CSE 490 GZ
Introduction to Data Compression
Winter 2004

Predictive Coding

Burrows-Wheeler Transform

Predictive Coding

- The next symbol can be statistically predicted from the past.
  - Code with context
  - Code the difference
  - Move to front, then code
- Goal of prediction
  - The prediction should make the distribution of probabilities of the next symbol as skewed as possible
  - After prediction there is no way to predict more so we are in the first order entropy model

Bad and Good Prediction

- From information theory – The lower the information the fewer bits are needed to code the symbol.

$$\text{inf}(a) = \log_2\left(\frac{1}{P(a)}\right)$$

- Examples:
  - P(a) = 1024/1024, inf(a) = .000977
  - P(a) = 1/2, inf(a) = 1
  - P(a) = 1/1024, inf(a) = 10

Entropy

- Entropy is the expected number of bit to code a symbol in the model with a_i having probability $P(a_i)$.

$$H = \sum_{i=1}^{m} P(a_i)\log_2\left(\frac{1}{P(a_i)}\right)$$

- Good coders should be close to this bound.
  - Arithmetic
  - Huffman
  - Golomb
  - Tunstall

PPM

- Prediction with Partial Matching
  - Cleary and Witten (1984)
  - Tries to find a good context to code the next symbol

good → he

context	a	e	i	r	s	y
the	0	0	5	7	4	7
he	10	1	7	10	9	7
e	12	2	10	15	10	10
<nil>	50	70	30	35	40	13

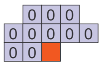
- Uses adaptive arithmetic coding for each context

JBIG

- Coder for binary images
  - documents
  - graphics
- Codes in scan line order using context from the same and previous scan lines.

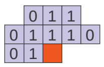
- Uses adaptive arithmetic coding with context

JBIG Example



next bit	0	1
frequency	100	10

$$H = \frac{10}{110}\log(\frac{110}{10}) + \frac{100}{110}\log(\frac{110}{100}) = .44$$



next bit	0	1
frequency	15	50

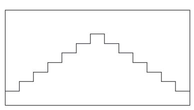
$$H = \frac{15}{65}\log(\frac{65}{15}) + \frac{50}{65}\log(\frac{65}{50}) = .78$$

Issues with Context

- Context dilution
  - If there are too many contexts then too few symbols are coded in each context, making them ineffective because of the zero-frequency problem.
- Context saturation
  - If there are too few contexts then the contexts might not be good as having more contexts.
- Wrong context
  - Again poor predictors.

Prediction by Differencing

- Used for Numerical Data
- Example: 2 3 4 5 6 7 8 7 6 5 4 3 2



- Transform to 2 1 1 1 1 1 1 –1 –1 –1 –1 –1 –1
  - much lower first-order entropy

General Differencing

- Let $x_1, x_2, ..., x_n$ be some numerical data that is correlated, that is x_i is near x_{i+1}
- Better compression can result from coding $x_1, x_2 - x_1, x_3 - x_2, ... , x_n - x_{n-1}$
- This idea is used in
  - signal coding
  - audio coding
  - video coding
- There are fancier prediction methods based on linear combinations of previous data, but these may require training.

Move to Front Coding

- Non-numerical data
- The data have a relatively small working set that changes over the sequence.
- Example: a b a b a a b c c b b c c c c b d b c c
- Move to Front algorithm
  - Symbols are kept in a list indexed 0 to m-1
  - To code a symbol output its index and move the symbol to the front of the list

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0

0	1	2	3
a	b	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1

0	1	2	3
a	b	c	d

↓

0	1	2	3
b	a	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1

0	1	2	3
b	a	c	d

↓

0	1	2	3
a	b	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1

0	1	2	3
a	b	c	d

↓

0	1	2	3
b	a	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1

0	1	2	3
b	a	c	d

↓

0	1	2	3
a	b	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1 0

0	1	2	3
a	b	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1 0 1

0	1	2	3
a	b	c	d

↓

0	1	2	3
b	a	c	d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1 0 1 2

0 1 2 3
b a c d

↓

0 1 2 3
c b a d

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1 0 1 2 0 1 0 1 0 00 1 3 1 2 0

0 1 2 3
c b d a

Example

- Example: a b a b a a b c c b b c c c c b d b c c
  0 1 1 1 1 0 1 2 0 1 0 1 0 00 1 3 1 2 0

Frequencies of {a, b, c, d}
a b c d
4 7 8 1

Frequencies of {0, 1, 2, 3}
0 1 2 3
8 9 2 1

Extreme Example

Input:
aaaaaaaaaabbbbbbbbbbccccccccccdddddddddd
Output
0000000000100000000020000000003000000000

Frequencies of a b c d
a b c d
10 10 10 10

Frequencies of 0 1 2 3
0 1 2 3
37 1 1 1

Burrows-Wheeler Transform

- Burrows-Wheeler, 1994
- BW Transform creates a representation of the data which has a small working set.
- The transformed data is compressed with move to front compression.
- The decoder is quite different from the encoder.
- The algorithm requires processing the entire string at once (it is not on-line).
- It is a remarkably good compression method.

