

Computer-Aided Reasoning for Software

Solver-Aided Languages

courses.cs.washington.edu/courses/cse507/17wi/

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Today

Last lecture

- Program synthesis

Today

- The next N years: solver-aided languages (?)

Reminders

- Next lecture: metasketches!
- Project presentations next Friday in class
 - 11 min per team: 8 min presentation + 3 min questions
- Project reports and prototypes due next Friday at 11:00pm

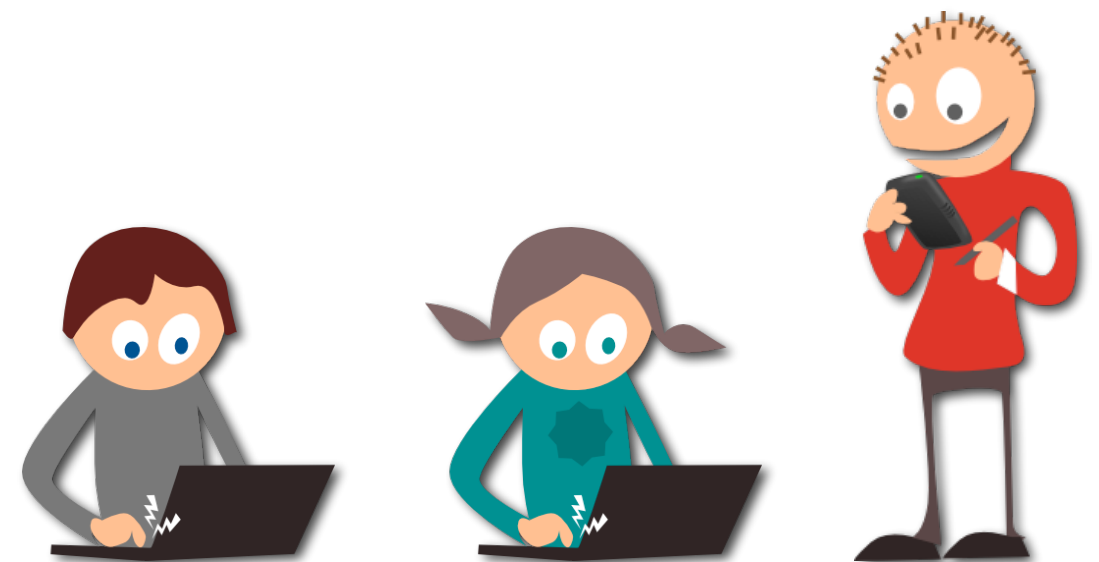


vision

a little programming for everyone

A little programming for **everyone**

Every knowledge worker wants to program ...



A little programming for **everyone**

Every knowledge worker wants to program ...

- spreadsheet data manipulation



A little programming for **everyone**

Every knowledge worker wants to program ...

- spreadsheet data manipulation
- models of cell fates



biologist



**social
scientist**

A little programming for **everyone**

Every knowledge worker wants to program ...

- spreadsheet data manipulation
- models of cell fates
- cache coherence protocols
- memory models



**hardware
designer**



biologist

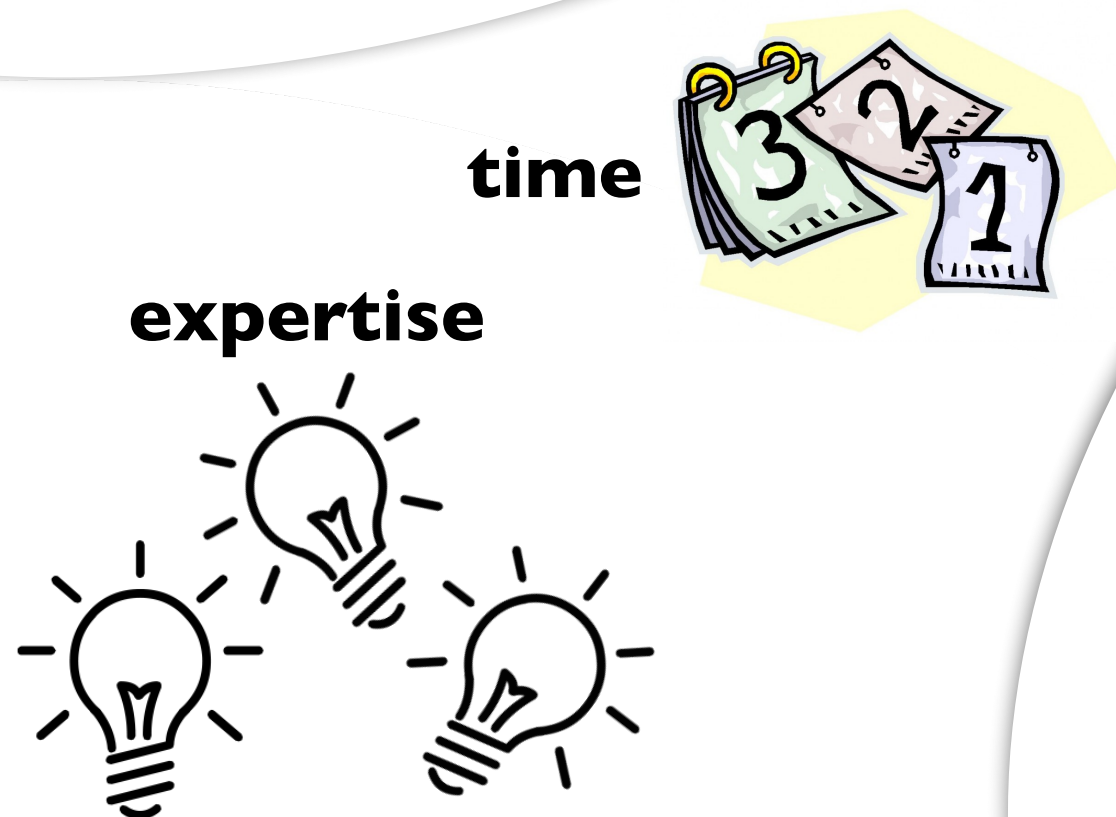


**social
scientist**

A little programming for **everyone**

Every knowledge worker wants to program ...

- ▶ spreadsheet data manipulation [Flashfill, **POPL'11**]
- ▶ models of cell fates [SBL, **POPL'13**]
- ▶ cache coherence protocols [Transit, **PLDI'13**]
- ▶ memory models [MemSAT, **PLDI'10**]



hardware
designer



biologist



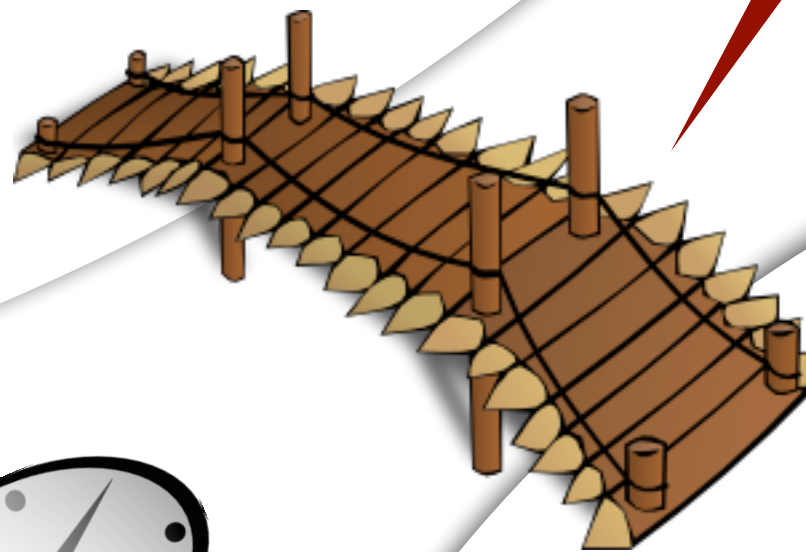
social
scientist

A **little** programming for everyone

We all want to build programs ...

- spreadsheet data manipulation
- models of cell fates
- cache coherence protocols
- memory models

solver-aided languages



less time



less expertise



**hardware
designer**



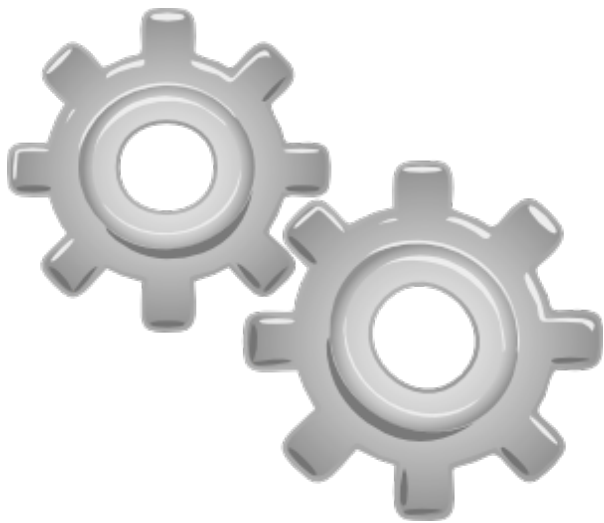
biologist



**social
scientist**

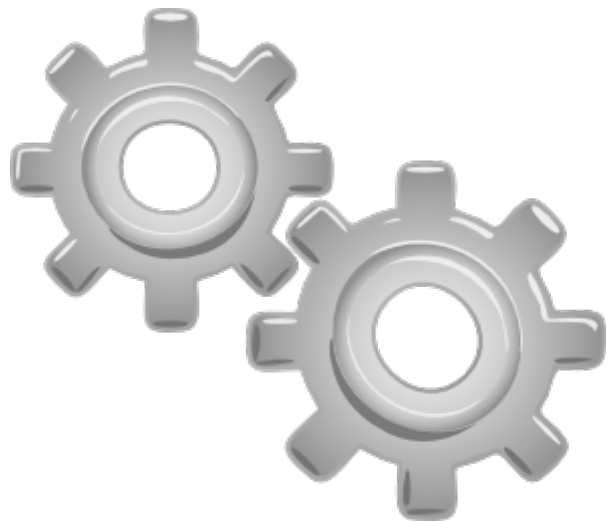
outline

solver-aided tools



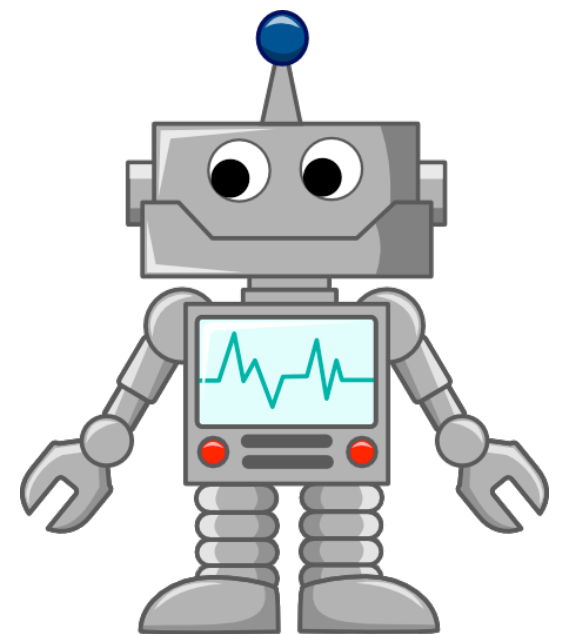
outline

solver-aided tools, languages



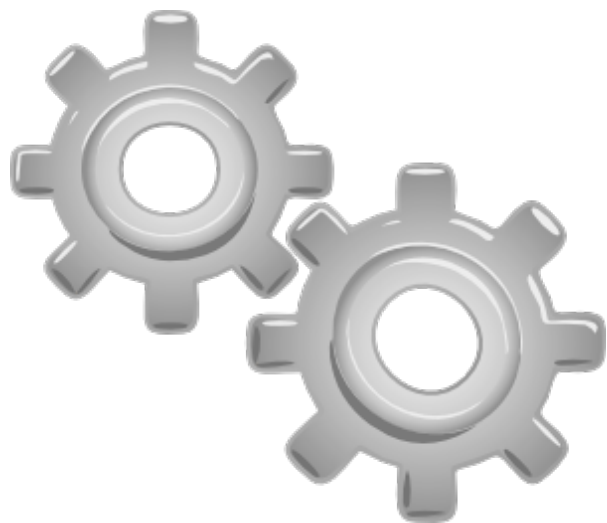
outline

solver-aided tools, languages, and applications



tools

solver-aided tools



Programming ...

specification

```
P(x) {  
  ...  
  ...  
}
```



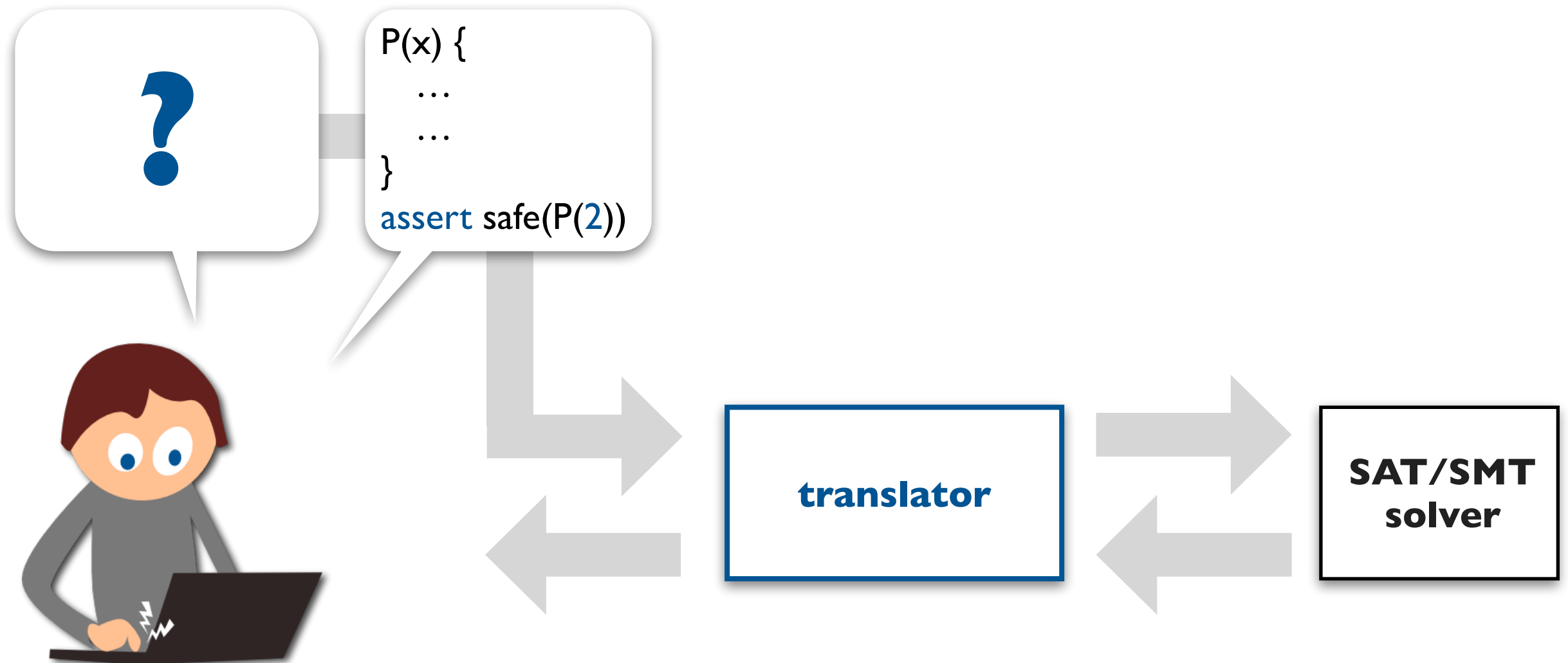
Programming ...

test case

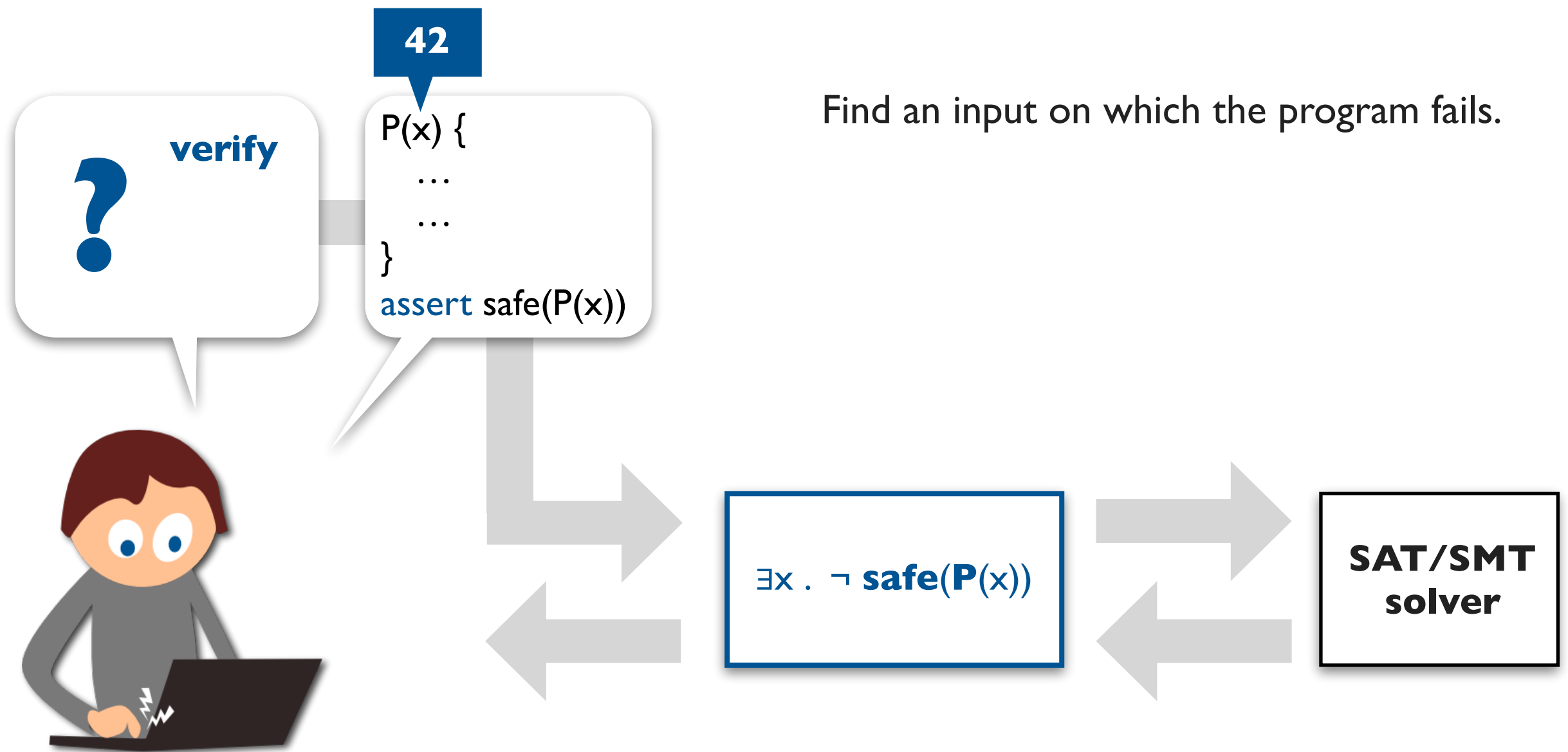
```
P(x) {  
  ...  
  ...  
}  
assert safe(P(2))
```



Programming with a solver-aided tool

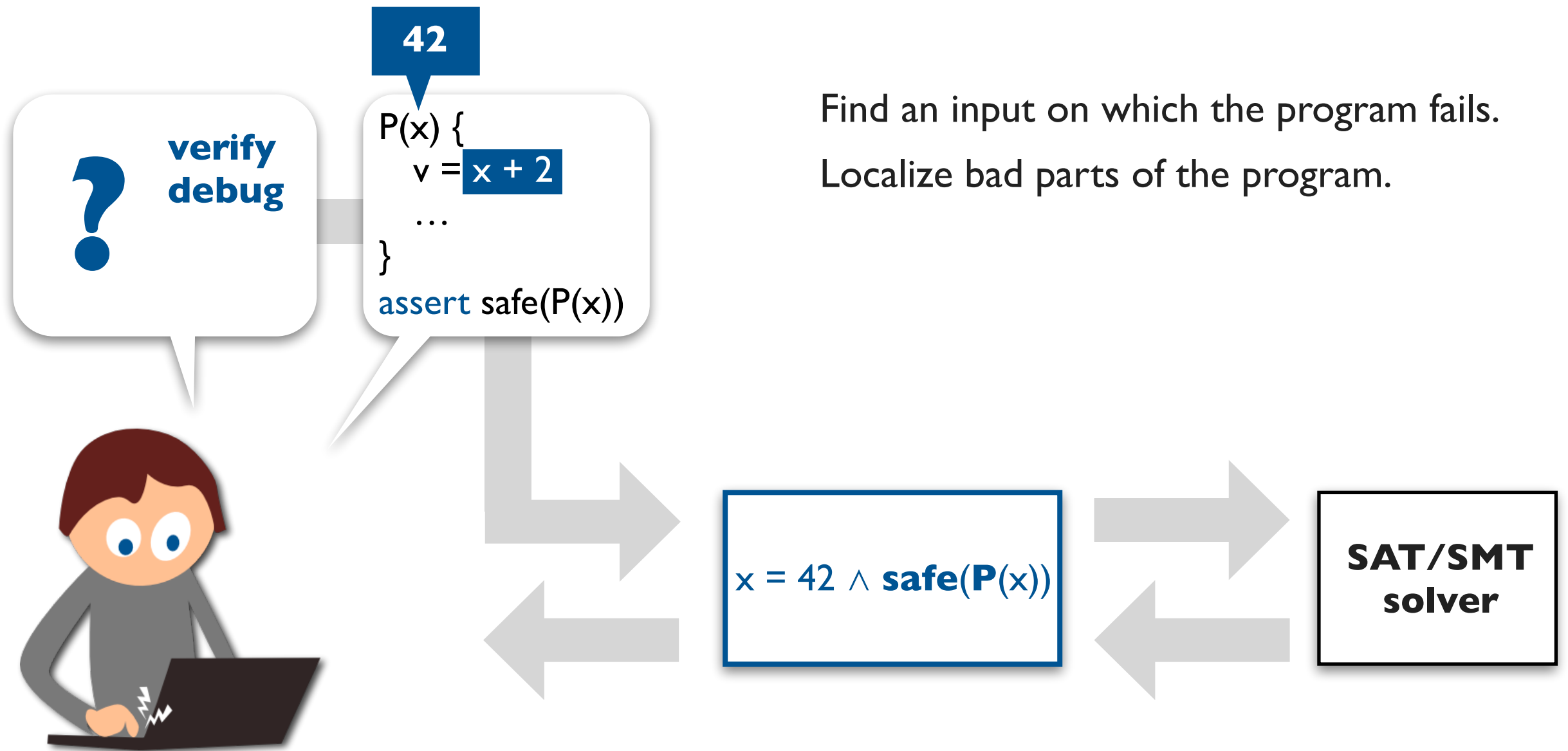


Programming with a solver-aided tool

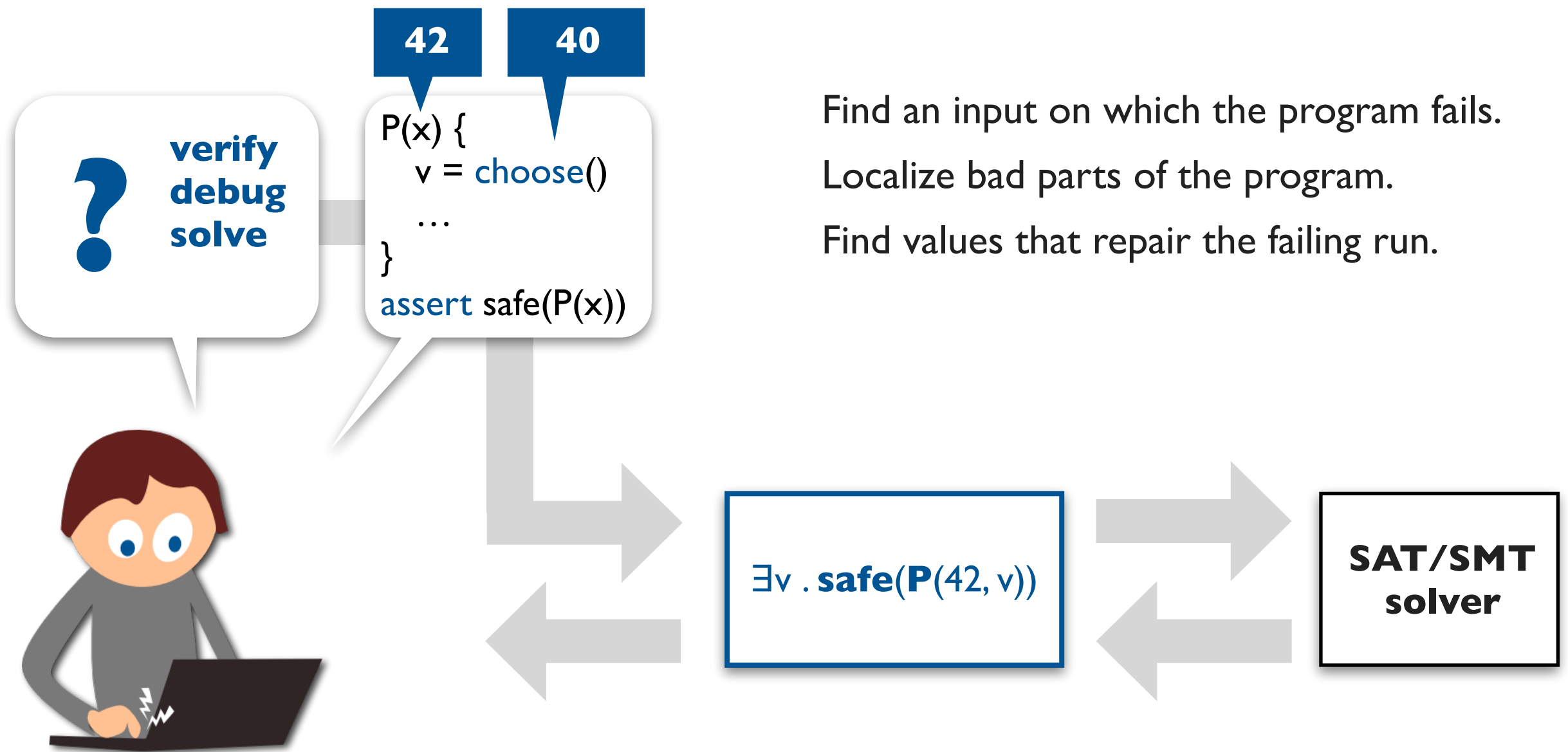


CBMC [Kroening et al., DAC'03]
Dafny [Leino, LPAR'10]
Miniatur [Vaziri et al., FSE'07]
Klee [Cadar et al., OSDI'08]

Programming with a solver-aided tool

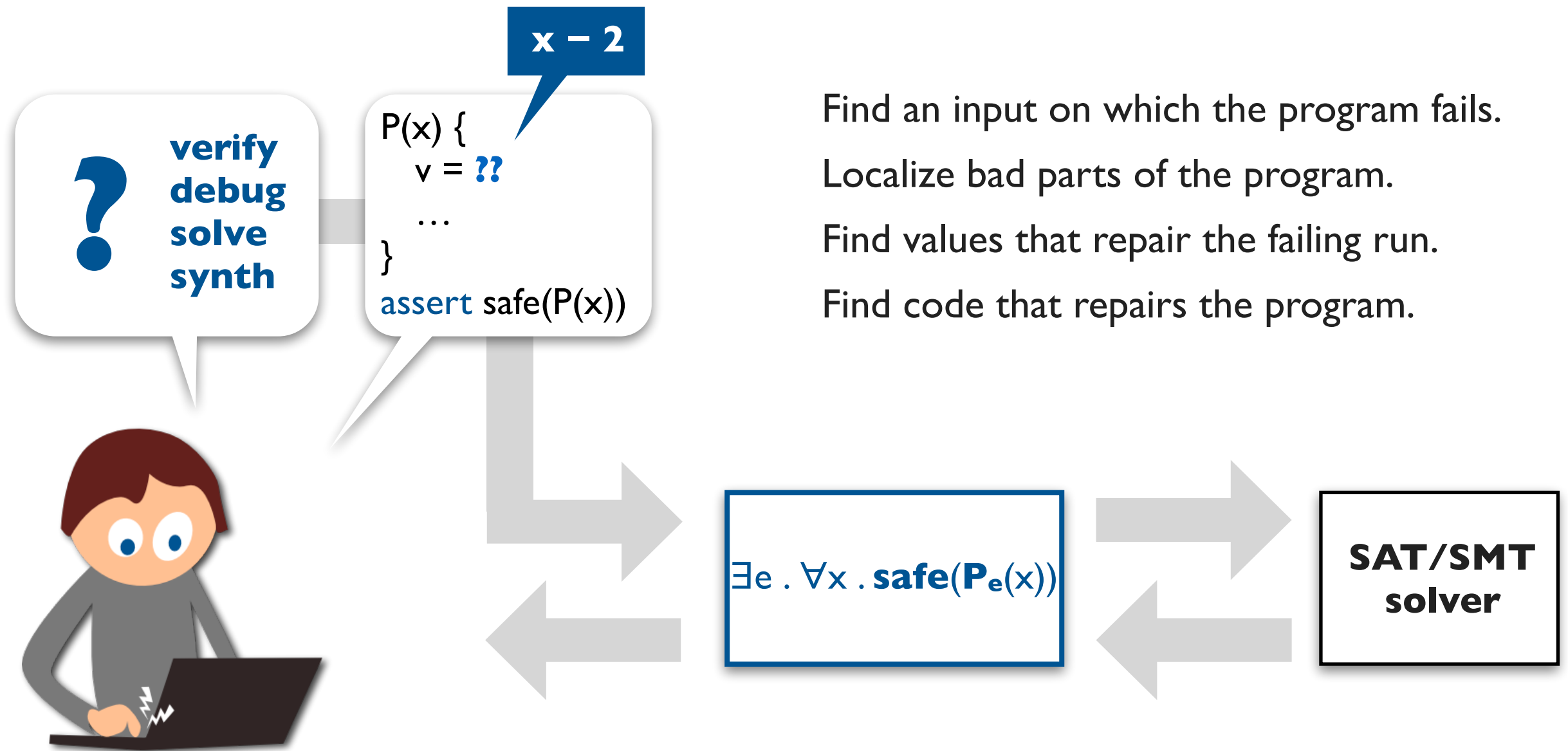


Programming with a solver-aided tool



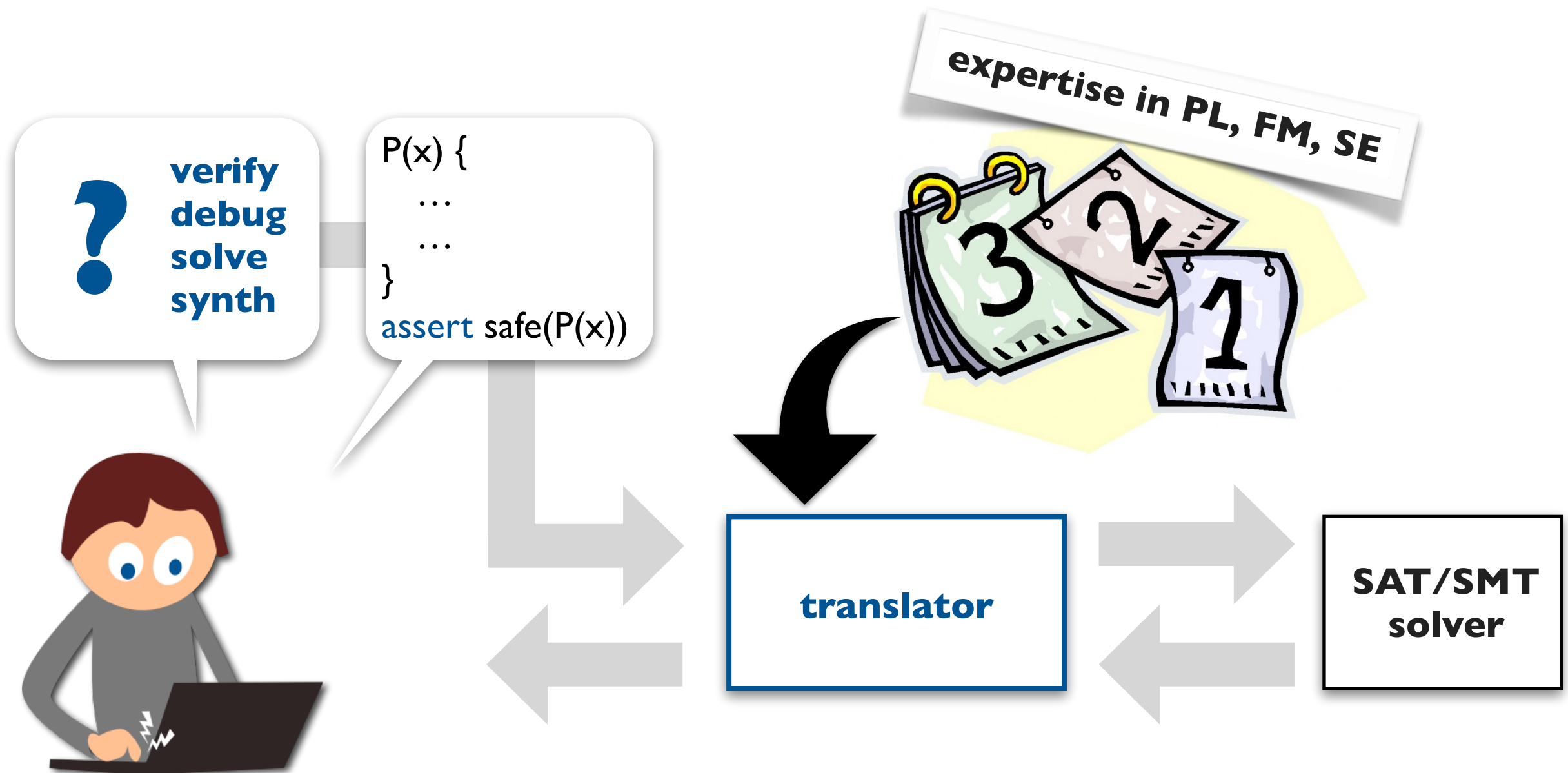
Kaplan [Koksal et al, POPL'12]
PBnJ [Samimi et al., ECOOP'10]
Squander [Milicevic et al., ICSE'11]

