

CSE 490 G
Introduction to Data Compression
Winter 2006

EBCOT
JPEG 2000

History

- Embedded Block Coding with Optimized Truncation (EBCOT)
 - Taubman – journal paper 2000
 - Algorithm goes back to 1998 or maybe earlier
 - Basis of JPEG 2000
- Embedded
 - Prefixes of the encoded bit stream are legal encodings at lower fidelity, like SPIHT and GTW
- Block coding
 - Entropy coding of blocks of bit planes, not block transform coding like JPEG.

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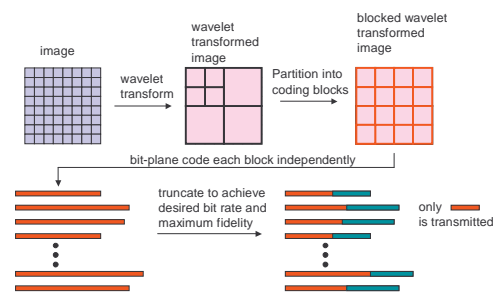
Features at a High Level

- SNR scalability (Signal to Noise Ratio)
 - Embedded code - The compressed bit stream can be truncated to yield a smaller compressed image at lower fidelity
 - Layered code – The bit stream can be partitioned into a base layer and enhancement layers. Each enhancement layer improves the fidelity of the image
- Resolution scalability
 - The lowest subband can be transmitted first yielding a smaller image at high fidelity.
 - Successive subbands can be transmitted to yield larger and larger images

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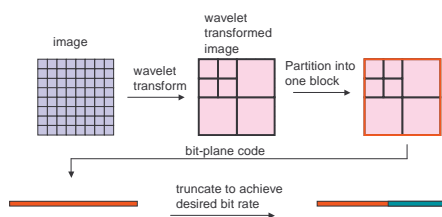
Block Diagram of Encoder



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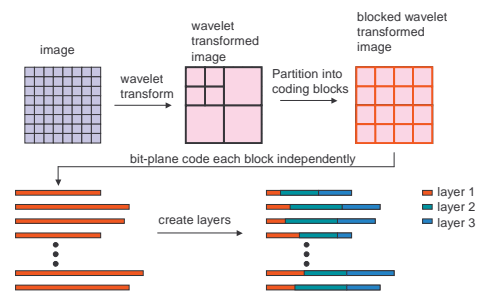
Extreme Case is Normal



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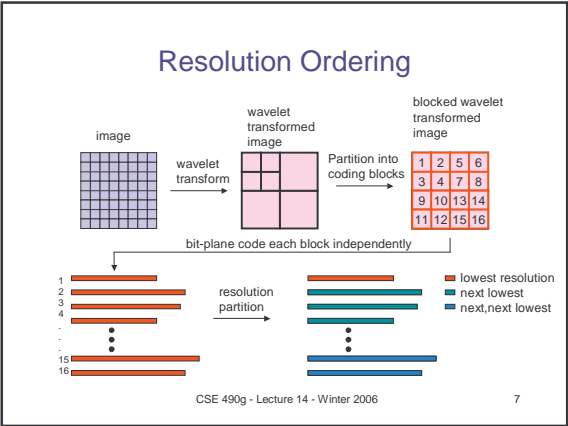
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Layering

