# CSE 490 GZ
# Introduction to Data Compression
Winter 2002

EBCOT

JPEG 2000

## History

- Embedded Block Coding with Optimized Truncation (EBCOT)
  - Taubman – journal paper 2000
  - Algorithm goes back to 1998 or maybe earlier
  - Basis of JPEG 2000
- Embedded
  - Prefixes of the encoded bit stream are legal encodings at lower fidelity, like SPIHT and GTW
- Block coding
  - Entropy coding of blocks of bit planes, not block transform coding like JPEG.

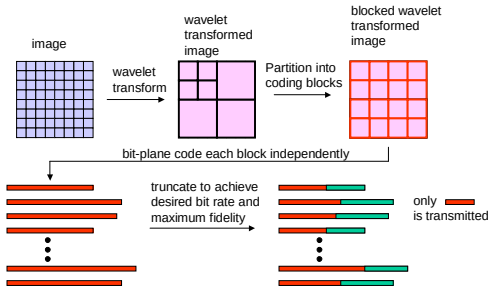


## Features at a High Level

- SNR scalability (Signal to Noise Ratio)
  - Embedded code - The compressed bit stream can be truncated to yield a smaller compressed image at lower fidelity
  - Layered code – The bit stream can be partitioned into a base layer and enhancement layers. Each enhancement layer improves the fidelity of the image
- Resolution scalability
  - The lowest subband can be transmitted first yielding a smaller image at high fidelity.
  - Successive subbands can be transmitted to yield larger and larger images

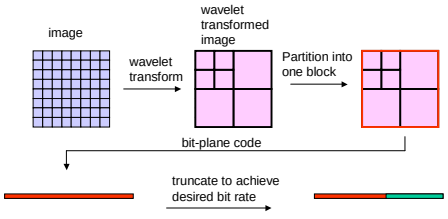


## Block Diagram of Encoder



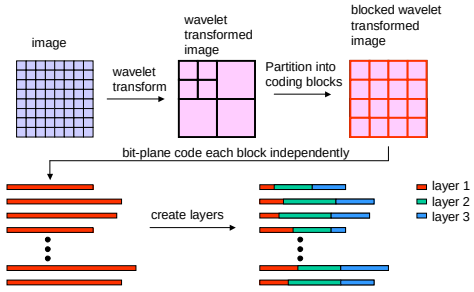



## Extreme Case is Normal

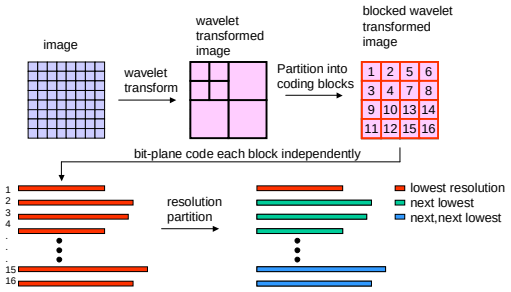





## Layering



## Resolution Ordering



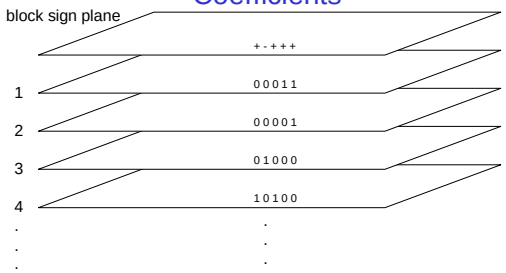



## Block Coding

- Assume we are in block k, and c(i,j) is a coefficient in block k.
- Divide c(i,j) into its sign s(i,j) and m(i,j) its magnitude.
- Quantize to $v(i,j) = \lfloor m(i,j)/q_k \rfloor$ where $q_k$ is the quantization step for block k.
- Example: c(i,j) = -10, $q_k$ = 3.
  - s(i,j) = 0
  - v(i,j) = 3