Programming with a solver-aided tool

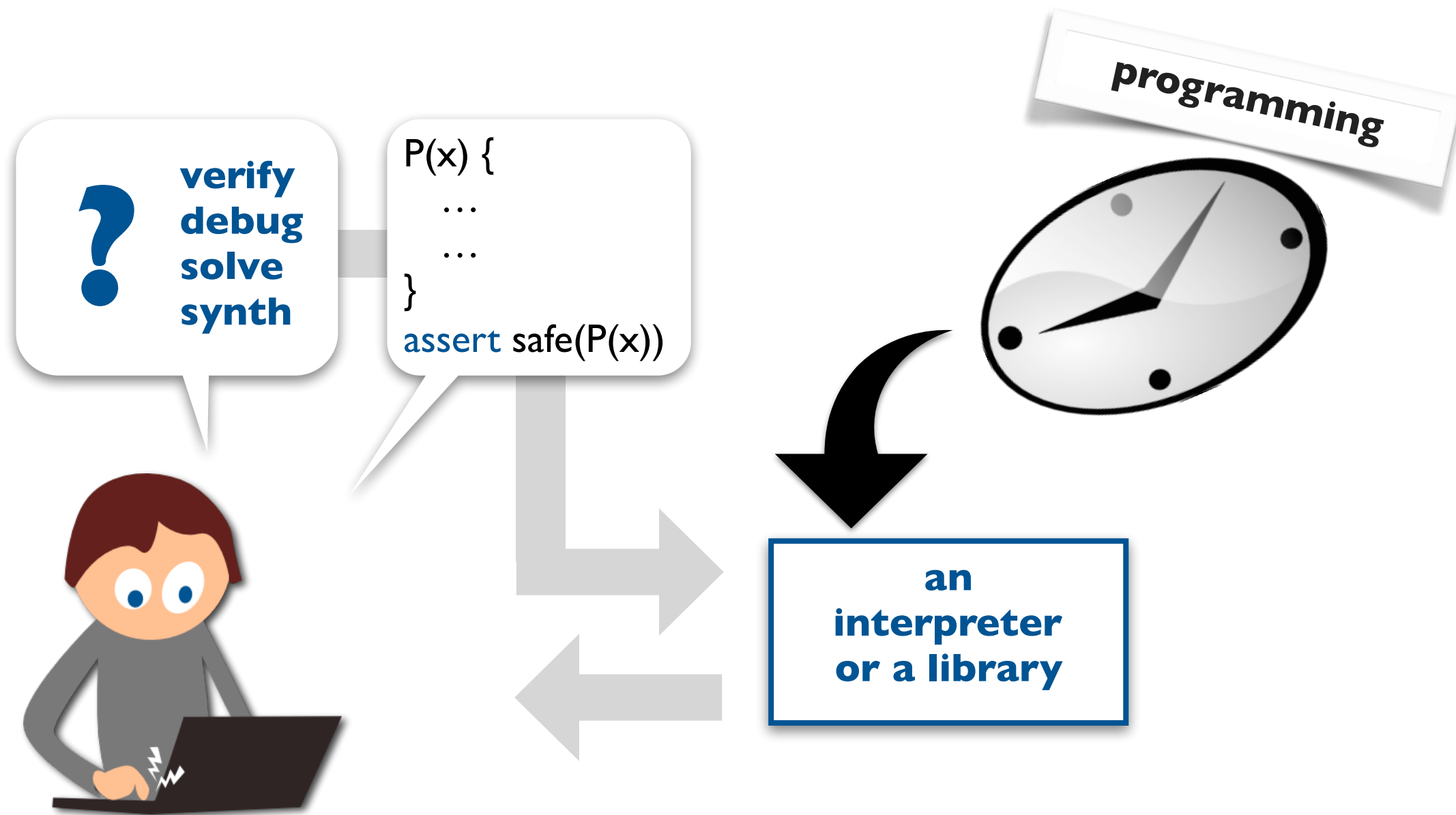


Sketch [Solar-Lezama et al., ASPLOS'06]
Comfusy [Kuncak et al., CAV'10]

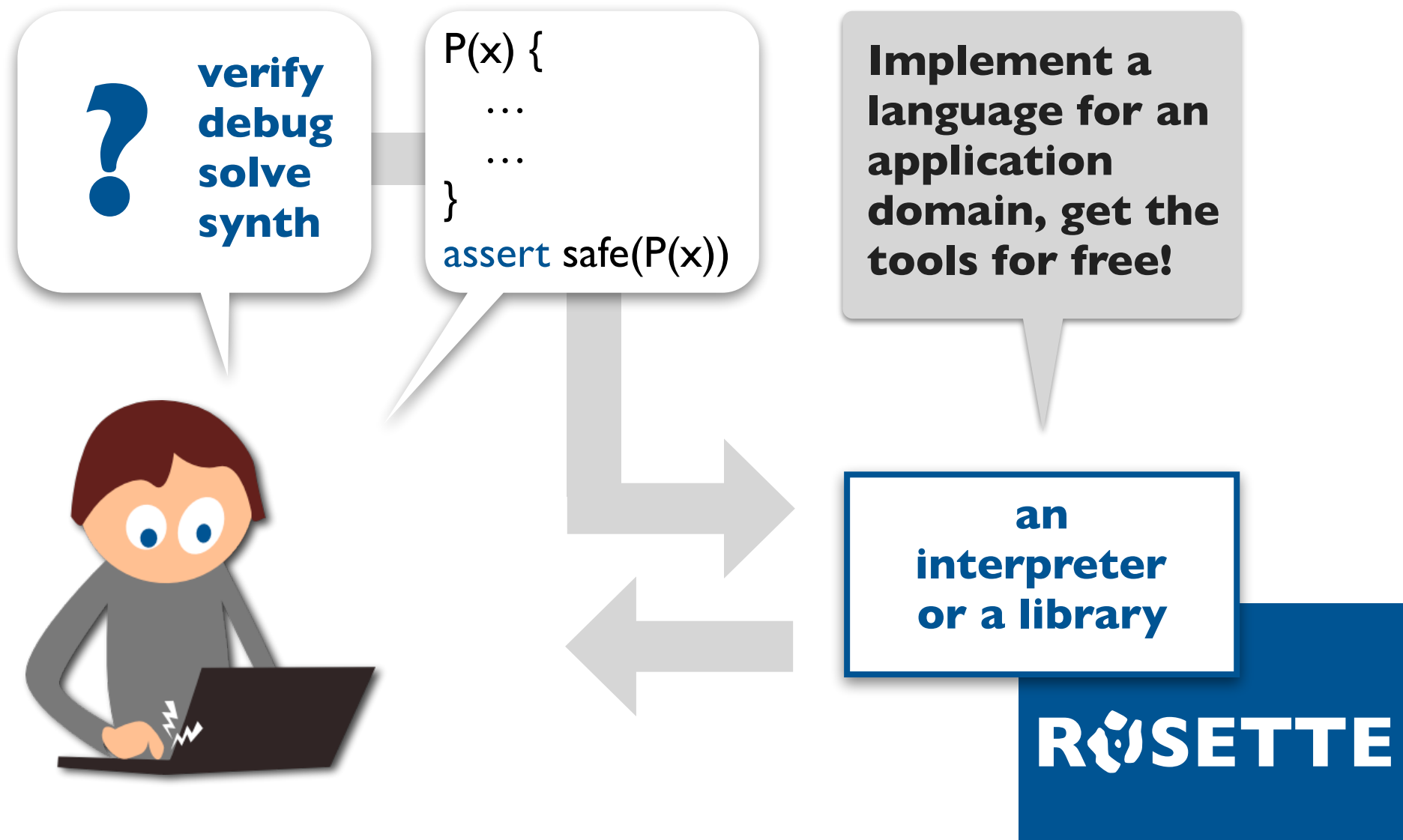
The standard (hard) way to build a tool



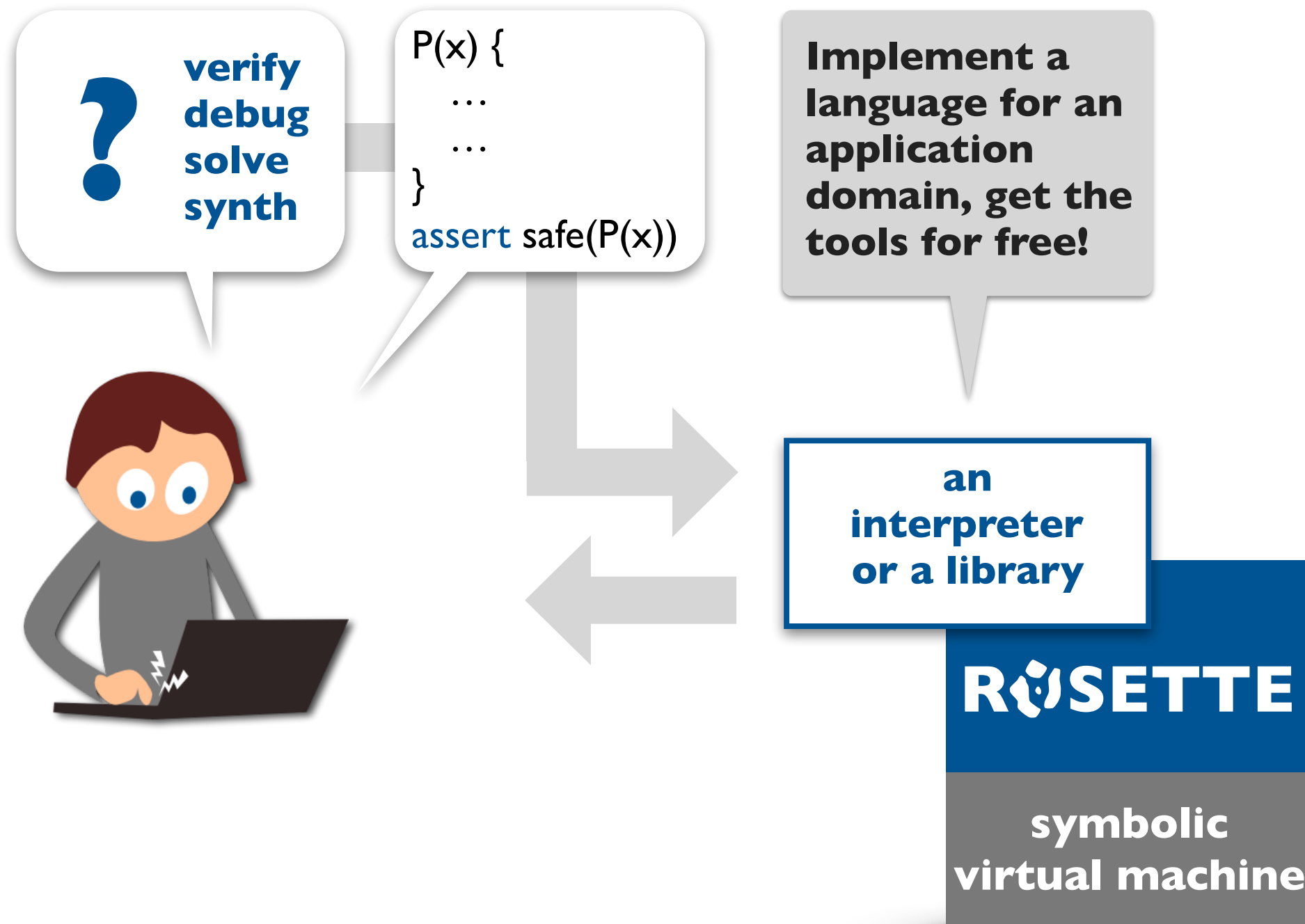
A new, easy way to build tools



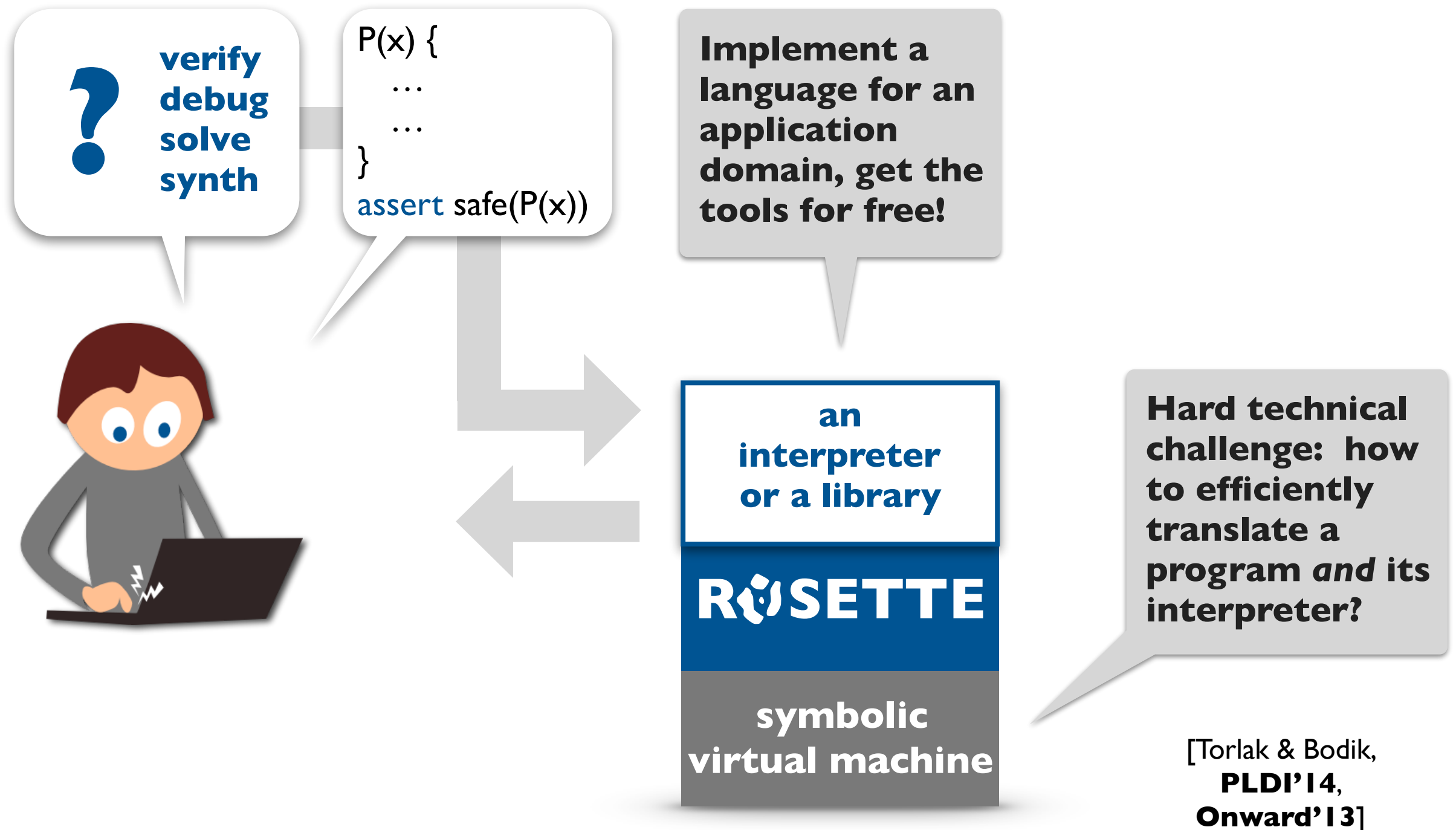
A new, easy way to build tools



A new, easy way to build tools



A new, easy way to build tools

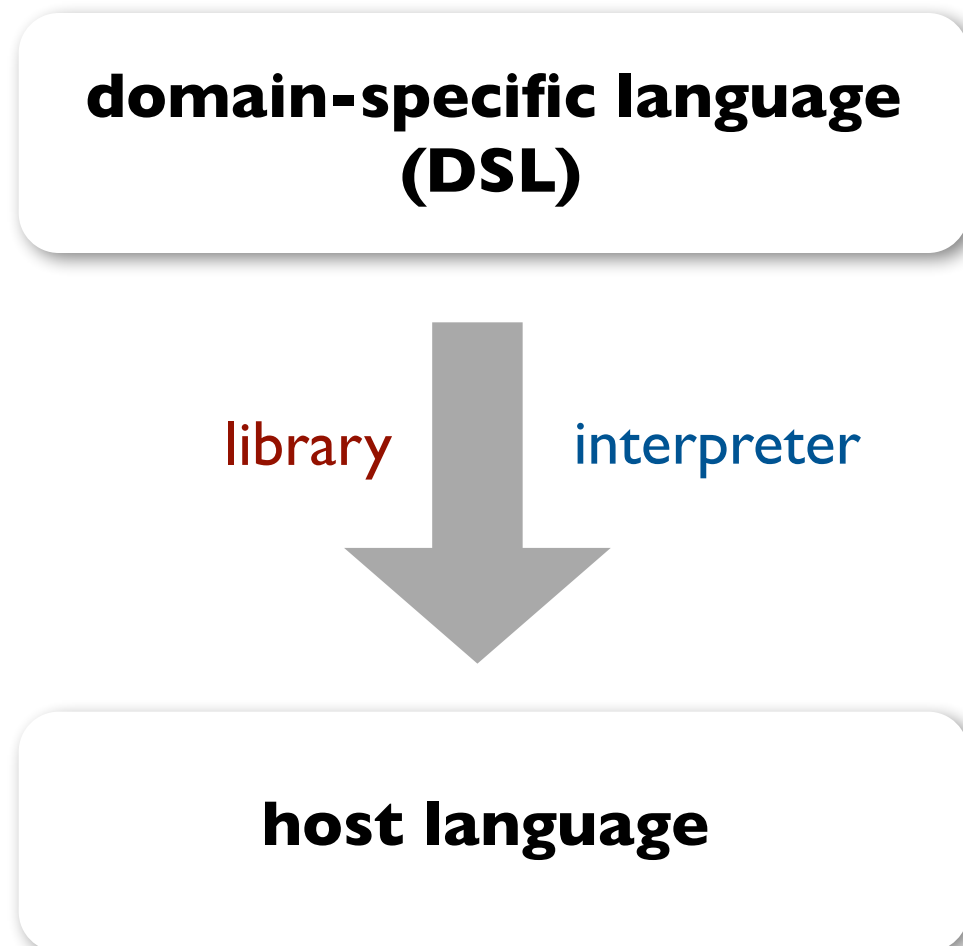


design

solver-aided languages



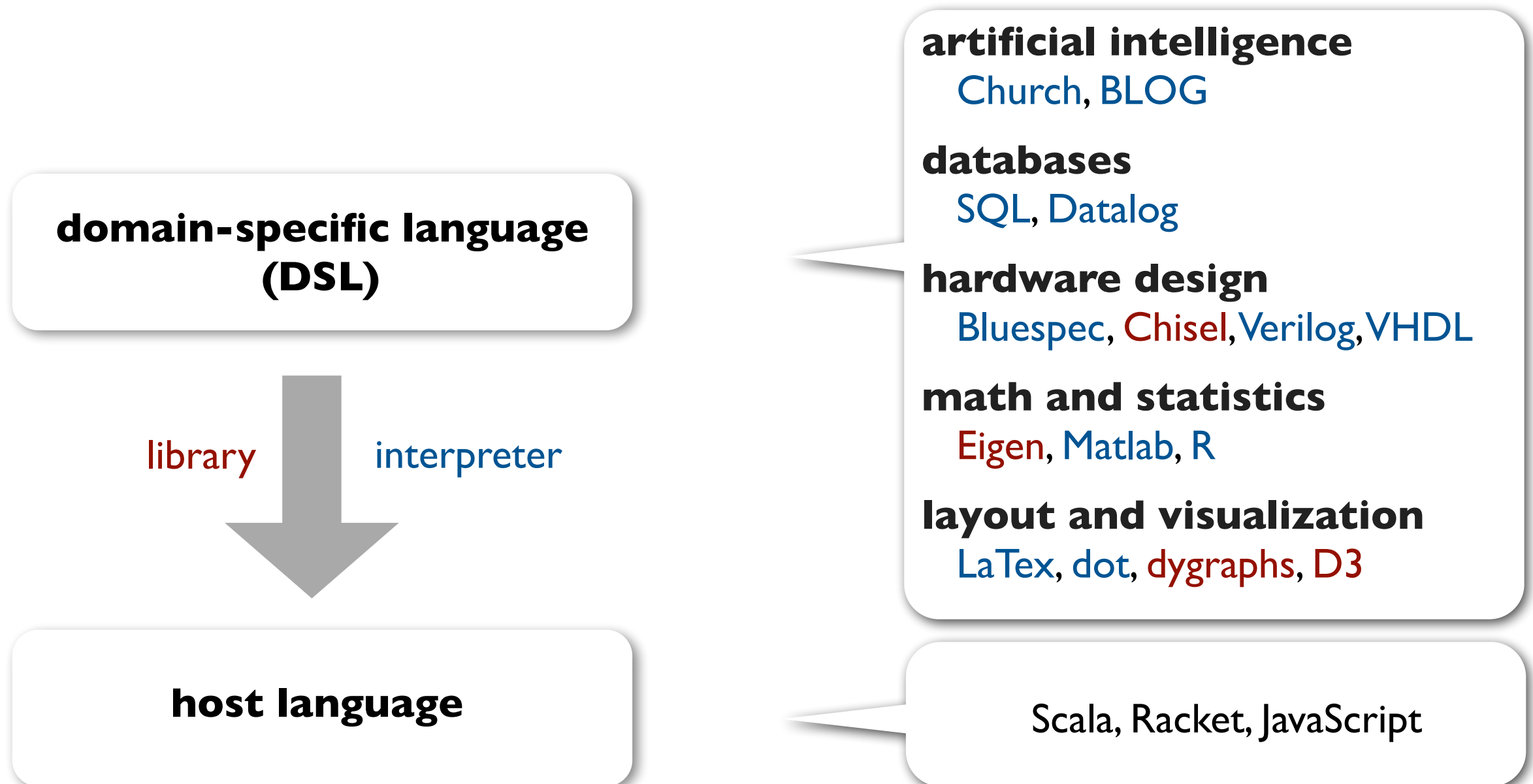
Layers of languages



A formal language that is specialized to a particular application domain and often limited in capability.

A high-level language for implementing DSLs, usually with meta-programming features.

Layers of languages



Layers of languages

**domain-specific language
(DSL)**

library

interpreter

host language

Eigen / Matlab

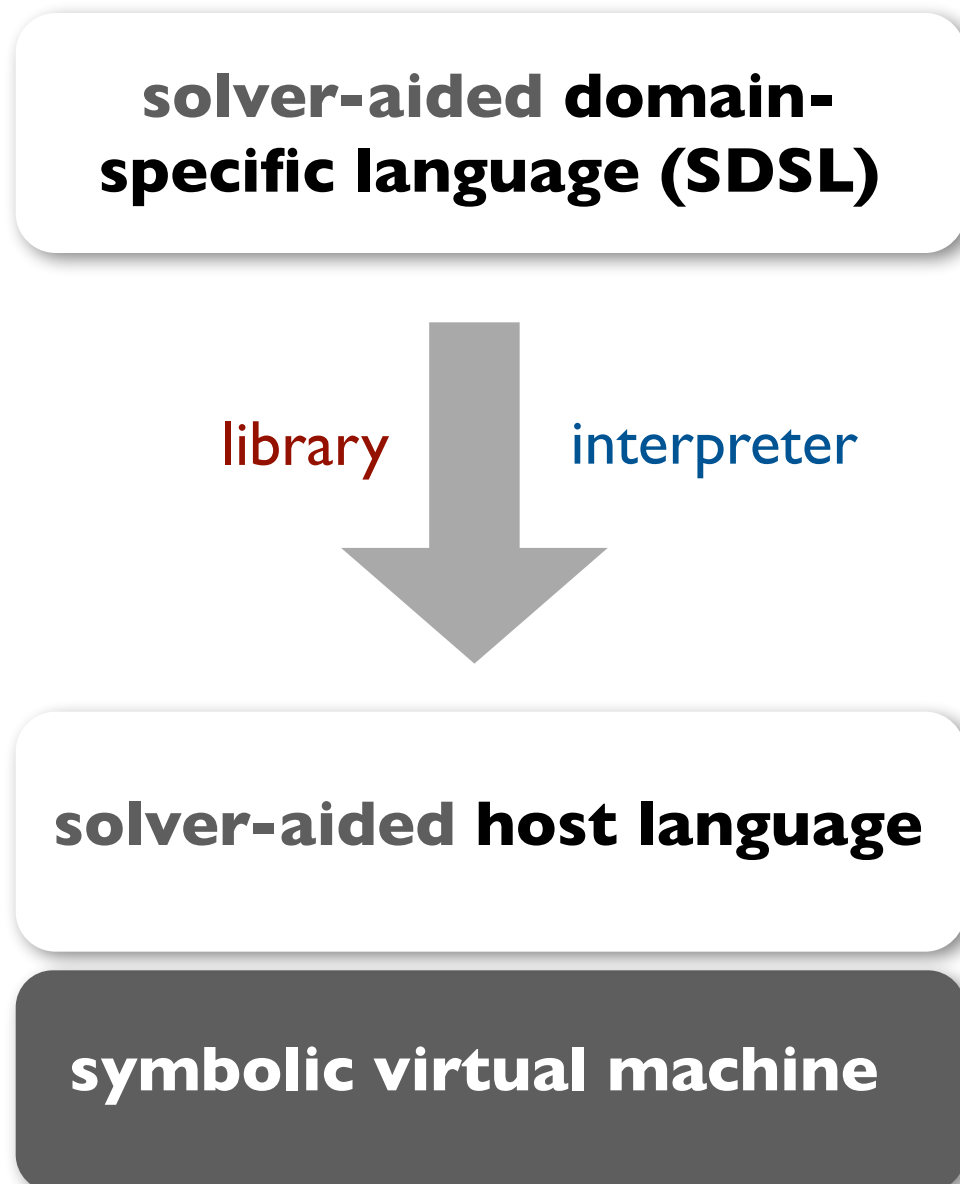
$C = A * B$

[associativity]

