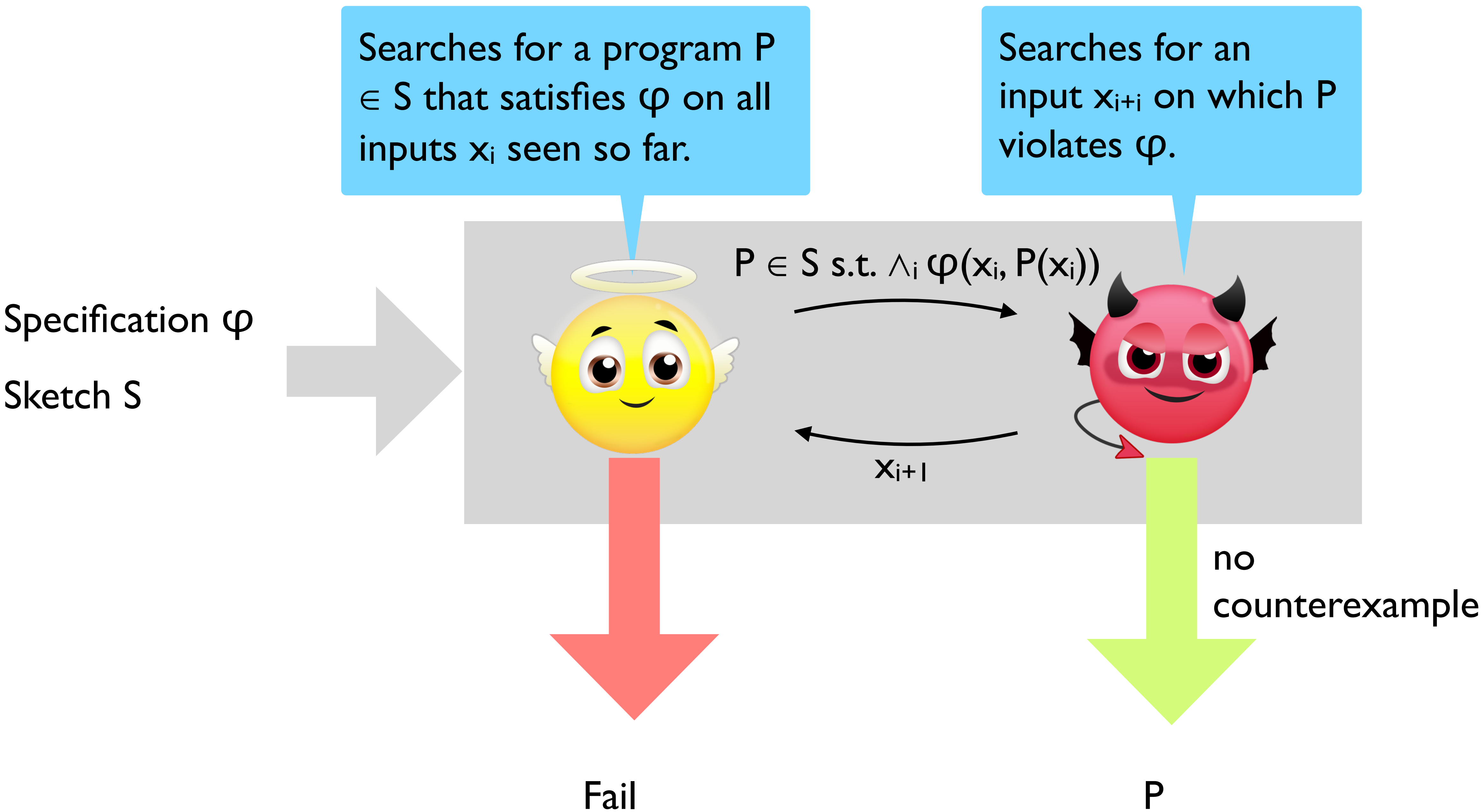

