Computer-Aided Reasoning for Software

CSE507

Combining Theories

courses.cs.washington.edu/courses/cse507/16sp/

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Today

Last lecture

- A survey of theory solvers and deciding $T_=$ with congruence closure

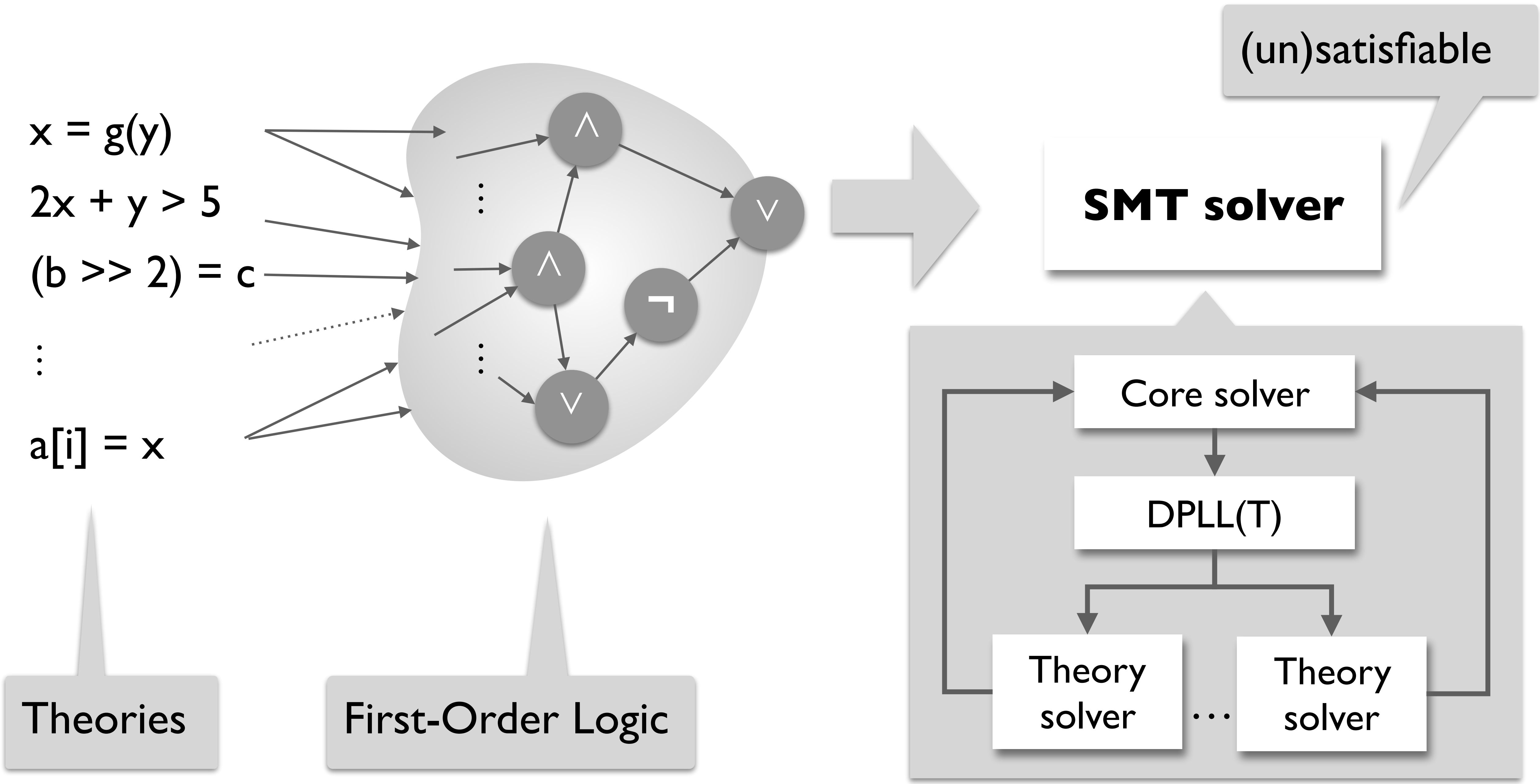
Today

- Deciding a combination of theories

Reminders

- HW1 is due by 11:00 pm
- HW2 is posted
  - Start early
  - Submit self-contained runnable code

Satisfiability Modulo Theories (SMT)



Combining theories with Nelson-Oppen

Σ_1-theory T_1
with axioms A_1

Theory solver

…

Σ_n-theory T_n
with axioms A_n

Theory solver

Combination solver

Theory $T_1 \cup \ldots \cup T_n$ with signature $\Sigma_1 \cup \ldots \cup \Sigma_n$ and axioms $A_1 \cup \ldots \cup A_n$

Combining theories with Nelson-Oppen

Σ_1-theory T_1
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Theory solver

Σ_2-theory T_2
with axioms A_2

Theory solver

We'll see how to combine two theories. Easy to generalize to n.

Combination solver

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We'll see how to combine two theories. Easy to generalize to n.

Combination solver

Theory $T_1 \cup T_2$ with signature $\Sigma_1 \cup \Sigma_2$ and axioms $A_1 \cup A_2$

The combination problem is undecidable for arbitrary (decidable) theories. It becomes decidable under **Nelson-Oppen restrictions**.