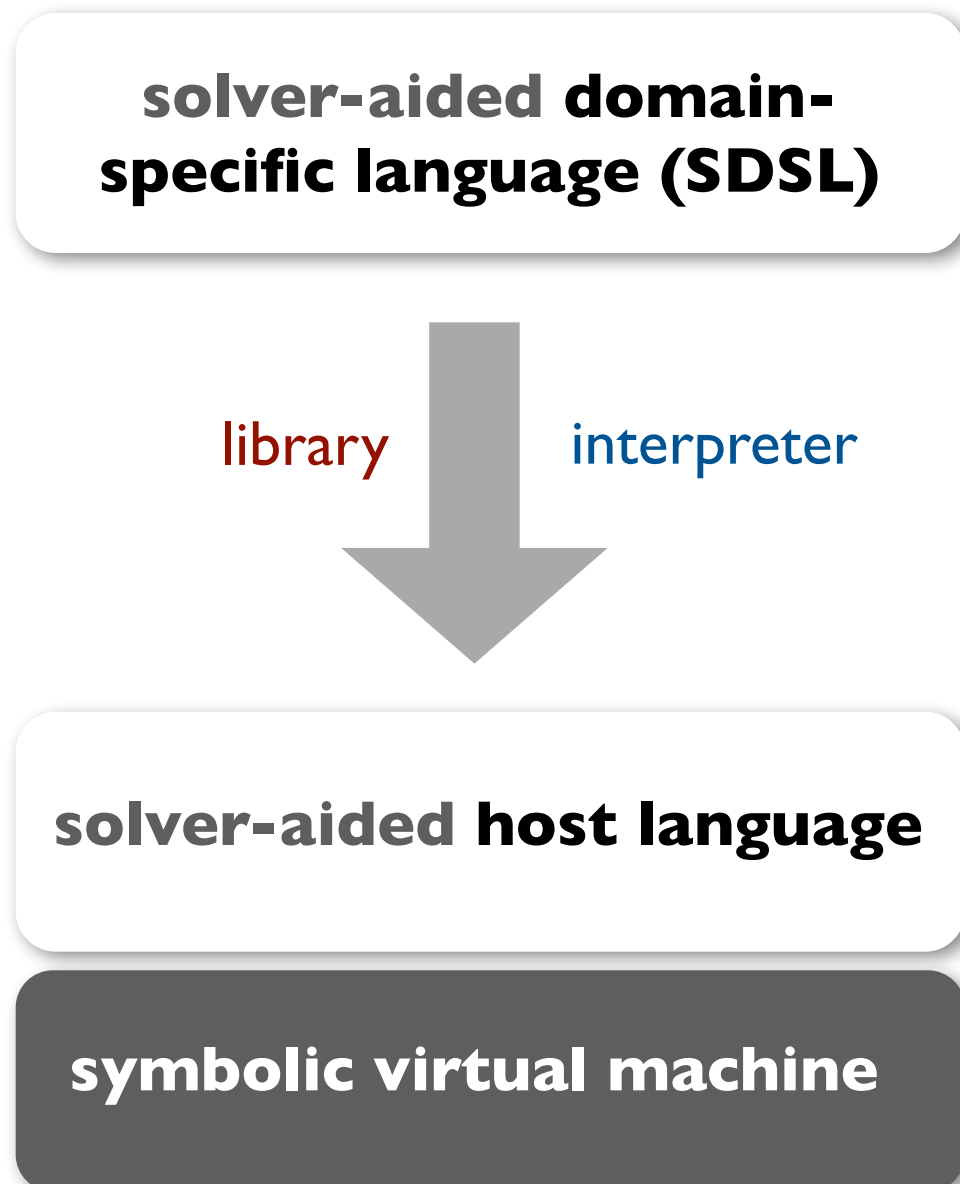
C / Java

```
for (i = 0; i < n; i++)  
  for (j = 0; j < m; j++)  
    for (k = 0; k < p; k++)  
      C[i][k] += A[i][j] * B[j][k]
```

Layers of solver-aided languages

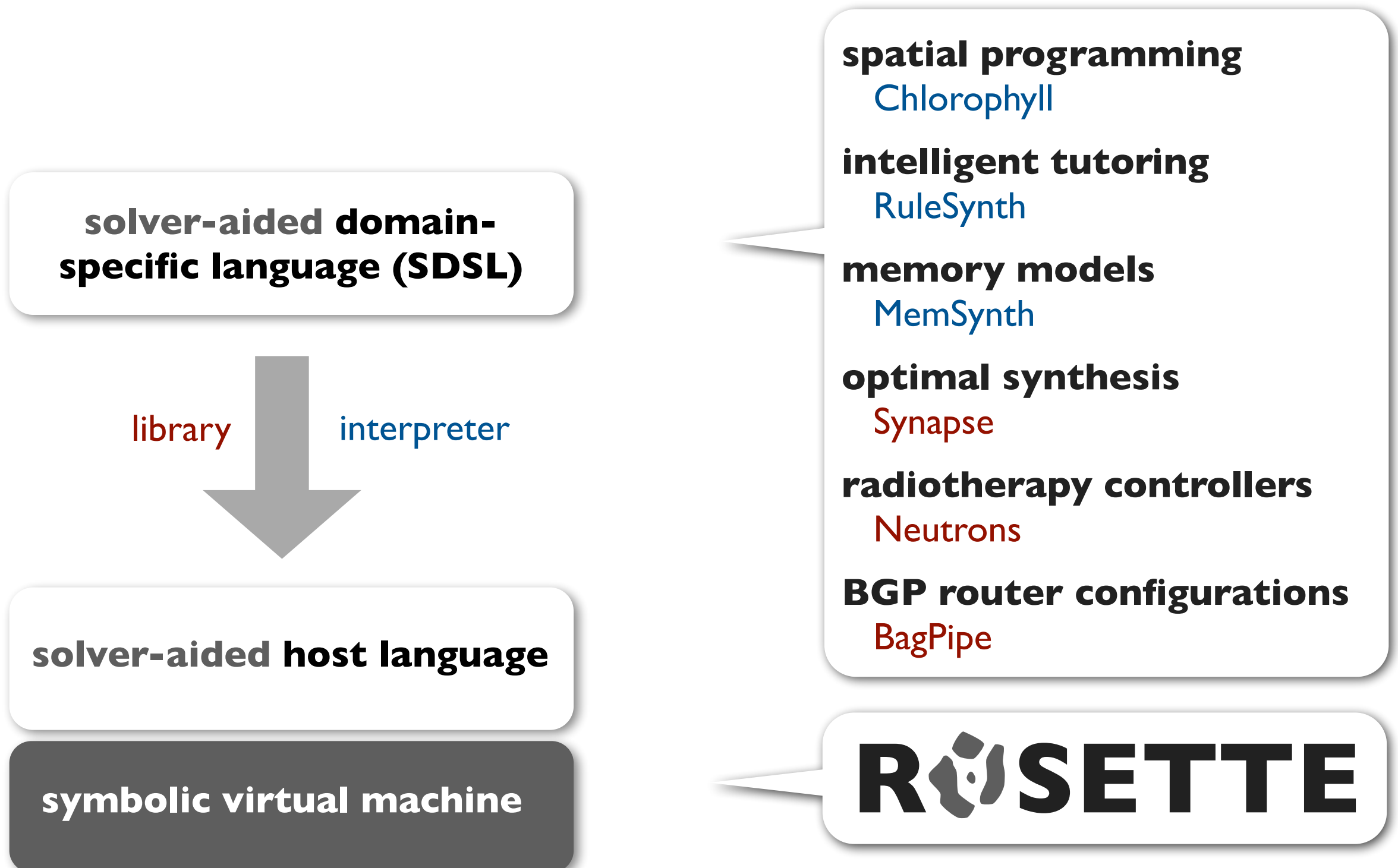


Layers of solver-aided languages



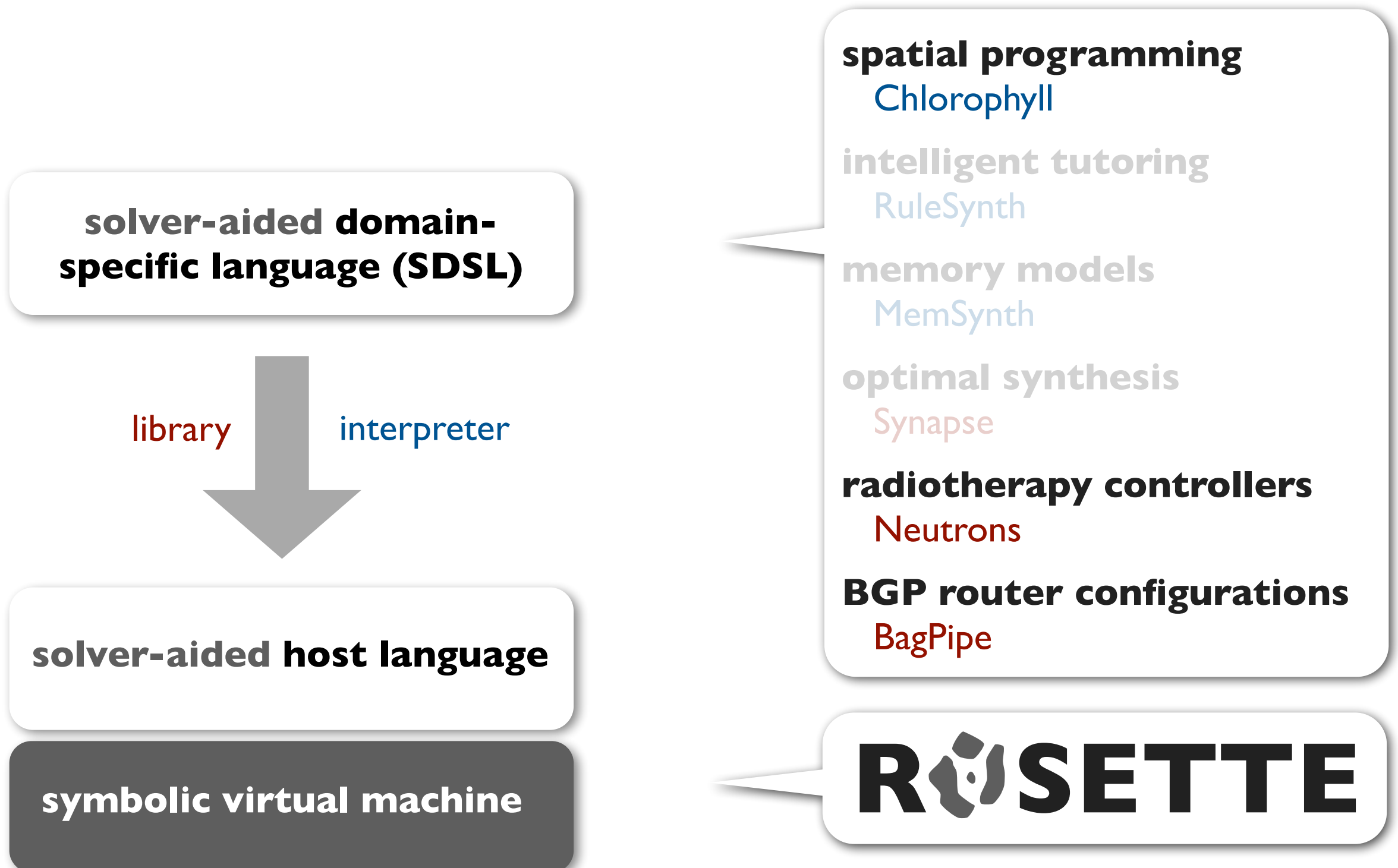
[Torlak & Bodik, **Onward'13**, **PLDI'14**]

Layers of solver-aided languages



[Torlak & Bodik, **Onward'13**, **PLDI'14**]

Layers of solver-aided languages



[Torlak & Bodik, **Onward'13**, **PLDI'14**]

Anatomy of a solver-aided host language

Modern descendent of
Scheme with macro-based
metaprogramming.



Racket

Anatomy of a solver-aided host language

```
(define-symbolic id type)
```

```
(assert expr)
```

```
(verify expr)
```

```
(debug [expr] expr)
```

```
(solve expr)
```

```
(synthesize [expr] expr)
```



ROSETTE

A tiny example SDSL

```
def bvmax(r0, r1) :  
    r2 = bvge(r0, r1)  
    r3 = bvneg(r2)  
    r4 = bvxor(r0, r2)  
    r5 = bvand(r3, r4)  
    r6 = bvxor(r1, r5)  
    return r6
```

BV: A tiny assembly-like language for writing fast, low-level library functions.

A tiny example SDSL

```
def bvmax(r0, r1) :  
    r2 = bvge(r0, r1)  
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    return r6
```

test debug
verify synth

BV: A tiny assembly-like language for writing fast, low-level library functions.

A tiny example SDSL

```
def bvmax(r0, r1) :  
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    return r6
```

test debug
verify synth

BV: A tiny assembly-like language for writing fast, low-level library functions.

- | | | |
|----|-------------|----------|
| 1. | interpreter | [10 LOC] |
| 2. | verifier | [free] |
| 3. | debugger | [free] |
| 4. | synthesizer | [free] |

A tiny example SDSL: RÖSETTE

```
def bvmax(r0, r1) :  
    r2 = bvge(r0, r1)  
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    r4 = bvxor(r0, r2)  
    r5 = bvand(r3, r4)  
    r6 = bvxor(r1, r5)  
    return r6
```

```
> bvmax(-2, -1)
```

A tiny example SDSL: **ROSETTE**

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
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```

```
> bvmax(-2, -1)
```

parse

```
(define bvmax  
  `((2 bvge 0 1)  
    (3 bvneg 2)  
    (4 bvxor 0 2)  
    (5 bvand 3 4)  
    (6 bvxor 1 5)))
```

A tiny example SDSL: SETTE

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
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  r4 = bvxor(r0, r2)  
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  return r6
```

```
> bvmax(-2, -1)
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```

(out opcode in ...)

A tiny example SDSL: **ROSETTE**

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```
> bvmax(-2, -1)
```

interpret

```
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    (6 bvxor 1 5)))
```

```
`(-2 -1)
```

```
(define (interpret prog inputs)  
  (make-registers prog inputs)  
  (for ([stmt prog])  
    (match stmt  
      [(list out opcode in ...)  
       (define op (eval opcode))  
       (define args (map load in))  
       (store out (apply op args))]))  
  (load (last)))
```

A tiny example SDSL:

ROSETTE

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0	-2
1	-1
2	
3	
4	
5	
6	

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A tiny example SDSL:

ROSETTE

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ROSETTE

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3	
4	
5	
6	

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A tiny example SDSL:

ROSETTE

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```

```
> bvmax(-2, -1)  
-1
```

interpret

```
(define bvmax  
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    (3 bvneg 2)  
    (4 bvxor 0 2)  
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```

0	-2
1	-1
2	0
3	0
4	-2
5	0
6	-1

```
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A tiny example SDSL:

ROSETTE

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```
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-1
```

```
(define bvmax  
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    (4 bvxor 0 2)  
    (5 bvand 3 4)  
    (6 bvxor 1 5)))
```

- ▶ pattern matching
- ▶ dynamic evaluation
- ▶ first-class & higher-order procedures
- ▶ side effects

```
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       (define op (eval opcode))  
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A tiny example SDSL: **ROSETTE**

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  r6 = bvxor(r1, r5)  
  return r6
```

```
> verify(bvmax, max)
```

query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(verify  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```


A tiny example SDSL: **ROSETTE**

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(r0, r2)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6
```

```
> verify(bvmax, max)
```

query

Creates two fresh symbolic constants of type number and binds them to variables n0 and n1.

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(verify  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```

A tiny example SDSL: ROSETTE

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
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  r6 = bvxor(r1, r5)  
  return r6
```

```
> verify(bvmax, max)
```

query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(verify  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```

Symbolic values can be used just like concrete values of the same type.

A tiny example SDSL: ROSETTE

```
def bvmax(r0, r1) :  
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  r3 = bvneg(r2)  
  r4 = bvxor(r0, r2)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6
```

```
> verify(bvmax, max)  
(0, -2)
```



query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(verify  
  (assert (= (interpret bvmax inputs)  
              (interpret max inputs))))
```

(*verify expr*) searches for a concrete interpretation of symbolic constants that causes *expr* to fail.

A tiny example SDSL: **ROSETTE**

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(r0, r2)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6
```

```
> verify(bvmax, max)  
(0, -2)
```

```
> bvmax(0, -2)  
-1
```

query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(verify  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```

A tiny example SDSL: RÖSETTE

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(r0, r2)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6
```

```
> debug(bvmax, max, (0, -2))
```

query

```
(define inputs (list 0 -2))  
(debug [input-register?]  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```

A tiny example SDSL: **RoSETTE**

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(r0, r2)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6
```

```
> debug(bvmax, max, (0, -2))
```

query

```
(define inputs (list 0 -2))  
(debug [input-register?]  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs))))
```

A tiny example SDSL: ROSETTE

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(??, ??)  
  r5 = bvand(r3, ??)  
  r6 = bvxor(??, ??)  
  return r6
```

```
> synthesize(bvmax, max)
```

query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(synthesize [inputs]  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs)))))
```

A tiny example SDSL: **ROSETTE**

```
def bvmax(r0, r1) :  
  r2 = bvge(r0, r1)  
  r3 = bvneg(r2)  
  r4 = bvxor(r0, r1)  
  r5 = bvand(r3, r4)  
  r6 = bvxor(r1, r5)  
  return r6  
  