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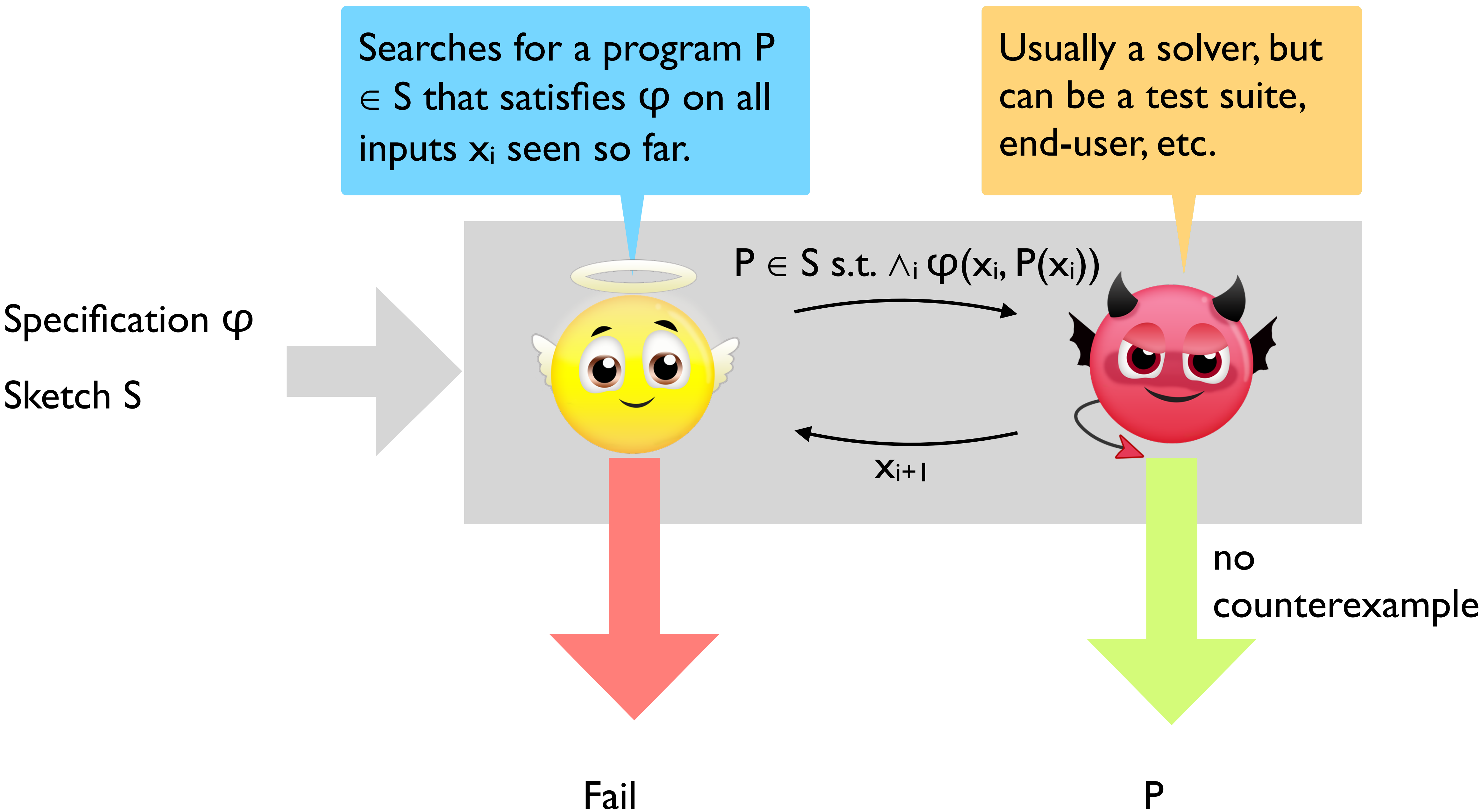