## Bit Planes of Normalized Quantized Coefficients




Quantized coefficients are normalized between –1 and 1



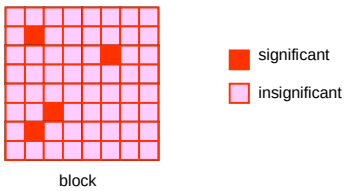
## Bit-Plane Coding of Blocks

- Sub-block significance coding
  - Some sub-blocks are declared insignificant
  - Significant sub-blocks must be coded
- Contexts are defined based on the previous bit-plane significance.
  - Zero coding (ZC) – 8 contexts
  - Run length coding (RLC) – 1 context
  - Sign coding (SC) – 5 contexts
  - Magnitude refinement coding (MR) – 3 contexts
- Code block in raster order using arithmetic coding



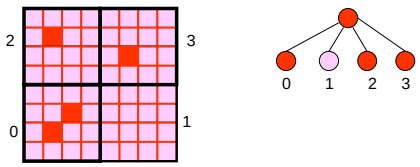
## Sub-Block Significance Coding

- Quad-tree organized group testing
- Block divided into 16x16 sub-blocks
- Identify in few bits the sub-blocks that are significant

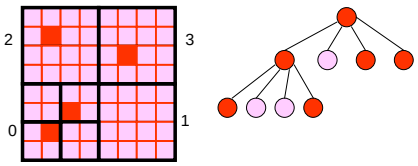





## Quad-Tree Subdivision

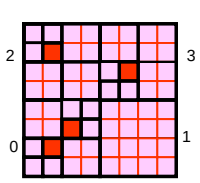


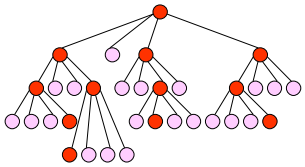
## Quad-Tree Subdivision






## Quad-Tree Subdivision Coding



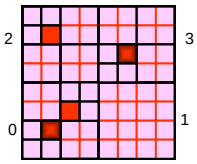
1110001001100001001010001100001000

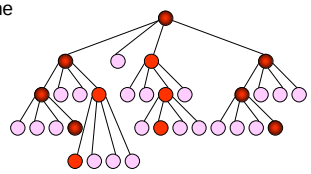
Depth-first code = 1 for significant
0 for insignificant



## Quad-Tree Subdivision Coding

known significant in last bit plane



111000100110000100101000110000100 0

Skip symbols that are already known:
1. nodes significant in previous bit plane
2. last child of significant parent of other children are insignificant



## ZC – Zero Coding


- LH is transposed so that it can be treated the same as HL. $(LH)^T$ has similar characteristics to HL.
- Each coefficient has its neighbors in the same subband


vertical neighbors
horizontal neighbors
diagonal neighbors



## ZC Contexts

- v = number of vertical neighbors significant in the previous bit-plane
- h = number of horizontal neighbors significant in the previous bit-plane
- d = number of diagonal neighbors significant in the previous bit-plane

$0 \le h, v \le 2$
$0 \le d \le 4$

higher labels mean more likely to be significant

| label | HL h | $(LH)^T$ v | LL d | HH d | HH h+v |
|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 |
| 2 | 0 | 0 | > 1 | 0 | > 1 |
| 3 | 0 | 1 | * | 1 | 0 |
| 4 | 0 | 2 | * | 1 | 1 |
| 5 | 1 | 0 | 0 | 1 | > 1 |
| 6 | 1 | 0 | > 0 | 2 | 0 |
| 7 | 1 | > 0 | * | 2 | > 0 |
| 8 | 2 | * | * | > 2 | * |



## Examples

significant in previous bit-plane

HL


h = 0
v = 2
d = 0
Context 4

HH


h = 0
v = 2
d = 0
Context 2

HL


h = 2
v = 0
d = 0
Context 8

HH


h = 2
v = 0
d = 0
Context 2

HL


h = 0
v = 0
d = 2
Context 2

HH


h = 0
v = 0
d = 2
Context 6

## RLC – Run Length Coding

- Looks for runs of 4 that are likely to be insignificant


- If all insignificant then code as a single symbol
- Main purpose – to lighten the load on the arithmetic coder.



