

CSE 484/M584: Computer Security (and Privacy)

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Admin

- Office hours are posted
 - They will start FRIDAY (none today or Thursday)
 - My office hours are Friday afternoons 230-330.
 - We have a LOT of office hours. Use them!
- Lab 1 will be posted tomorrow/tonight.
 - 1a will be due next Wednesday (Section should get you 90% of the way done)



Threat Modeling

Threat Modeling

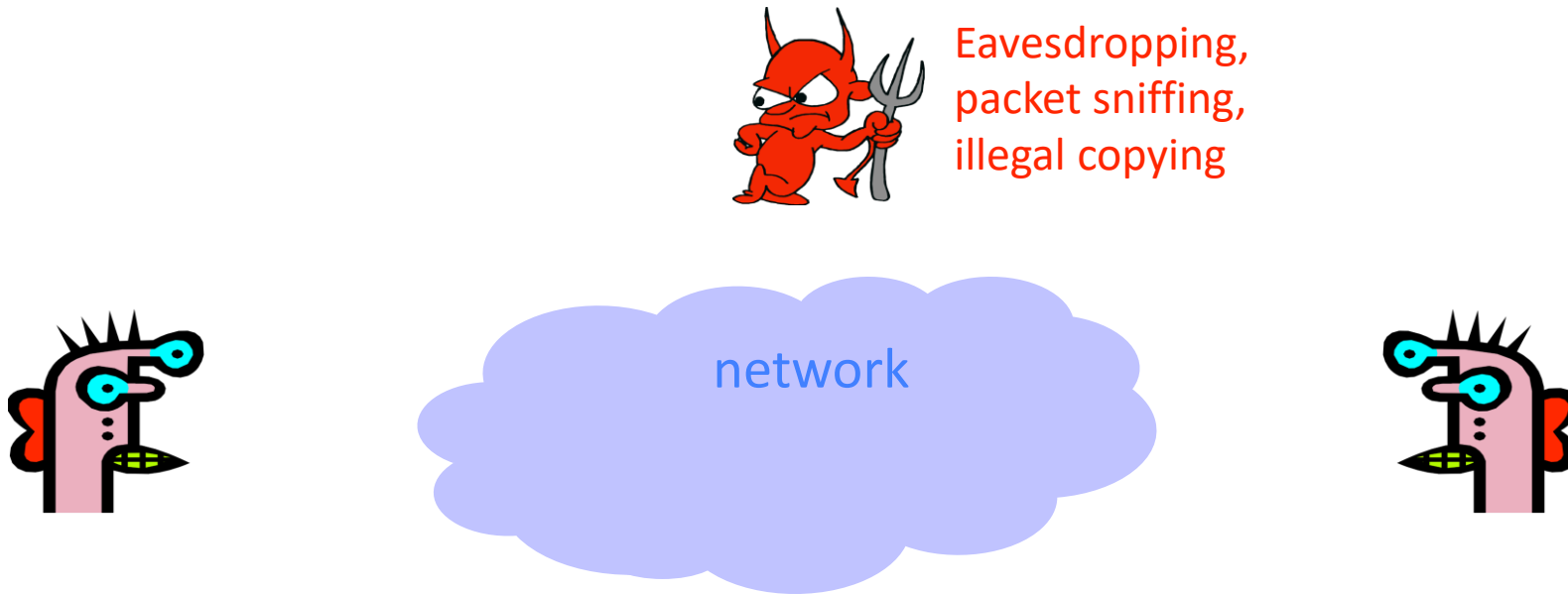
- **Assets**: What are we trying to protect?
- **Adversaries**: Who might try to attack, and why?
- **Vulnerabilities**: How might the system be weak?
- **Threats**: What actions might an adversary take to exploit vulnerabilities?
- **Risk**: How important are assets? How likely is exploit?
- **Possible Defenses**: What mitigations do we have available?
- Not “traditional” threat modeling, but important (both in general, and to help better understand the system prior to threat modeling):
 - **Benefits**: Who might the system benefit, and how?
 - **Harms**: Who might the system harm, and how?

