

Theorem [Fundamental Theorem of Statistical Learning]

Let H be a hypothesis class, from $\mathcal{X}$ to $\{0,1\}$,
and the loss function is a 0-1 loss, then the
following are equivalent.

1. H has the uniform convergence property
2. Any ERM rule is a successful agnostic PAC learner of H .
3. H is agnostic PAC learnable.
4. H is PAC learnable
5. Any ERM rule is a successful PAC learner for H
6. H has a finite VC-dimension.

We already learned that

1 $\rightarrow$ 2	(lecture 3)
2 $\rightarrow$ 3	trivial
3 $\rightarrow$ 4	trivial
2 $\rightarrow$ 5	trivial
2 $\rightarrow$ 6	: $n \rightarrow \infty$ not needed - theorem.

We are left to prove 6 $\rightarrow$ 1

Theorem [Fundamental theorem, Quantitative version]

$VCdim(H) = d$, then

1. H has uniform convergence property with
$$C_1 \cdot \frac{d \log \frac{1}{\epsilon \delta}}{\epsilon^2} \leq m_H^{uc}(\epsilon, \delta) \leq C_2 \cdot \frac{d \log \frac{1}{\epsilon \delta}}{\epsilon^2}$$
2. H is agnostic PAC learnable with
$$C_1 \cdot \frac{d \log \frac{1}{\epsilon \delta}}{\epsilon^2} \leq m_H(\epsilon, \delta) \leq C_2 \cdot \frac{d \log \frac{1}{\epsilon \delta}}{\epsilon^2}$$
3. H is PAC learnable with
$$C_1 \cdot \frac{d \log \frac{1}{\epsilon \delta}}{\epsilon^2} \leq m_H(\epsilon, \delta) \leq C_2 \cdot \frac{d \log \frac{1}{\epsilon} + \log \frac{1}{\delta}}{\epsilon^2}$$

proof later.

proof strategy for fundamental theorem of statistical learning.

① [Sauer's lemma]

If $\text{VCdim}(\mathcal{H}) = d$, then for any $C \subseteq \mathcal{X}$,

the effective size of $\mathcal{H}$ restricted to C , i.e., $|\mathcal{H}_C|$ is only $O(|C|^d) \ll 2^{|C|}$
polynomial exponential
worst case
big-O notation

$|\mathcal{H}_C| = O(f(|C|, d))$ means that $\exists c > 0$ s.t. $\leq c f(|C|, d)$ for large enough d .

Definition (Restriction of $\mathcal{H}$ to C)

$\mathcal{H}$ is class of functions $\mathcal{X} \rightarrow \{0, 1\}$, and $C = (C_1, \dots, C_m) \subseteq \mathcal{X}$,

the Restriction of $\mathcal{H}$ to C

$$\mathcal{H}_C \triangleq \{ \tilde{h} : C \rightarrow \{0, 1\} \mid h \in \mathcal{H} \}$$

$$C_i \mapsto h(C_i)$$

equivalently we can represent $\tilde{h}$ by a vector $(h(C_1), \dots, h(C_m))$.

② If $|\mathcal{H}_C|$ grows polynomially in $|C|$, then we have Uniform Convergence Property.

Definition (Uniform Convergence)

A hypothesis class $\mathcal{H}$ has the uniform convergence property,

if $\exists m_{\mathcal{H}}^{UC}$ s.t. for every ϵ, δ and every $\mathcal{D}$, if $m \geq m_{\mathcal{H}}^{UC}(\epsilon, \delta)$

the S is ϵ -representative with probability at least $1 - \delta$.

$$\uparrow |L_D(h) - L_S(h)| < \epsilon, \forall h \in \mathcal{H}.$$

Definition [Growth Function]

The Growth Function $T_{\mathcal{H}} : \mathbb{N} \rightarrow \mathbb{N}$, is defined as

$$T_{\mathcal{H}}(m) = \max_{C \subseteq \mathcal{X} : |C|=m} |\mathcal{H}_C|$$

which is the number of different functions from C of size m to $\{0, 1\}$ that can be obtained by restricting $\mathcal{H}$ to C .

* If $m \leq \text{VCdim}(\mathcal{H}) = d$, then by definition of VCdim, $T_{\mathcal{H}}(m) = 2^m$.
largest m s.t. $T_{\mathcal{H}}(m) = 2^m$.

* For $m > d$, m grows polynomially in m .

① Lemma [Sauer-Shelah-Perles]

If $\text{VCdim}(\mathcal{H}) = d < \infty$, then for all m , $T_{\mathcal{H}}(m) \leq \sum_{i=0}^d \binom{m}{i}$.