## SC – Sign Coding

$$hs = \begin{cases} 0 & \text{if horizontal neighbors are both insignificant or of opposite sign} \\ 1 & \text{if at least one horizontal neighbor is positive} \\ -1 & \text{if at least one horizontal neighbor is negative} \end{cases}$$

$$vs = \begin{cases} 0 & \text{if vertical neighbors are both insignificant or of opposite sign} \\ 1 & \text{if at least one vertical neighbor is positive} \\ -1 & \text{if at least one vertical neighbor is negative} \end{cases}$$

| | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|
| hs | 1 | 1 | 1 | 0 | 0 | 0 | -1 | -1 | -1 |
| vs | 1 | 0 | -1 | 1 | 0 | -1 | 1 | 0 | -1 |
| sign prediction | 1 | 1 | 1 | -1 | 1 | 1 | -1 | -1 | -1 |
| label | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 |



## MR – Magnitude Refinement

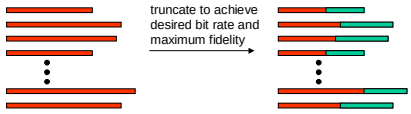
- This is the refinement pass.
- Define t = 0 if first refinement bit, t = 1 otherwise.

| label | t | h + v |
|---|---|---|
| 0 | 0 | 0 |
| 1 | 0 | > 0 |
| 2 | 1 | * |



## Bit Allocation

- How do we truncate the encoded blocks to to achieve a desired bit rate and get maximum fidelity






## Basic Set Up

- Encoded block k can be truncated to $n_k$ bits.
- Total Bit Rate

$$\sum_k n_k$$

- Distortion attributable to block k is

$$D_k^{n_k} = w_k^2 \sum_{(i,j)\in B_k} (c^{n_k}(i,j) - c(i,j))^2$$

where $w_k$ is the "weight" of the basis vectors for block k and $c^{n_k}(i,j)$ is the recovered coefficients from $n_k$ bits of block k.



## Bit Allocation as an Optimization Problem

- Input: Given m embedded codes and a bit rate target R
- Output: Find truncation values $n_k$, $1 \le k \le m$, such that

$$D = \sum_k D_k^{n_k} \quad \text{is minimized and}$$

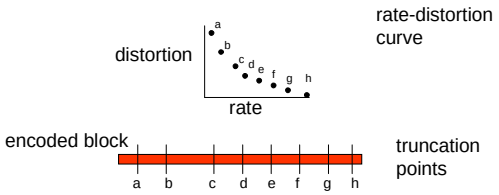
$$\sum_k n_k \le R$$

## Facts about Bit Allocation

- It is an NP-hard problem generally
- There are fast approximate algorithms that work well in practice
  - Lagrange multiplier method
  - Multiple choice knapsack method

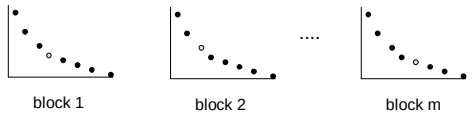


## Rate-Distortion Curve






## Picture of Bit Allocation




Pick one point from each curve so that the sum of the x values is bounded by R and the sum of the y values is minimized.

Good approximate algorithms exist because the curves are almost convex.



## Notes on EBCOT

- EBCOT is quite complicated with many features.
- JPEG 2000 based on EBCOT but differs to improve compression and decompression time.
- EBCOT has
  - resolution scalability
  - SNR scalability
  - quantization
  - bit allocation
  - arithmetic coding with context and adaptivity
  - group testing (quad trees)
  - sign and refinement bit contexts
  - lots of engineering



## Notes on Wavelet Compression

- Wavelets appear to be excellent for image compression
  - No blocking artifacts
  - Wavelet coding techniques abound and are very effective
- Some of the wavelet coding techniques can apply to block transforms.
- Newest generation of image compressor use wavelets, JPEG 2000.