What's *Security*, Anyway?

- Common general security goals: “CIA”
 - Confidentiality
 - Integrity
 - Availability
- Or the extension: CPIAAU (Parkerian Hexad)
 - Control
 - Authenticity
 - Utility

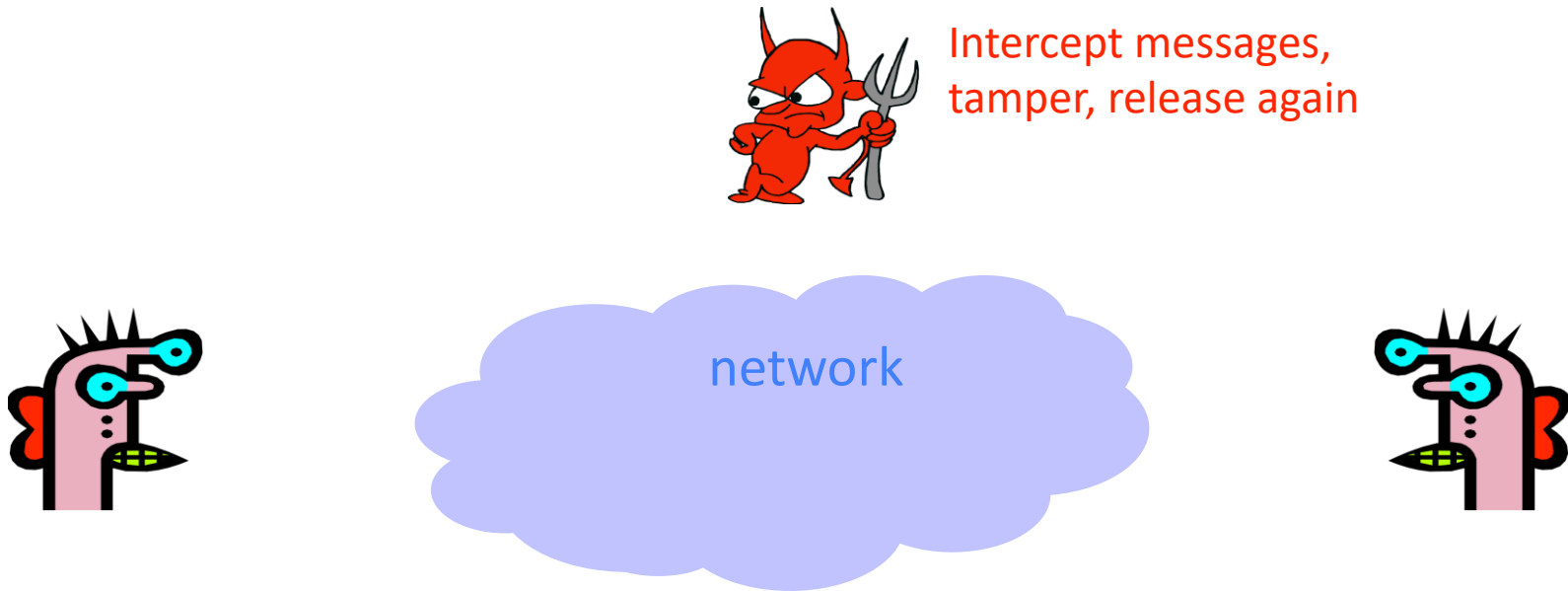
Confidentiality (Privacy)

- Confidentiality is **concealment of information**.



