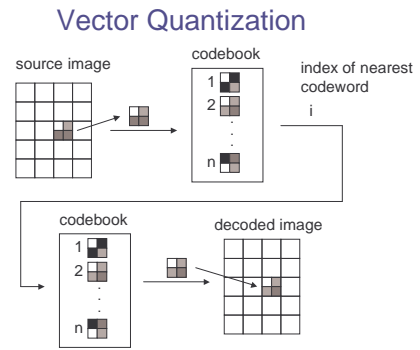


CSE 490 GZ
Introduction to Data Compression
Winter 2004

Vector Quantization



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Vectors

- An $a \times b$ block can be considered to be a vector of dimension ab .

$$\text{block } \begin{bmatrix} w & x \\ y & z \end{bmatrix} = (w, x, y, z) \text{ vector}$$

- Nearest means in terms of Euclidian distance or Euclidian squared distance. Both equivalent.

$$\text{Distance} = \sqrt{(w_1 - w_2)^2 + (x_1 - x_2)^2 + (y_1 - y_2)^2 + (z_1 - z_2)^2}$$

$$\text{Squared Distance} = (w_1 - w_2)^2 + (x_1 - x_2)^2 + (y_1 - y_2)^2 + (z_1 - z_2)^2$$

- Squared distance is easier to calculate.

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Vector Quantization Facts

- The image is partitioned into $a \times b$ blocks.
- The codebook has n representative $a \times b$ blocks called codewords, each with an index.
- Compression with fixed length codes is

$$\frac{\log_2 n}{ab} \text{ bpp}$$

- Example: $a = b = 4$ and $n = 1,024$
 - compression is $10/16 = .63$ bpp
 - compression ratio is $8 : .63 = 12.8 : 1$
- Better compression with entropy coding of indices

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Examples



4 x 4 blocks
.63 bpp

4 x 8 blocks
.31 bpp

8 x 8 blocks
.16 bpp

Codebook size = 1,024

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Scalar vs. Vector

- Pixels within a block are correlated.
 - This tends to minimize the number of codewords needed to represent the vectors well.
- More flexibility.
 - Different size blocks
 - Different size codebooks

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Encoding and Decoding

- Encoding:
 - Scan the $a \times b$ blocks of the image. For each block find the nearest codeword in the codebook and output its index.
 - Nearest neighbor search.
- Decoding:
 - For each index output the codeword with that index into the destination image.
 - Table lookup.

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The Codebook

- Both encoder and decoder must have the same codebook.
- The codebook must be useful for many images and be stored someplace.
- The codebook must be designed properly to be effective.
- Design requires a representative training set.
- These are major drawbacks to VQ.

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Codebook Design Problem

- Input: A training set X of vectors of dimension d and a number n . ($d = a \times b$ and n is number of codewords)
- Output: n codewords $c(0), c(1), \dots, c(n-1)$ that minimizes the distortion.

$$D = \sum_{x \in X} \|x - c(\text{index}(x))\|^2 \quad \text{sum of squared distances}$$

where $\text{index}(x)$ is the index of the nearest codeword to x .

$$\|(x_0, x_1, \dots, x_{d-1})\|^2 = x_0^2 + x_1^2 + \dots + x_{d-1}^2 \quad \text{squared norm}$$

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GLA

- The Generalized Lloyd Algorithm (GLA) extends the Lloyd algorithm for scalars.
 - Also called LBG after inventors Linde, Buzo, Gray (1980)
- It can be very slow for large training sets.

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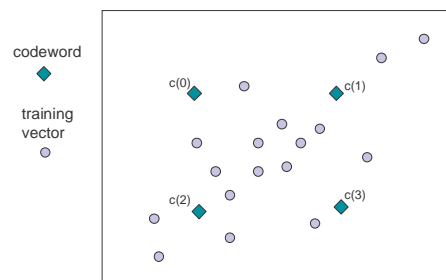
GLA