Nelson-Oppen restrictions

T_1 and T_2 can be combined when

- Both are decidable, quantifier-free conjunctive fragments
- Equality (=) is the only symbol in the intersection of their signatures: $\Sigma_1 \cap \Sigma_2 = \{ = \}$
- Both are **stably infinite**

Nelson-Oppen restrictions

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- Both are **stably infinite**

A theory T is stably infinite if for every satisfiable Σ_T-formula F, there is a T-model that satisfies F and that has a universe of infinite cardinality.

Examples of (non-)stably infinite theories

Σ_T : $\{ a, b, = \}$

A_T: $\forall x . x = a \vee x = b$

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✗

Fixed width bit vectors (T_{bv})

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Equality and uninterpreted functions ($T_=$)

Fixed width bit vectors (T_{bv}) ✗

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Equality and uninterpreted functions ($T_=$) ✓

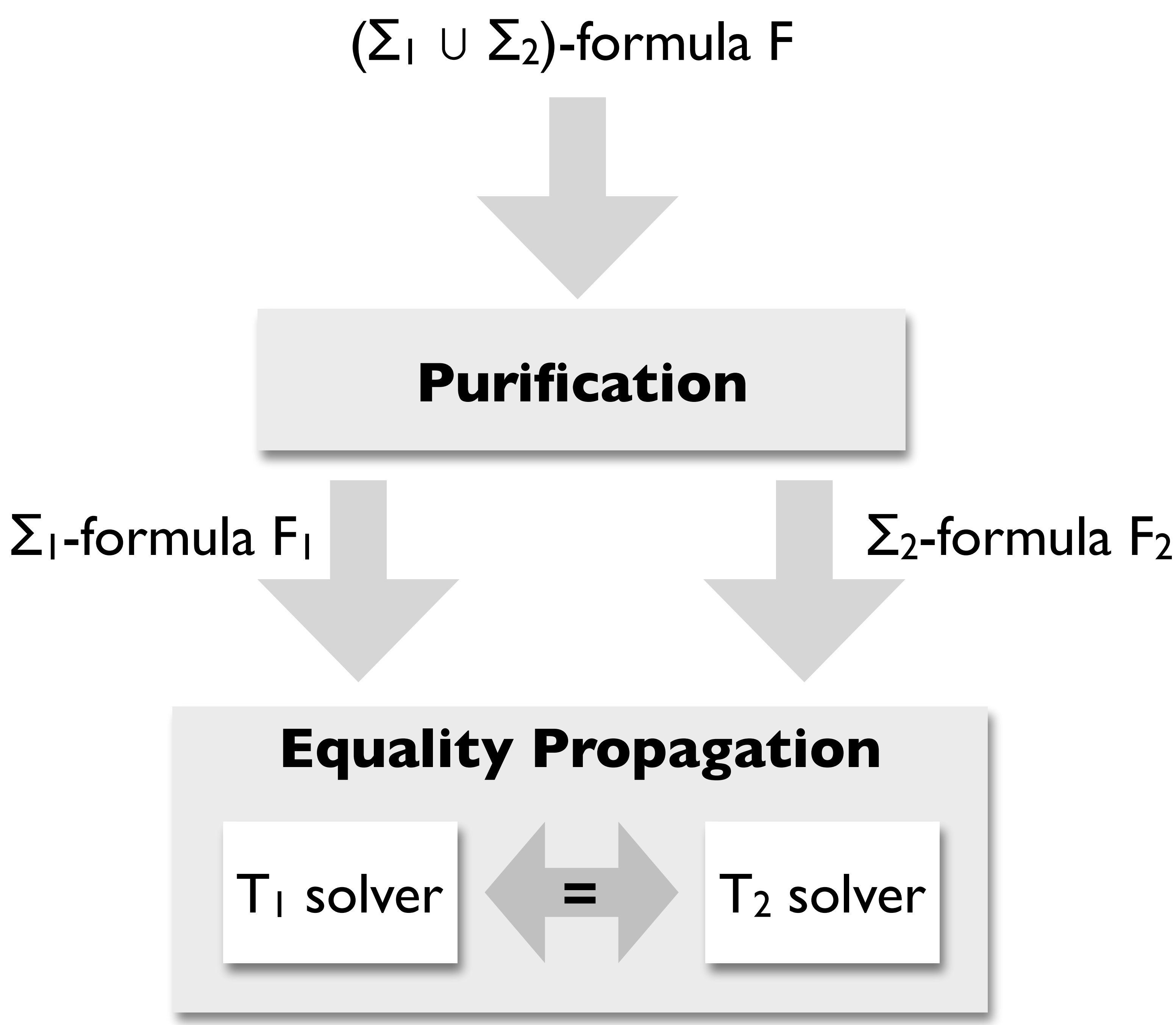
Fixed width bit vectors (T_{bv}) ✗

Linear real arithmetic (T_R) ✓

Arrays (T_A) ✓

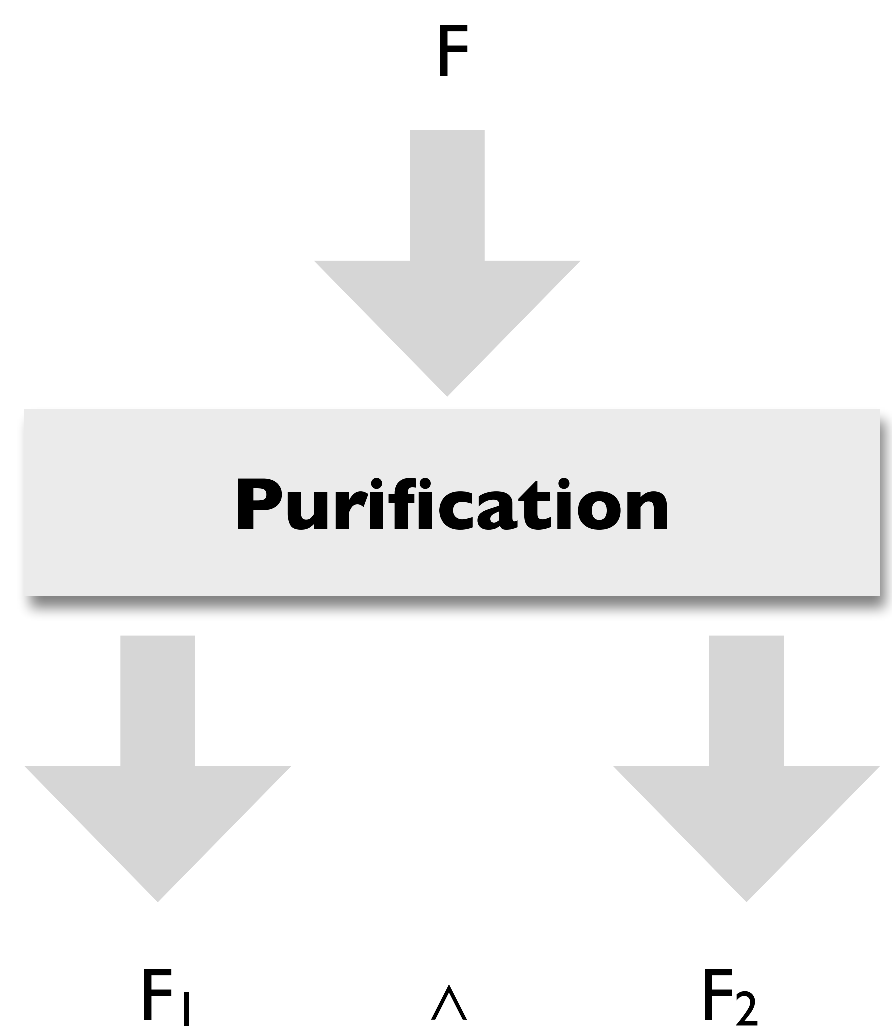
Linear integer arithmetic (T_R) ✓

Overview of Nelson-Oppen



Overview of purification

Transforms a $(\Sigma_1 \cup \Sigma_2)$-formula F into an equisatisfiable formula $F_1 \wedge F_2$ with F_1 in T_1 and F_2 in T_2



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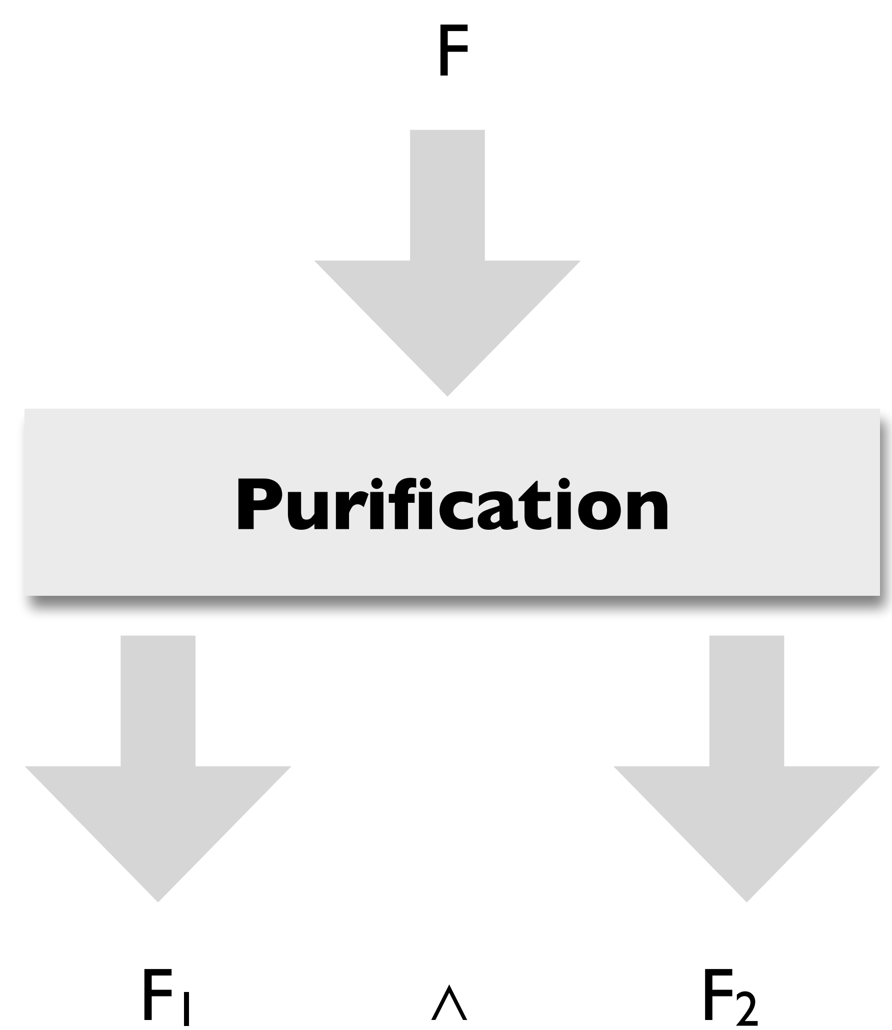