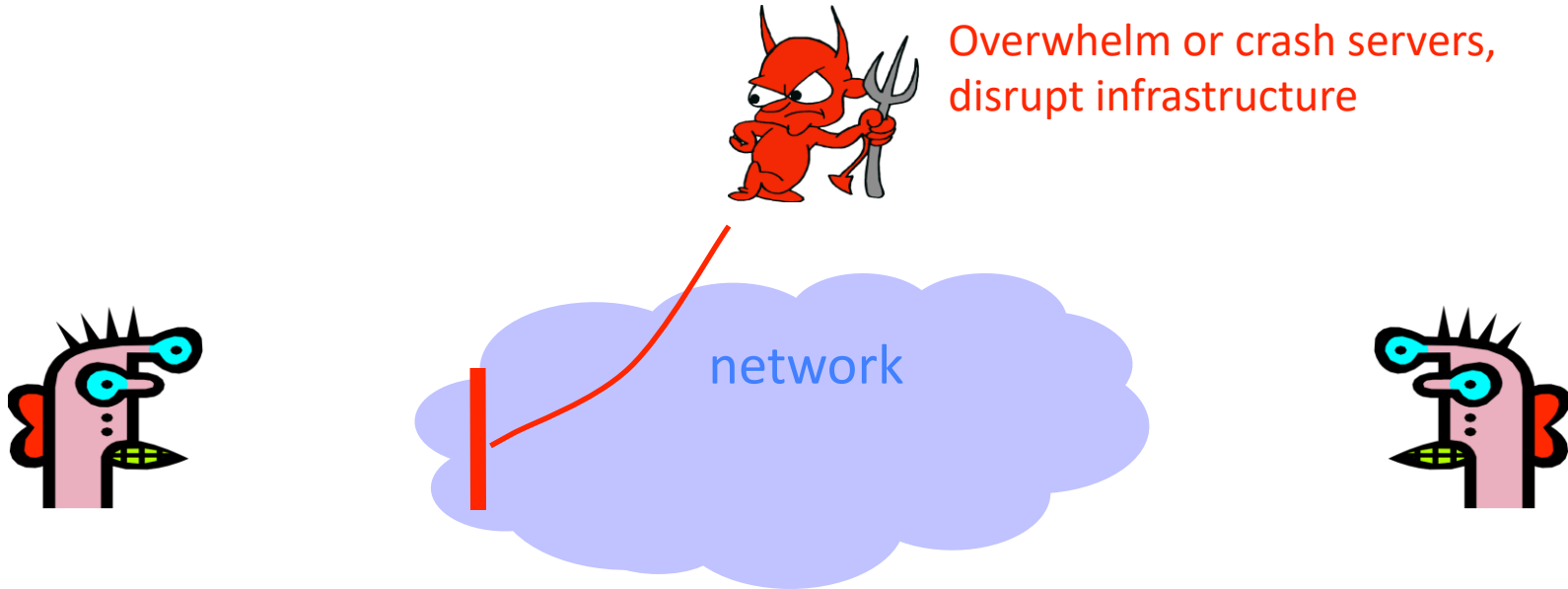
Integrity

- Integrity is prevention of unauthorized changes.