Choose a training set X and small error tolerance $\epsilon > 0$.
 Choose start codewords $c(0), c(1), \dots, c(n-1)$
 Compute $X(j) := \{x : x \text{ is a vector in } X \text{ closest to } c(j)\}$
 Compute distortion D for $c(0), c(1), \dots, c(n-1)$
 Repeat
 Compute new codewords
 $c'(j) := \text{round}\left(\frac{1}{|X(j)|} \sum_{x \in X(j)} x\right)$ (centroid)
 Compute $X'(j) := \{x : x \text{ is a vector in } X \text{ closest to } c'(j)\}$
 Compute distortion D' for $c'(0), c'(1), \dots, c'(n-1)$
 if $|(D - D')/D| < \epsilon$ then quit
 else $c := c'$; $X := X'$; $D := D'$
 End(repeat)

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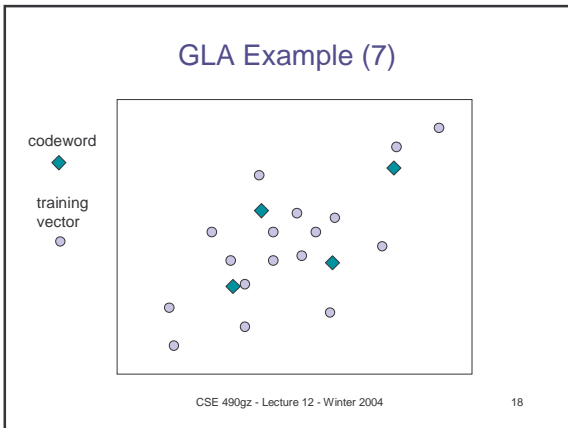
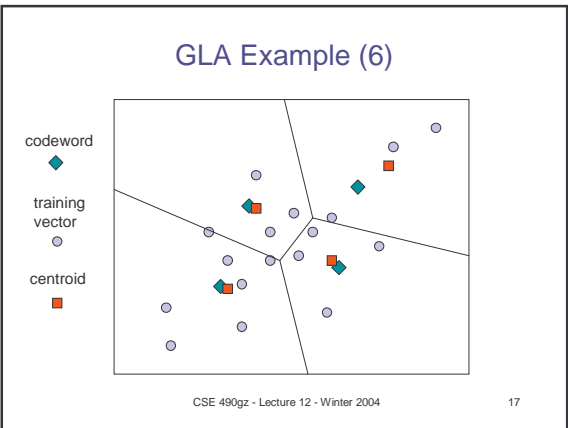
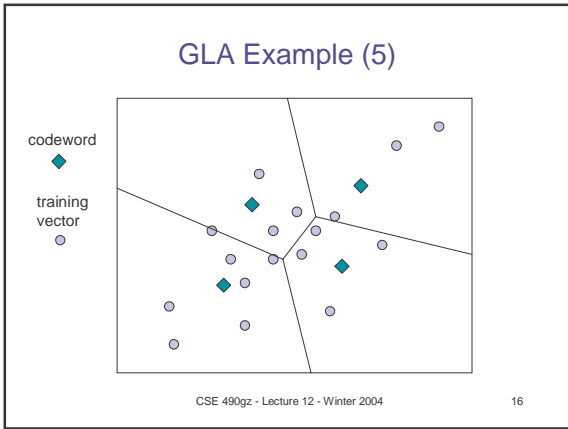
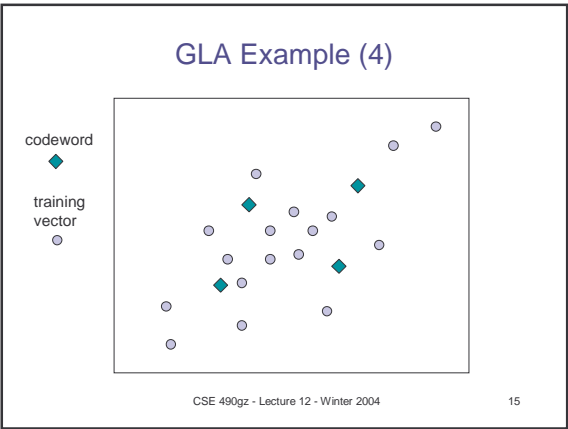
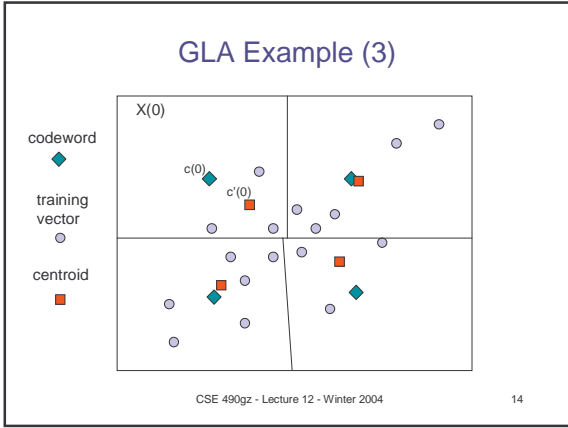
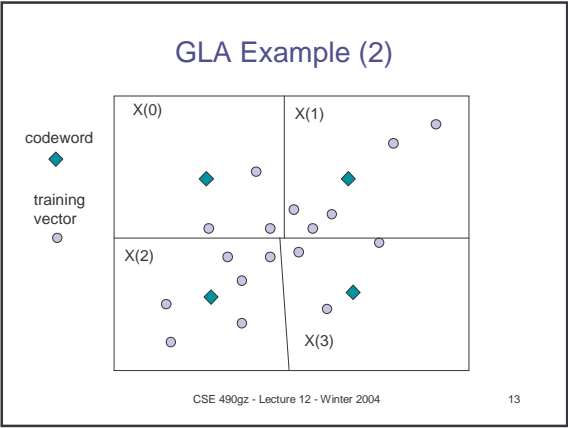
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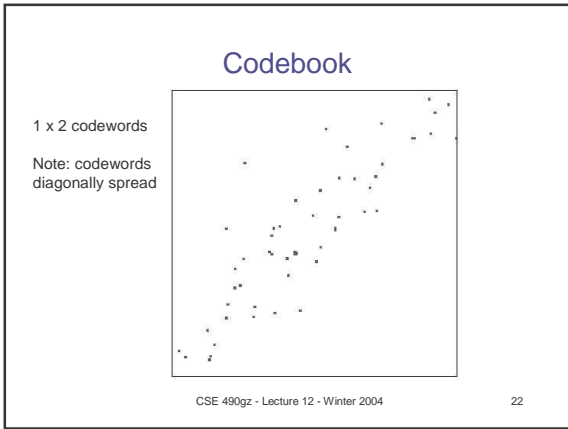
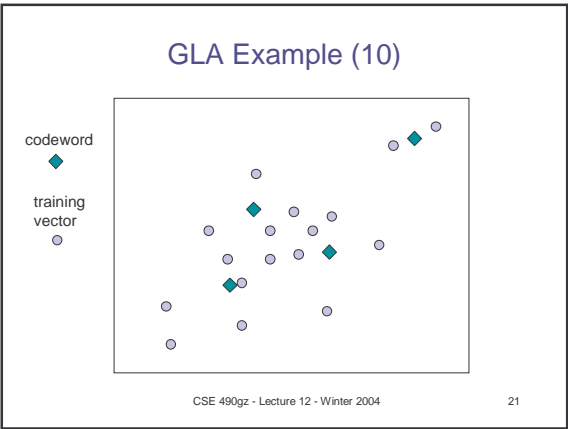
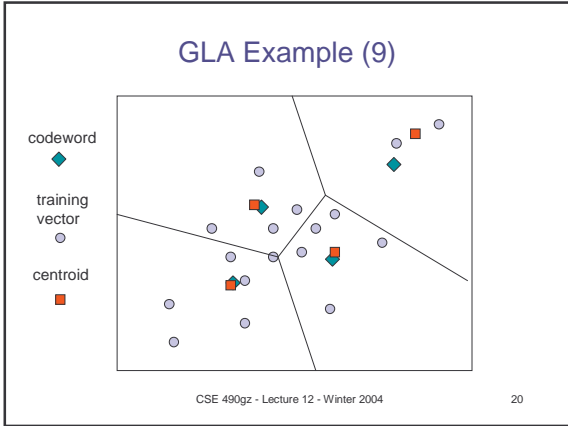
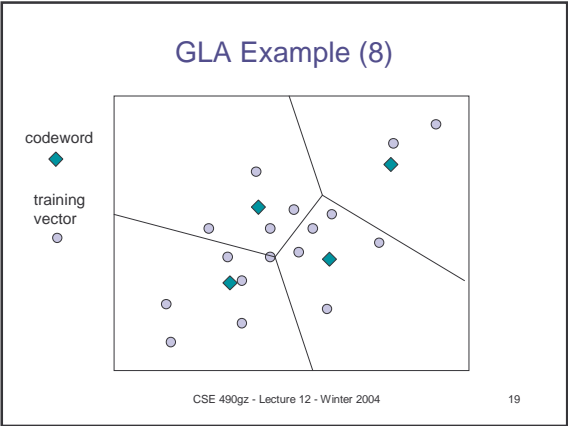
GLA Example (1)