$\rightarrow d \geq m$, $T_{\mathcal{H}}(m) \leq \sum_{i=0}^m \binom{m}{i} = 2^m \leftarrow \text{exp in } m$
 m choose i , $\frac{m!}{i!(m-i)!}$

$\rightarrow d < m$, $T_{\mathcal{H}}(m) \leq \left(\frac{em}{d}\right)^d \leftarrow \text{polynomial in } m$
 $\uparrow$ Serfling's formula.

② Theorem [for uniform convergence] for every $D, \delta, h \in \mathbb{N}$

$$|L_D(h) - L_S(h)| \leq \frac{4 + \sqrt{\log(C_H(2m))}}{\delta \cdot \sqrt{m/2}} \quad \text{w.p. } 1-\delta.$$

Proof of fundamental theorem of statistical learning.

using ① Sauer's Lemma, wlog from ②, for $m > d$

$$|L_D(h) - L_S(h)| \leq \frac{4 + \sqrt{2d \log(2em/d)}}{\delta \sqrt{m/2}}$$

for simplicity let $\sqrt{2d \log(2em/d)} \geq 4$,

$$\leq \frac{2}{\delta} \sqrt{\frac{2d \log(2em/d)}{m}} \stackrel{\text{want}}{\leq} \varepsilon$$

$$\rightarrow \frac{2d \log(2em/d)}{(\delta \varepsilon)^2} \leq m$$

$$\rightarrow m \geq 4 \cdot \frac{2d}{(\delta \varepsilon)^2} \log(2em/d) \stackrel{\text{is sufficient}}{=} \frac{4d \log(2em/d)}{(\delta \varepsilon)^2}$$

this gives a loose bound on $m_{\mathcal{H}}(\varepsilon, \delta)$, but sufficient to prove PAC learnability.

Agnostic

Proof of ② > claim: $\mathbb{E}_S \left[\sup_{h \in \mathcal{H}} |L_D(h) - L_S(h)| \right] \leq \frac{4 + \sqrt{\log(C_H(2m))}}{\sqrt{m/2}}$

Markov's inequality: $\mathbb{P} \left(\sup_{h \in \mathcal{H}} |L_D(h) - L_S(h)| > \varepsilon \right) \leq \frac{\mathbb{E} \left[\sup_{h \in \mathcal{H}} |L_D(h) - L_S(h)| \right]}{\varepsilon}$
for $\varepsilon \geq \frac{4 + \sqrt{\log(C_H(2m))}}{\sqrt{m/2} \cdot \delta} \leq \delta$

$$\mathbb{E}_S \left[\sup_{h \in \mathcal{H}} |L_D(h) - L_S(h)| \right] = \mathbb{E}_{S=\{(x_i, y_i)\}_{i=1}^m} \left[\sup_{h \in \mathcal{H}} \left| \mathbb{E}_{S'=\{(x_i, y_i')\}_{i=1}^m} [L_S'(h)] - L_S(h) \right| \right]$$

Jensen's inequality
Convex f,

$$\leq \mathbb{E}_S \left[\sup_{h \in \mathcal{H}} \mathbb{E}_{S'} |L_S'(h) - L_S(h)| \right]$$

$$\mathbb{E}[f(\bar{x})] \geq f(\mathbb{E}[\bar{x}])$$

$$\leq \mathbb{E}_S \left[\mathbb{E}_{S'} \left[\sup_{h \in \mathcal{H}} |L_S'(h) - L_S(h)| \right] \right]$$

$$\checkmark \quad \text{max of linear f's}$$

$$= \mathbb{E}_{S, S'} \left[\sup_{h \in \mathcal{H}} \left| \frac{1}{m} \sum_{i=1}^m (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \right| \right] \quad (*)$$

$$L_h(\bar{x}_i') - L_h(\bar{x}_i) \stackrel{\text{def}}{=} b_i \cdot (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \quad , b_i \in \{-1, 1\}$$

$$\bar{x}_i, \bar{x}_i' \text{ are iid} \stackrel{\text{def}}{=} b_i \cdot (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \quad , b_i \sim \begin{matrix} \text{w.p.} \\ 1/2 \end{matrix}$$

$$\Rightarrow (*) = \mathbb{E}_{b_i \in \{-1, 1\}} \left[\mathbb{E}_{S, S'} \left[\sup_{h \in \mathcal{H}} \left| \frac{1}{m} \sum_{i=1}^m b_i (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \right| \right] \right]$$

linearity of expectation

$$\Rightarrow \mathbb{E}_{S, S'} \left[\mathbb{E}_{b_i \in \{-1, 1\}} \left[\sup_{h \in \mathcal{H}} \left| \frac{1}{m} \sum_{i=1}^m b_i (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \right| \right] \right]$$