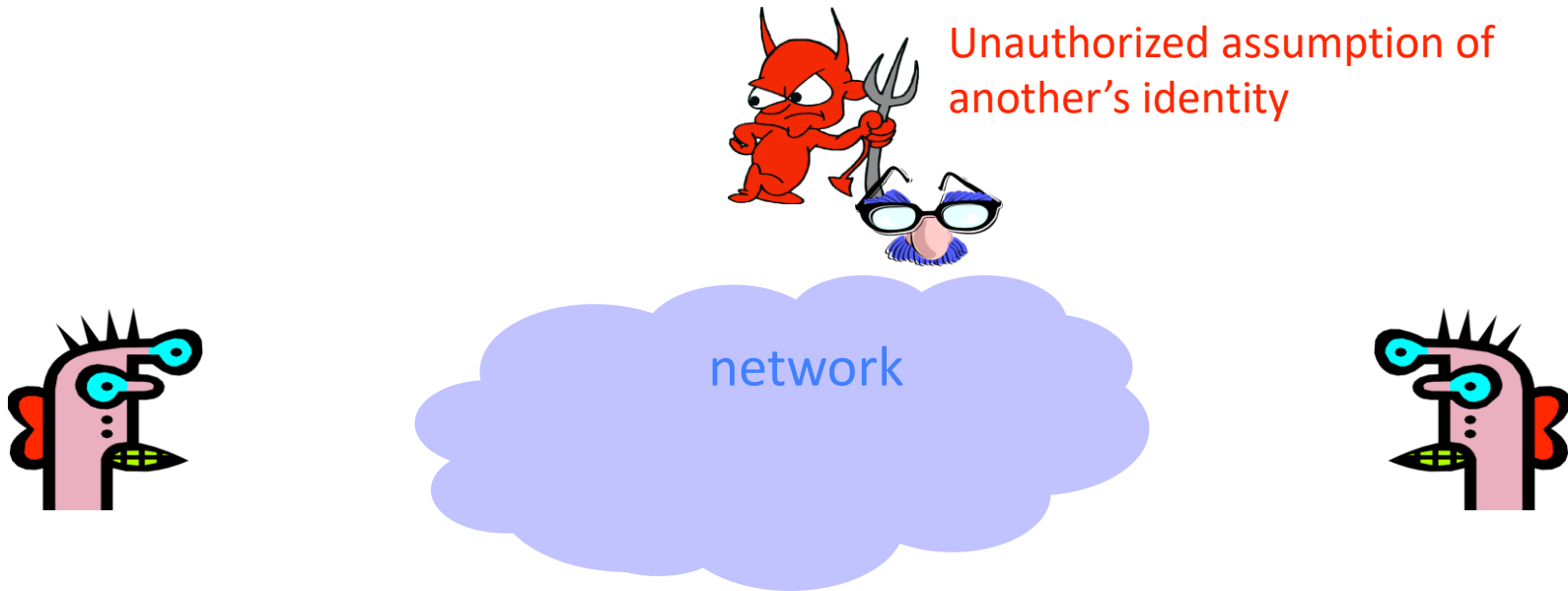
Availability

- Availability is **ability to use information or resources.**



Authenticity

- Authenticity is **knowing who you're talking to**.

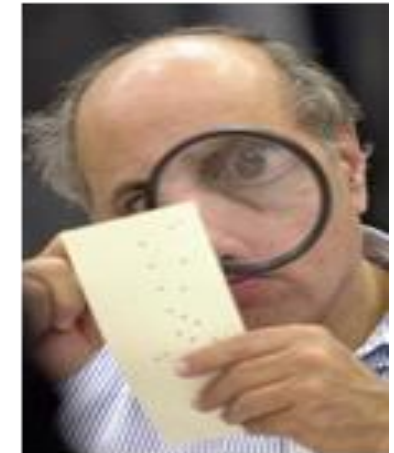


Threat Modeling

- There's no such thing as perfect security
 - But, attackers have limited resources
 - **Make them pay unacceptable costs / take on unacceptable risks to succeed!**
- Defining security per context: identify assets, adversaries, motivations, threats, vulnerabilities, risk, possible defenses

Threat Modeling Example: Electronic Voting

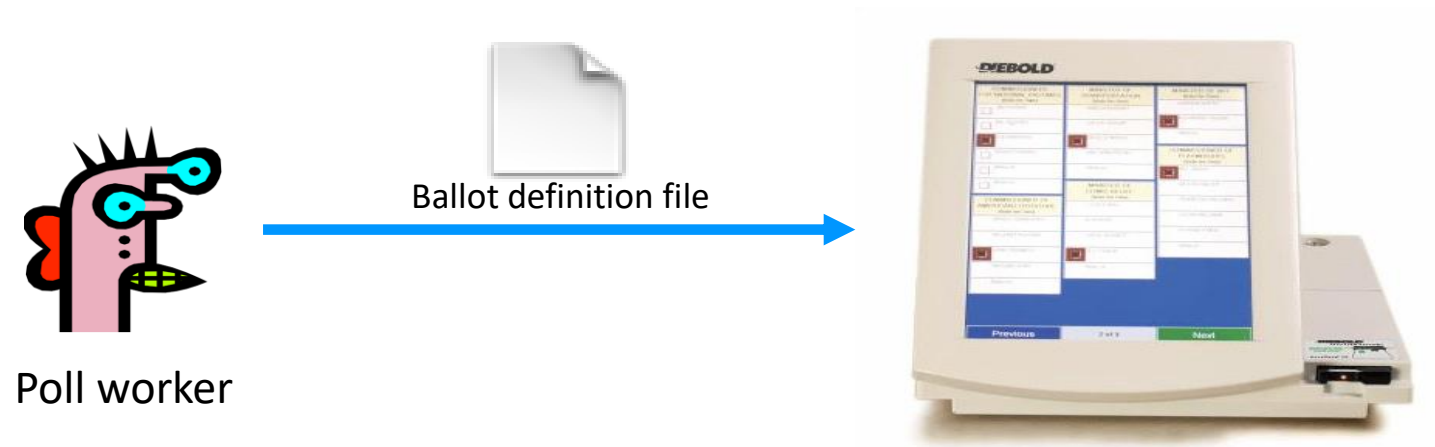
- Popular replacement to traditional paper ballots