> synthesize(bvmax, max)
```



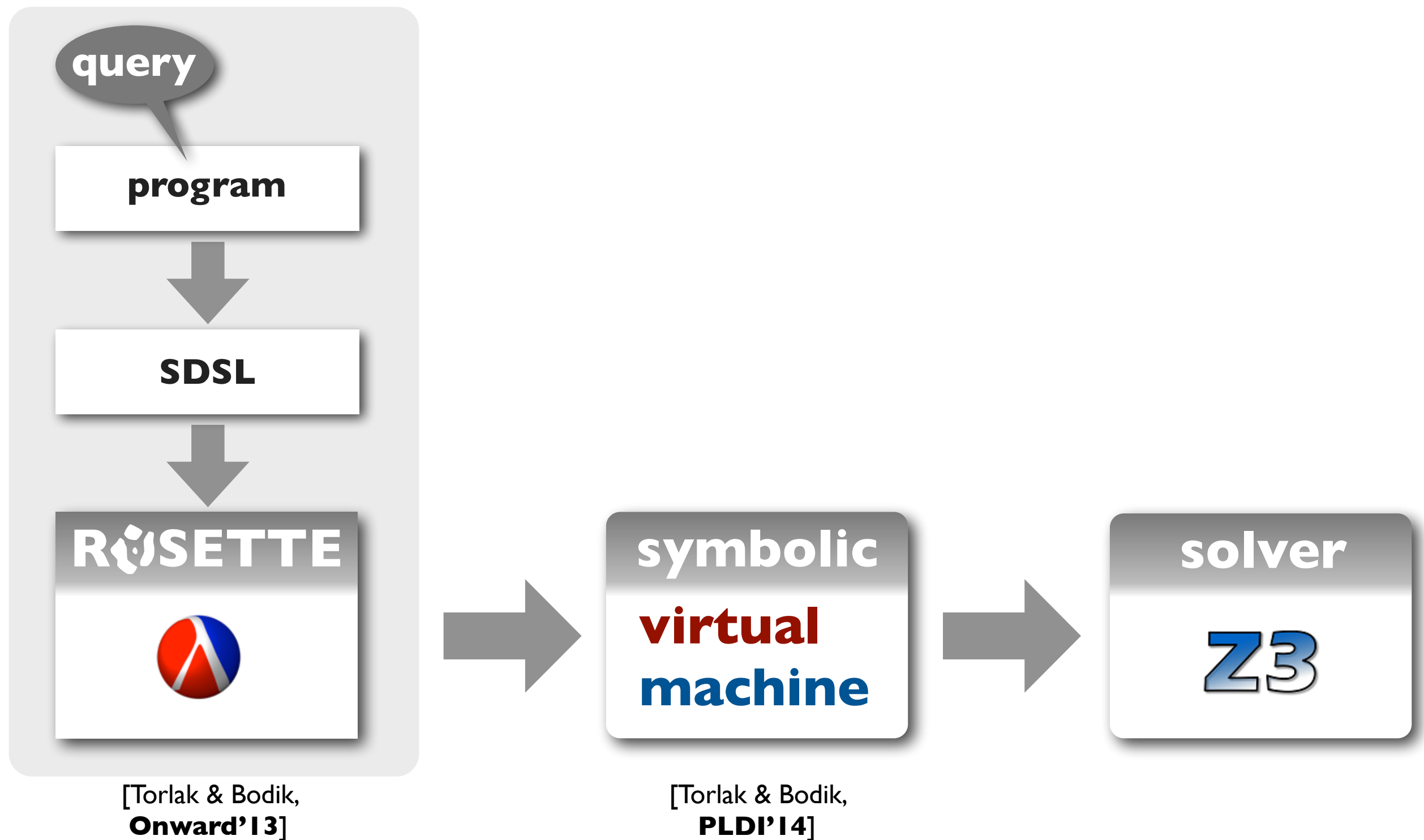
query

```
(define-symbolic n0 n1 integer?)  
(define inputs (list n0 n1))  
(synthesize [inputs]  
  (assert (= (interpret bvmax inputs)  
             (interpret max inputs)))))
```

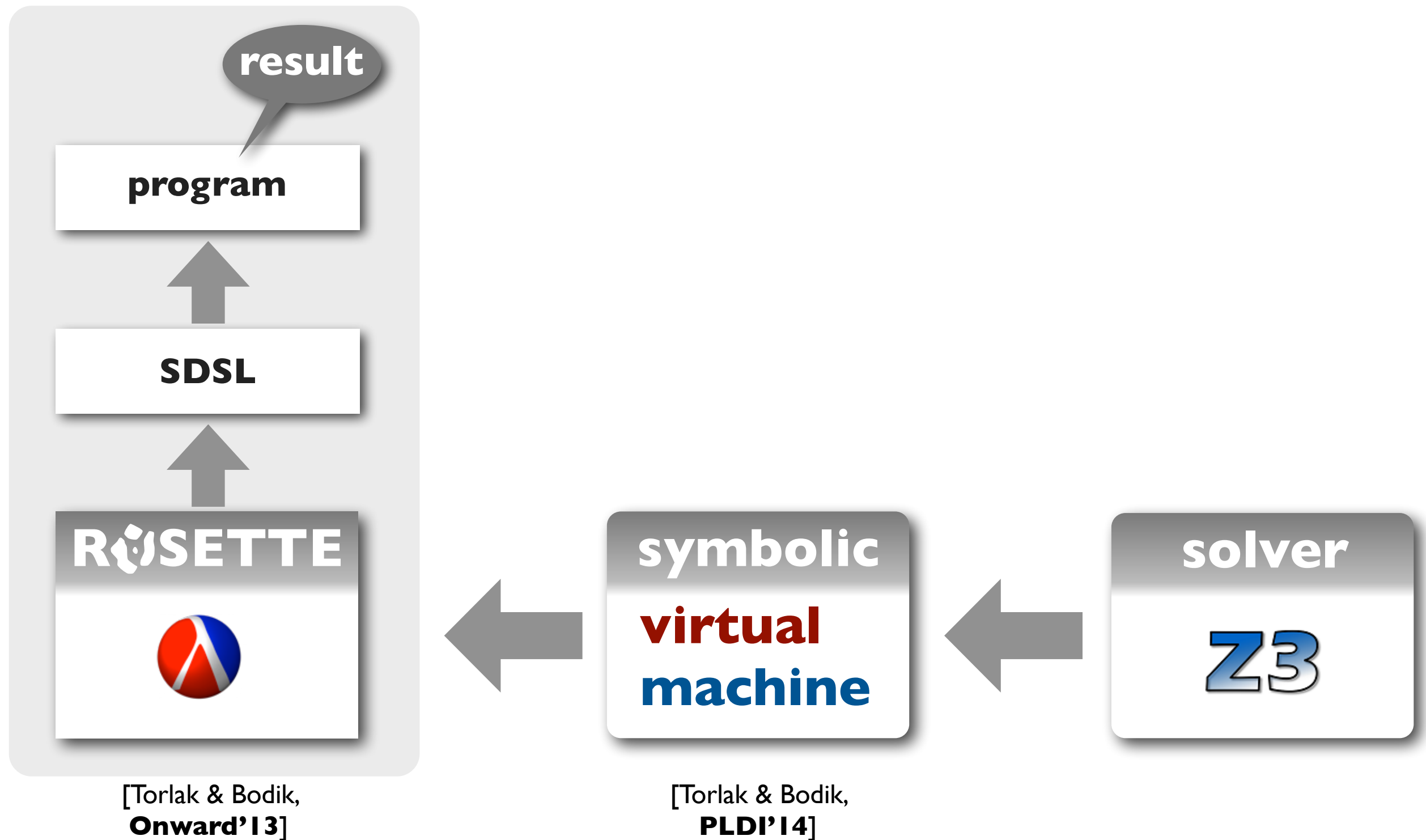

symbolic virtual machine (SVM)



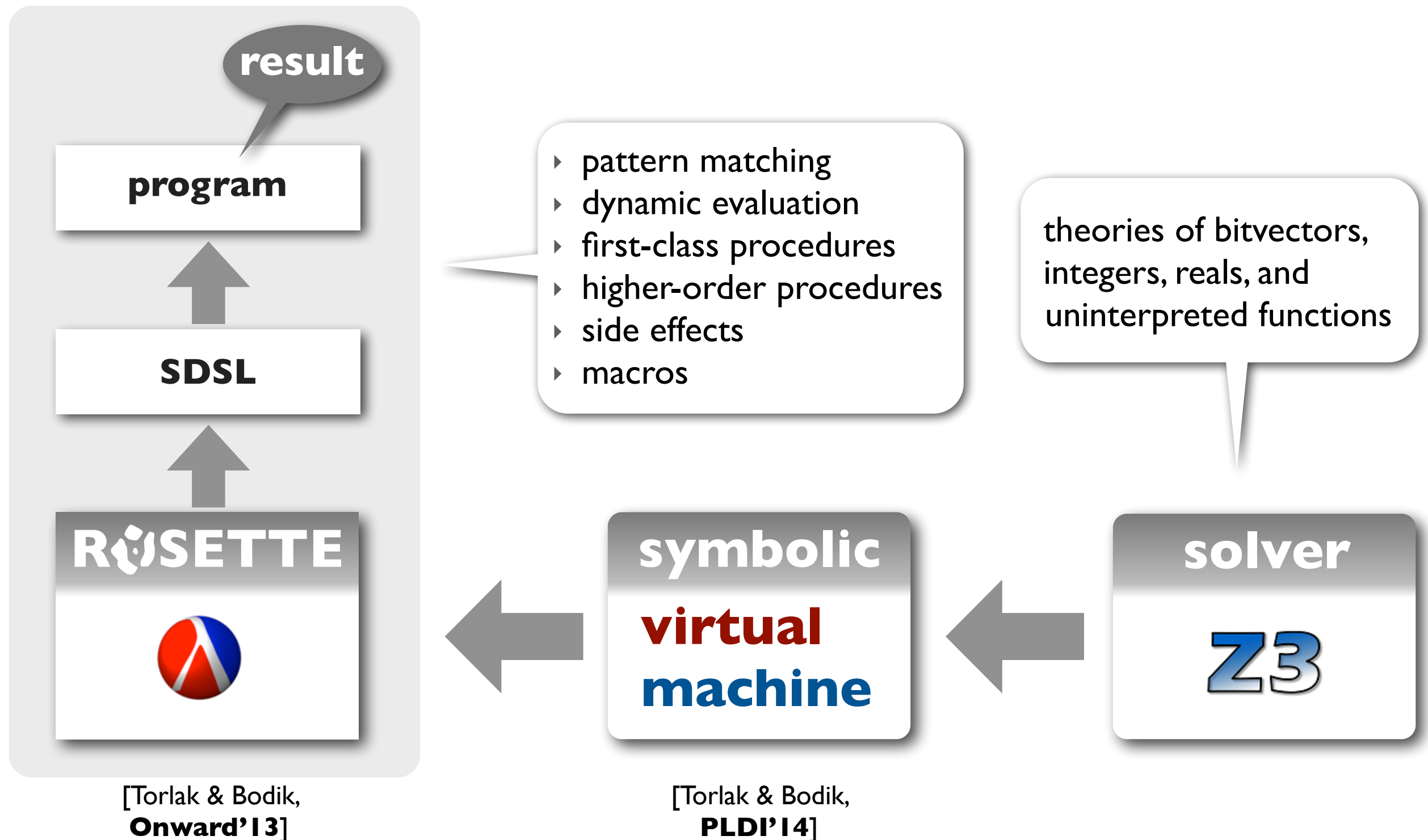
How it all works: a big picture view



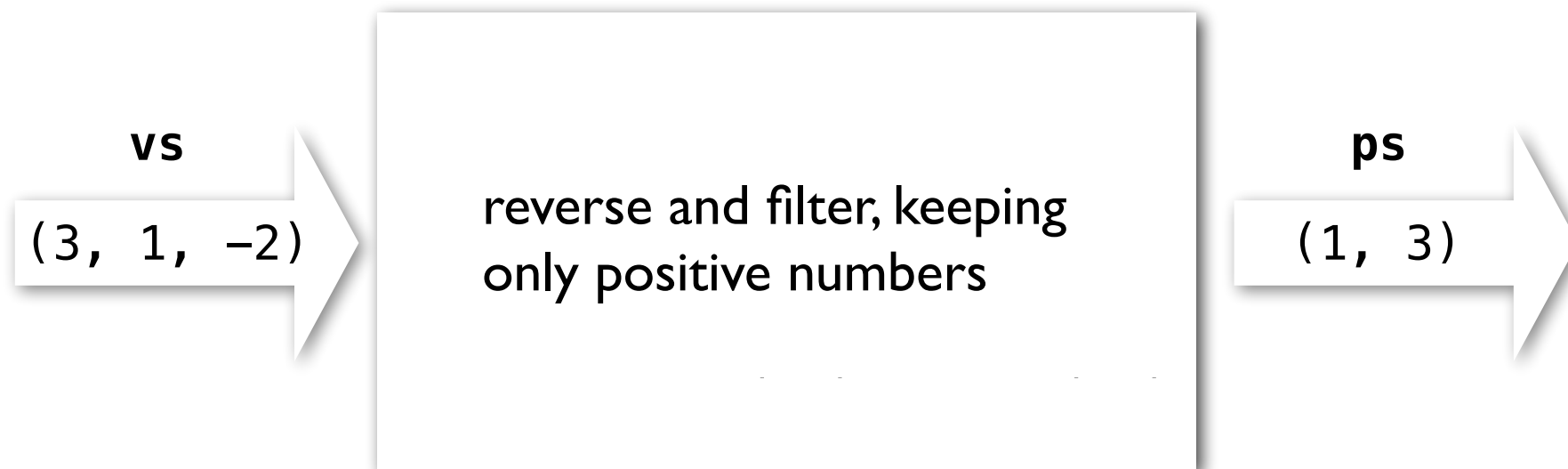
How it all works: a big picture view



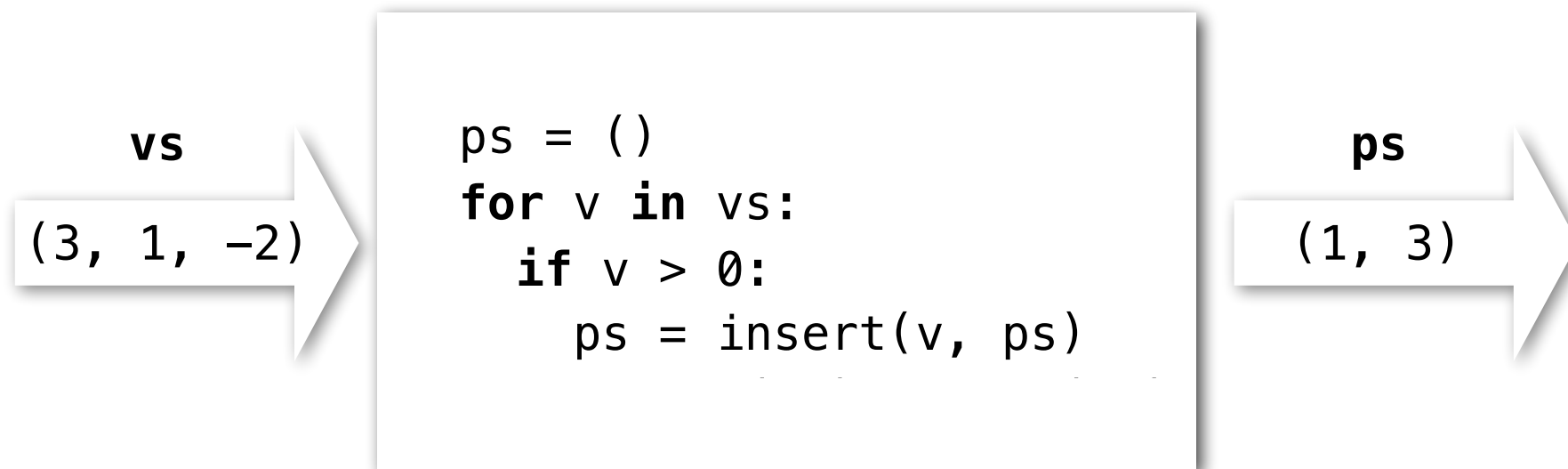
How it all works: a big picture view



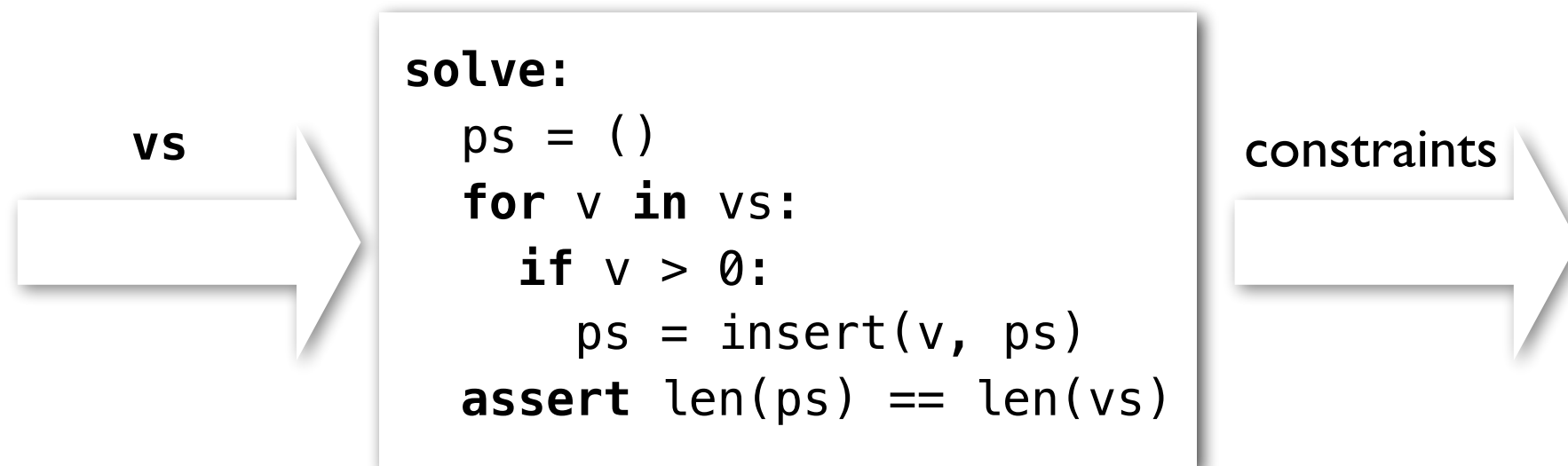
Translation to constraints by example



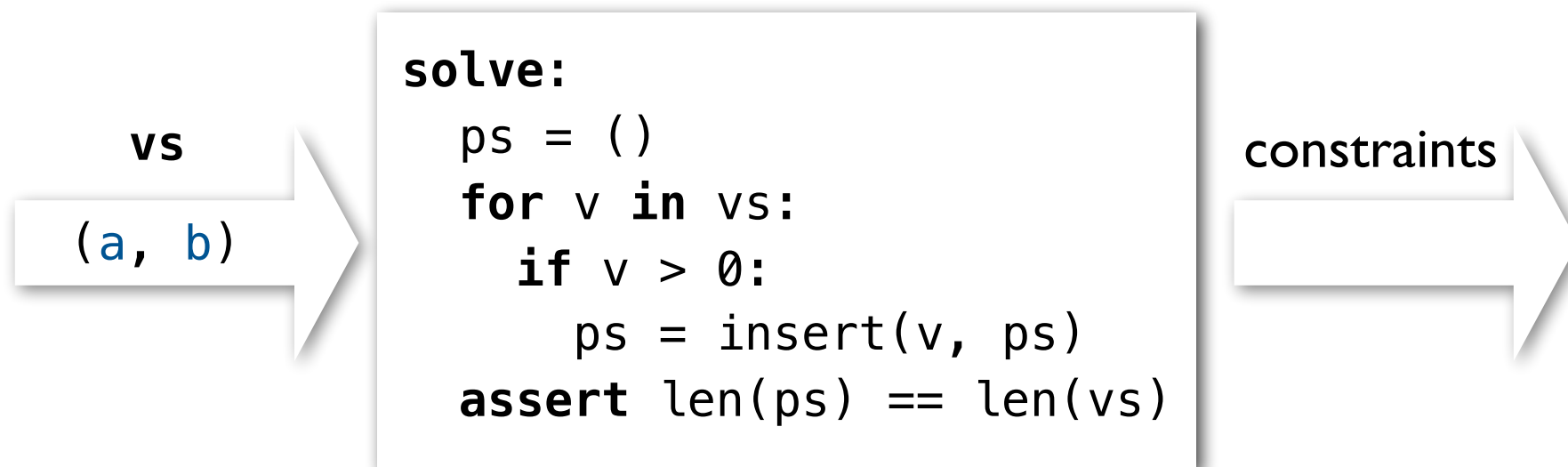
Translation to constraints by example



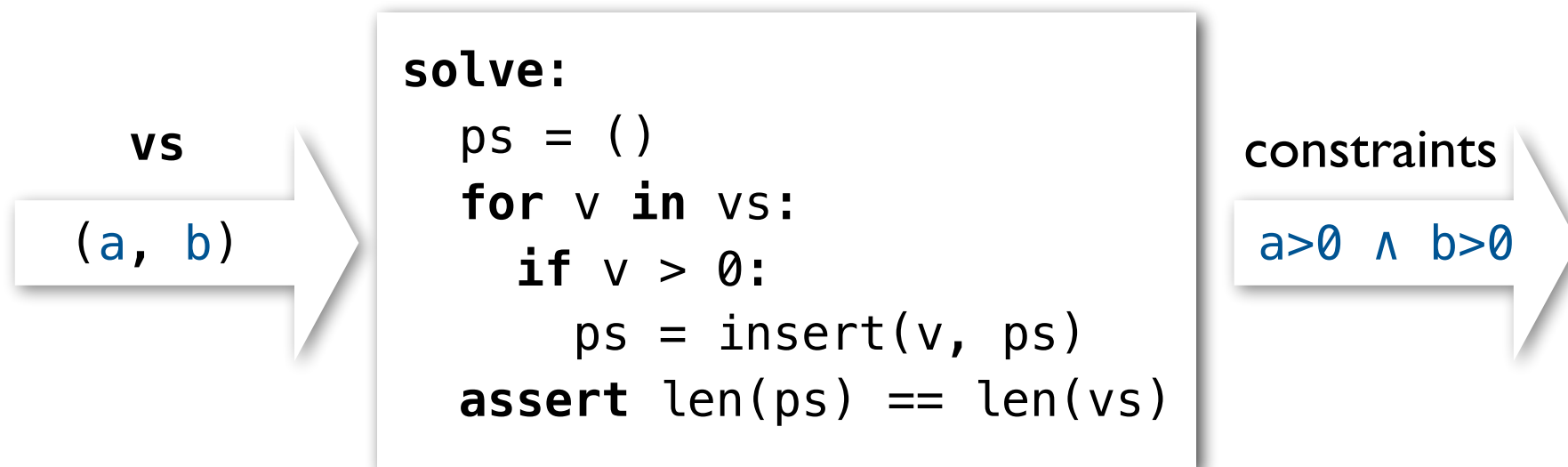
Translation to constraints by example



Translation to constraints by example



Translation to constraints by example

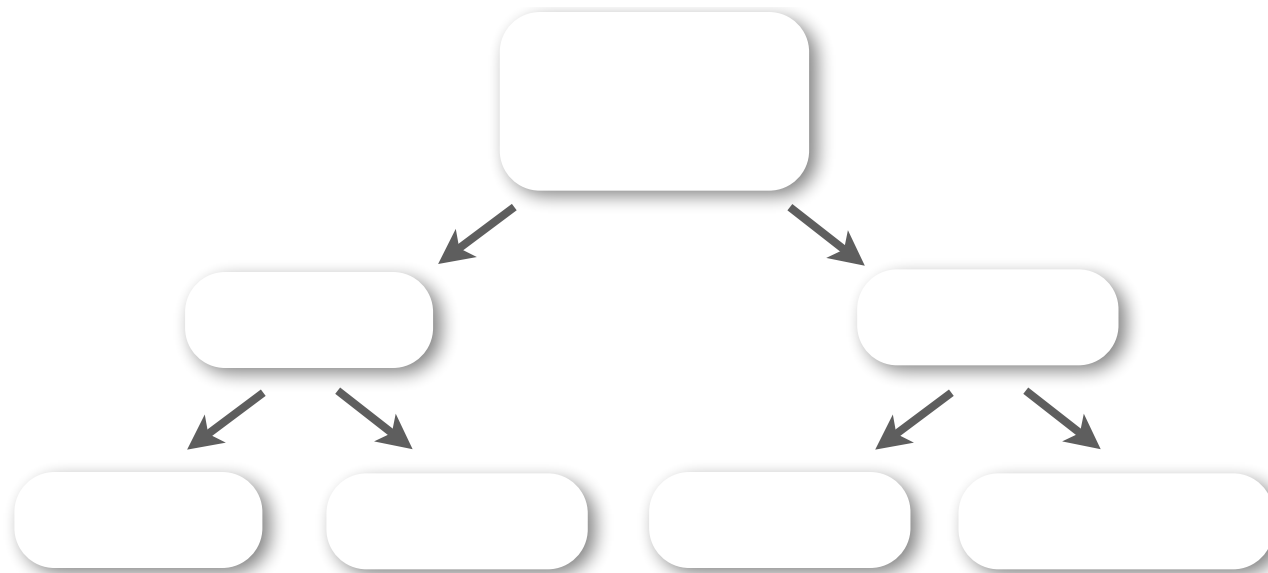


Design space of precise symbolic encodings

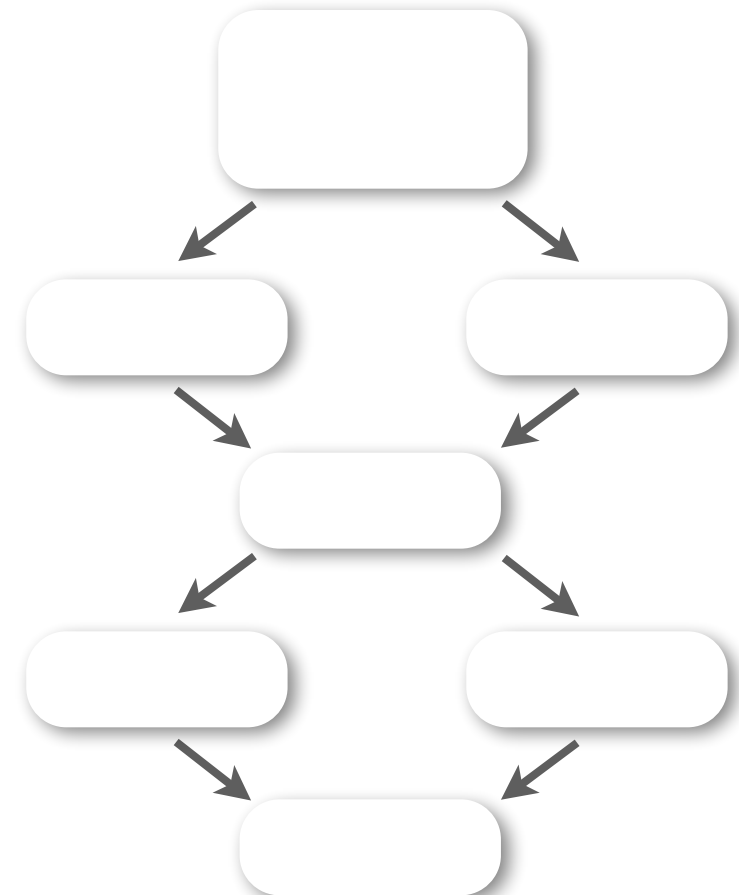
solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking

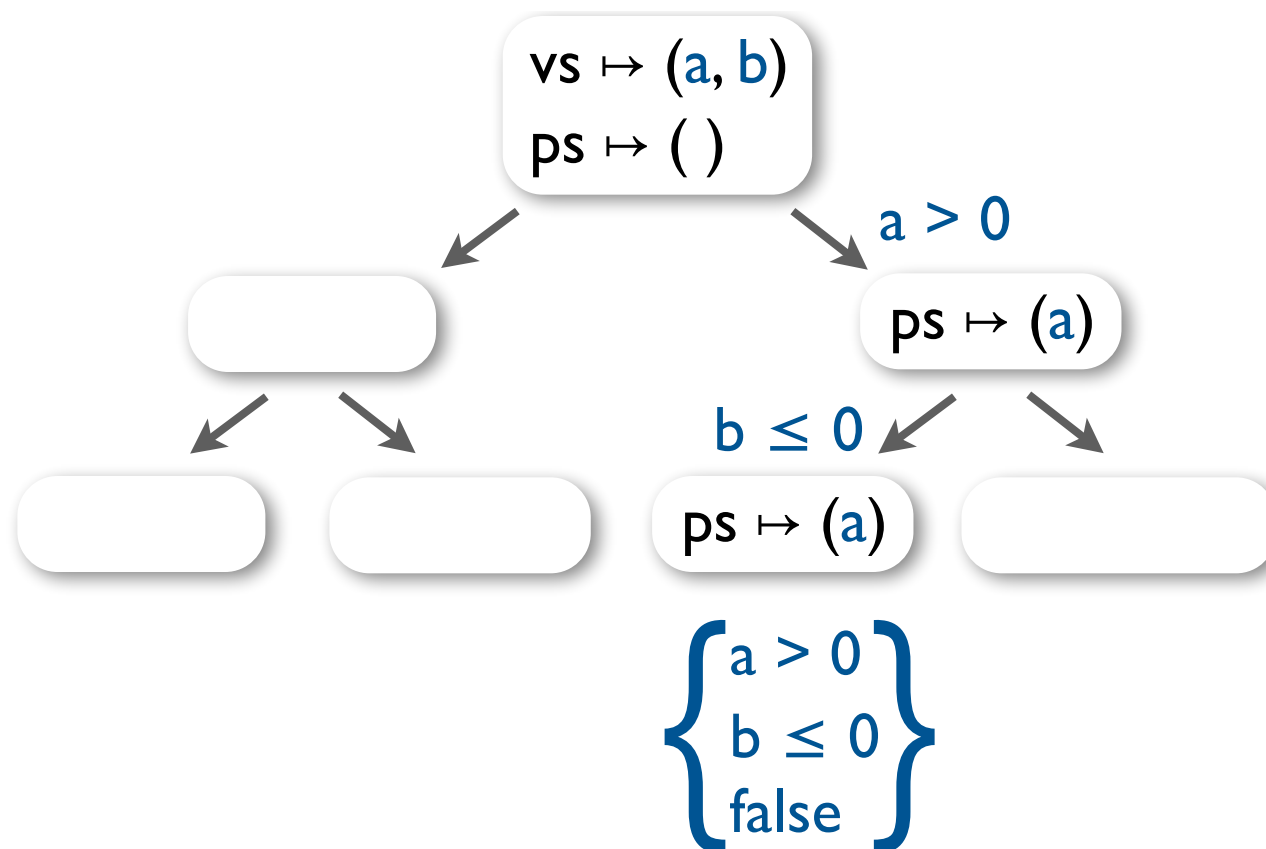


Design space of precise symbolic encodings

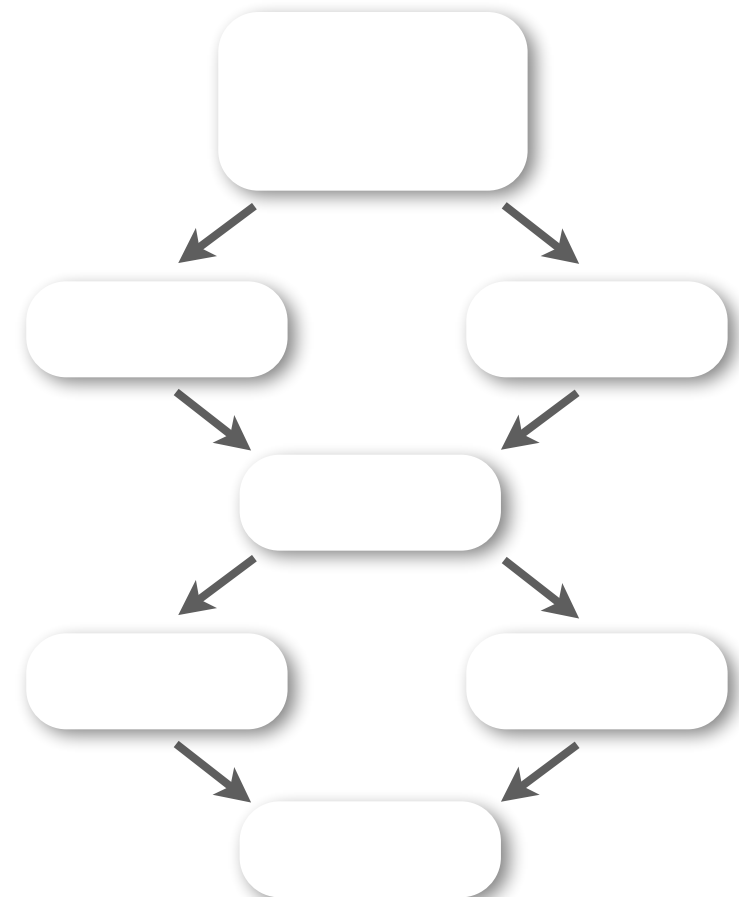
solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking

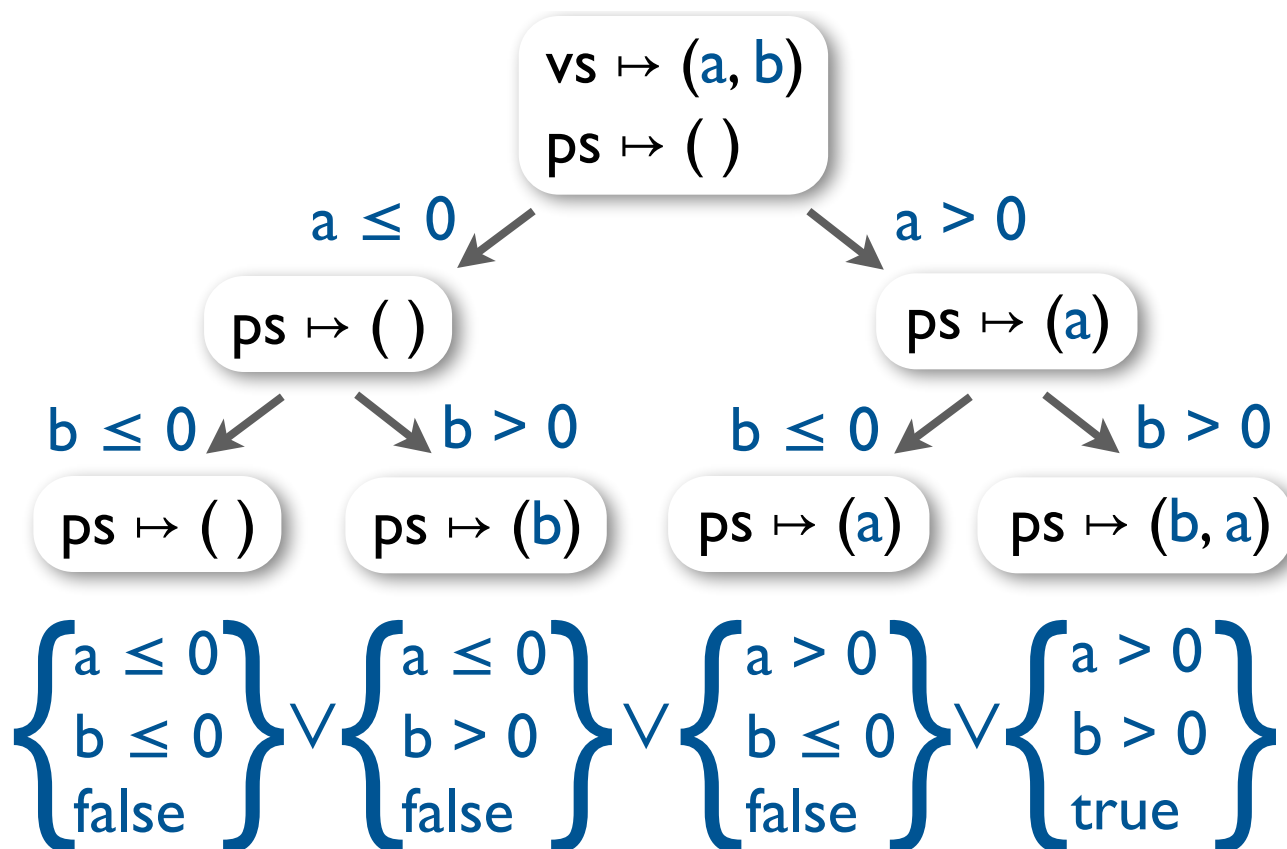


Design space of precise symbolic encodings

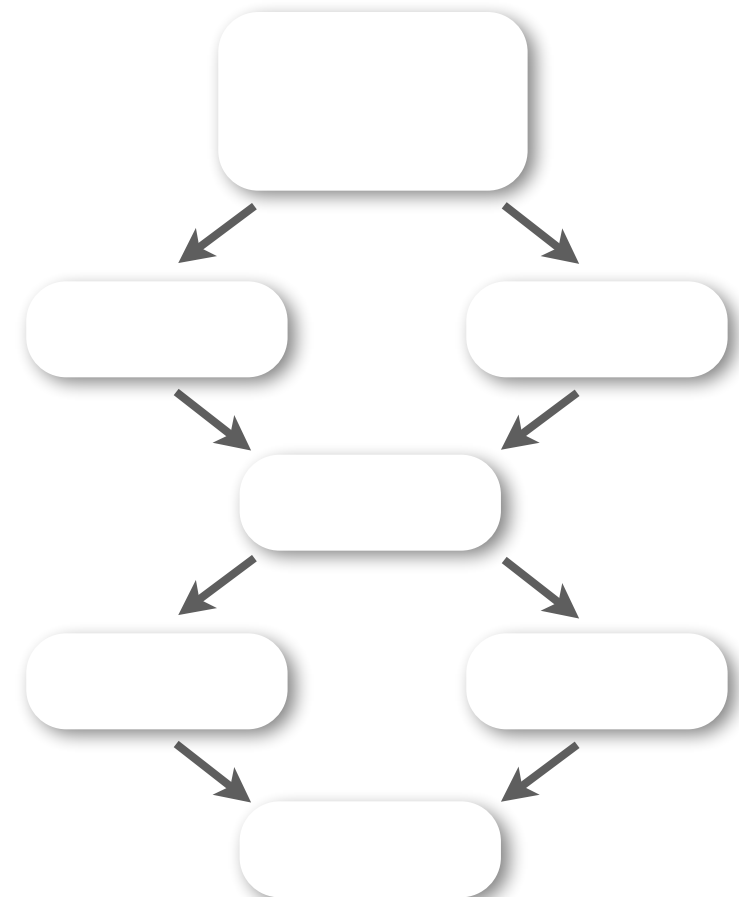
solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking

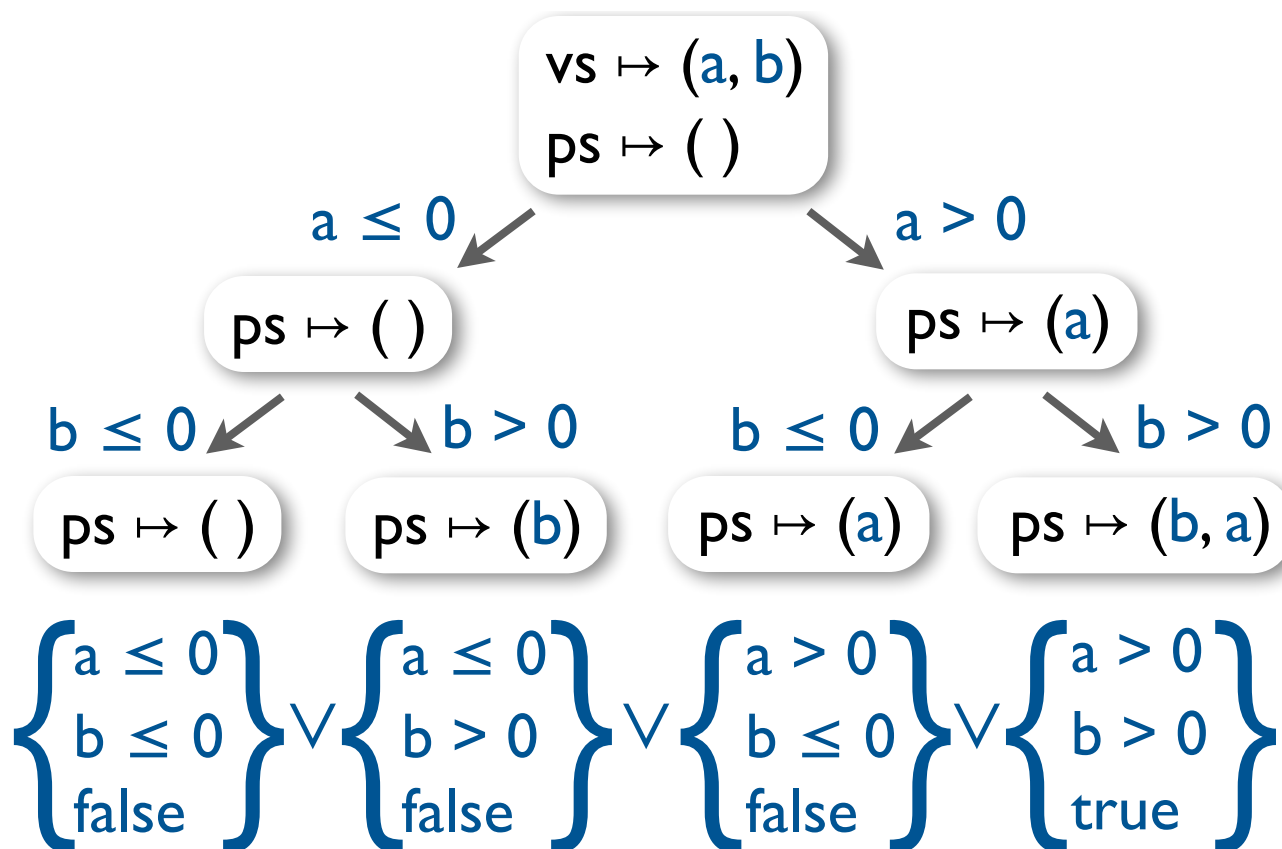


Design space of precise symbolic encodings

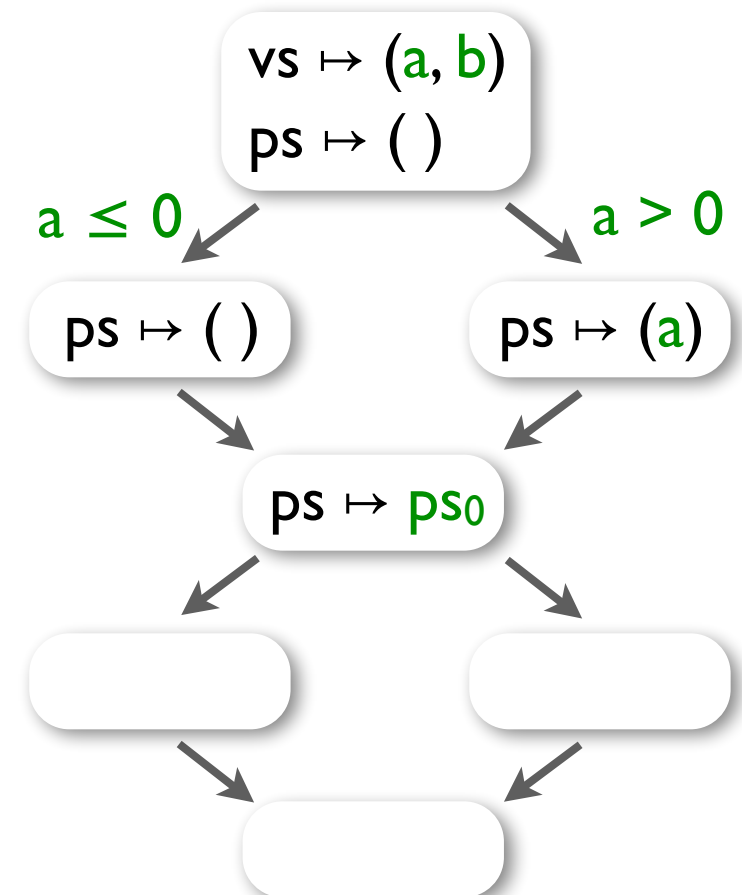
solve:

```
ps = ()
for v in vs:
    if v > 0:
        ps = insert(v, ps)
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking