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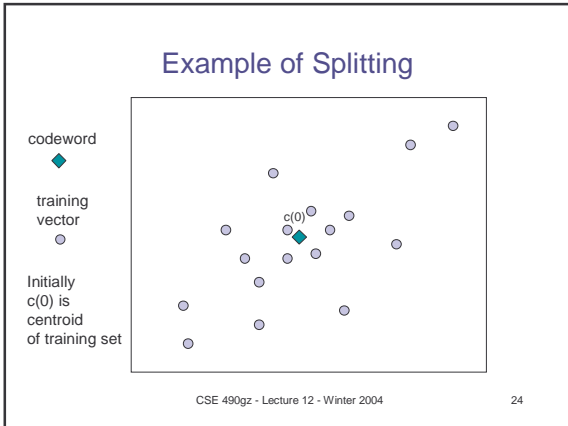


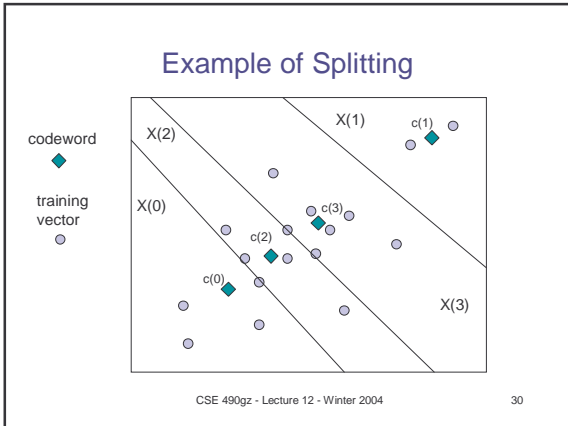
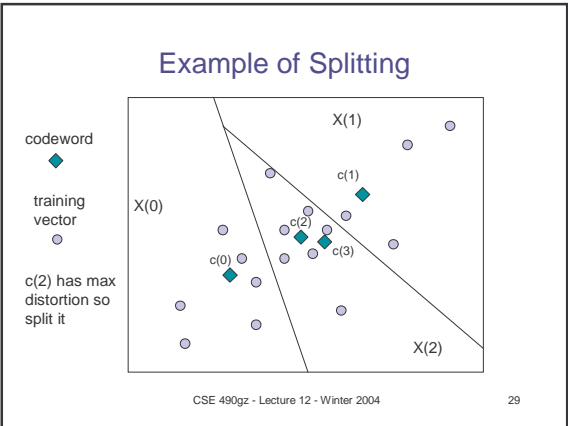
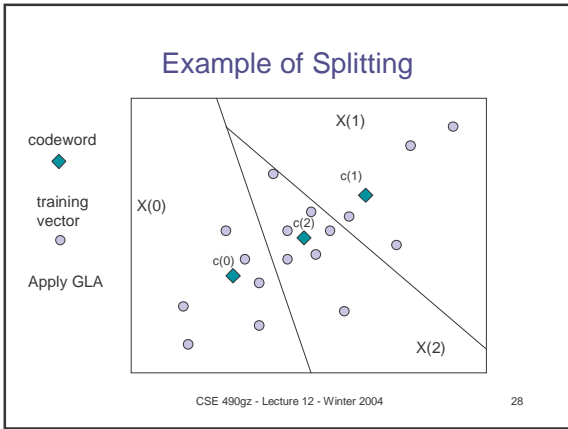
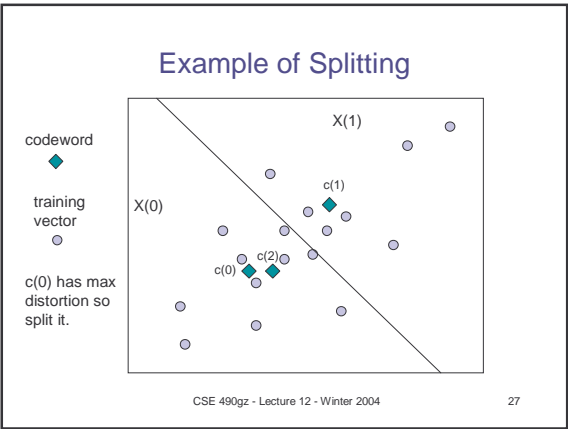
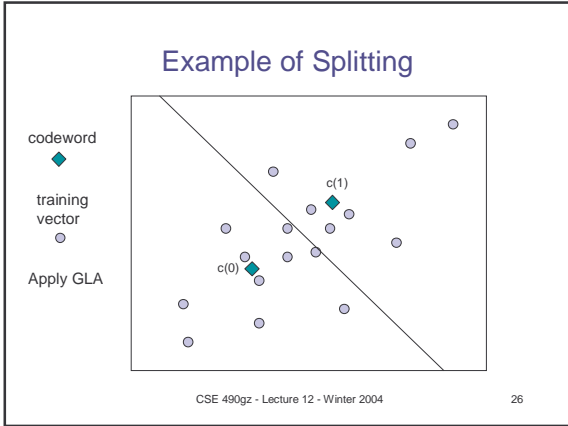
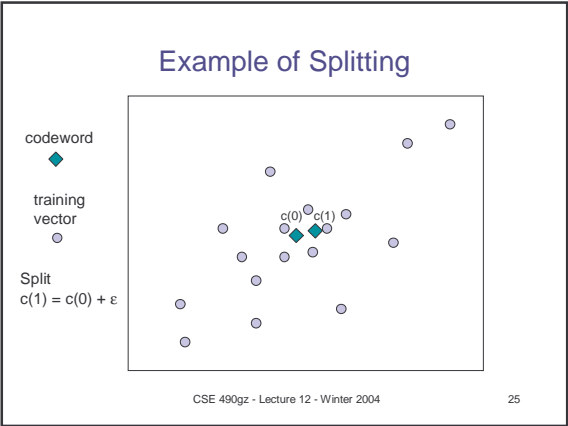


Codeword Splitting

- It is possible that a chosen codeword represents no training vectors, that is, $X(j)$ is empty.
 - **Splitting** is an alternative codebook design algorithm that avoids this problem.
- Basic Idea
 - Select codeword $c(j)$ with the greatest distortion.
$$D(j) = \sum_{x \in X(j)} \|x - c(j)\|^2$$
 - Split it into two codewords then do the GLA.

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GLA Advice

- Time per iteration is dominated by the partitioning step, which is m nearest neighbor searches where m is the training set size.
 - Average time per iteration $O(m \log n)$ assuming d is small.
- Training set size.
 - Training set should be at least 20 training vectors per code word to get reasonable performance.
 - To small a training set results in “over training”.
- Number of iterations can be large.

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Encoding

- Naive method.
 - For each input block, search the entire codebook to find the closest codeword.
 - Time $O(T n)$ where n is the size of the codebook and T is the number of blocks in the image.
 - Example: $n = 1024$, $T = 256 \times 256 = 65,536$ (2×2 blocks for a 512×512 image)
 $nT = 1024 \times 65536 = 2^{26} \approx 67$ million distance calculations.
- Faster methods are known for doing “Full Search VQ”. For example, k-d trees.
 - Time $O(T \log n)$

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