Before we get into the system itself

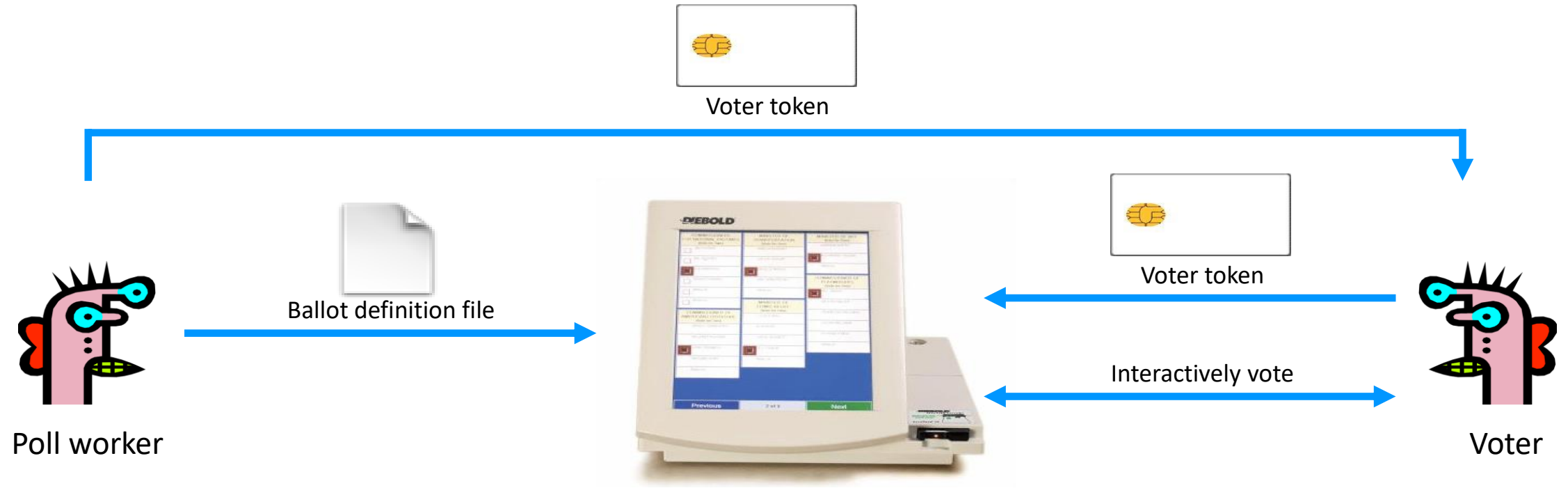
- Think about how in-person electronic voting must work
 - Machines
 - Votes
 - Voters
 - Etc.
- Lightning round threat modeling!
 - Assets?
 - Adversaries?
 - Risks (What might happen if assets are compromised?)
 - Defenses?