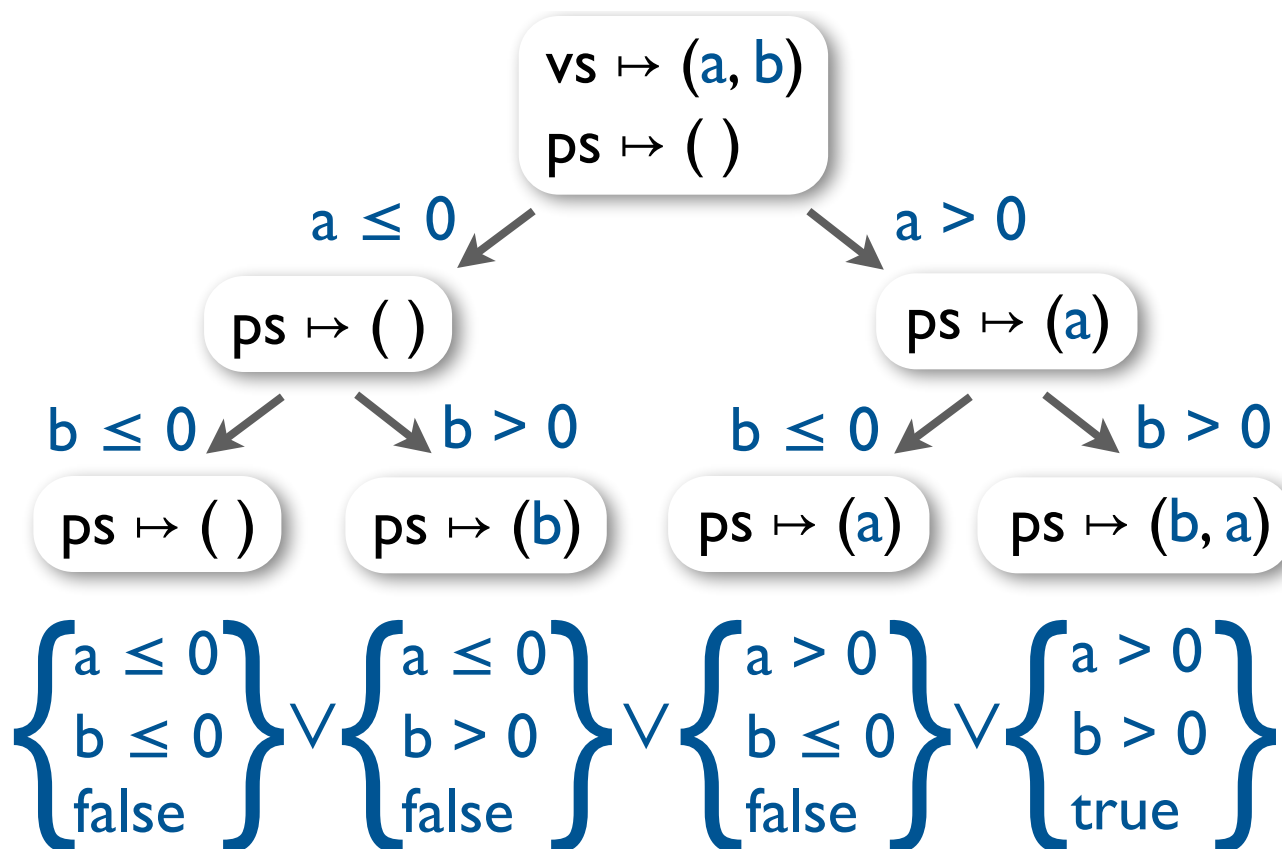
$ps_0 = \text{ite}(a > 0, (a), ())$

Design space of precise symbolic encodings

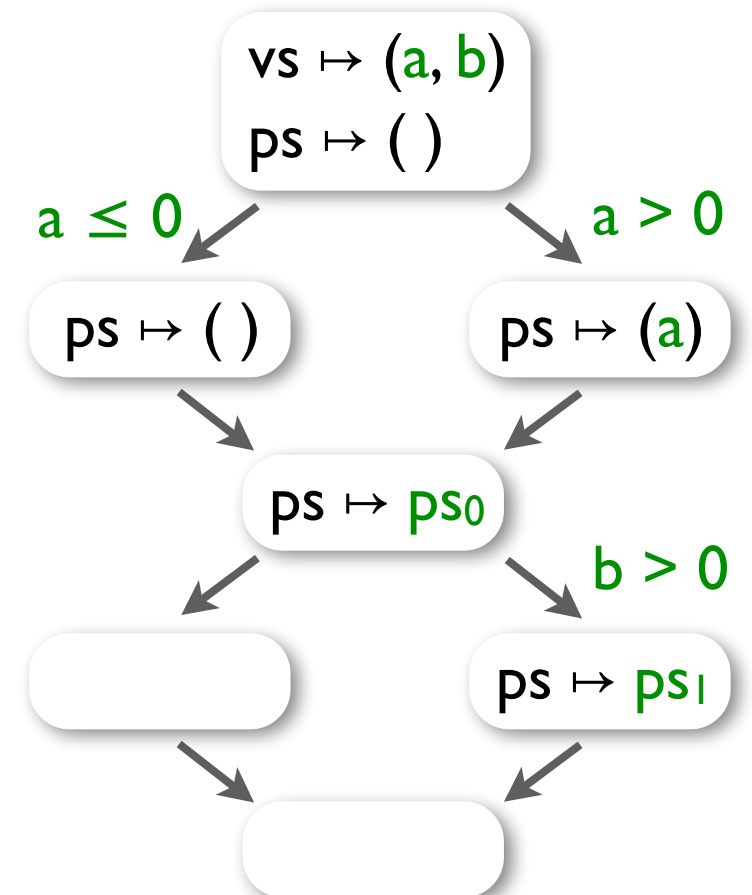
solve:

```
ps = ()
for v in vs:
    if v > 0:
        ps = insert(v, ps)
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking



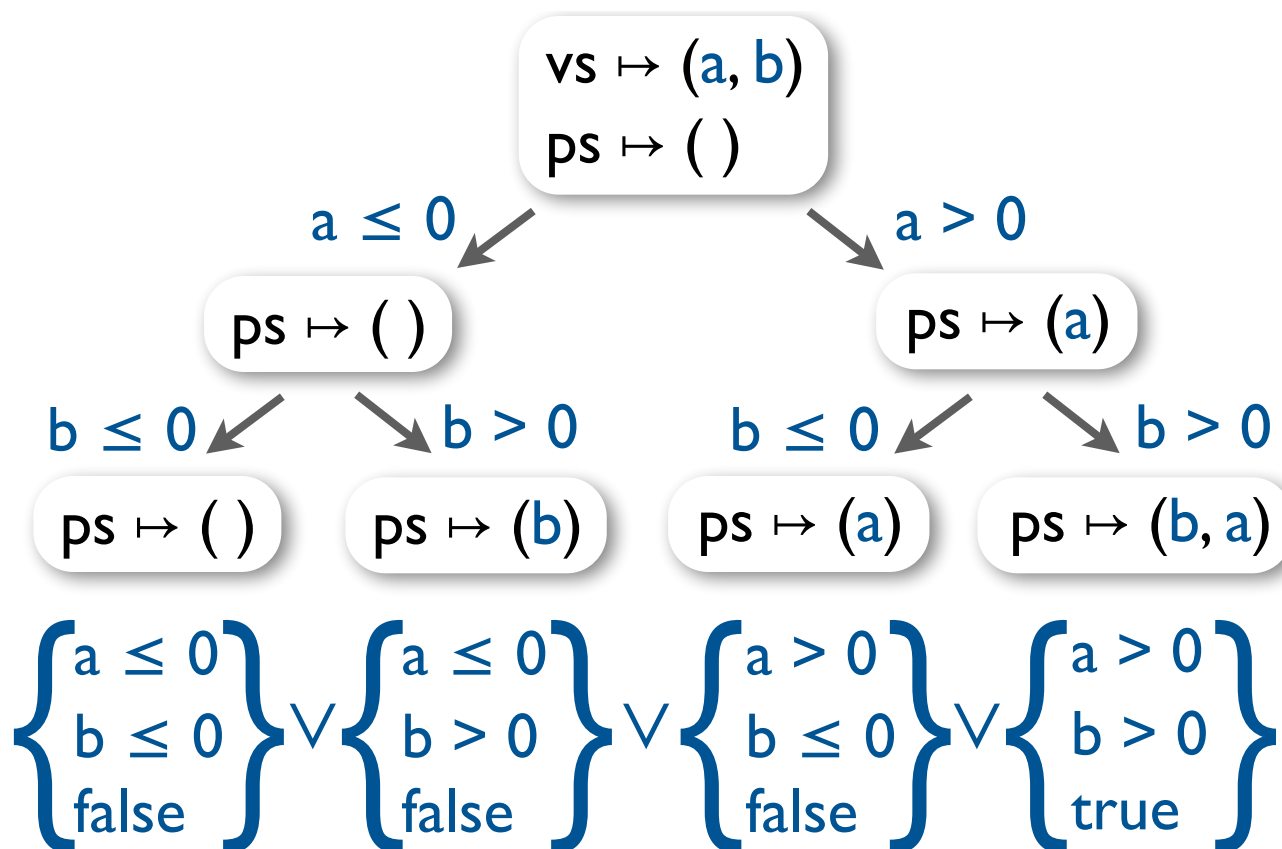
ps₀ = ite(a > 0, (a), ())
ps₁ = insert(b, ps₀)

Design space of precise symbolic encodings

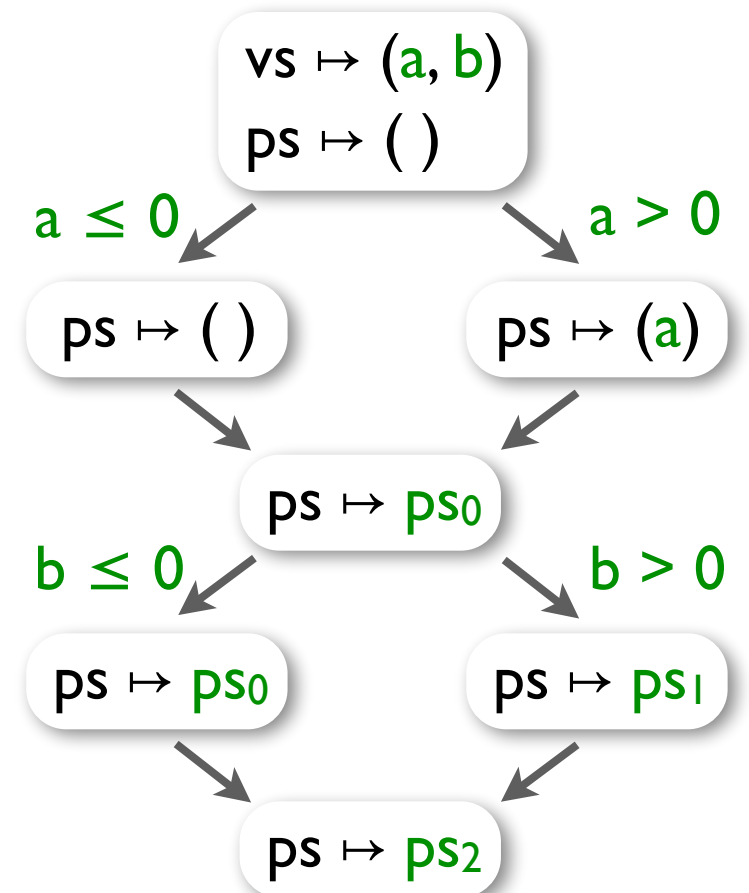
solve:

```
ps = ()
for v in vs:
    if v > 0:
        ps = insert(v, ps)
assert len(ps) == len(vs)
```

symbolic execution



bounded model checking

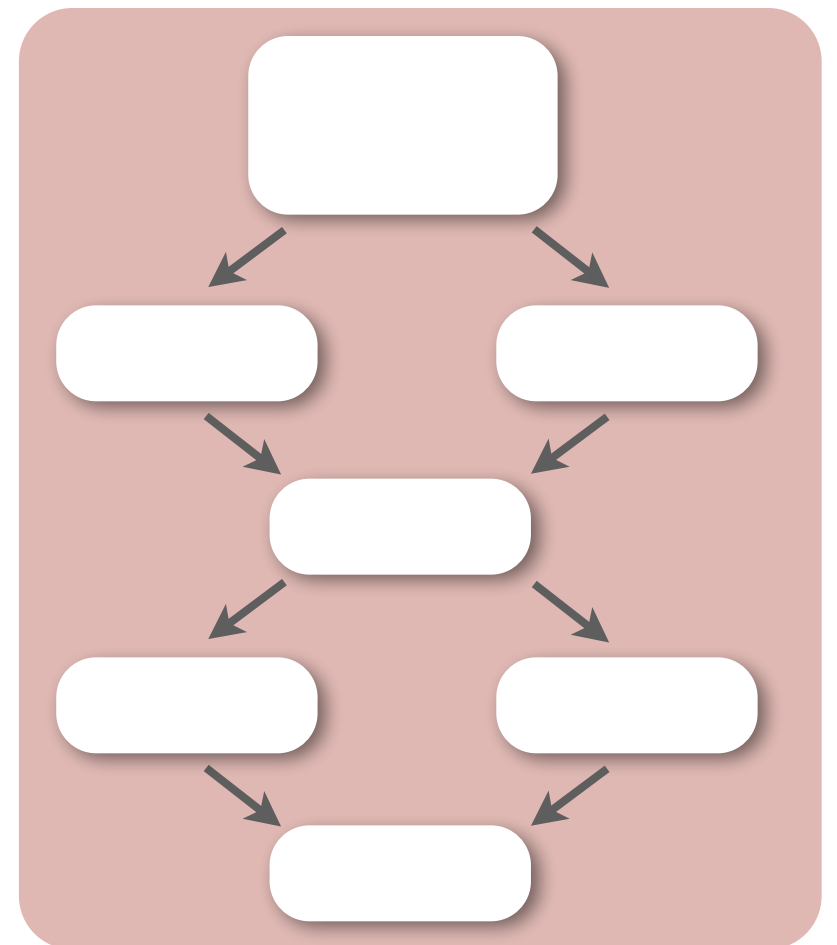


```
ps0 = ite(a > 0, (a), ())
ps1 = insert(b, ps0)
ps2 = ite(b > 0, ps0, ps1)
assert len(ps2) = 2
```

A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```



$\left\{ \begin{array}{l} a > 0 \\ b > 0 \\ \text{true} \end{array} \right\}$



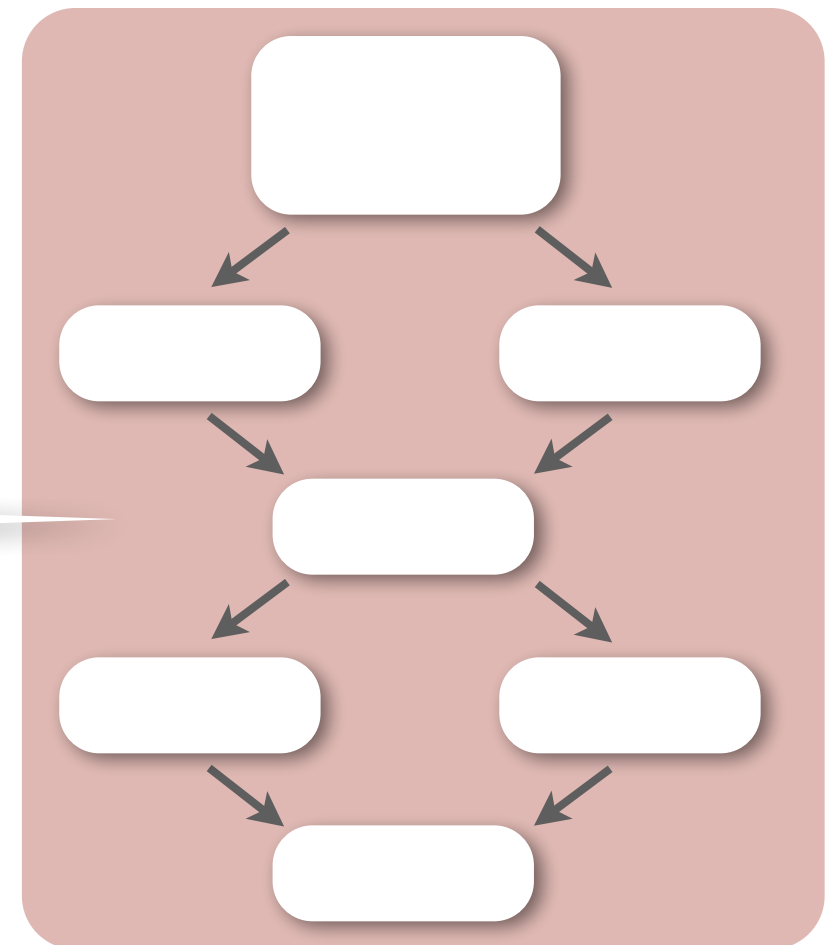
A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Merge values of

- primitive types: **symbolically**
- immutable types: **structurally**
- all other types: via **unions**


$$\left\{ \begin{array}{l} a > 0 \\ b > 0 \\ \text{true} \end{array} \right\}$$

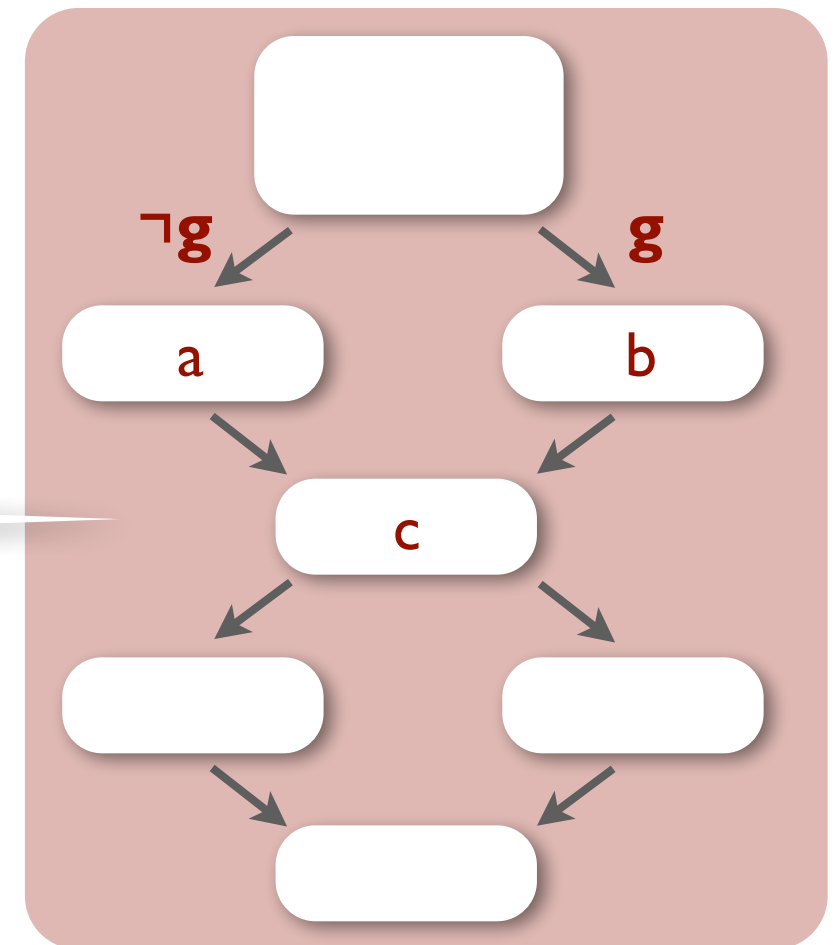

A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Merge values of

- primitive types: **symbolically**
- immutable types: structurally
- all other types: via **unions**



$\left\{ \begin{array}{l} a > 0 \\ b > 0 \\ \text{true} \end{array} \right\}$



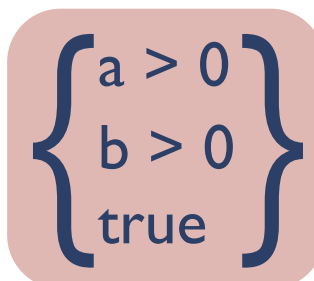
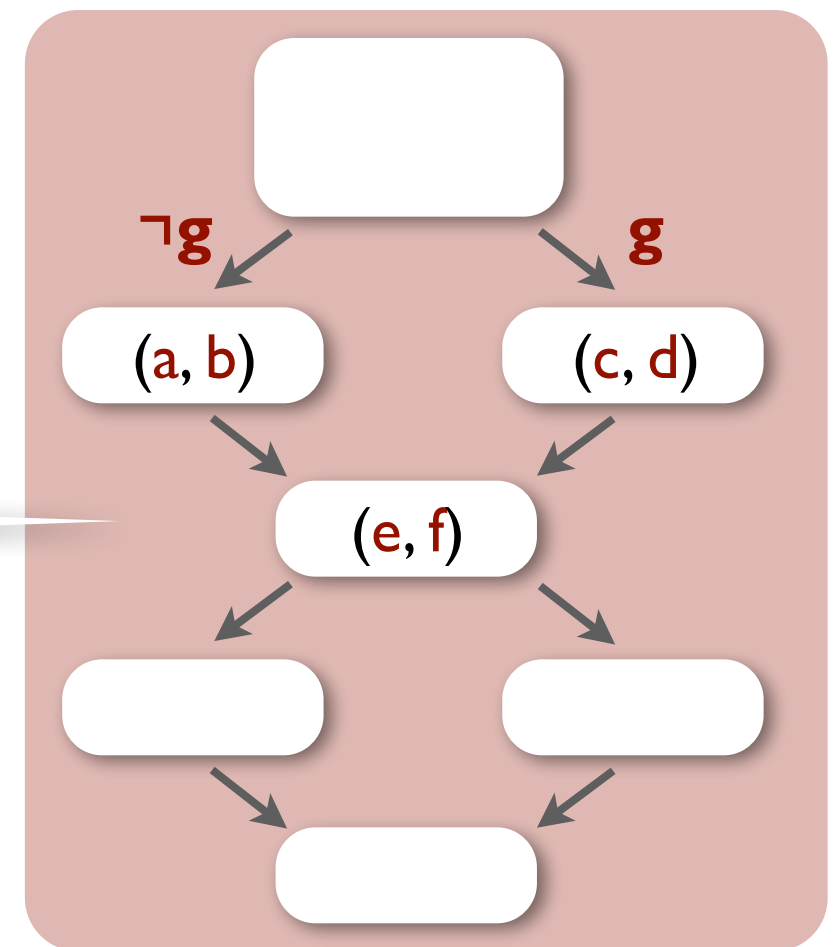
A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Merge values of

- primitive types: *symbolically*
- immutable types: *structurally*
- all other types: via *unions*



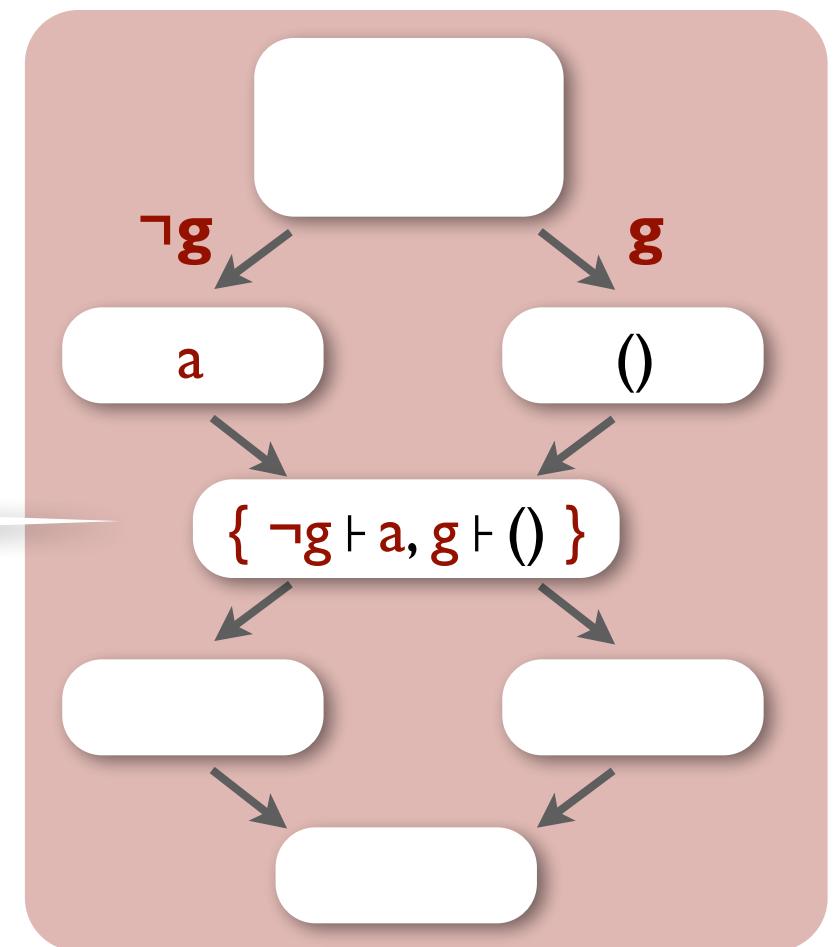
A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Merge values of

- primitive types: **symbolically**
- immutable types: **structurally**
- all other types: **via unions**



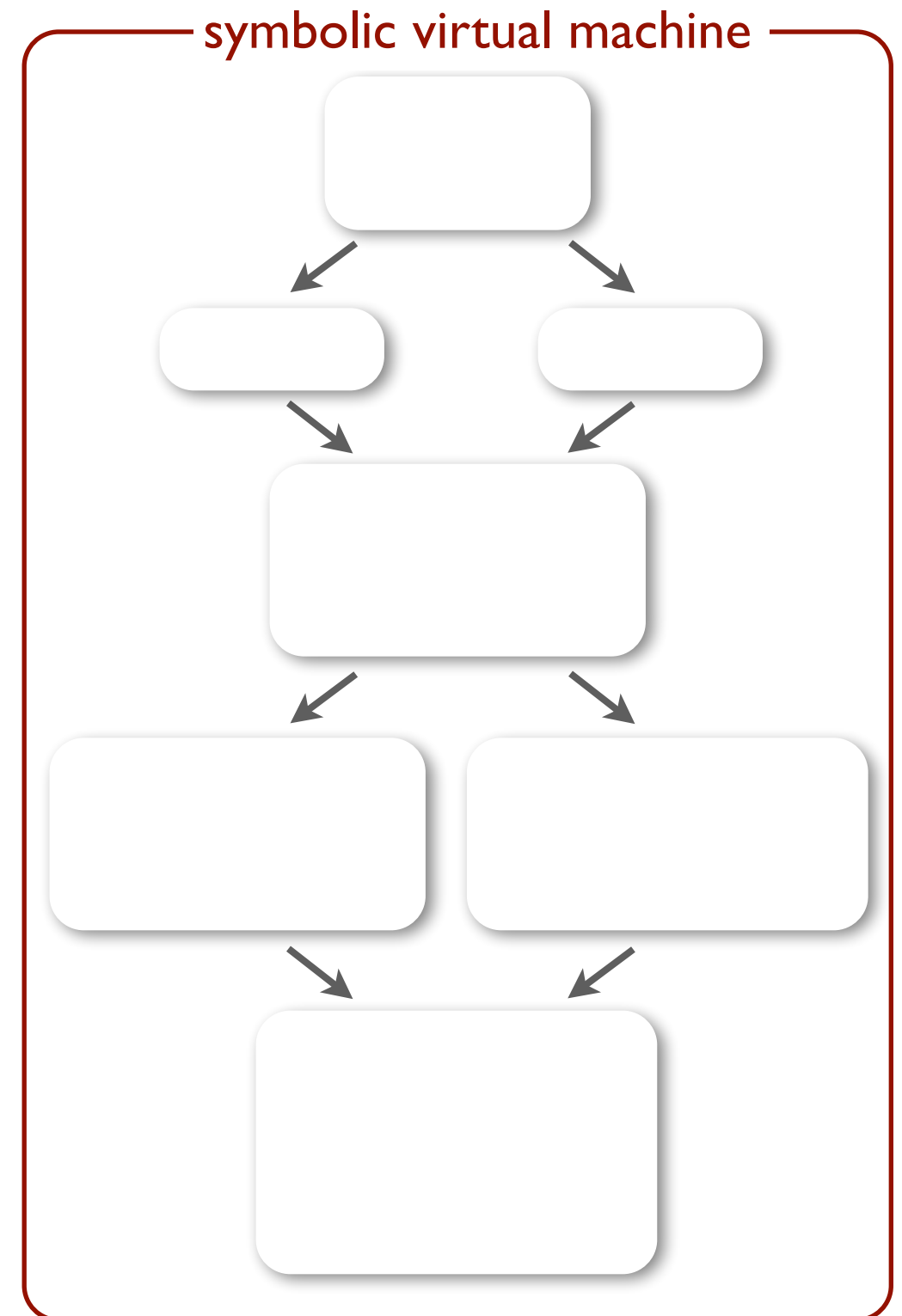
$\left\{ \begin{array}{l} a > 0 \\ b > 0 \\ \text{true} \end{array} \right\}$



A new design: type-driven state merging

solve:

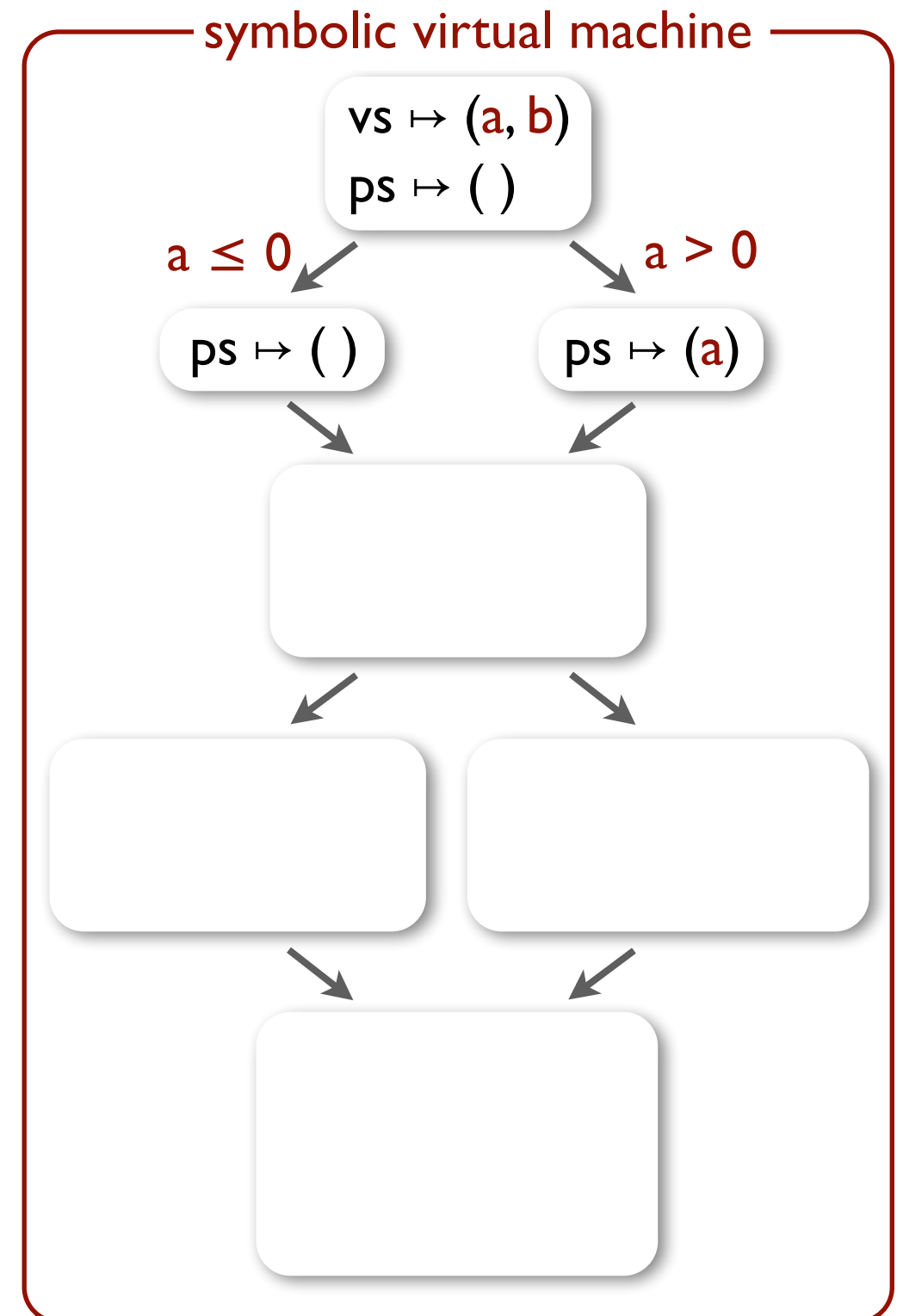
```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```



