

# CSE 490 GZ Introduction to Data Compression Winter 2002

## Adaptive Huffman Coding

## Adaptive Huffman Coding

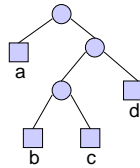
- One pass
- During the pass calculate the frequencies
- Update the Huffman tree accordingly
  - Coder – new Huffman tree computed after transmitting the symbol
  - Decoder – new Huffman tree computed after receiving the symbol
- Symbol set and their basic codes must be known ahead of time.
- Need NYT (not yet transmitted symbol) to indicate a new leaf is needed in the tree.

CSE 490gz - Lecture 3 - Winter 2002

2

## Optimal Tree Numbering

- a : 5, b : 2, c : 1, d : 3

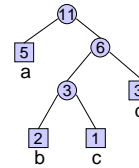


CSE 490gz - Lecture 3 - Winter 2002

3

## Weight the Nodes

- a : 5, b : 2, c : 1, d : 3

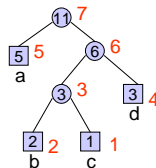


CSE 490gz - Lecture 3 - Winter 2002

4

## Number the Nodes

- a : 5, b : 2, c : 1, d : 3



Number the nodes as they are removed from the priority queue.

CSE 490gz - Lecture 3 - Winter 2002

5

## Adaptive Huffman Principle

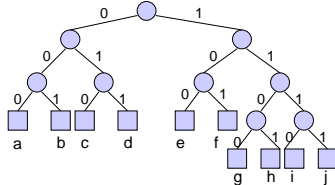
- In an optimal tree for  $n$  symbols there is a numbering of the nodes  $y_1 < y_2 < \dots < y_{2n-1}$  such that their corresponding weights  $x_1, x_2, \dots, x_{2n-1}$  satisfy:
  - $x_1 \leq x_2 \leq \dots \leq x_{2n-1}$
  - siblings are numbered consecutively
- And **vice versa**
  - That is, if there is such a numbering then the tree is optimal. We call this the **node number invariant**.

CSE 490gz - Lecture 3 - Winter 2002

6

## Initialization

- Symbols  $a_1, a_2, \dots, a_m$  have a basic prefix code, used when symbols are first encountered.
- Example: a, b, c, d, e, f, g, h, i, j

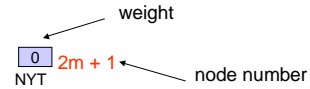


CSE 490gz - Lecture 3 - Winter 2002

7

## Initialization

- The tree will encode up to  $m + 1$  symbols including NYT.
- We reserve numbers 1 to  $2m + 1$  for node numbering.
- The initial Huffman tree consists of a single node



CSE 490gz - Lecture 3 - Winter 2002

8

## Coding Algorithm

- If a new symbol is encountered then output the code for NYT followed by the fixed code for the symbol. Add the new symbol to the tree.
- If an old symbol is encountered then output its code.
- Update the tree to preserve the node number invariant.

CSE 490gz - Lecture 3 - Winter 2002

9

## Decoding Algorithm

- Decode the symbol using the current tree.
- If NYT is encountered then use the fixed code to decode the symbol. Add the new symbol to the tree.
- Update the tree to preserve the node number invariant.

CSE 490gz - Lecture 3 - Winter 2002

10

## Updating the Tree

- Let  $y$  be leaf (symbol) with current weight  $x$ .\*
- If  $y$  the root update  $x$  by 1, otherwise,
- Exchange  $y$  with the largest numbered node with the same weight (unless it is the parent).\*\*
- Update  $x$  by 1
- Let  $y$  be the parent with its weight  $x$  and go to 2.

\*We never update the weight of NYT

\*\* This exchange will preserve the node number invariant

CSE 490gz - Lecture 3 - Winter 2002

11

## Example

- aabcdad in alphabet  $\{a, b, \dots, j\}$



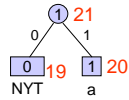
output = 000  
fixed code

CSE 490gz - Lecture 3 - Winter 2002

12

### Example

- aabcdad



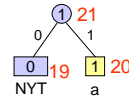
output = 000

CSE 490gz - Lecture 3 - Winter 2002

13

### Example

- aabcdad



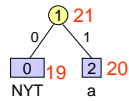
output = 0001

CSE 490gz - Lecture 3 - Winter 2002

14

### Example

- aabcdad



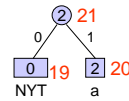
output = 0001

CSE 490gz - Lecture 3 - Winter 2002

15

### Example

- aabcdad



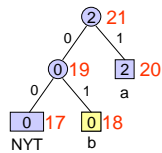
NYT  
fixed code for b  
output = 00010001