Pre-Election



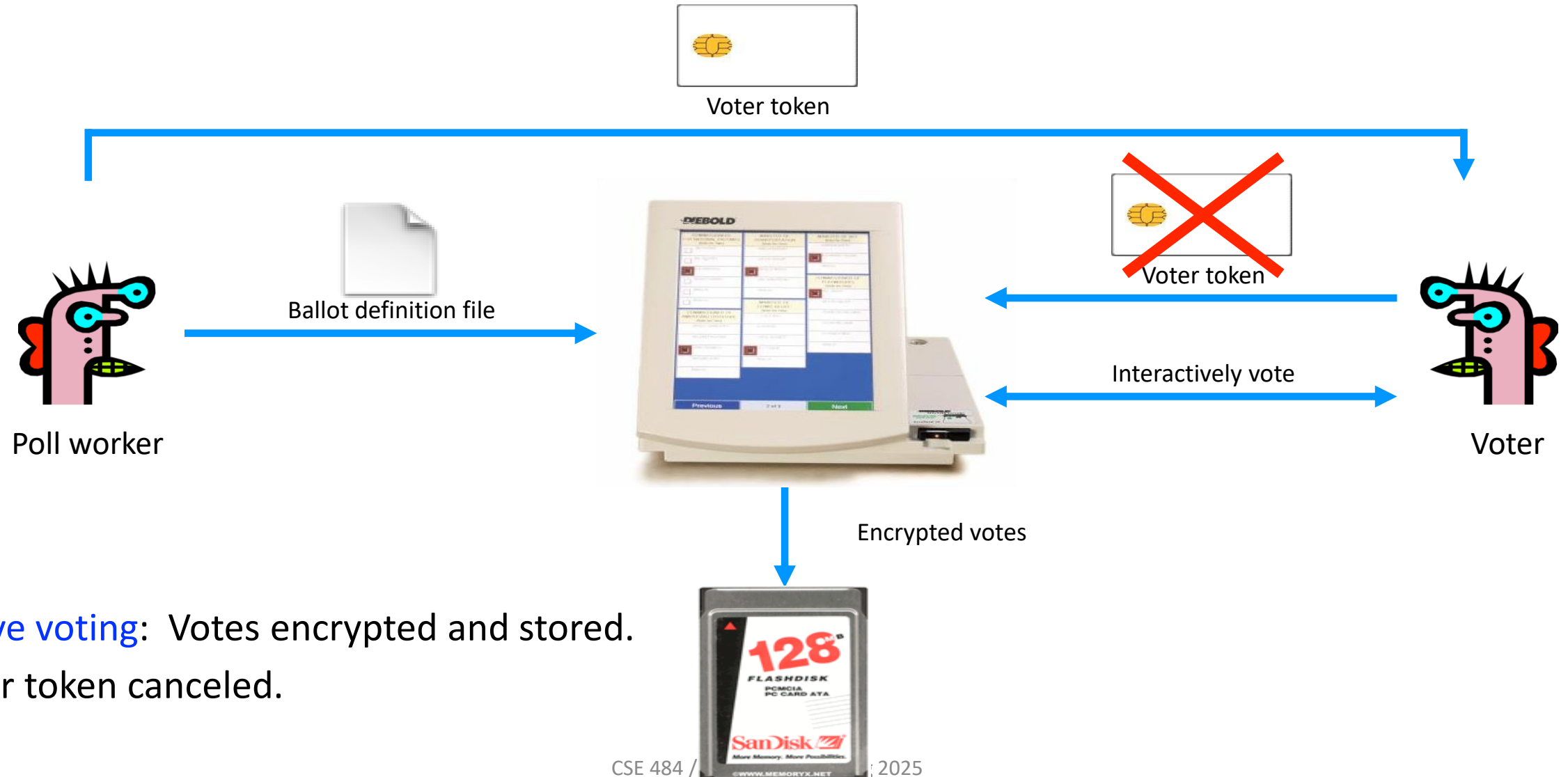
Pre-election: Poll workers load “ballot definition files” on voting machine.