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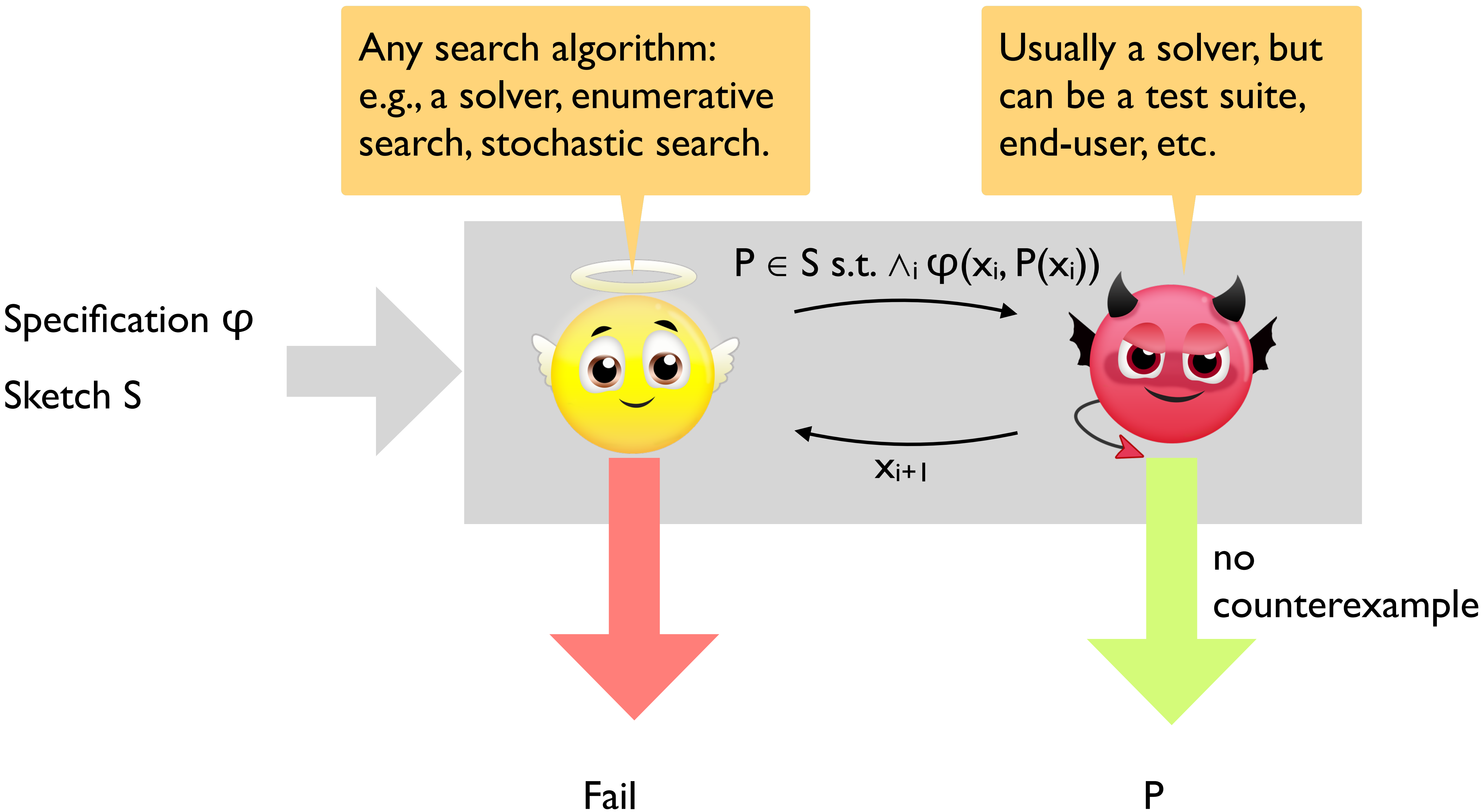
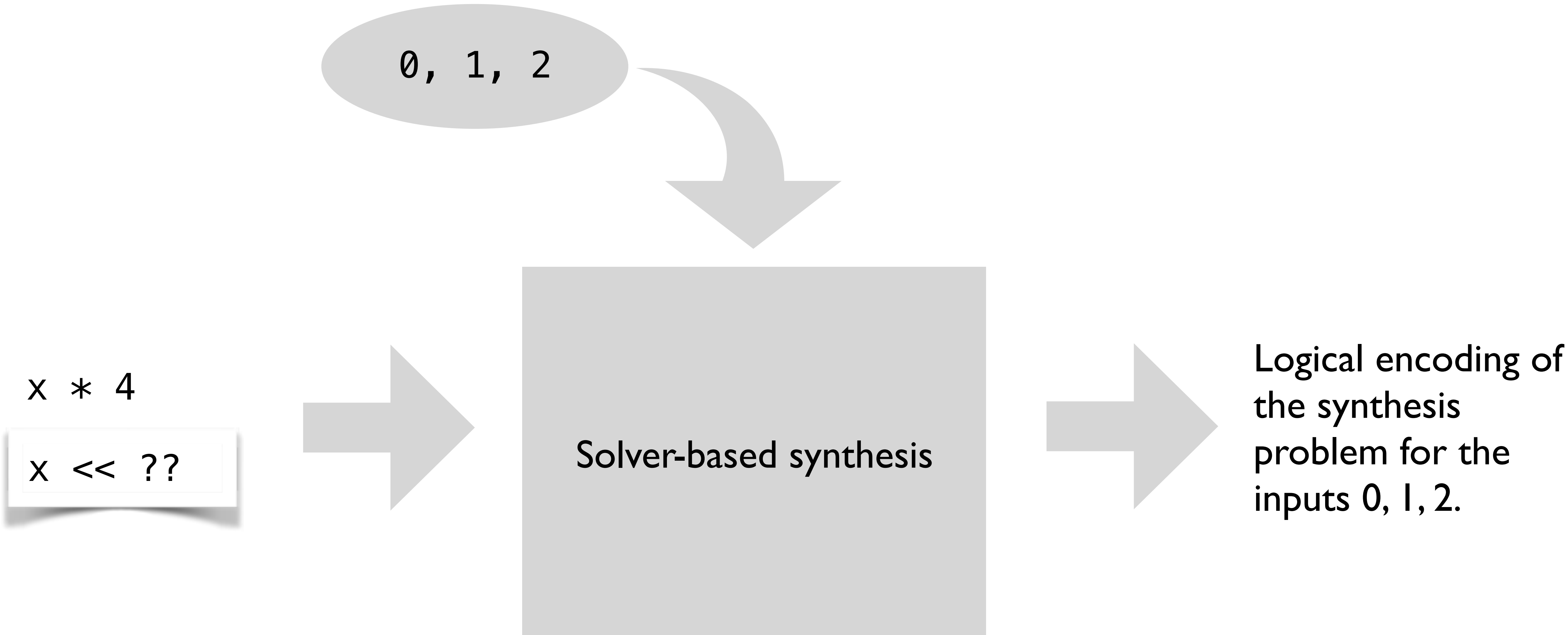