S, S' is fixed, only $\mathcal{H}_C$ matters $C = S \cup S', |C| \leq 2m$

$$= \mathbb{E}_C \left[\max_{h \in \mathcal{H}_C} \left| \frac{1}{m} \sum_{i=1}^m b_i (L_h(\bar{x}_i') - L_h(\bar{x}_i)) \right| \mid S, S' \right]$$

$$\theta_i, \mathbb{E} \theta_i = 0, -1 \leq \theta_i \leq 1$$

$$\text{Hoeffding's inequality} \Rightarrow \mathbb{P} \left(\left| \frac{1}{m} \sum_{i=1}^m \theta_i \right| > p \right) \leq 2 \cdot e^{-\frac{2mp^2}{4}}$$

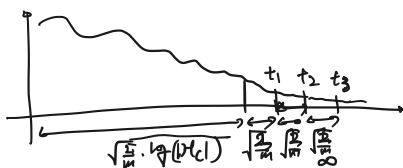
$$\text{union bound} \Rightarrow \mathbb{P} \left(\max_{h \in \mathcal{H}_C} \left| \frac{1}{m} \sum_{i=1}^m \theta_i \right| > p \right) \leq 2|\mathcal{H}_C| \cdot \mathbb{P} \left(\left| \frac{1}{m} \sum_{i=1}^m \theta_i \right| > p \right) \leq 2|\mathcal{H}_C| \cdot e^{-\frac{2mp^2}{4}}$$

$$(*) \Rightarrow \mathbb{E} \left[\max_{h \in \mathcal{H}_C} \left| \frac{1}{m} \sum_{i=1}^m L_h(\bar{x}_i') - L_h(\bar{x}_i) \right| \right] \leq \frac{4 + \sqrt{\log(|\mathcal{H}_C|)}}{\sqrt{m/2}} \leq \frac{4 + \sqrt{\log(C_H(2m))}}{\sqrt{m/2}}$$

$$(*) \Rightarrow \mathbb{P}(\theta > p) \leq 2|\mathcal{H}_C| \cdot e^{-\frac{2p^2}{4}} \Rightarrow \mathbb{E}[\theta] \leq \frac{4 + \sqrt{\log(|\mathcal{H}_C|)}}{\sqrt{m/2}}$$

Lemma A.4 book

proof
PDF



$$\mathbb{E}[\theta] = \int_0^1 \theta \cdot f(\theta) d\theta \leq \sum_{i=1}^m t_i \mathbb{P}(\theta > t_{i-1}) + \frac{\sqrt{2 \cdot \log(|\mathcal{H}_C|)}}{\sqrt{m}} - (i-1 + \sqrt{\log(|\mathcal{H}_C|)})^2$$

$$\leq 2\sqrt{\frac{2}{m}} \cdot |\mathcal{H}_C| \cdot \sum_{i=1}^m \sqrt{\eta(\mathcal{H}_C) + i} e + \dots$$

$$\leq 2\sqrt{\frac{2}{m}} + \sqrt{\frac{2}{m}} \eta(2|\mathcal{H}_C|)$$

proof of Sauer's Lemma

we will show that $\forall \mathcal{H}$

def of $VC(\mathcal{H}) = d$

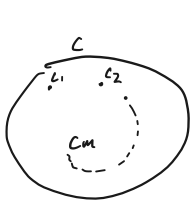
$$|\mathcal{H}_C| \stackrel{(*)}{\leq} |\{B \subseteq C : \mathcal{H} \text{ shatters } B\}| \leq \sum_{i=0}^d \binom{m}{i}$$

because largest set that can be shattered is d size.

proof of $(*)$ by induction

- $m=1$: $d=0 \rightarrow |\mathcal{H}_C| \leq 1$. trivial since $d < 1 \rightarrow$ does not shatter when $m=1$.
 $d=1 \rightarrow |\mathcal{H}_C| \leq 2 = 2^m$

- suppose $(*)$ holds for sets C' of size m , n.t.s. $(*)$ holds for set size m .



$$\mathcal{H}_C = \begin{matrix} c_1, c_2, \dots, c_m \\ 0 \\ 1 \\ 1 \\ 0, 1 \end{matrix}$$

$\forall B \subseteq C$ that $\mathcal{H}$ shatters B .

$$|\mathcal{H}_C| = |\mathcal{Y}_0| + |\mathcal{Y}_1|$$

$\mathcal{Y}_0$ is $\mathcal{H}$ on C' , $\mathcal{Y}_1$ is $\mathcal{H}$ on $C \setminus C'$

$c_i \in B, \mathcal{H} \text{ shatters } B$