Repeat until fix point:

- If f is in T_i and t is not, and u is fresh:
  $F[f(\ldots, t, \ldots)] \rightsquigarrow F[f(\ldots, u, \ldots)] \wedge u = t$
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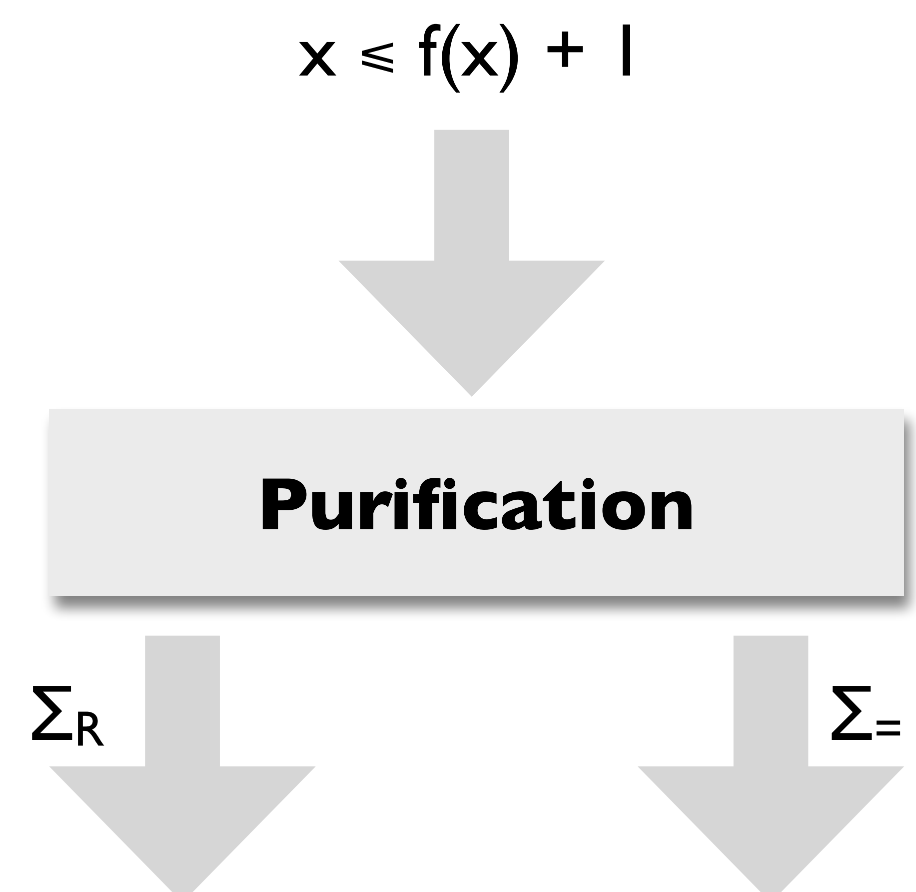

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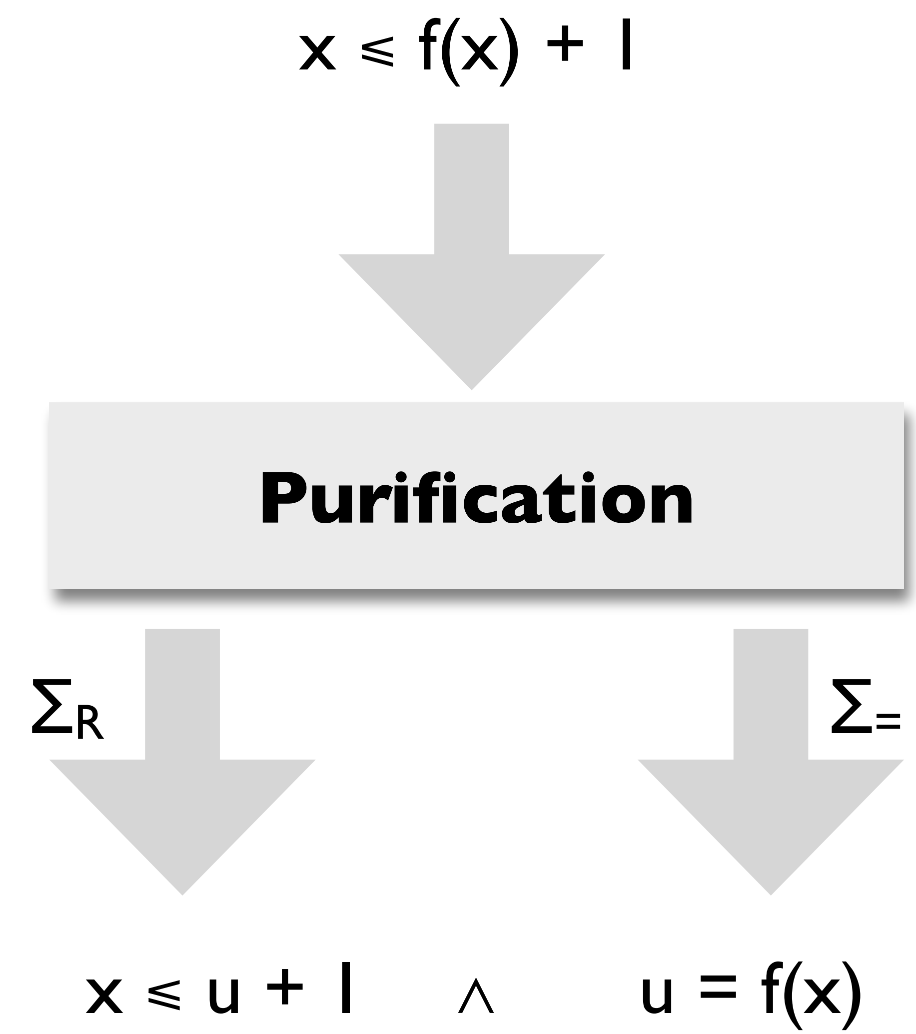