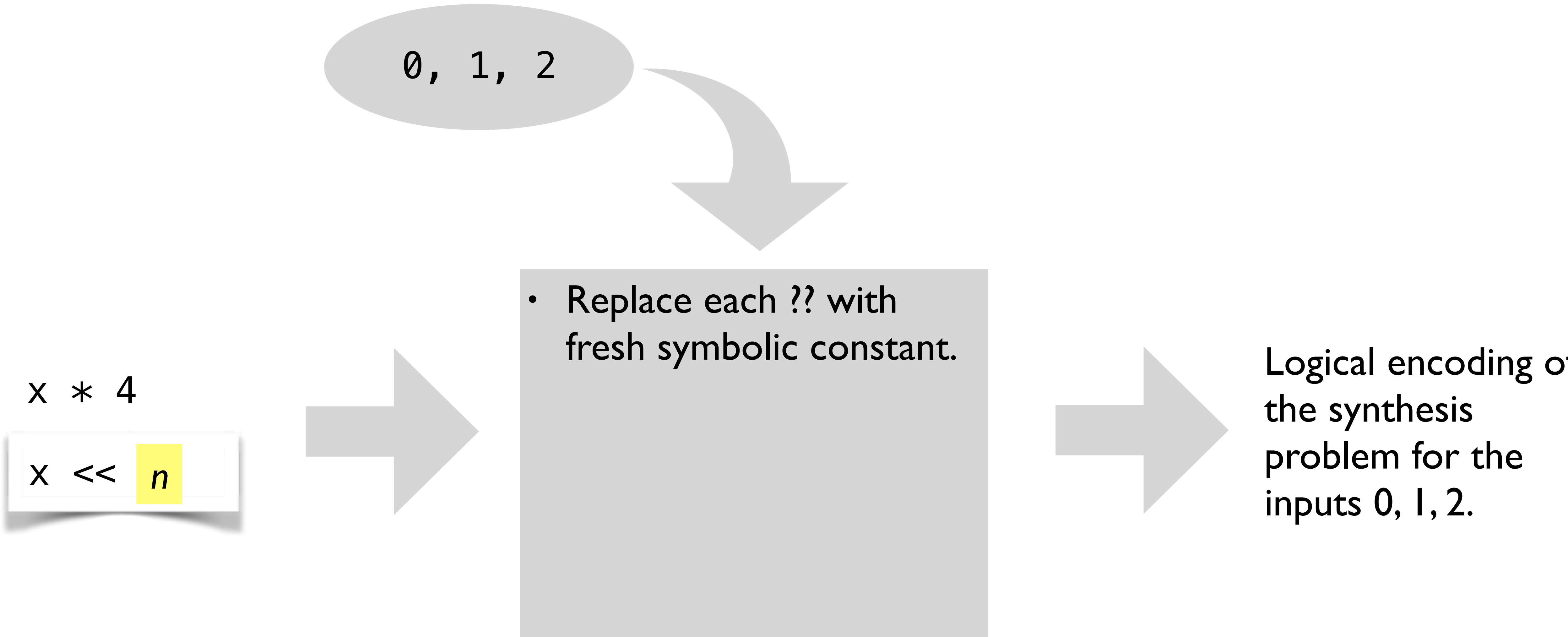