A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```



A new design: type-driven state merging

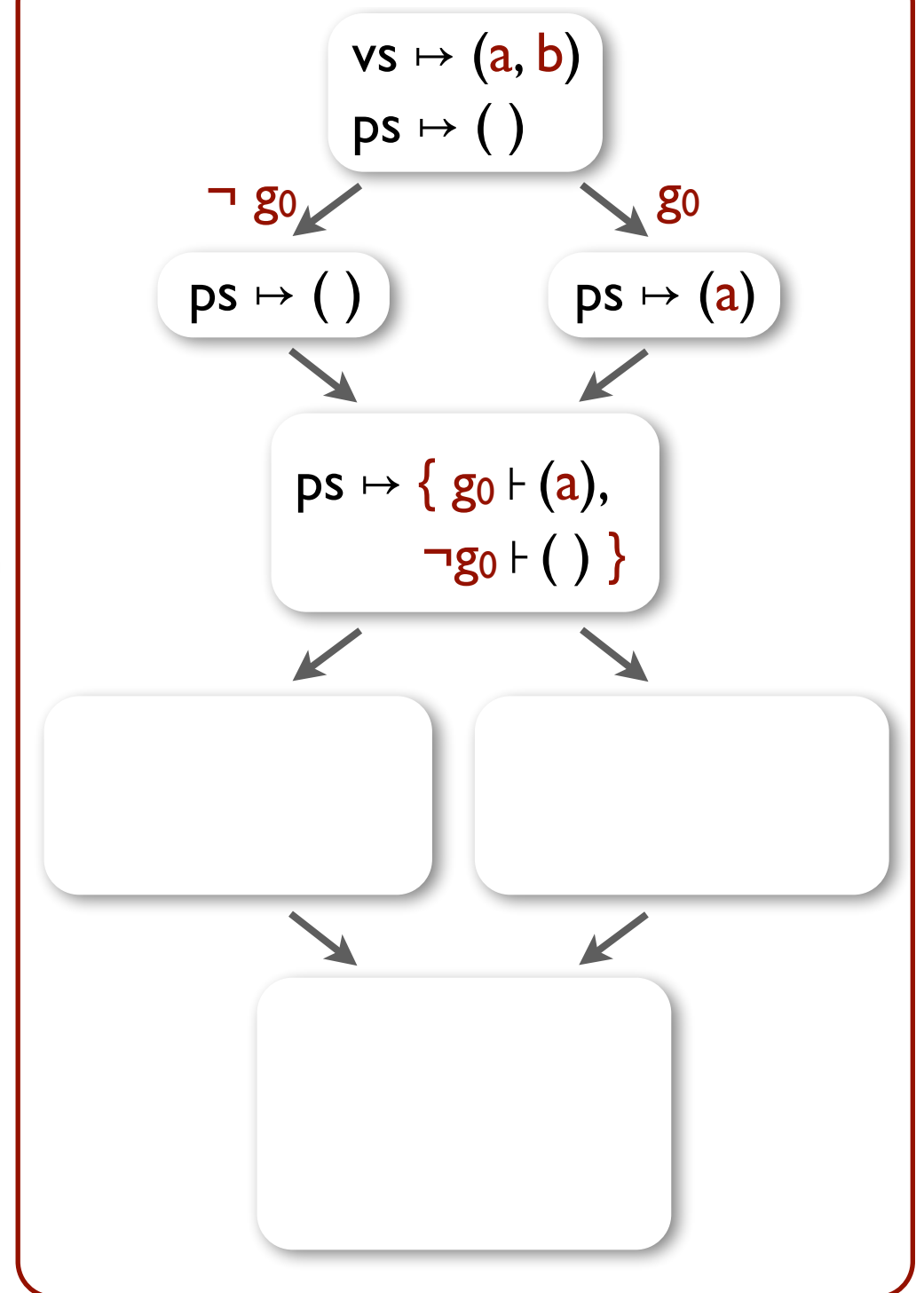
solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Symbolic union: a set of guarded values, with disjoint guards.

$g_0 = a > 0$

symbolic virtual machine



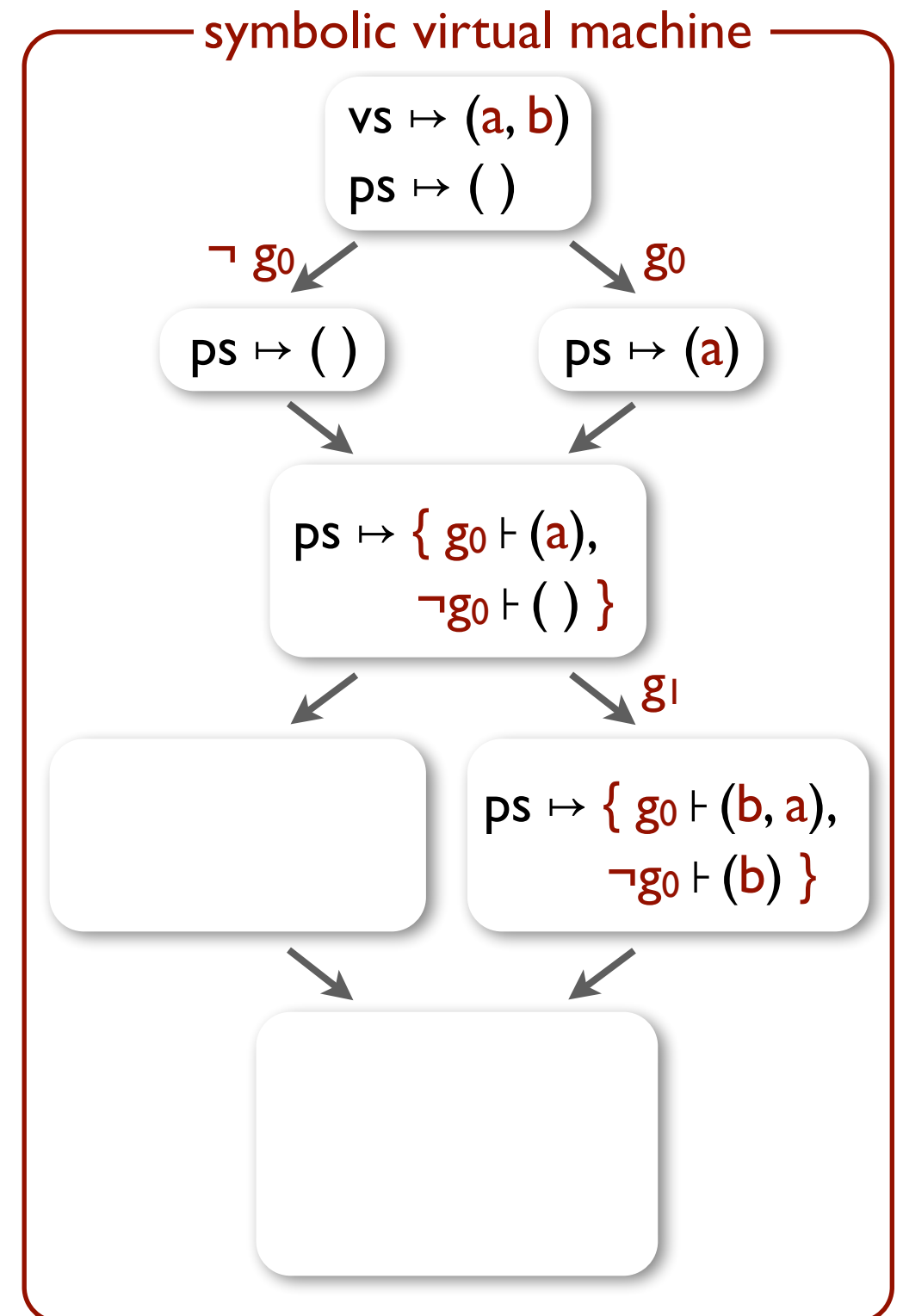
A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

Execute insert
concretely on all
lists in the union.

$g_0 = a > 0$
 $g_1 = b > 0$

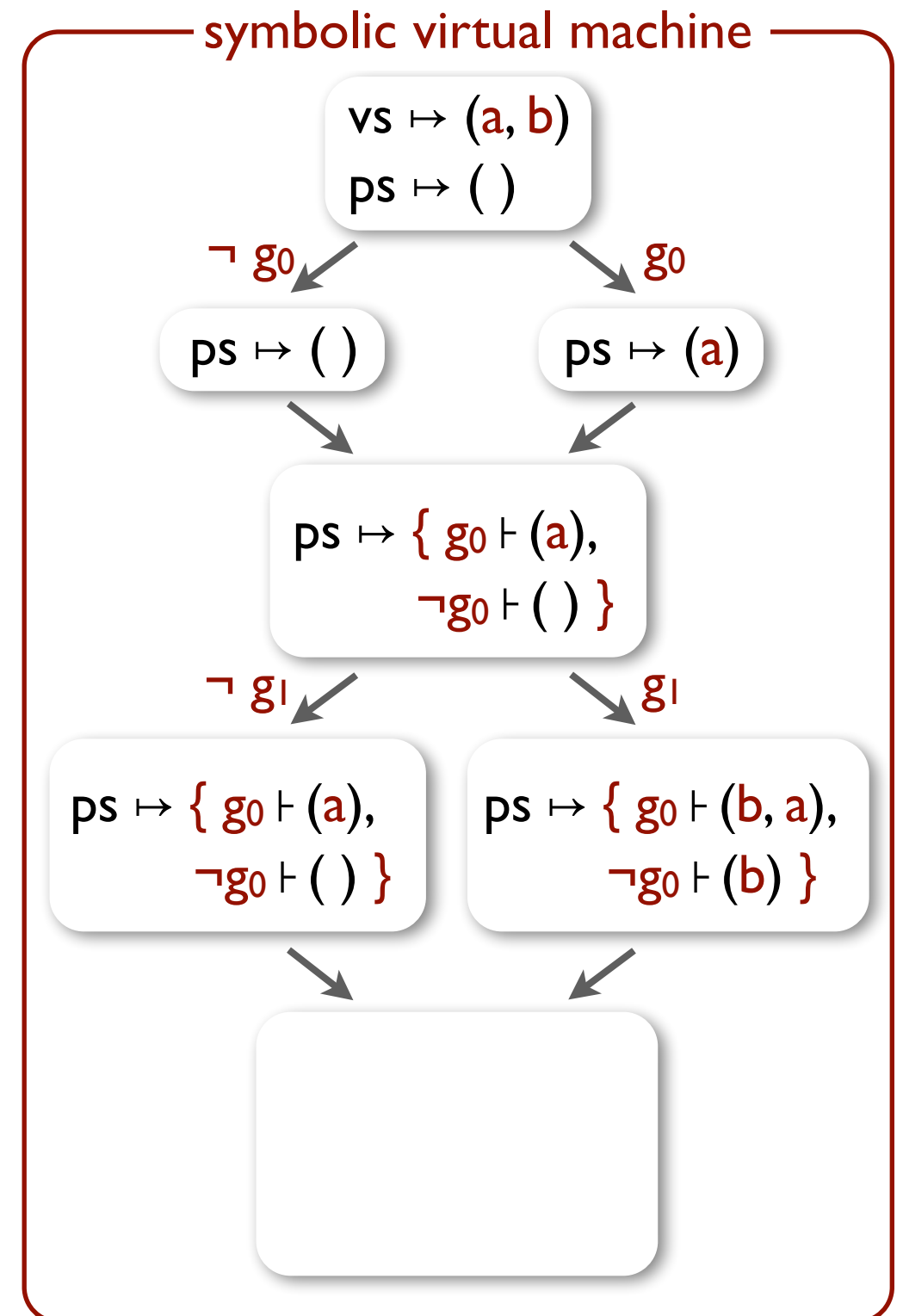


A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

$g_0 = a > 0$
 $g_1 = b > 0$

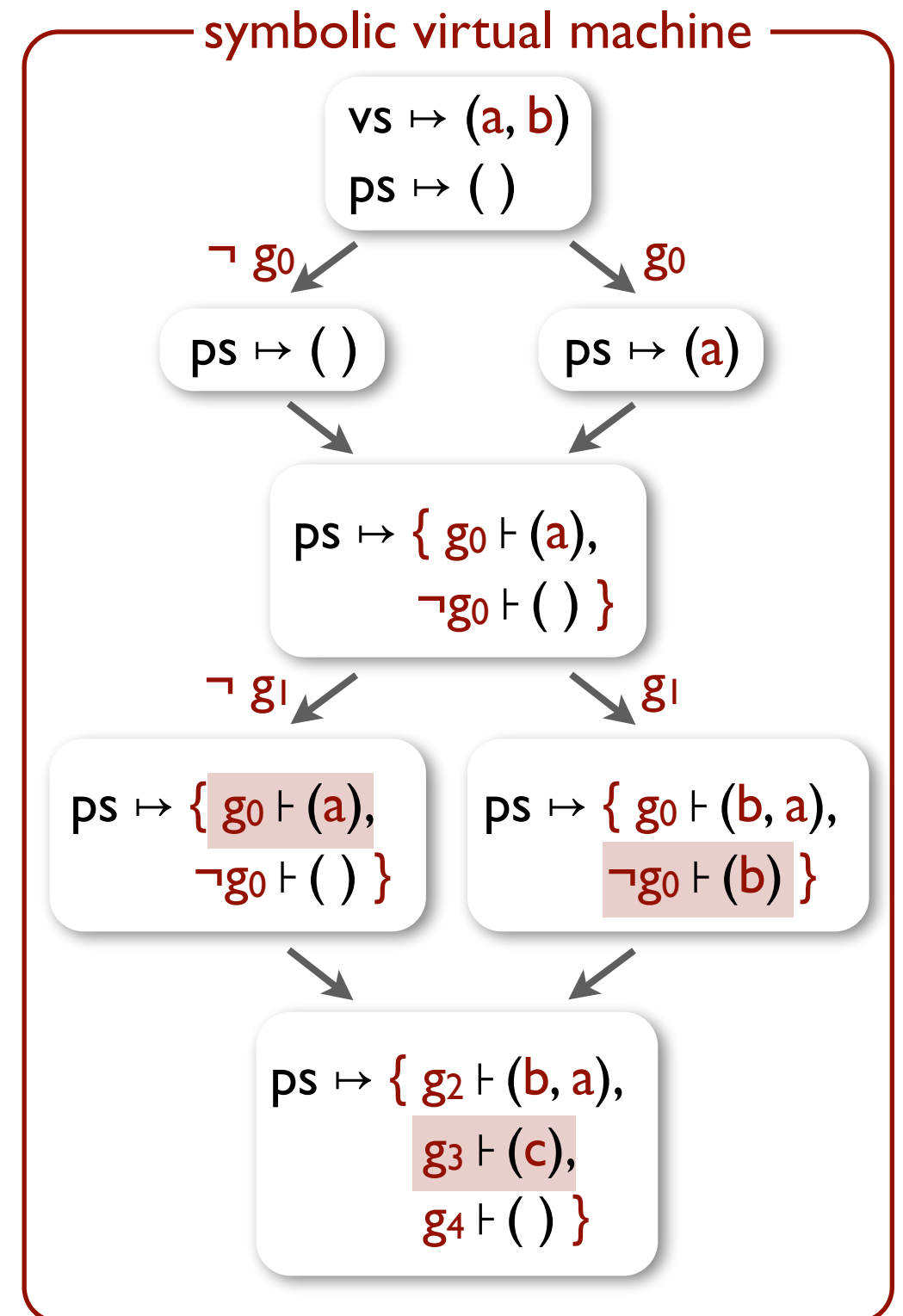


A new design: type-driven state merging

solve:

```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

$g_0 = a > 0$
 $g_1 = b > 0$
 $g_2 = g_0 \wedge g_1$
 $g_3 = \neg(g_0 \Leftrightarrow g_1)$
 $g_4 = \neg g_0 \wedge \neg g_1$
 $c = \text{ite}(g_1, b, a)$



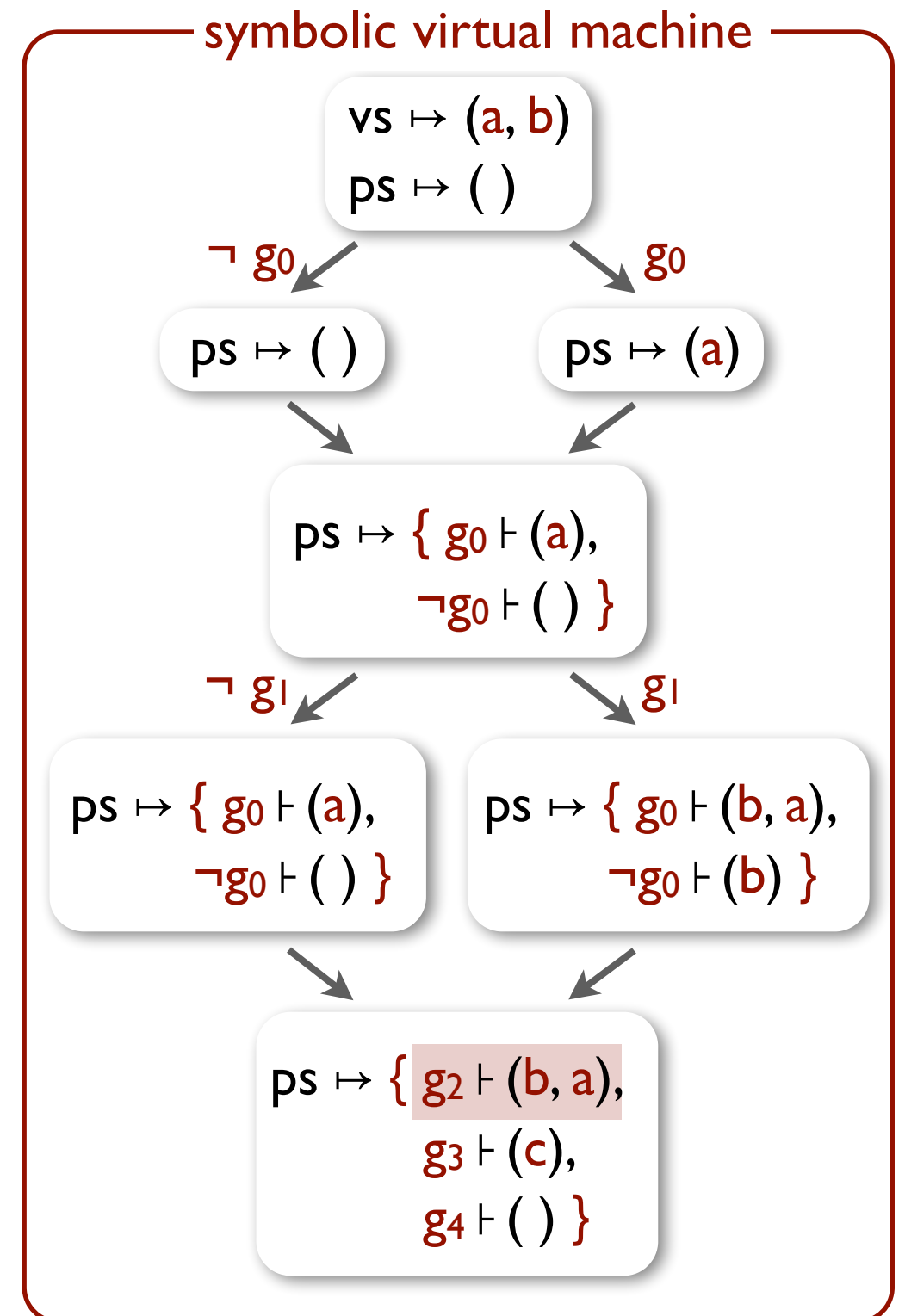
A new design: type-driven state merging

solve:

```
ps = ()
for v in vs:
    if v > 0:
        ps = insert(v, ps)
assert len(ps) == len(vs)
```

Evaluate `len` concretely on all lists in the union; assertion true only on the list guarded by g_2 .

```
g0 = a > 0
g1 = b > 0
g2 = g0 ∧ g1
g3 = ¬(g0 ⇔ g1)
g4 = ¬g0 ∧ ¬g1
c = ite(g1, b, a)
assert g2
```



A new design: type-driven state merging

solve:

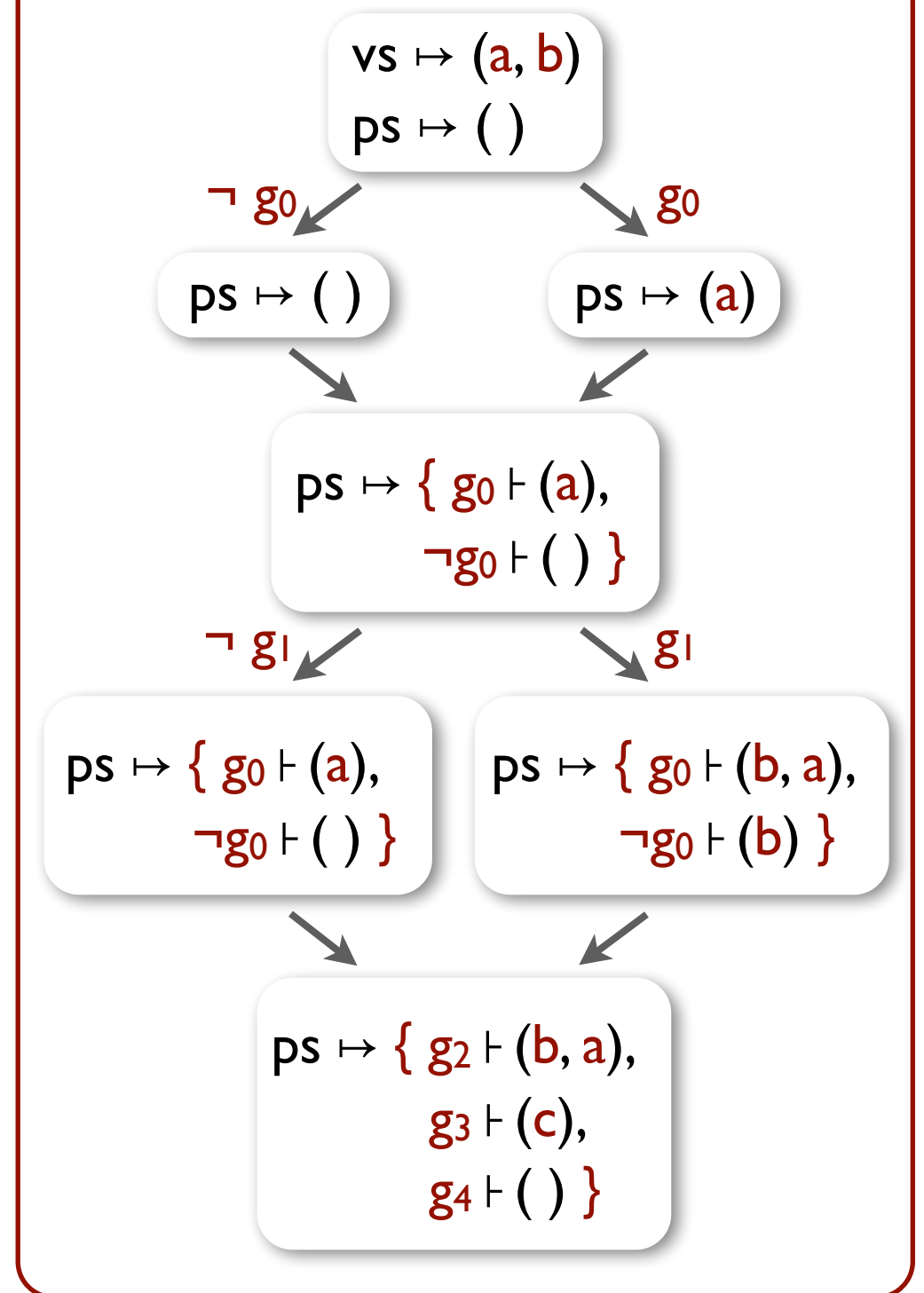
```
ps = ()  
for v in vs:  
    if v > 0:  
        ps = insert(v, ps)  
assert len(ps) == len(vs)
```

polynomial encoding

concrete evaluation

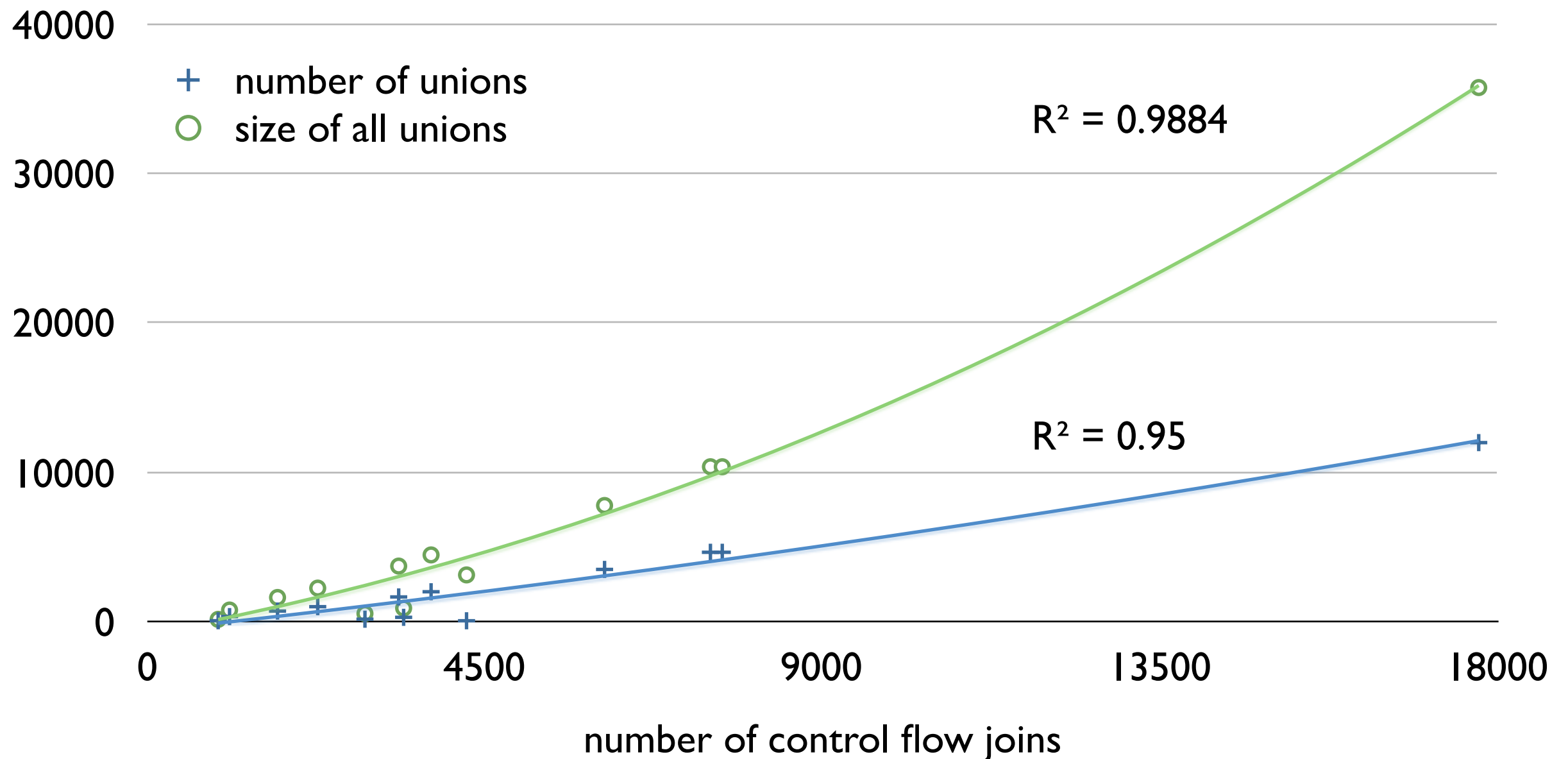
```
g0 = a > 0  
g1 = b > 0  
g2 = g0 ∧ g1  
g3 = ¬(g0 ⇔ g1)  
g4 = ¬g0 ∧ ¬g1  
c = ite(g1, b, a)  
assert g2
```

symbolic virtual machine



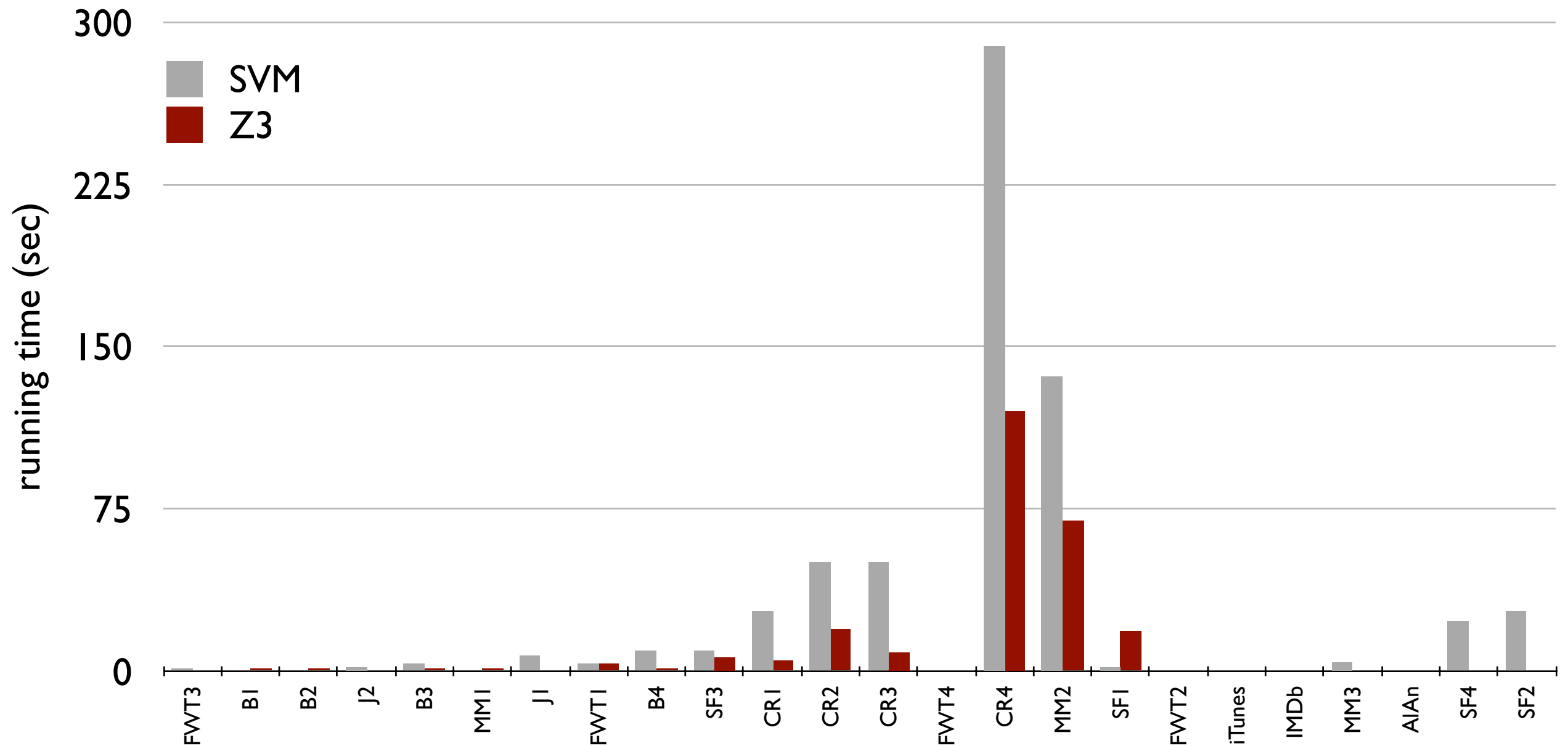
Effectiveness of type-driven state merging

Merging performance for verification and synthesis queries in SynthCL, WebSynth and IFC programs



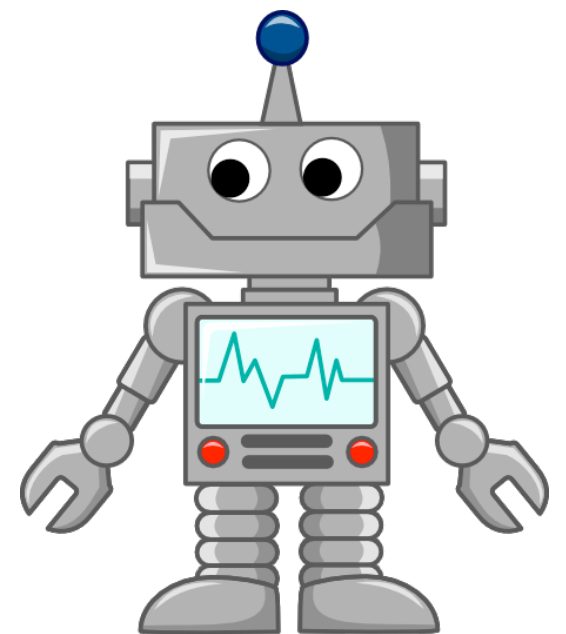
Effectiveness of type-driven state merging

SVM and solving time for verification and synthesis queries in SynthCL, WebSynth and IFC programs



apps

solver-aided programming for everyone



Chlorophyll: ultra low-power computing

Instructions/Second vs Power

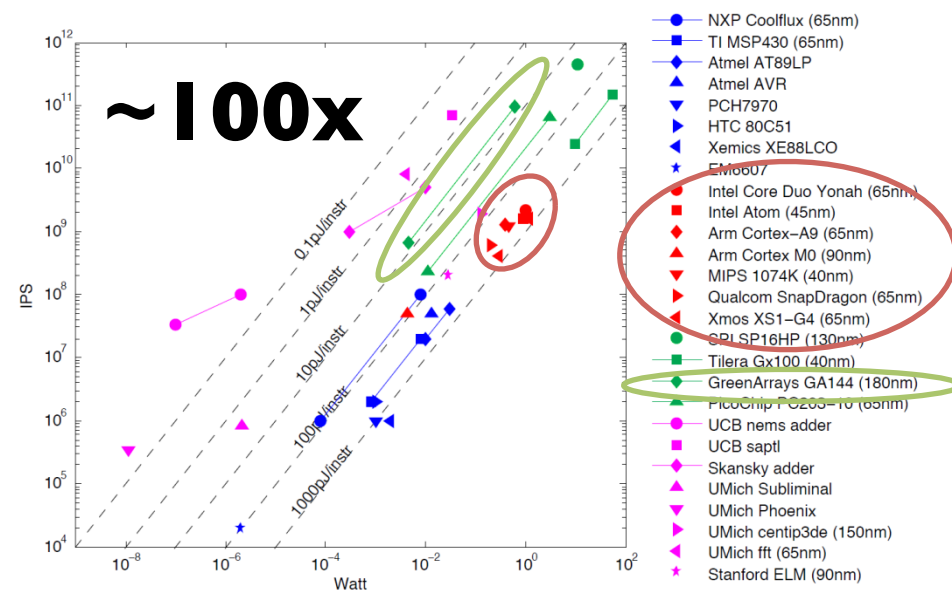
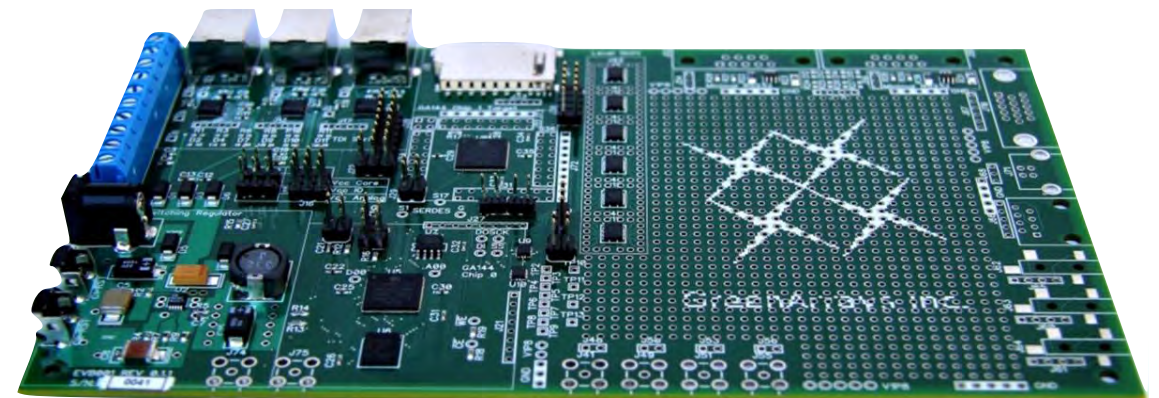


Figure by Per Ljung

GreenArrays GA144 Processor



Chlorophyll: ultra low-power computing

GreenArrays GA144 Processor

- ▶ Stack-based 18-bit architecture
- ▶ 32 instructions
- ▶ 8 x 18 array of asynchronous cores
- ▶ No shared resources (cache, memory)
- ▶ Limited communication, neighbors only
- ▶ < 300 byte memory per core

Manual program partitioning:
break programs up into a pipeline
with a few operations per core.

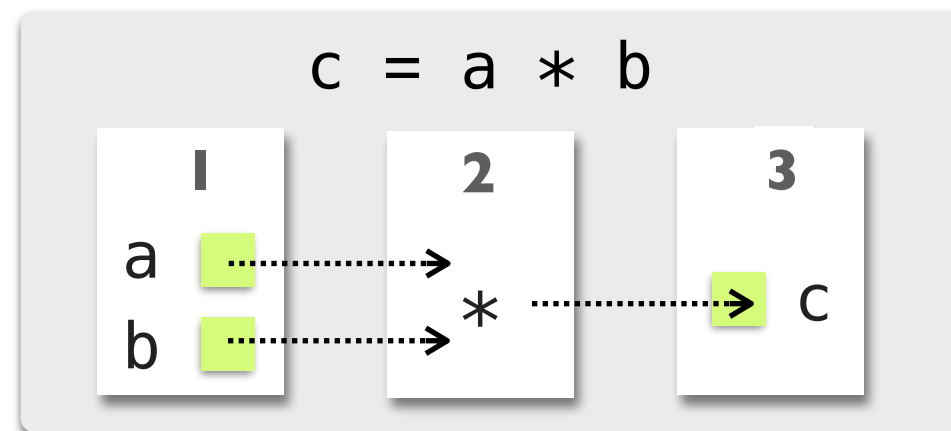


Drawing by Mangpo Phothilimthana

Chlorophyll: ultra low-power computing

GreenArrays GA144 Processor

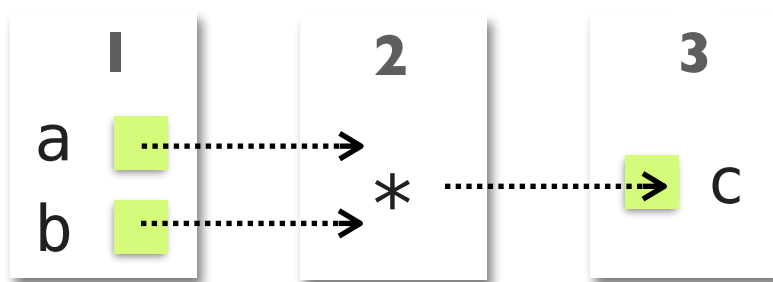
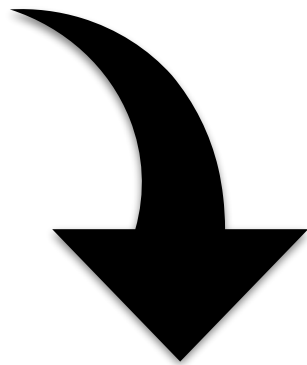
- ▶ Stack-based 18-bit architecture
- ▶ 32 instructions
- ▶ 8 x 18 array of asynchronous cores
- ▶ No shared resources (cache, memory)
- ▶ Limited communication, neighbors only
- ▶ < 300 byte memory per core



Drawing by Mangpo Phothilimthana

Chlorophyll: ultra low-power computing