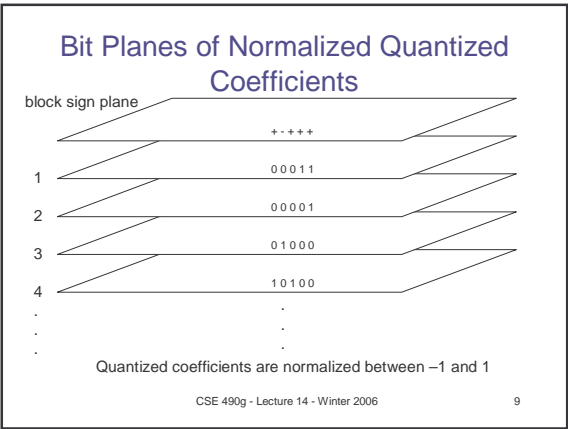


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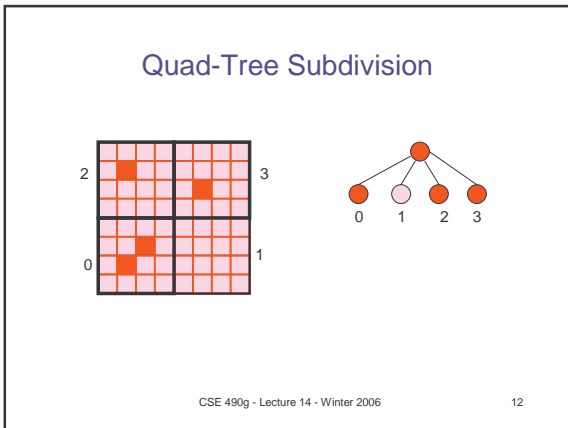
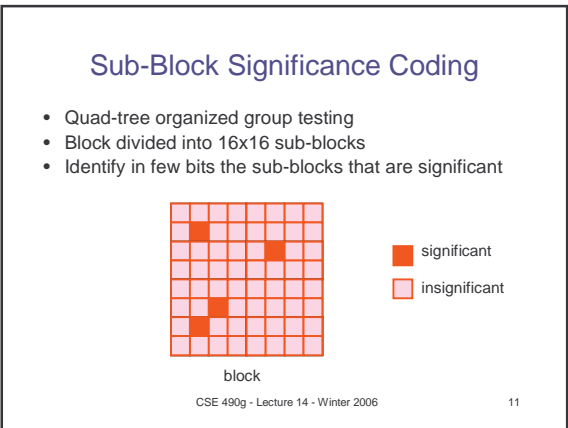
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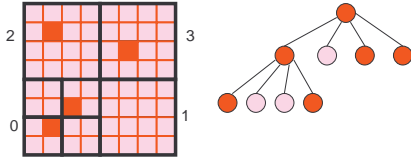
- ### Block Coding
- Assume we are in block k , and $c(i,j)$ is a coefficient in block k .
 - Divide $c(i,j)$ into its sign $s(i,j)$ and $m(i,j)$ its magnitude.
 - Quantize to $v(i,j) = \lfloor m(i,j)/q_k + .5 \rfloor$ where q_k is the quantization step for block k .
 - Example: $c(i,j) = -10$, $q_k = 3$.
 - $s(i,j) = 0$
 - $v(i,j) = \text{floor}(-10/3 + .5) = -2$
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- ### Bit-Plane Coding of Blocks
- Sub-block significance coding (like group testing)
 - Some sub-blocks are declared insignificant
 - Significant sub-blocks must be coded
 - Contexts are defined based on the previous bit-plane significance.
 - Zero coding (ZC) – 9 contexts
 - Run length coding (RLC) – 1 context
 - Sign coding (SC) – 5 contexts
 - Magnitude refinement coding (MR) – 3 contexts
 - Block coded in raster order using arithmetic coding
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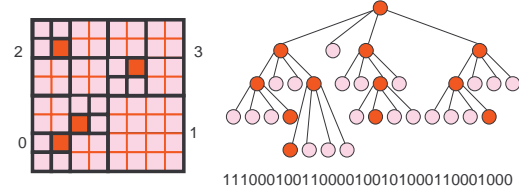
Quad-Tree Subdivision



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Quad-Tree Subdivision Coding



111000100110000100101000110001000

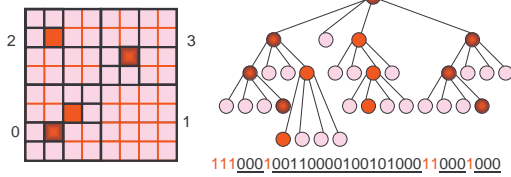
Depth-first code = 1 for significant
0 for insignificant

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Quad-Tree Subdivision Coding

known significant in last bit plane



111000100110000100101000110001000

- Skip symbols that are already known:
- nodes significant in previous bit plane
 - last child of significant parent of other children are insignificant

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ZC – Zero Coding



- LH is transposed so that it can be treated the same as HL. $(LH)^T$ has similar characteristics to HL.
- Each coefficient has its neighbors in the **same** subband



- vertical neighbors
- horizontal neighbors
- diagonal neighbors

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ZC Contexts

- v = number of vertical neighbors significant in the previous bit-plane
- h = number of horizontal neighbors significant in the previous bit-plane
- d = number of diagonal neighbors significant in the previous bit-plane

$0 \leq h, v \leq 2$
 $0 \leq d \leq 4$

higher labels mean more likely to be significant

label	HL	(LH) ^T	LL	HH
	h	v	d	h+v
0	0	0	0	0
1	0	0	1	1
2	0	0	>1	>1
3	0	1	*	1
4	0	2	*	1
5	1	0	0	>1
6	1	0	>0	2
7	1	>0	*	>0
8	2	*	*	>2