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Another purification example

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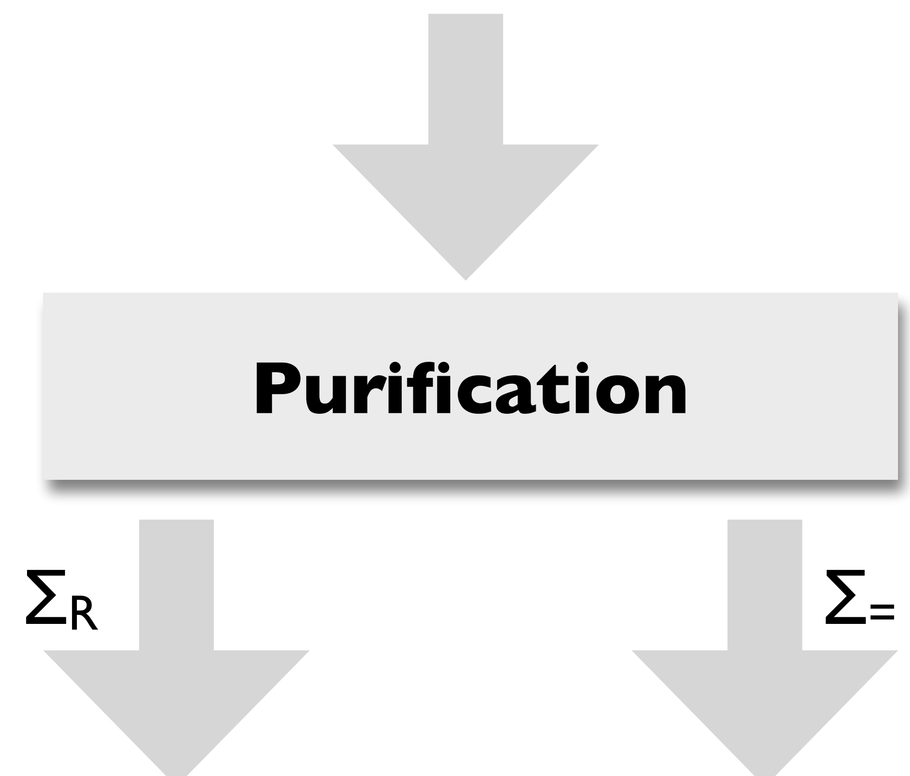
$f(x + g(y)) \leqslant g(a) + f(b)$



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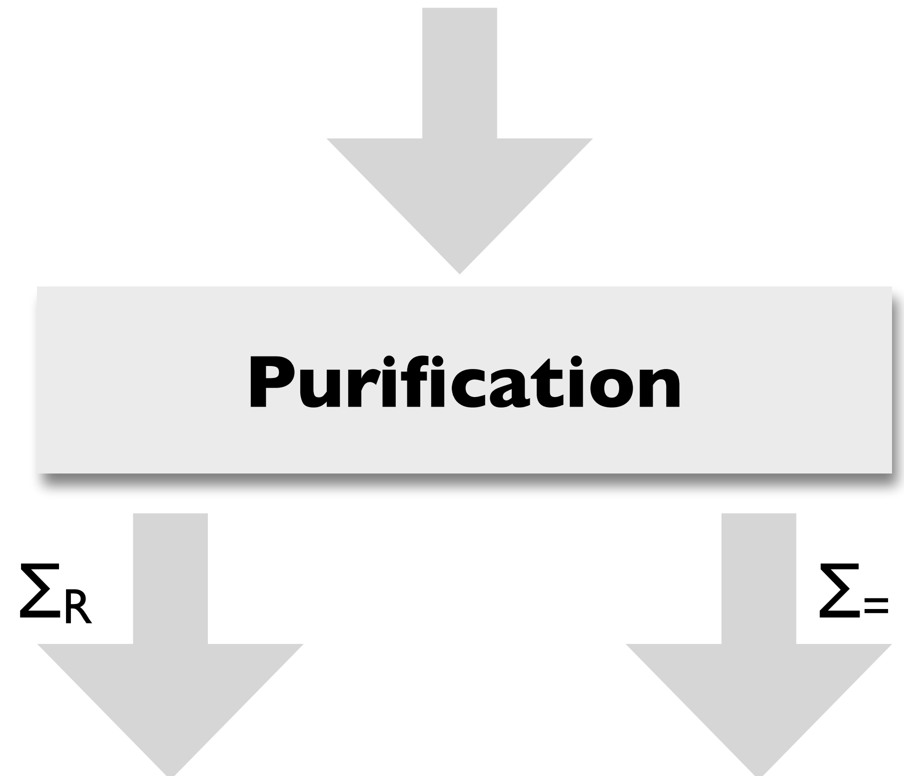

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$f(x + u_1) \leqslant u_2 + u_3$

Purification

Σ_R

$\Sigma_=$

$u_1 = g(y)$
$u_2 = g(a)$
$u_3 = f(b)$

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$f(u_4) \leqslant u_2 + u_3$

Purification

Σ_R

$u_4 = x + u_1$

$\Sigma_=$

$u_1 = g(y)$
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$u_3 = f(b)$
$u_5 = f(u_4)$

Shared and local constants

A constant is *shared* if it occurs in both F_1 and F_2, and it is *local* otherwise.