```
int a, b;  
int c = a * b;
```

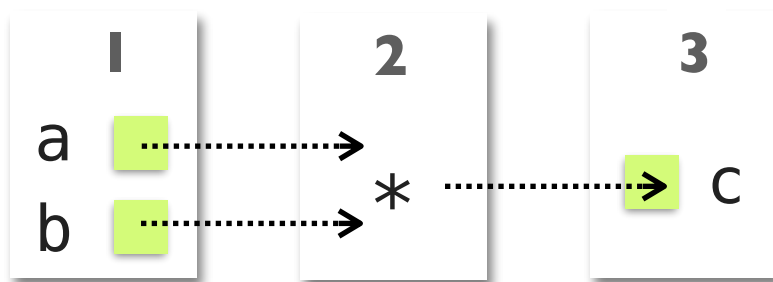


Synthesizes placement of code and data onto cores, by type-checking a program sketch in a C-like DSL.

Chlorophyll: ultra low-power computing

```
int@1 a, b;  
int@3 c = a *@2 b;
```

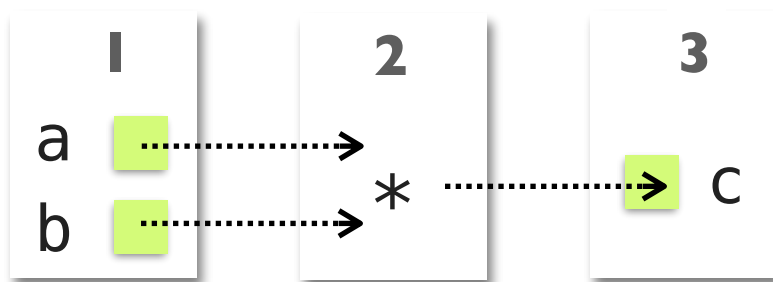
Synthesizes placement of
code and data onto cores, by
type-checking a program
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Chlorophyll: ultra low-power computing

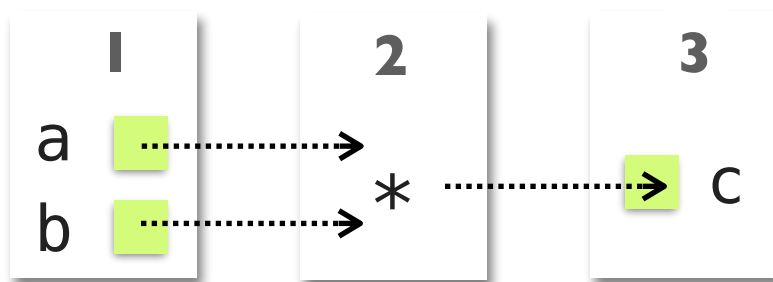
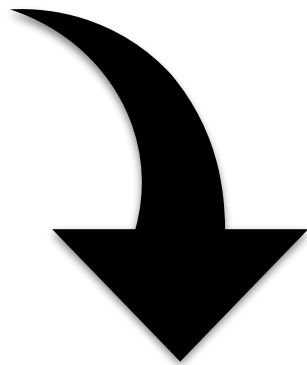
```
int@?? a, b;  
int@?? c = a *@?? b;
```

Synthesizes placement of
code and data onto cores, by
type-checking a program
sketch in a C-like DSL.

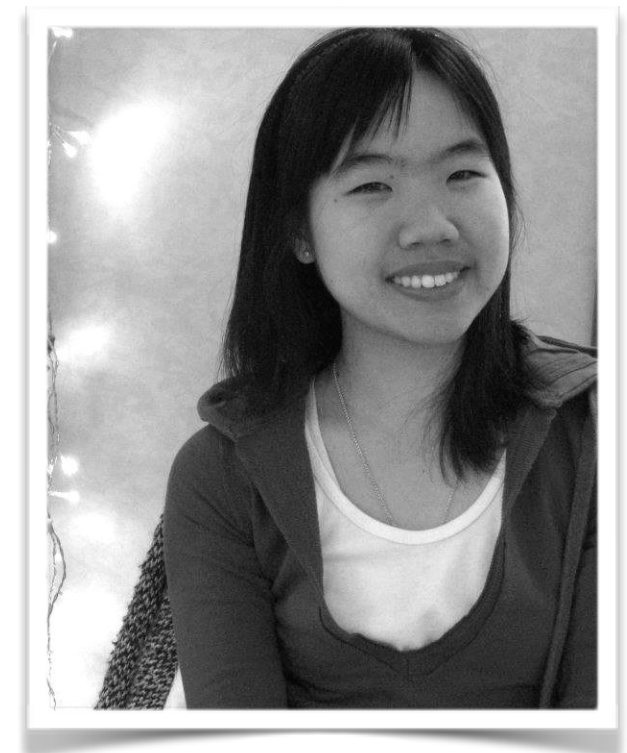


Chlorophyll: ultra low-power computing

```
int@?? a, b;  
int@?? c = a *@?? b;
```



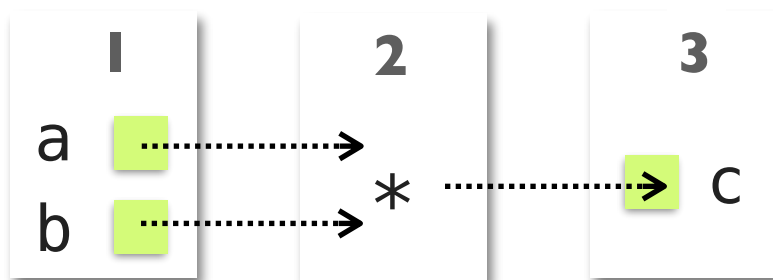
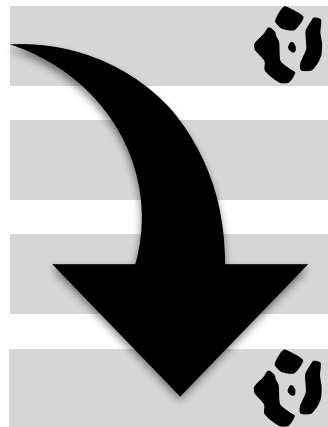
**Built by a first-year
grad in a few weeks**



Phitchaya Mangpo Phothilimthana

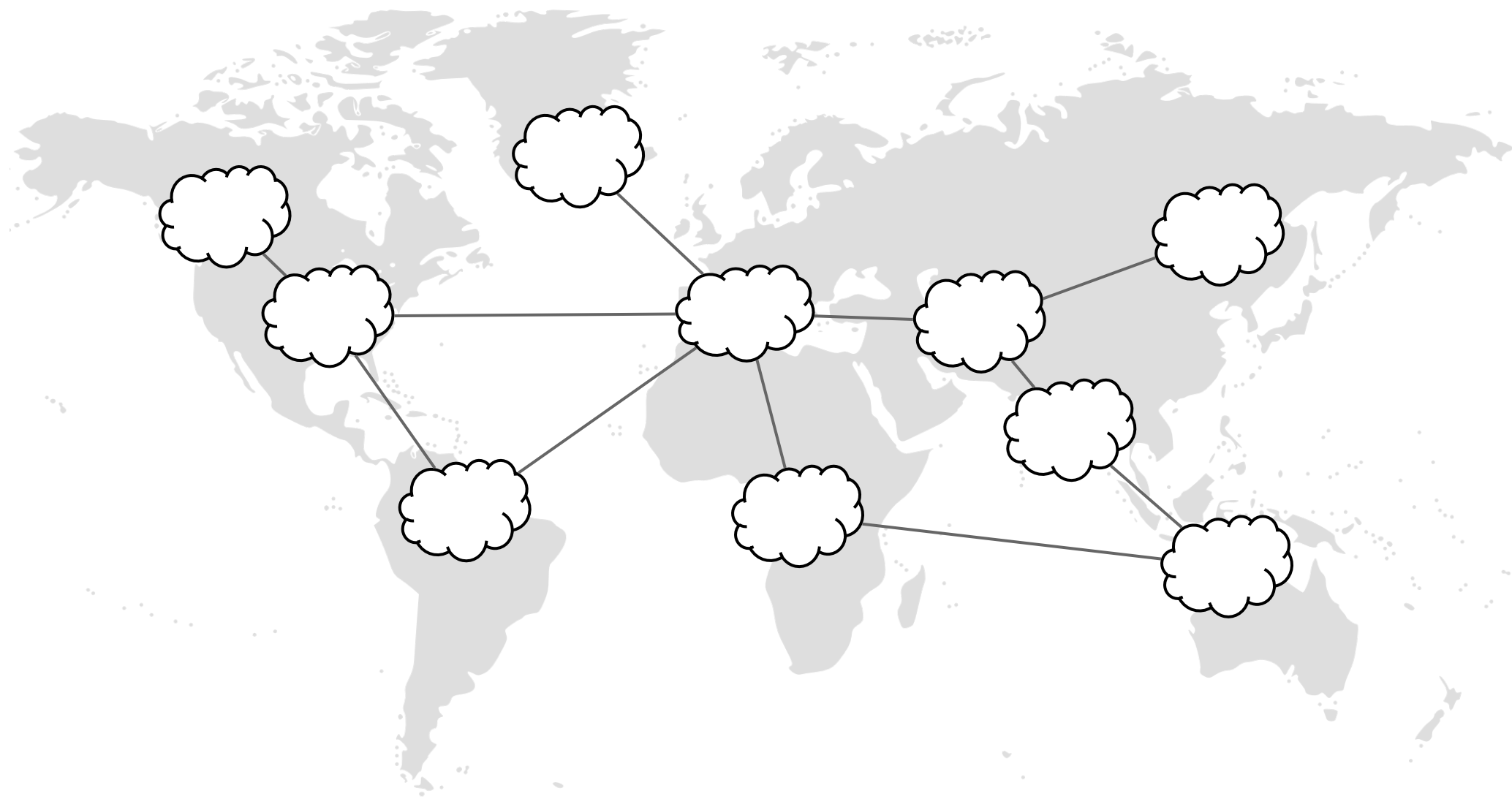
Chlorophyll: ultra low-power computing

