Ok, we've analysed that for a given H

Uniform convergence

(Agnostic) PAC learnability

FIRM being a la Agnostic) PAC learner

VCdim (24) is bounded

We equivalent. ERM often isn't efficient, certainly in the agnostic setting. But it is efficient for realizeable linear separaters: Him-s Pick well s.t. squ(w, xi> +b) = yi \ (xi)yi)=S Vidim ( H 1 - d + 1 Shaffureable set:  $x_0 = 0$  for  $y_0$  Pick  $b = y_0 = 0$   $\langle w, x_0 \rangle + b = b = y_0 \rangle$   $x_1 = e_1$   $\langle w_1, x_1 \rangle + b$   $x_2 = e_3$   $y_3 = 0$   $x_4 = e_4$   $y_5 = y_1 \wedge y_1 \wedge y_2 \wedge y_3 \wedge y_4 \wedge y_5 \wedge$ LB. = w; +b= yi Take and set [C]? In Rd. [Let's ignore WTS it Connot be shattered by Hlind b and just [find a set of labels we con't induce!] argue about Since [C]? It!, C is not linearly I the bias-less =) I ato St. Ohes? P= {i: a; > 0 } N= 3; aj < 03 > Neither Con be empty. Why? Saixi = Slajlxj

[If on 1370, the other can't be...

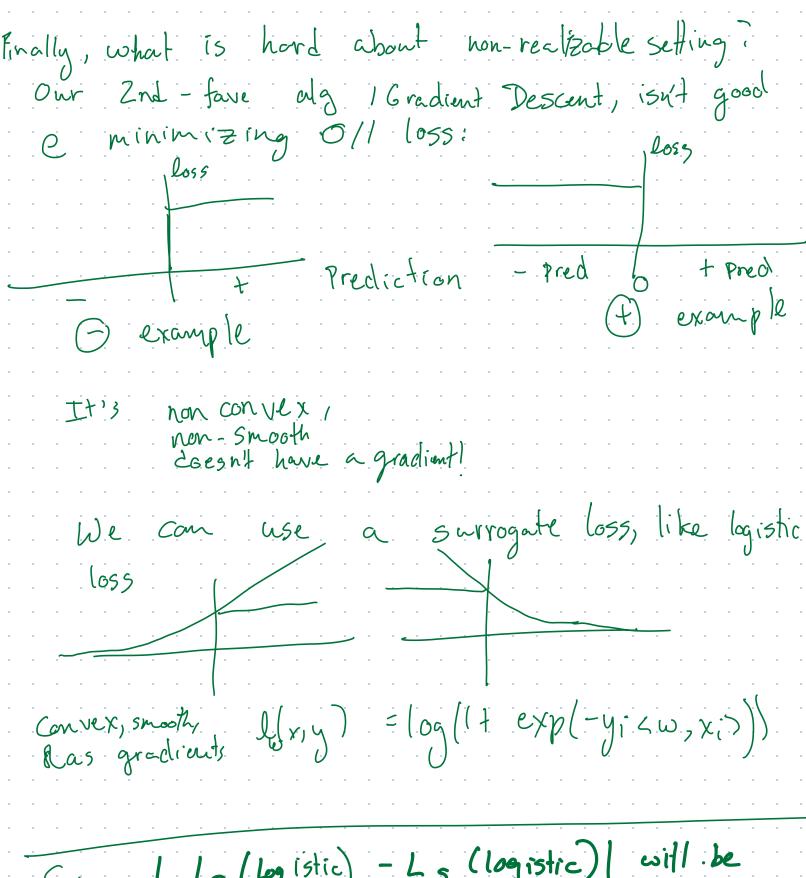
2+0 so at least

one is nonempter

Suppose  $\omega$  induces labels  $sgn(a_i)$  for  $\chi_i$ .

This means  $a_i < \omega, \chi_i > 0$   $\forall i$ :  $0 \le \underbrace{\leq}_{i \in P} a_i < \omega, \chi_i > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > = (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i, \omega) = \underbrace{\leq}_{i \in P} a_i \times_i, \omega > (\underbrace{\leq}_{i \in P} a_i \times_i,$ A contradiction? So C can't be shattered Since 5 gn (x, w) by Alin-d This means ERM on m2 C. d+logts samples

PAC-learns linear separators in Rd. [ When not realizable, it's amputationally hard is ] [Simple Alg: use Linear Programming! Other things a 130 work. If our data is not only realizable (separchie) but also Mas a margin min width Then we can learn tewer samples is the margin samples = argmin //vl/ st Y, (w, x; ) = ] Vie[m]  $||x|| \leq |x| + ||x|| \leq |x|$ 



So, [ Lp (log istic) - Ls (logistic) | will be small [ Uniform Convergence!) but no gnar that 0-1-Ls (logistic) will be small even if  $\exists \ \omega \ \omega$ . small 0-1 Ls what's another set of methods for arguing about good  $L_p(h)$ ?  $L_p(h) = |L_p(h) - L_s(h)| + L_s(h)$ ?