Synthesizing programs with a solver



[Solar-Lezama et al, ASPLOS'06]

Synthesizing programs with a solver



[Solar-Lezama et al, ASPLOS'06]

Synthesizing programs with a solver

`0, 1, 2`

```
x * 4
```

```
x << n
```

- Replace each ?? with fresh symbolic constant.
- Translate the resulting problem to constraints w.r.t. the current inputs.

$(0 << n = 0) \wedge$

$(1 << n = 4) \wedge$

$(2 << n = 8)$

[Solar-Lezama et al, ASPLOS'06]

Synthesizing programs with a solver

`0, 1, 2`

`x * 4`

`x << n`

- Replace each ?? with fresh symbolic constant.
- Translate the resulting problem to constraints w.r.t. the current inputs.
- If SAT, convert the model to a program P.

$(0 \ll n = 0) \land$

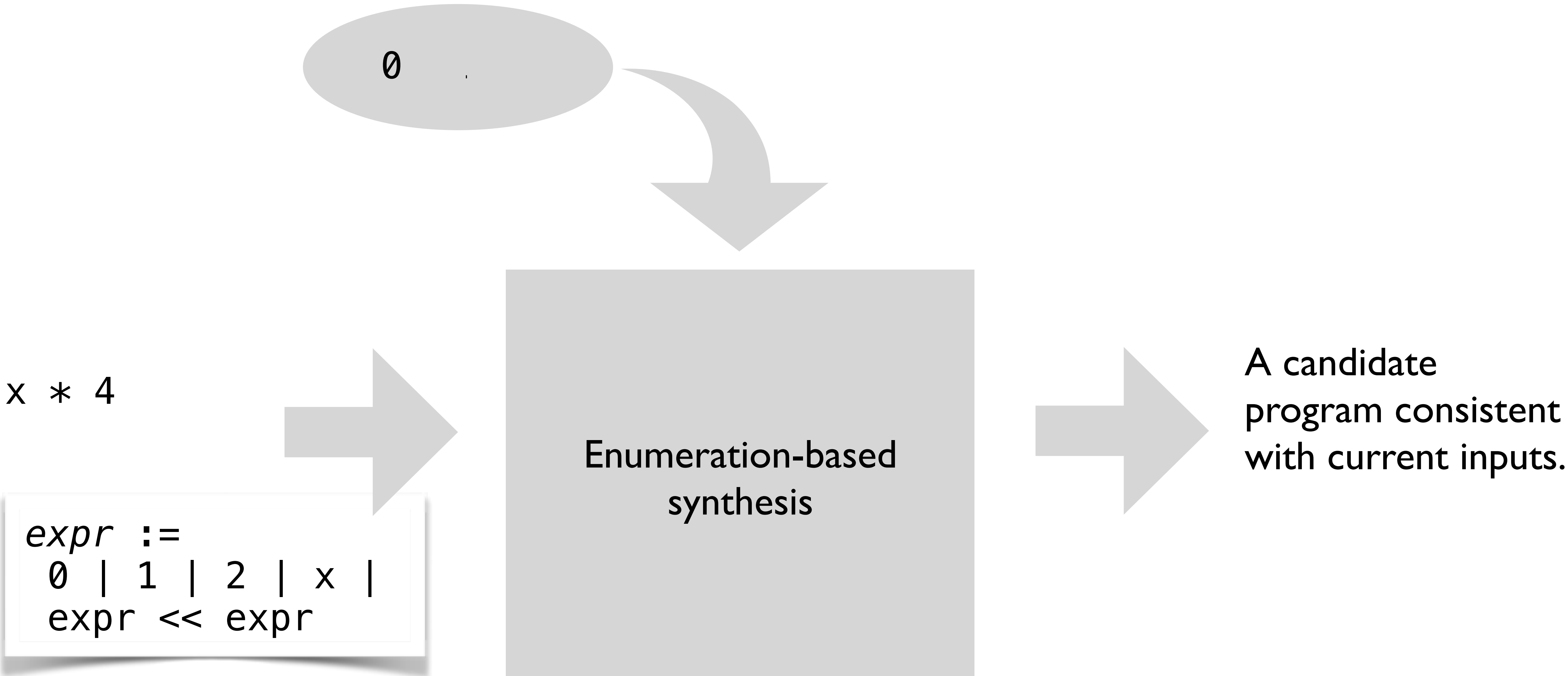
$(1 \ll n = 4) \land$

$(2 \ll n = 8)$

`x << 2`

[Solar-Lezama et al, ASPLOS'06]

Synthesizing programs with enumerative search



[Udupa et al, PLDI'13]

Synthesizing programs with enumerative search

0

x * 4

```
expr :=
 0 | 1 | 2 | x |
 expr << expr
```

- Iteratively construct all programs of size *K* until one is consistent with the current inputs.
- If two programs produce the same output on all current inputs, keep just one of the two.

A candidate program consistent with current inputs.

[Udupa et al, PLDI'13]

Synthesizing programs with enumerative search

0

x * 4

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K=*1*: 0

[Udupa et al, PLDI'13]

Synthesizing programs with enumerative search

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0, 1
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x * 4
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[Udupa et al, PLDI'13]

Synthesizing programs with enumerative search

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x * 4
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expr :=
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```

- Iteratively construct all programs of size *K* until one is consistent with the current inputs.
- If two programs produce the same output on all current inputs, keep just one of the two.