```
int@?? a, b;  
int@?? c = a *@?? b;
```

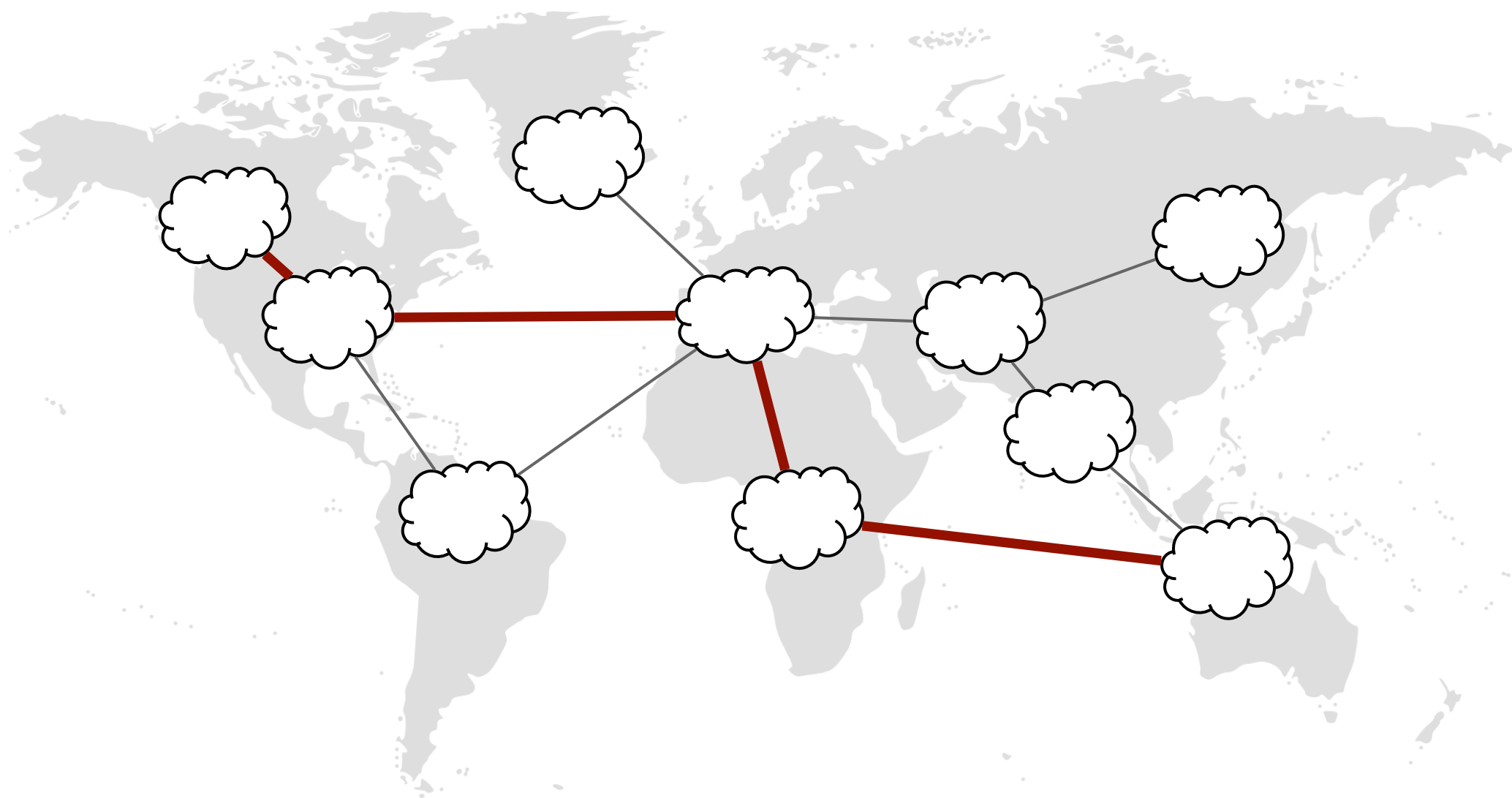


[Phothilimthana et al.,
PLDI'14]

Bagpipe: verifying BGP router configurations

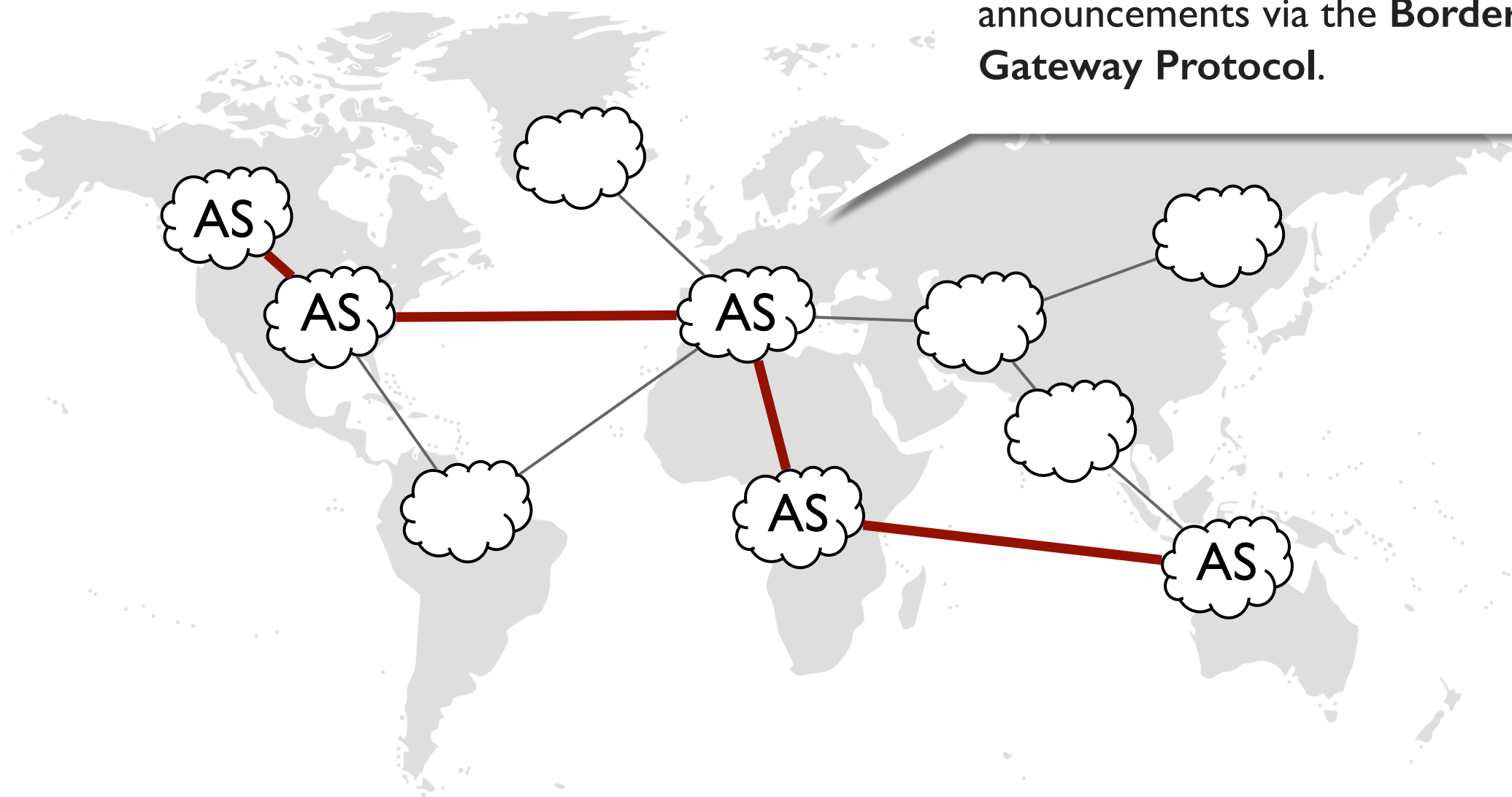


Bagpipe: verifying BGP router configurations



Bagpipe: verifying BGP router configurations

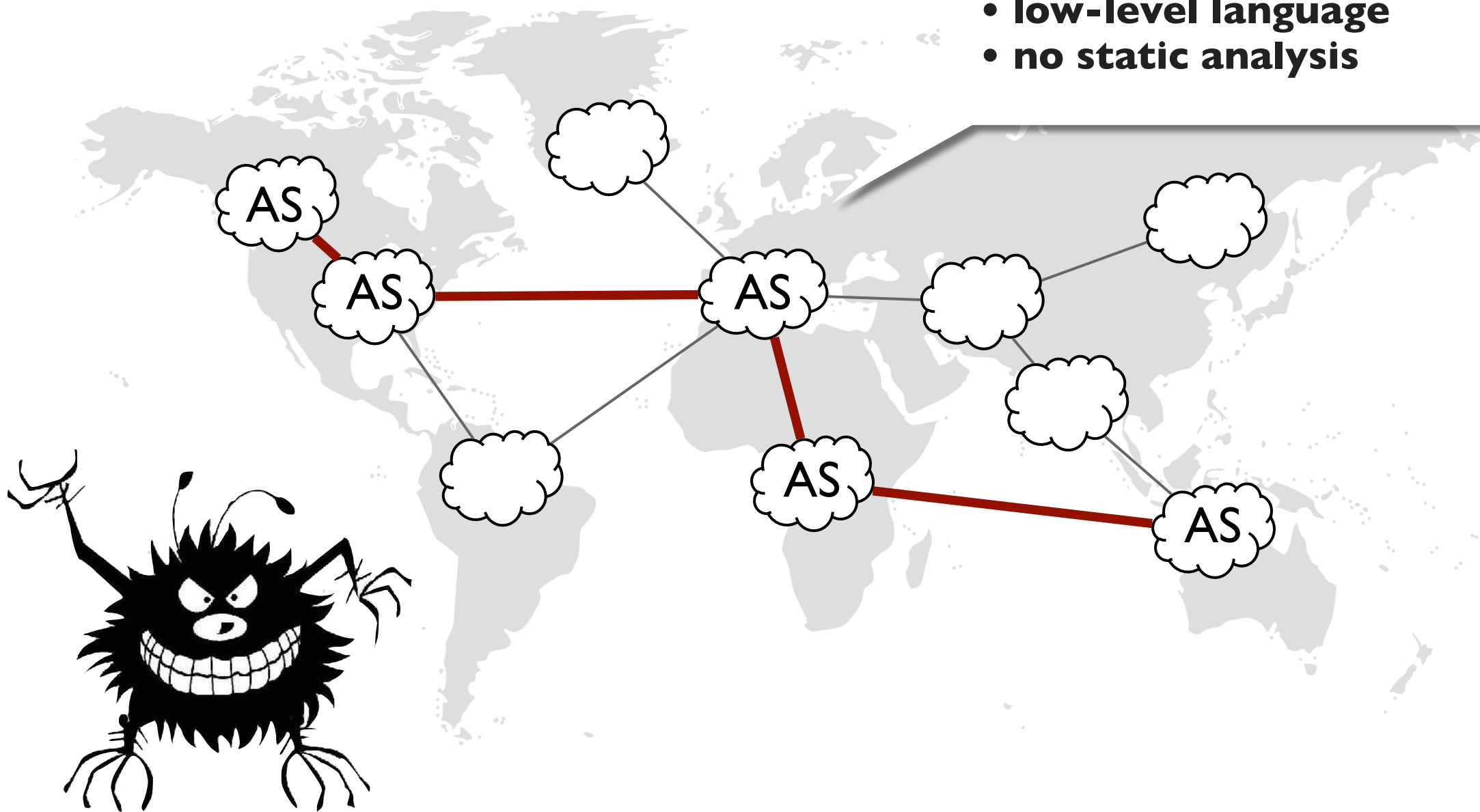
Autonomous systems communicate routing information by sending announcements via the **Border Gateway Protocol**.



Bagpipe: verifying BGP router configurations

Configuring BGP is tricky

- **distributed system**
- **low-level language**
- **no static analysis**



Bagpipe: verifying BGP router configurations

BGP configuration property



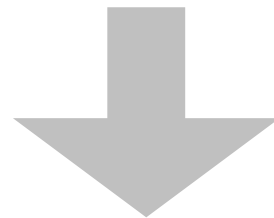
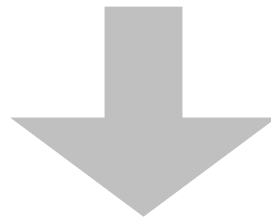
A BGP interpreter
implemented in
Rosette.



policy violation

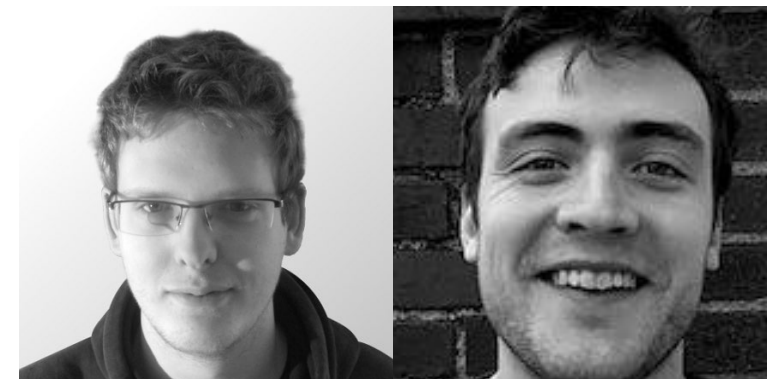
Bagpipe: verifying BGP router configurations

BGP configuration property



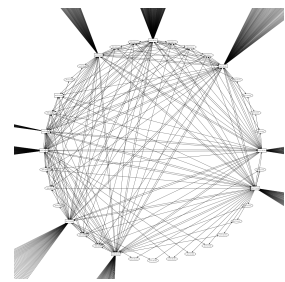
policy violation

Built by two grads in a few weeks



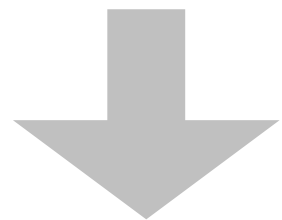
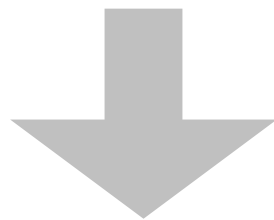
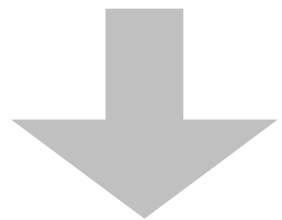
Konstantin Weitz and Doug Woos

Bagpipe: verifying BGP router configurations

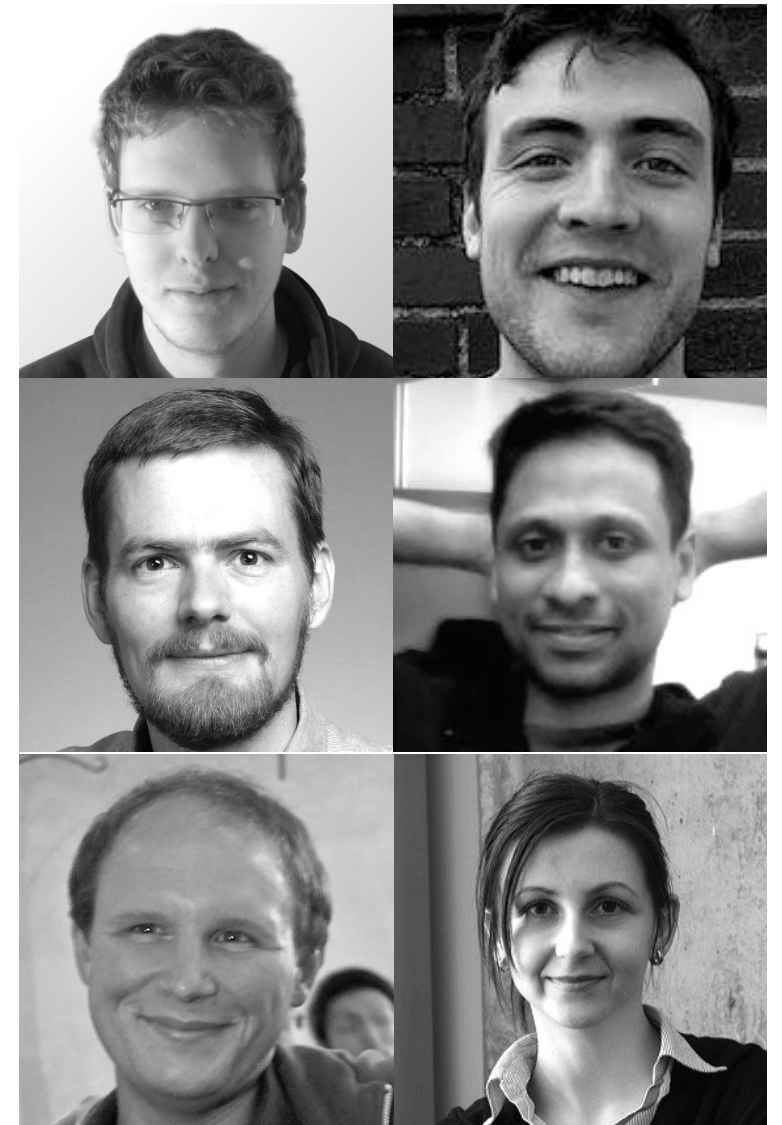


Internet2

**private
announcements
are not leaked**



route leaks!

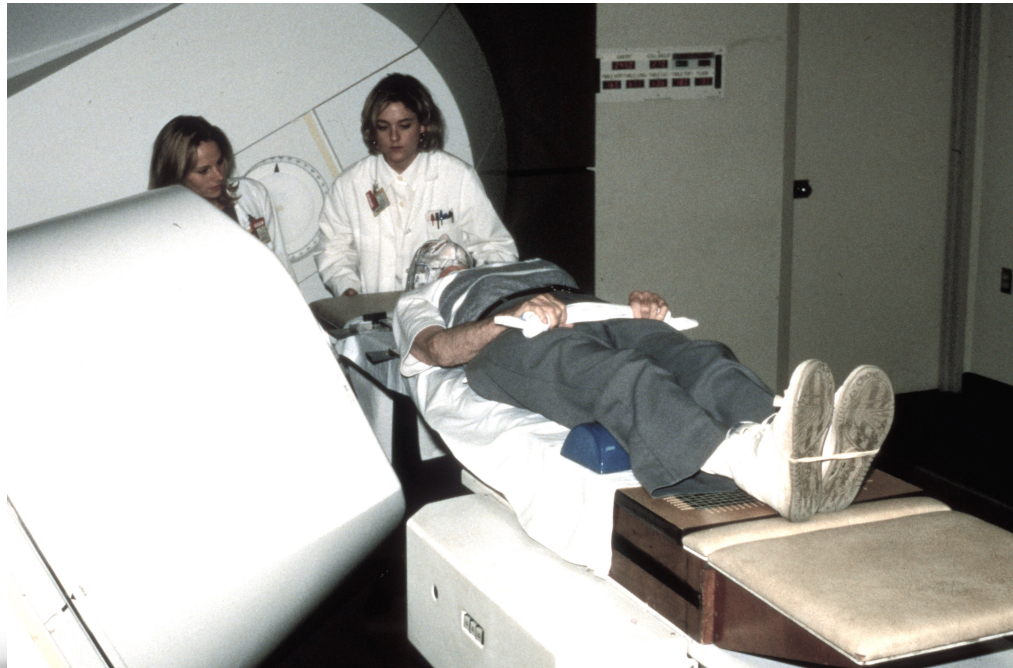


[Weitz et al., **OOPSLA'16**]

Neutrons: verifying a radiotherapy system

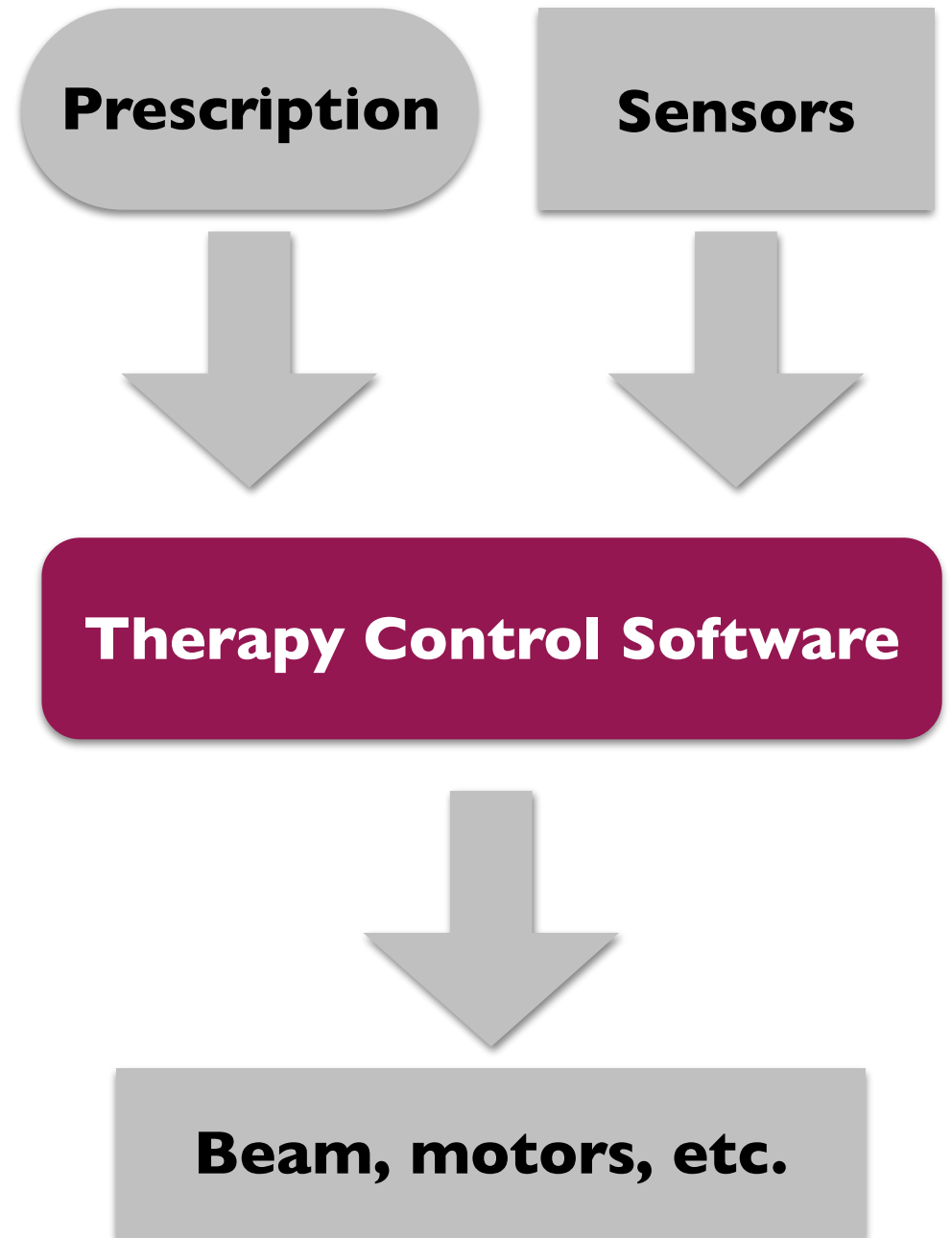
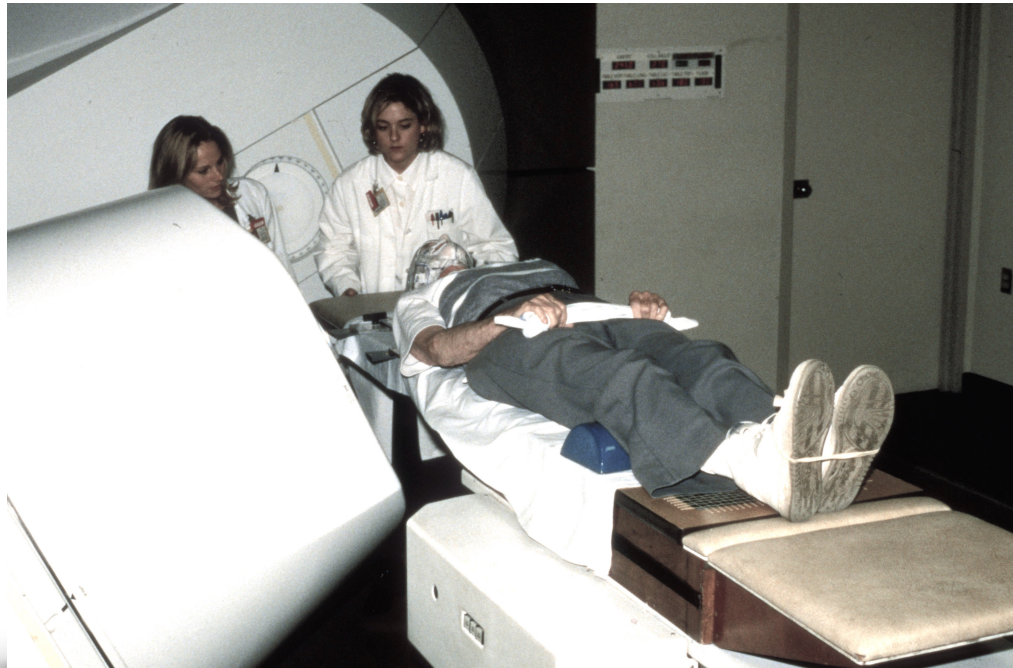
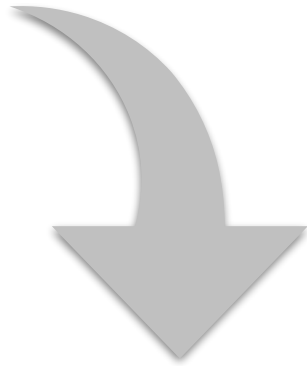
Clinical Neutron Therapy System (CNTS) at UW

- 30 years of incident-free service.
- Controlled by custom software, built by CNTS engineering staff.
- Third generation of Therapy Control software built recently.

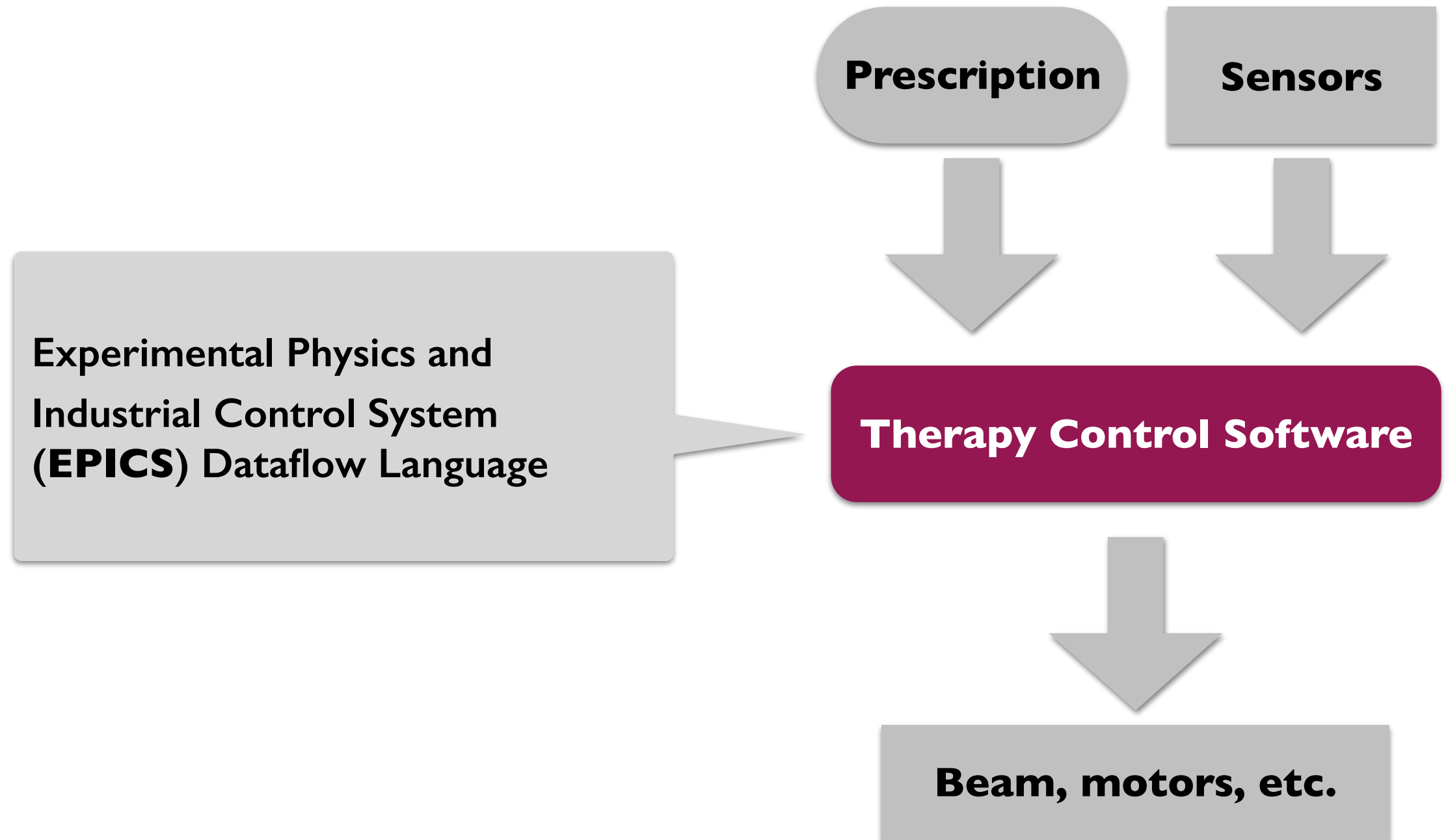


Neutrons: verifying a radiotherapy system

Clinical Neutron Therapy System (CNTS) at UW



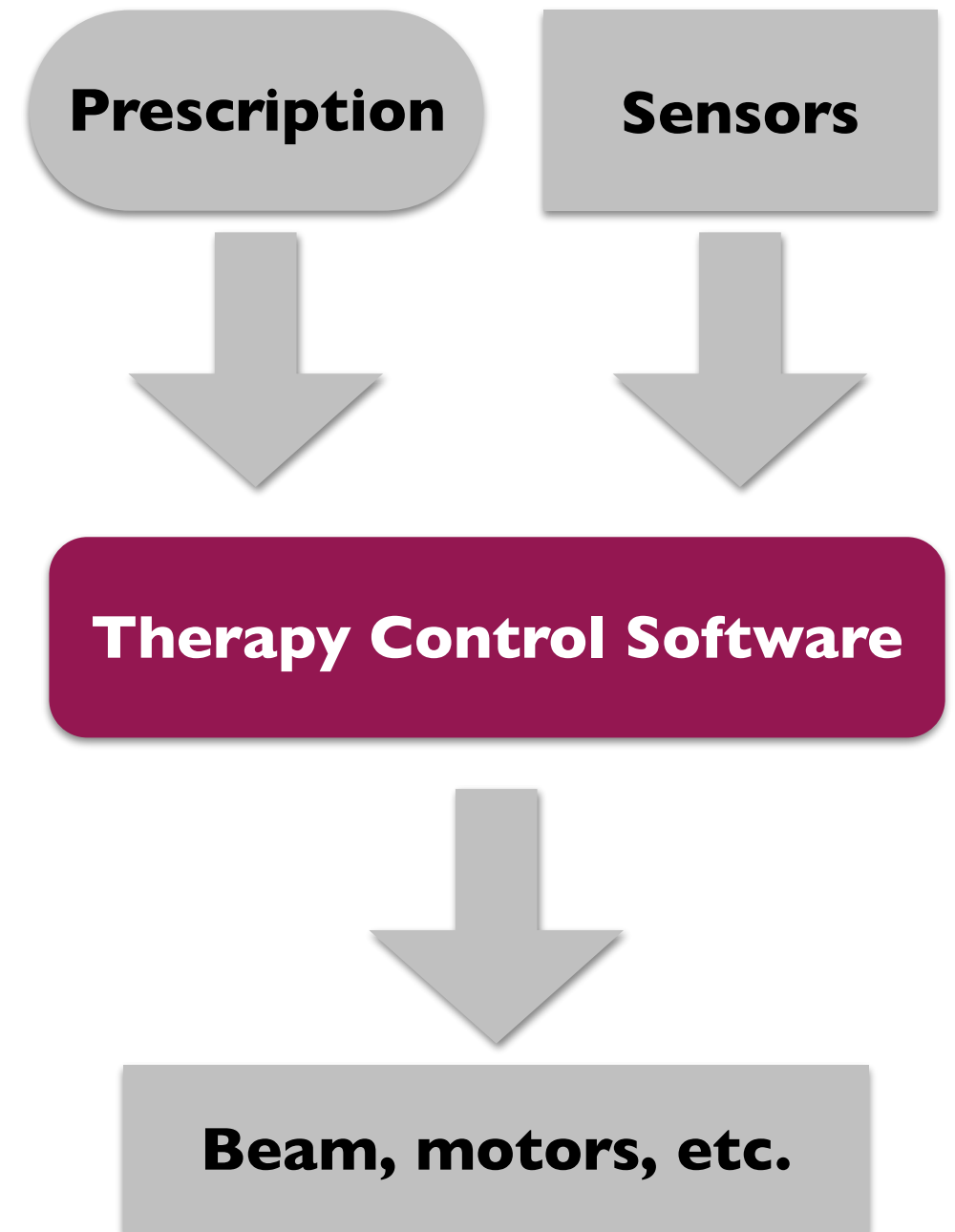
Neutrons: verifying a radiotherapy system



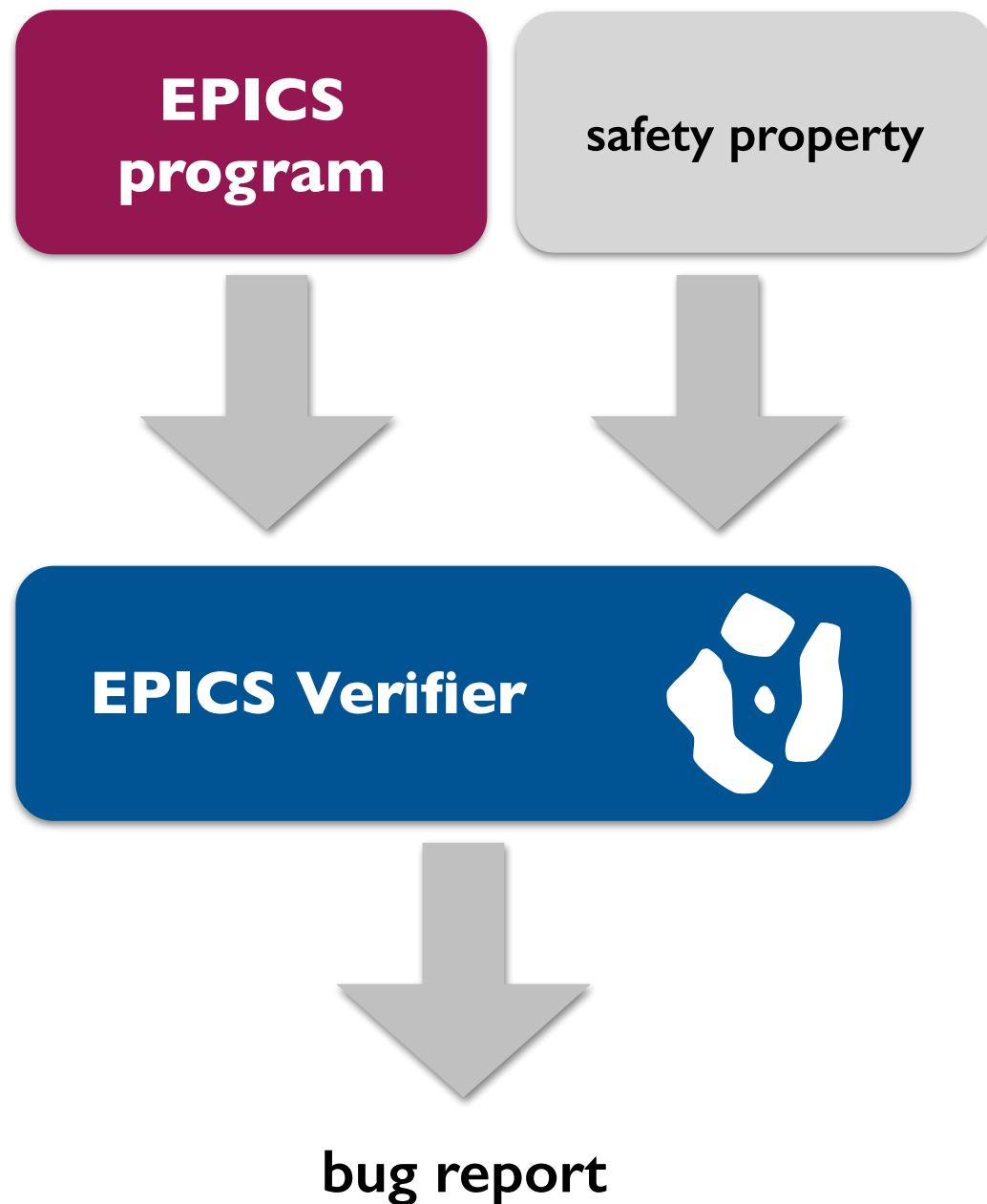
Neutrons: verifying a radiotherapy system

EPICS documentation / semantics

The Maximize Severity attribute is one of NMS (Non-Maximize Severity), MS (Maximize Severity), MSS (Maximize Status and Severity) or MSI (Maximize Severity if Invalid). It determines whether alarm severity is propagated across links. If the attribute is MSI only a severity of `INVALID_ALARM` is propagated; settings of MS or MSS propagate all alarms that are more severe than the record's current severity. For input links the alarm severity of the record referred to by the link is propagated to the record containing the link. For output links the alarm severity of the record containing the link is propagated to the record referred to by the link. If the severity is changed the associated alarm status is set to `LINK_ALARM`, except if the attribute is MSS when the alarm status will be copied along with the severity.



Neutrons: verifying a radiotherapy system

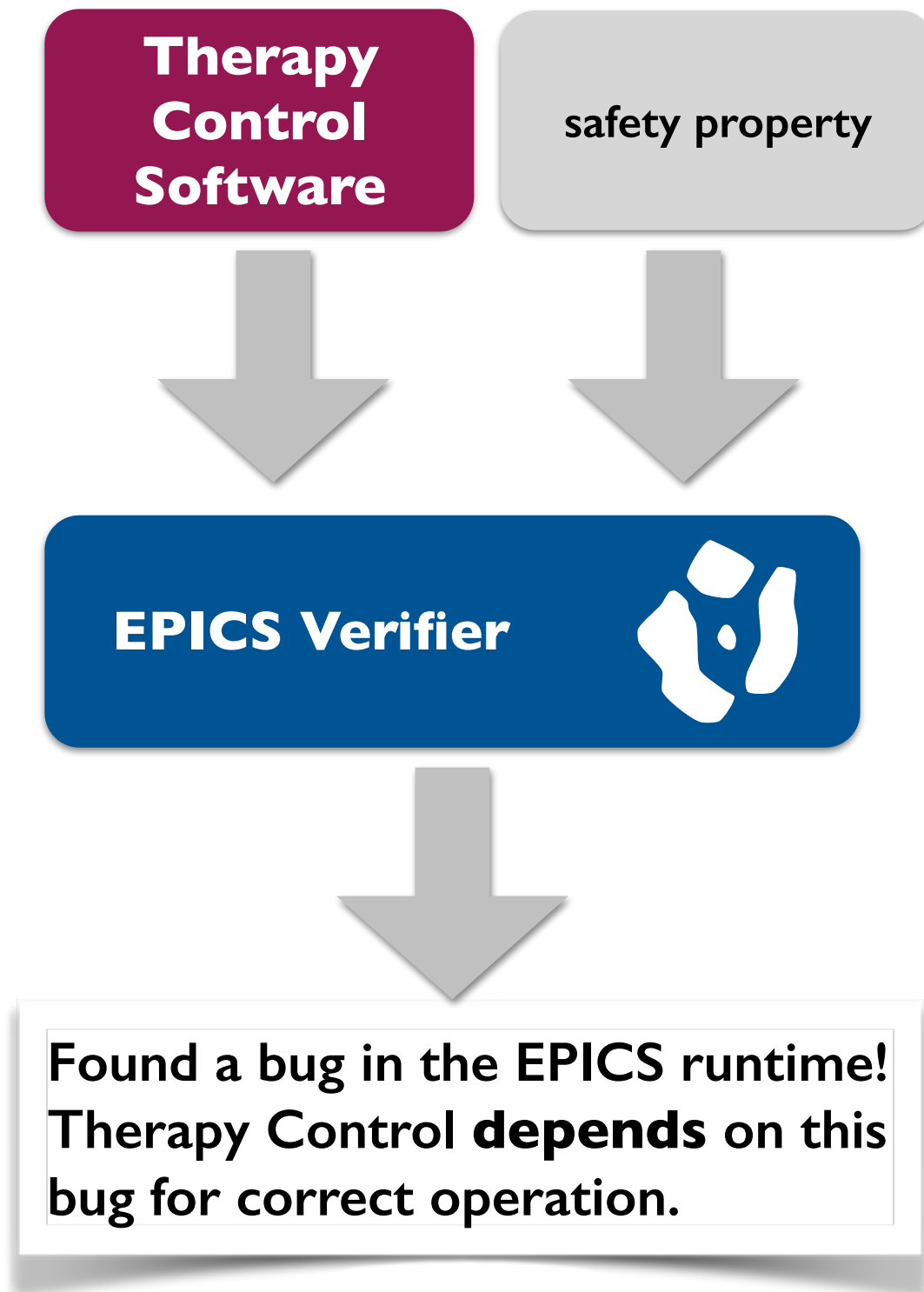


Built by a 2nd year grad in a few days!



Calvin Loncaric

Neutrons: verifying a radiotherapy system



[Pernsteiner et al., CAV'16]



Thanks for a great quarter!