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$u_3 = f(b)$
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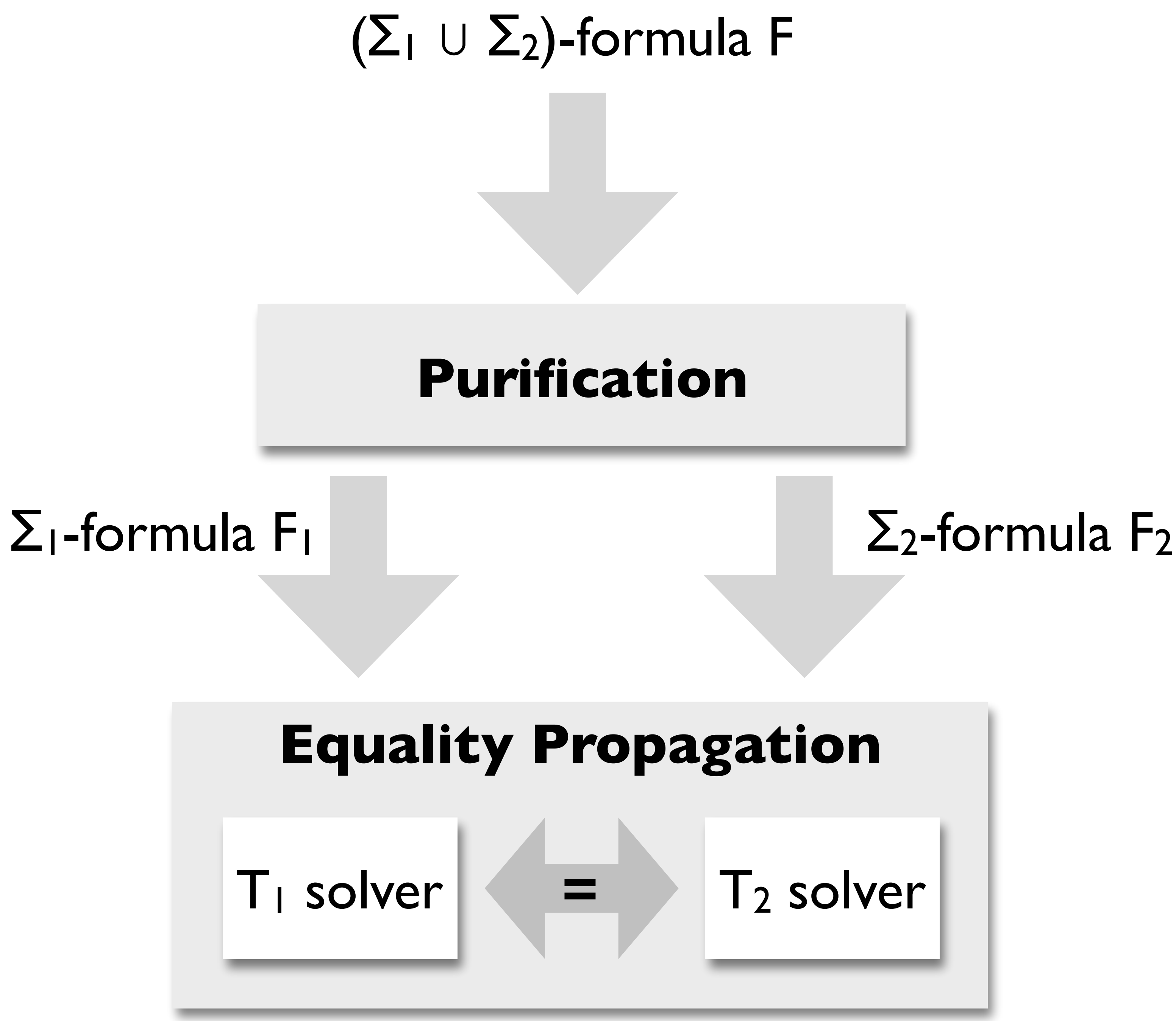
Shared: {u1, u2, u3, u4, u5}
Local: {x, y, a, b}