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Examples

significant in previous bit-plane



$h = 0$
 $v = 2$
 $d = 0$

Context 4



$h = 0$
 $v = 2$
 $d = 0$

Context 2



$h = 2$
 $v = 0$
 $d = 0$

Context 8



$h = 2$
 $v = 0$
 $d = 0$

Context 2



$h = 0$
 $v = 0$
 $d = 2$

Context 2



$h = 0$
 $v = 0$
 $d = 2$

Context 6

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RLC – Run Length Coding

- Looks for runs of 4 that are likely to be insignificant

?	?	?	?
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- If all insignificant then code as a single symbol
- Main purpose – to lighten the load on the arithmetic coder.

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SC – Sign Coding

$$hs = \begin{cases} 0 & \text{if horizontal neighbors are both insignificant or of opposite sign} \\ 1 & \text{if at least one horizontal neighbor is positive} \\ -1 & \text{if at least one horizontal neighbor is negative} \end{cases}$$

$$vs = \begin{cases} 0 & \text{if vertical neighbors are both insignificant or of opposite sign} \\ 1 & \text{if at least one vertical neighbor is positive} \\ -1 & \text{if at least one vertical neighbor is negative} \end{cases}$$

	1	1	1	0	0	0	-1	-1
hs	1	0	-1	1	0	-1	1	0
vs	1	1	1	-1	1	1	-1	-1
sign prediction	4	3	2	1	0	1	2	3
label								

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MR – Magnitude Refinement

- This is the refinement pass.
- Define $t = 0$ if first refinement bit, $t = 1$ otherwise.

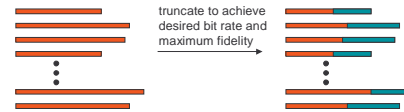
label	t	h + v
0	0	0
1	0	> 0
2	1	*

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Bit Allocation

- How do we truncate the encoded blocks to achieve a desired bit rate and get maximum fidelity



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Basic Set Up

- Encoded block k can be truncated to n_k bits.
- Total Bit Rate

$$\sum_k n_k$$

- Distortion attributable to block k is

$$D_k^{n_k} = w_k^2 \sum_{(i,j) \in B_k} (c^{n_k}(i,j) - c(i,j))^2$$

where w_k is the “weight” of the basis vectors for block k and $c^{n_k}(i,j)$ is the recovered coefficients from n_k bits of block k .

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Bit Allocation as an Optimization Problem

- Input: Given m embedded codes and a bit rate target R
- Output: Find truncation values n_k , $1 \leq k \leq m$, such that

$$D = \sum_k D_k^{n_k} \text{ is minimized and}$$

$$\sum_k n_k \leq R$$

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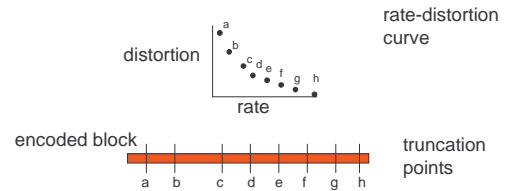
Facts about Bit Allocation

- It is an NP-hard problem generally
- There are fast approximate algorithms that work well in practice
 - Lagrange multiplier method
 - Multiple choice knapsack method

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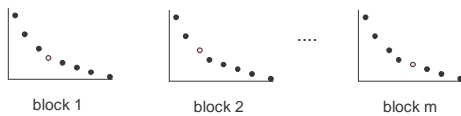
Rate-Distortion Curve



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Picture of Bit Allocation



Pick one point from each curve so that the sum of the x values is bounded by R and the sum of the y values is minimized.

Good approximate algorithms exist because the curves are almost convex.

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Notes on EBCOT

- EBCOT is quite complicated with many features.
- JPEG 2000 based on EBCOT but differs to improve compression and decompression time.
- EBCOT has
 - resolution scalability
 - SNR scalability
 - quantization
 - bit allocation
 - arithmetic coding with context and adaptivity
 - group testing (quad trees)
 - sign and refinement bit contexts
 - lots of engineering

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Notes on Wavelet Compression

- Wavelets appear to be excellent for image compression
 - No blocking artifacts
 - Wavelet coding techniques abound and are very effective
- Some of the wavelet coding techniques can apply to block transforms.
- Newest generation of image compressor use wavelets, JPEG 2000.

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