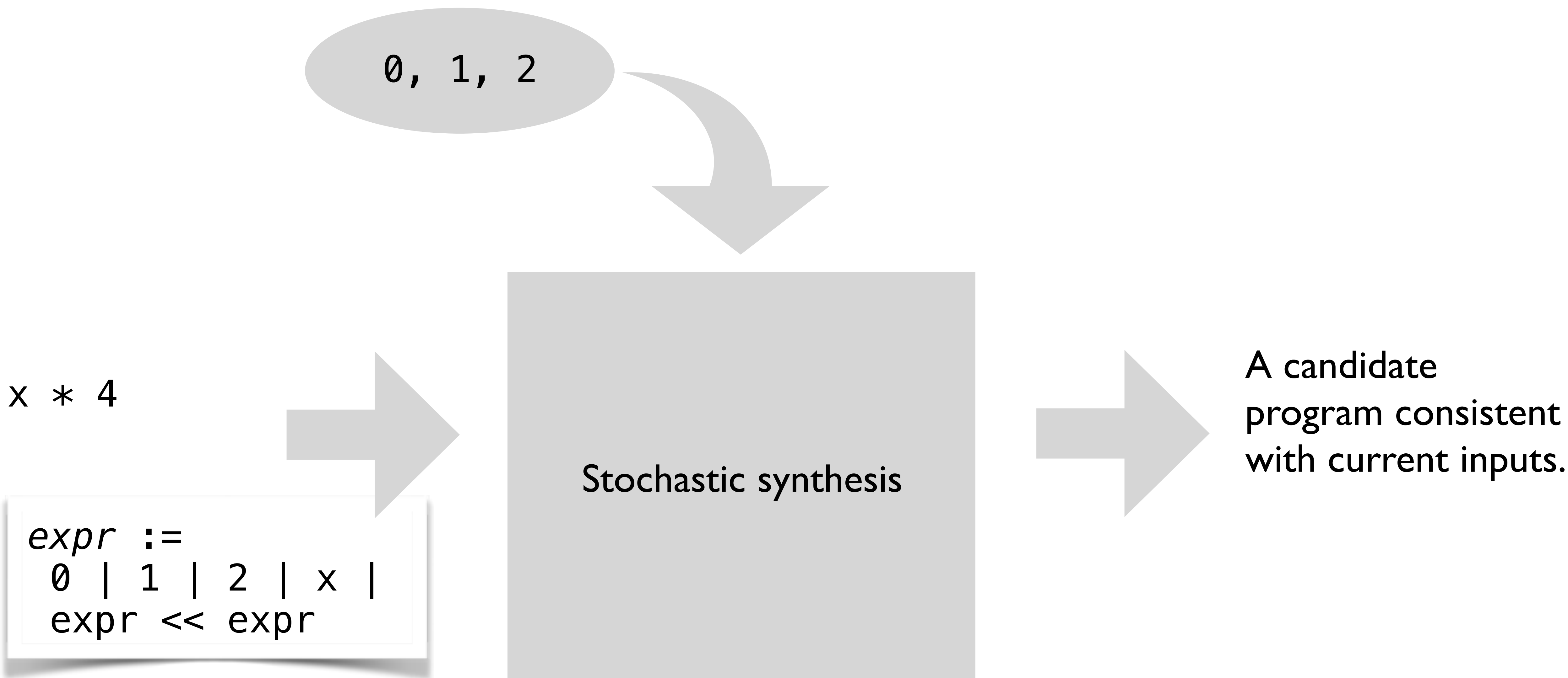
K=*1*: 0, 1, 2, x

K=2: 1 << 2, 2 << 2,

x << 1, **x << 2**

[Udupa et al, PLDI'13]

Synthesizing programs with stochastic search



[Schkufza et al, ASPLOS'13]

Synthesizing programs with stochastic search

```
0, 1, 2
```

```
x * 4
```

```
expr :=
 0 | 1 | 2 | x |
 expr << expr
```

- Use Metropolis-Hastings to sample expressions.
- Mutate the current candidate program and keep the mutation with probability proportional to its correctness w.r.t. the current inputs.

A candidate program consistent with current inputs.

[Schkufza et al, ASPLOS'13]

Summary

Today

- Deductive synthesis with axioms and E-graphs
- Inductive synthesis with solvers, enumeration, and stochastic search

Next lecture

- Solver-aided languages