CSE 490gz - Lecture 3 - Winter 2002

16

### Example

- aabcdad



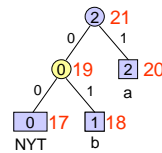
output = 00010001

CSE 490gz - Lecture 3 - Winter 2002

17

### Example

- aabcdad



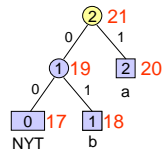
output = 00010001

CSE 490gz - Lecture 3 - Winter 2002

18

### Example

- aab**cd**ad



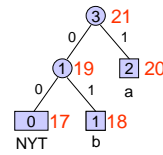
output = 00010001

CSE 490gz - Lecture 3 - Winter 2002

19

### Example

- aab**cd**ad



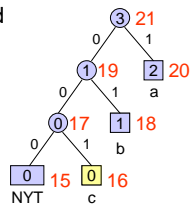
NYT → fixed code for c  
output = 0001000100010

CSE 490gz - Lecture 3 - Winter 2002

20

### Example

- aab**cd**ad



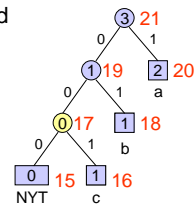
output = 0001000100010

CSE 490gz - Lecture 3 - Winter 2002

21

### Example

- aab**cd**ad



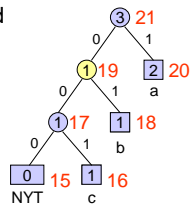
output = 0001000100010

CSE 490gz - Lecture 3 - Winter 2002

22

### Example

- aab**cd**ad



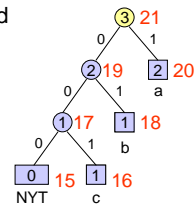
output = 0001000100010

CSE 490gz - Lecture 3 - Winter 2002

23

### Example

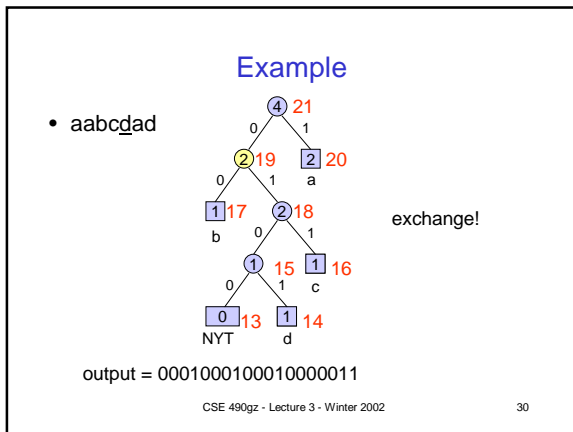
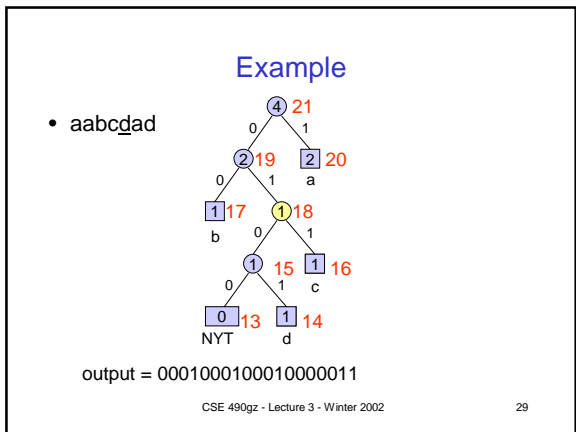
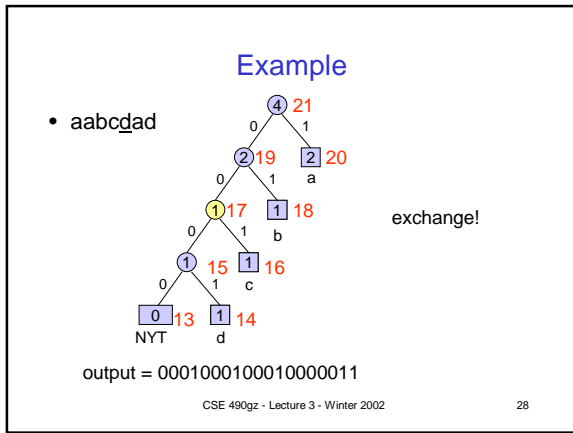
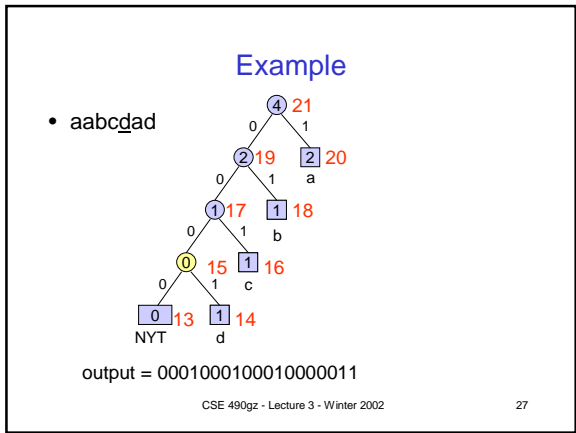
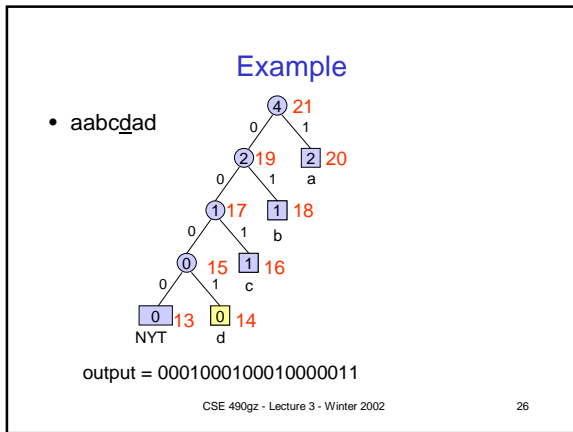
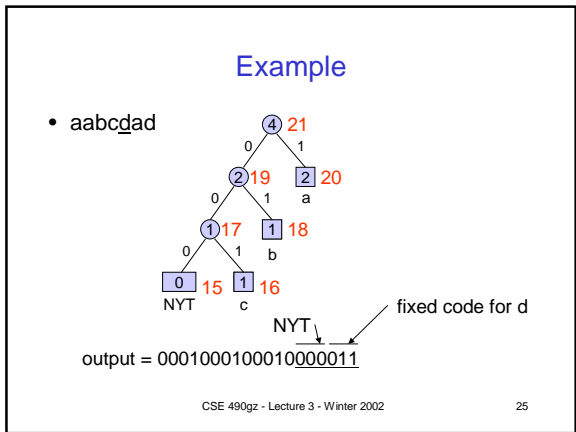
- aab**cd**ad