Encoding Example

- abracadabra

1. Create all cyclic shifts of the string.

```
0   abracadabra
1   bracadabraa
2   racadabraab
3   acadabraabr
4   cadabraabra
5   adabraabrac
6   dabraabraca
7   abraabracad
8   braabracada
9   raabracadab
10  aabracadabr
```

Encoding Example

2. Sort the strings alphabetically in to array A

```
0   abracadabra         A 0   aabracadabr
1   bracadabraa           1   abraabracad
2   racadabraab           2   abracadabra
3   acadabraabr           3   acadabraabr
4   cadabraabra    -->    4   adabraabrac
5   adabraabrac           5   braabracada
6   dabraabraca           6   bracadabraa
7   abraabracad           7   cadabraabra
8   braabracada           8   dabraabraca
9   raabracadab           9   raabracadab
10  aabracadabr           10  racadabraab
```


Encoding Example

3. L = the last column

```
A 0  aabracadabr
  1  abraabracad      L = rdarcaaaabb
  2  abracadabra
  3  acadabraabr
  4  adabraabrac
  5  braabracada
  6  bracadabraa
  7  cadabraabra
  8  dabraabraca
  9  raabracadab
  10 racadabraab
```


Encoding Example

4. Transmit X the index of the input in A and L (using move to front coding).

```
A 0  aabracadabr
  1  abraabracad      L = rdarcaaaabb
  2  abracadabra      X = 2
  3  acadabraabr
  4  adabraabrac
  5  braabracada
  6  bracadabraa
  7  cadabraabra
  8  dabraabraca
  9  raabracadab
  10 racadabraab
```


Why BW Works

- Ignore decoding for the moment.
- The prefix of each shifted string is a context for the last symbol.
  - The last symbol appears just before the prefix in the original.
- By sorting similar contexts are adjacent.
  - This means that the predicted last symbols are similar.

Decoding Example

- We first decode assuming some information. We then show how compute the information.
- Let A^s be A shifted by 1

```
A 0  aabracadabr     As 0  raabracadab
  1  abraabracad        1  dabraabraca
  2  abracadabra        2  aabracadabr
  3  acadabraabr        3  racadabraab
  4  adabraabrac        4  cadabraabra
  5  braabracada        5  abraabracad
  6  bracadabraa        6  abracadabra
  7  cadabraabra        7  acadabraabr
  8  dabraabraca        8  adabraabrac
  9  raabracadab        9  braabracada
  10 racadabraab        10 bracadabraa
```


Decoding Example

- Assume we know the mapping T[i] is the index in A^s of the string i in A.
- T = [2 5 6 7 8 9 10 4 1 0 3]

```
A 0  aabracadabr     As 0  raabracadab
  1  abraabracad        1  dabraabraca
  2  abracadabra        2  aabracadabr
  3  acadabraabr        3  racadabraab
  4  adabraabrac        4  cadabraabra
  5  braabracada        5  abraabracad
  6  bracadabraa        6  abracadabra
  7  cadabraabra        7  acadabraabr
  8  dabraabraca        8  adabraabrac
  9  raabracadab        9  braabracada
  10 racadabraab        10 bracadabraa
```

Decoding Example

- Let F be the first column of A, it is just L, sorted.

```
F = 0 1 2 3 4 5 6 7 8 9 10
    a a a a a b b c d r r

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1 0 3
```

- Follow the pointers in T in F to recover the input starting with X.

Decoding Example

```
F = 0 1 2 3 4 5 6 7 8 9 10
    a a a a a b b c d r r

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1 0 3
```

a

Decoding Example

```
F = 0 1 2 3 4 5 6 7 8 9 10
    a a a a a b b c d r r

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1 0 3
```

ab

Decoding Example

```
F = 0 1 2 3 4 5 6 7 8 9 10
    a a a a a b b c d r r

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1 0 3
```

abr

Decoding Example

- Why does this work?
- The first symbol of A[T[i]] is the second symbol of A[i] because $A^s[T[i]] = A[i]$.

	A	T		A^s
0	aabracadabr	2	0	raabracadab
1	abraabracad	5	1	dabraabraca
2	abracadabra	6	2	aabracadabr
3	acadabraabr	7	3	racadabraab
4	adabraabrac	8	4	cadabraabra
5	braabracada	9	5	abraabracad
6	bracadabraa	10	6	abracadabra
7	cadabraabra	4	7	acadabraabr
8	dabraabraca	1	8	adabraabrac
9	raabracadab	0	9	braabracada
10	racadabraab	3	10	bracadabraa

Decoding Example

- How do we compute F and T from L and X?
  F is just L sorted

```
    0 1 2 3 4 5 6 7 8 9 10
F = a a a a a b b c d r r
L = r d a r c a a a a b b
```

Note that L is the first column of A^s and A^s is in the same order as A.

If i is the k-th x in F then T[i] is the k-th x in L.

Decoding Example

```
     0 1 2 3 4 5 6 7 8 9 10
F =  a a a a a b b c d r r

L =  r d a r c a a a a b b

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8
```


Decoding Example

```
     0 1 2 3 4 5 6 7 8 9 10
F =  a a a a a b b c d r r

L =  r d a r c a a a a b b

T=  0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10
```


Decoding Example

```
     0 1 2 3 4 5 6 7 8 9 10
F =  a a a a a b b c d r r

L =  r d a r c a a a a b b

T=  0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4
```


Decoding Example

```
     0 1 2 3 4 5 6 7 8 9 10
F =  a a a a a b b c d r r

L =  r d a r c a a a a b b

T=  0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1
```


Decoding Example

```
     0 1 2 3 4 5 6 7 8 9 10
F =  a a a a a b b c d r r

L =  r d a r c a a a a b b

T = 0 1 2 3 4 5 6  7 8 9 10
    2 5 6 7 8 9 10 4 1 0 3
```


Notes on BW

- Alphabetic sorting does not need the entire cyclic shifted inputs.
  - Sort the indices of the string
  - Most significant symbols first radix sort works
- There are high quality practical implementations
  - Bzip
  - Bzip2 (seems to be w/o patents)

Encoding Exercise

Encode the string abababababababab = $(ab)^8$
1. Find L and X
2. Do move-to-front coding of L.
3. Estimate the length of the code using first order entropy.

Decoding Exercise

Decode L = baaaaaba, X = 6
1. First Compute F and T
2. Use those to decode.