Overview of Nelson-Oppen



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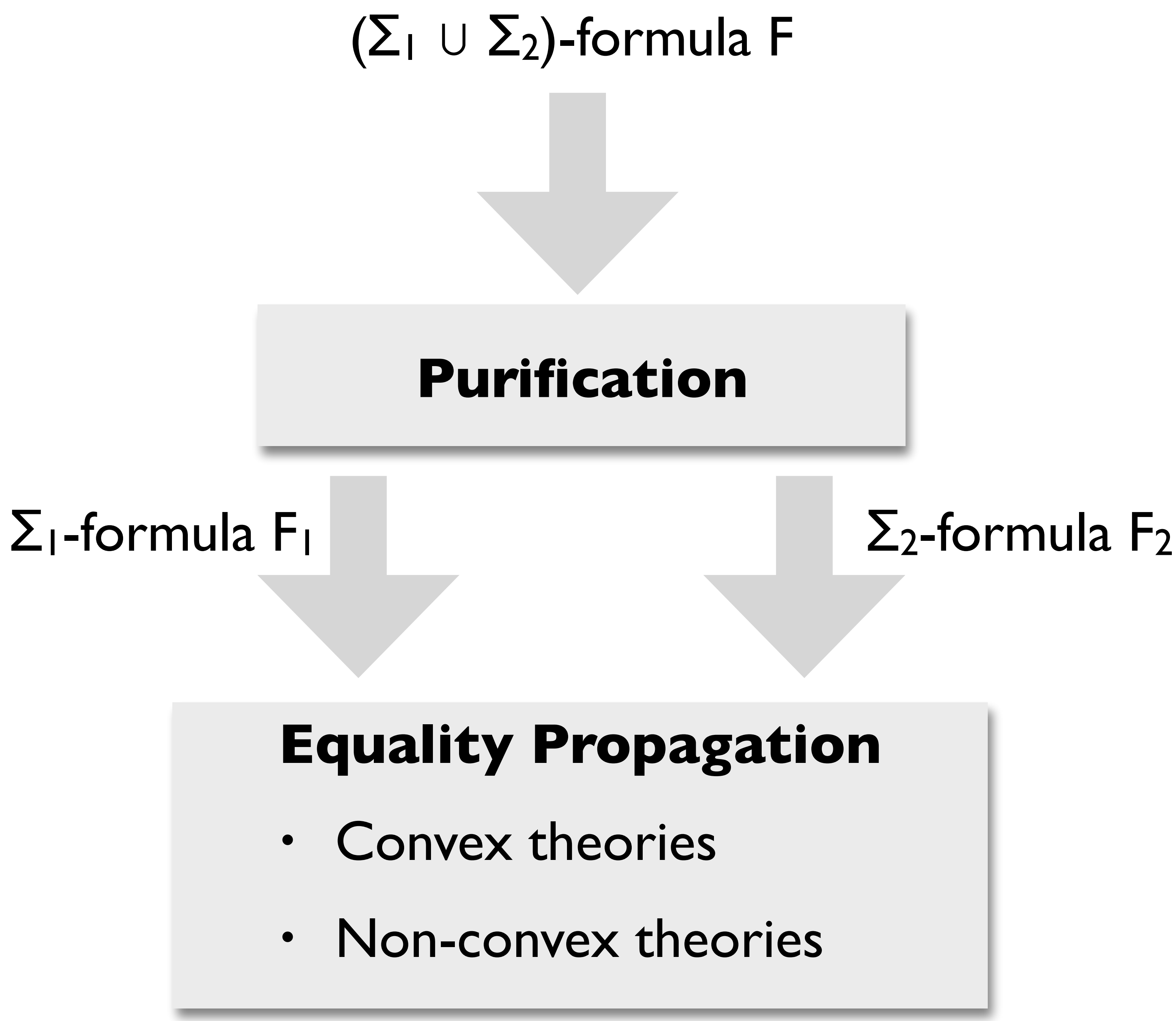

Convex theories

A theory T is *convex* if for every conjunctive formula F, the following holds:

If $F \Rightarrow x_1 = y_1 \vee \ldots \vee x_n = y_n$ for a finite $n > 1$,

then $F \Rightarrow x_i = y_i$ for some $i \in \{1, \ldots, n\}$.

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If F implies a disjunction of equalities, then it also implies at least one of the equalities.

Examples of (non-)convex theories

Linear arithmetic over integers (T_Z)

Examples of (non-)convex theories

Linear arithmetic over integers (T_Z) ✗

$1 \leq x \wedge x \leq 2 \Rightarrow x = 1 \vee x = 2$ but

not $1 \leq x \wedge x \leq 2 \Rightarrow x = 1$

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Equality and uninterpreted functions ($T_=$) ✓

Linear real arithmetic (T_R) ✓

Nelson-Oppen for convex theories

NELSON-OPPEN-CONVEX(F)

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Nelson-Oppen for convex theories

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Is F satisfiable if both F_1 and F_2 are satisfiable? **No**:

$x = 1 \wedge 2 = x + y \wedge f(x) \neq f(y)$

Nelson-Oppen for convex theories

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Nelson-Oppen for convex theories: example

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$$f(f(x) - f(y)) \neq f(z) \wedge x \leq y \wedge y + z \leq x \wedge 0 \leq z$$

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$w = u - v$	
$x = y \wedge$	$x = y \wedge$
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$x = y \wedge$ $u = v \wedge$ $w = z \wedge$	$x = y \wedge$ $u = v \wedge$ $w = z \wedge$ UNSAT

This doesn't work for non-convex theories ...

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$$1 \leq x \wedge x \leq 2 \wedge$$
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Σ_Z	$\Sigma_=$
$1 \leq x \wedge$ $x \leq 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \neq f(z_1) \wedge$ $f(x) \neq f(z_2)$

This doesn't work for non-convex theories ...

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2. Run T_1-solver on F_1 and T_2-solver on F_2 and return UNSAT if either is unsatisfiable
3. If there are shared constants x and y such that $F_i \Rightarrow x = y$ but F_j does not
    1. $F_j \leftarrow F_j \wedge x = y$
    2. Go to step 2.
4. Return SAT

$$1 \leq x \wedge x \leq 2 \wedge$$
$$f(x) \neq f(1) \wedge f(x) \neq f(2)$$

$1 \leq x \wedge$ $x \leq 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \neq f(z_1) \wedge$ $f(x) \neq f(z_2)$
SAT	SAT
Σ_Z	$\Sigma_=$

This doesn't work for non-convex theories …

NELSON-OPPEN-CONVEX(F)

1. Purify F into $F_1 \wedge F_2$
2. Run T_1-solver on F_1 and T_2-solver on F_2 and return UNSAT if either is unsatisfiable
3. If there are shared constants x and y such that $F_i \Rightarrow x = y$ but F_j does not
    1. $F_j \leftarrow F_j \wedge x = y$
    2. Go to step 2.
4. Return SAT

If T is non-convex, it may imply a disjunction of equalities without implying any single equality.