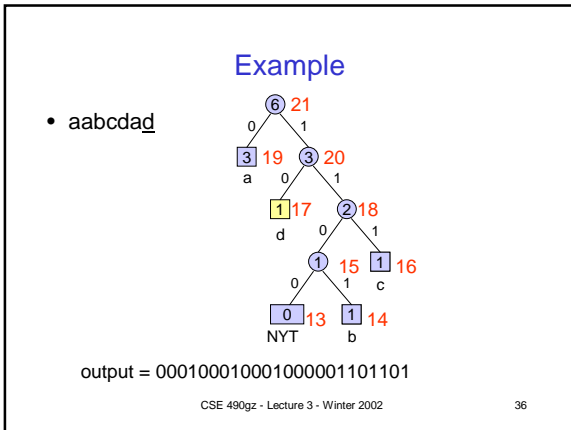
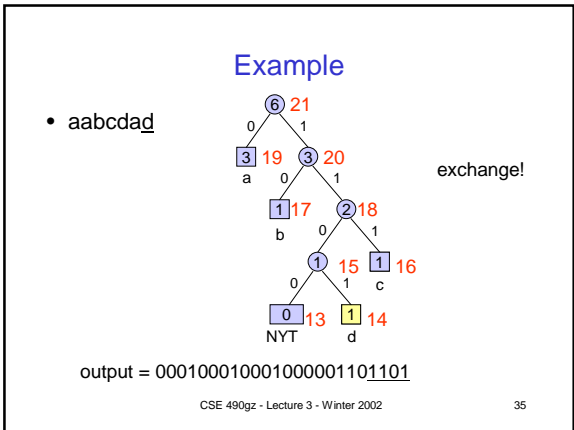
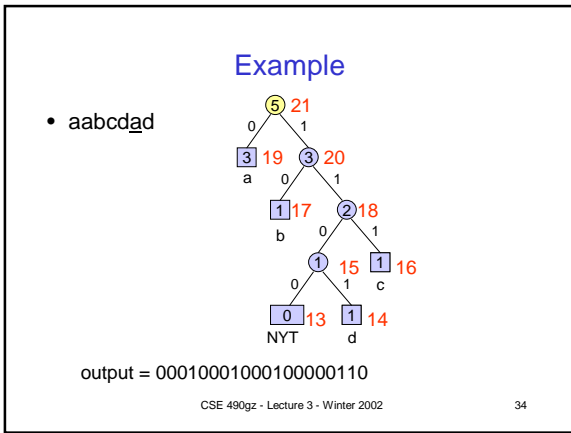
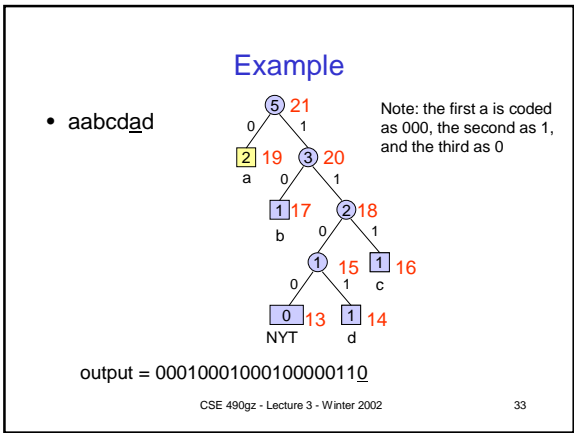
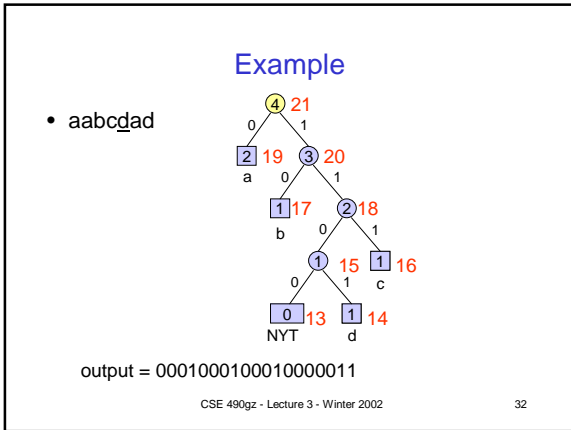
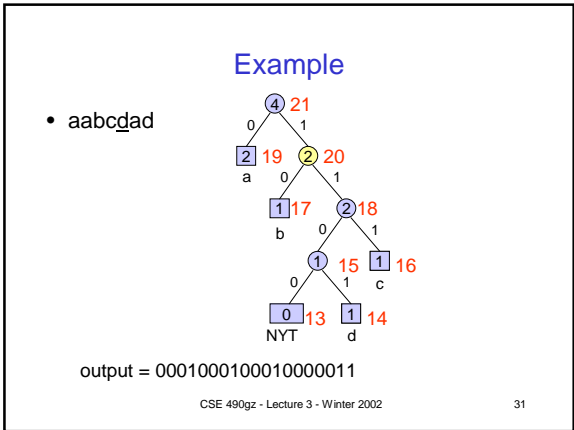


output = 0001000100010

CSE 490gz - Lecture 3 - Winter 2002

24





Example

- aab~~c~~dad

output = 000100010001000001101101

CSE 490gz - Lecture 3 - Winter 2002 37

Example

- aab~~c~~dad

output = 000100010001000001101101

CSE 490gz - Lecture 3 - Winter 2002 38

Example

- aab~~c~~dad

output = 000100010001000001101101

CSE 490gz - Lecture 3 - Winter 2002 39

Data Structure for Adaptive Huffman

1. Fixed code table
2. Binary tree with parent pointers
3. Table of pointers nodes in tree
4. Doubly linked list to rank the nodes

CSE 490gz - Lecture 3 - Winter 2002 40

In Class Exercise

- Decode using adaptive Huffman coding assuming the following fixed code

- 00110000

CSE 490gz - Lecture 3 - Winter 2002 41

Example of Huffman

- Statistical compression algorithm
- Prefix code
- Fixed-to-variable rate code
- Optimization to create a best code
- Symbol merging
- Context
- Adaptive coding
- Decoder and encoder behave almost the same
- Need for data structures and algorithms

CSE 490gz - Lecture 3 - Winter 2002 42