We have to propagate disjunctions as well as individual equalities. Which disjunctions? How do we propagate disjunctions to theory solvers which reason only about conjunctions?

Nelson-Oppen for non-convex theories

NELSON-OPPEN(F)

1. Purify F into $F_1 \wedge F_2$
2. Run T_1-solver on F_1 and T_2-solver on F_2 and return UNSAT if either is unsatisfiable
3. If there are shared constants x and y such that $F_i \Rightarrow x = y$ but F_j does not
    1. $F_j \leftarrow F_j \wedge x = y$
    2. Go to step 2.
4. If $F_i \Rightarrow x_1 = y_1 \vee \ldots \vee x_n = y_n$ but F_j does not, then if NELSON-OPPEN($F_i \wedge F_j \wedge x_k = y_k$) outputs SAT for any k, return SAT. Otherwise, return UNSAT.
5. Return SAT

Nelson-Oppen for non-convex theories

NELSON-OPPEN(F)

1. Purify F into $F_1 \wedge F_2$
2. Run T_1-solver on F_1 and T_2-solver on F_2 and return UNSAT if either is unsatisfiable
3. If there are shared constants x and y such that $F_i \Rightarrow x = y$ but F_j does not
    1. $F_j \leftarrow F_j \wedge x = y$
    2. Go to step 2.
4. If $F_i \Rightarrow x_1 = y_1 \vee \ldots \vee x_n = y_n$ but F_j does not, then if NELSON-OPPEN($F_i \wedge F_j \wedge x_k = y_k$) outputs SAT for any k, return SAT. Otherwise, return UNSAT.
5. Return SAT

Propagate a *minimal* disjunction.

Nelson-Oppen for non-convex theories: example

$$1 \leq x \wedge x \leq 2 \wedge$$
$$f(x) \neq f(1) \wedge f(x) \neq f(2)$$

Nelson-Oppen for non-convex theories: example

$$1 \leq x \wedge x \leq 2 \wedge$$
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$1 \leq x \wedge$ $x \leq 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \neq f(z_1) \wedge$ $f(x) \neq f(z_2)$
Σ_Z	$\Sigma_=$

Nelson-Oppen for non-convex theories: example

$$1 \leq x \land x \leq 2 \land$$
$$f(x) \neq f(1) \land f(x) \neq f(2)$$

$1 \leq x \land$ $x \leq 2 \land$ $z_1 = 1 \land$ $z_2 = 2$	$f(x) \neq f(z_1) \land$ $f(x) \neq f(z_2)$
$(x=z_1 \lor x=z_2) \land$ Σ_Z	$\Sigma_=$

Nelson-Oppen for non-convex theories: example

$1 \le x \wedge x \le 2 \wedge$

$f(x) \ne f(1) \wedge f(x) \ne f(2)$

$1 \le x \wedge$ $x \le 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \ne f(z_1) \wedge$ $f(x) \ne f(z_2)$
$(x=z_1 \vee x=z_2) \wedge$	
Σ_Z	$\Sigma_=$

$1 \le x \wedge$ $x \le 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \ne f(z_1) \wedge$ $f(x) \ne f(z_2)$
$x = z_1$	$x = z_1 \wedge$ UNSAT

Nelson-Oppen for non-convex theories: example

$$1 \le x \wedge x \le 2 \wedge$$
$$f(x) \ne f(1) \wedge f(x) \ne f(2)$$

$1 \le x \wedge$ $x \le 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \ne f(z_1) \wedge$ $f(x) \ne f(z_2)$
$(x=z_1 \vee x=z_2) \wedge$ Σ_Z	 $\Sigma_=$

$1 \le x \wedge$ $x \le 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \ne f(z_1) \wedge$ $f(x) \ne f(z_2)$
$x = z_1$	$x = z_1 \wedge$ UNSAT

$1 \le x \wedge$ $x \le 2 \wedge$ $z_1 = 1 \wedge$ $z_2 = 2$	$f(x) \ne f(z_1) \wedge$ $f(x) \ne f(z_2)$
$x = z_2$	$x = z_2 \wedge$ UNSAT

Soundness and completeness of Nelson-Oppen

If the theories T_1 and T_2 satisfy Nelson-Open restrictions, then the combination procedure returns UNSAT for a formula F in $T_1 \cup T_2$ iff F is unsatisfiable modulo $T_1 \cup T_2$.

Complexity of Nelson-Oppen

If decision procedures for convex theories T_1 and T_2 have polynomial time complexity, so does their Nelson-Oppen combination.

If decision procedures for non-convex theories T_1 and T_2 have NP time complexity, so does their Nelson-Oppen combination.

Summary

Today

- Sound and complete procedure for a combination of restricted theories
- Stably infinite, conjunctive, quantifier-free with signatures that are disjoint except for =

Next lecture

- Deciding satisfiability of arbitrary boolean combinations of quantifier-free first-order formulas