Active Voting

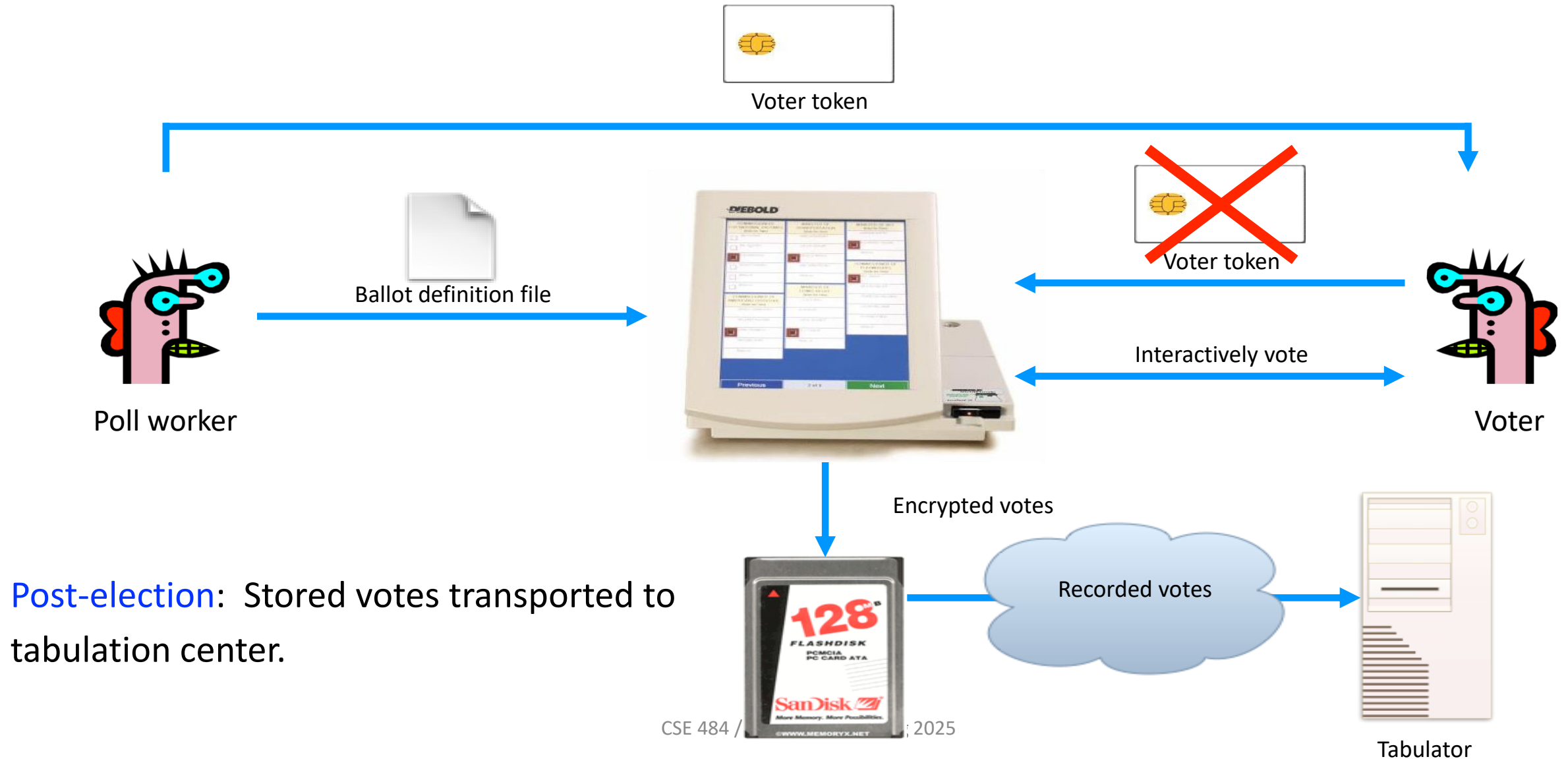


Active voting: Voters obtain **single-use** tokens from poll workers. Voters use tokens to **activate machines** and vote.

Active Voting



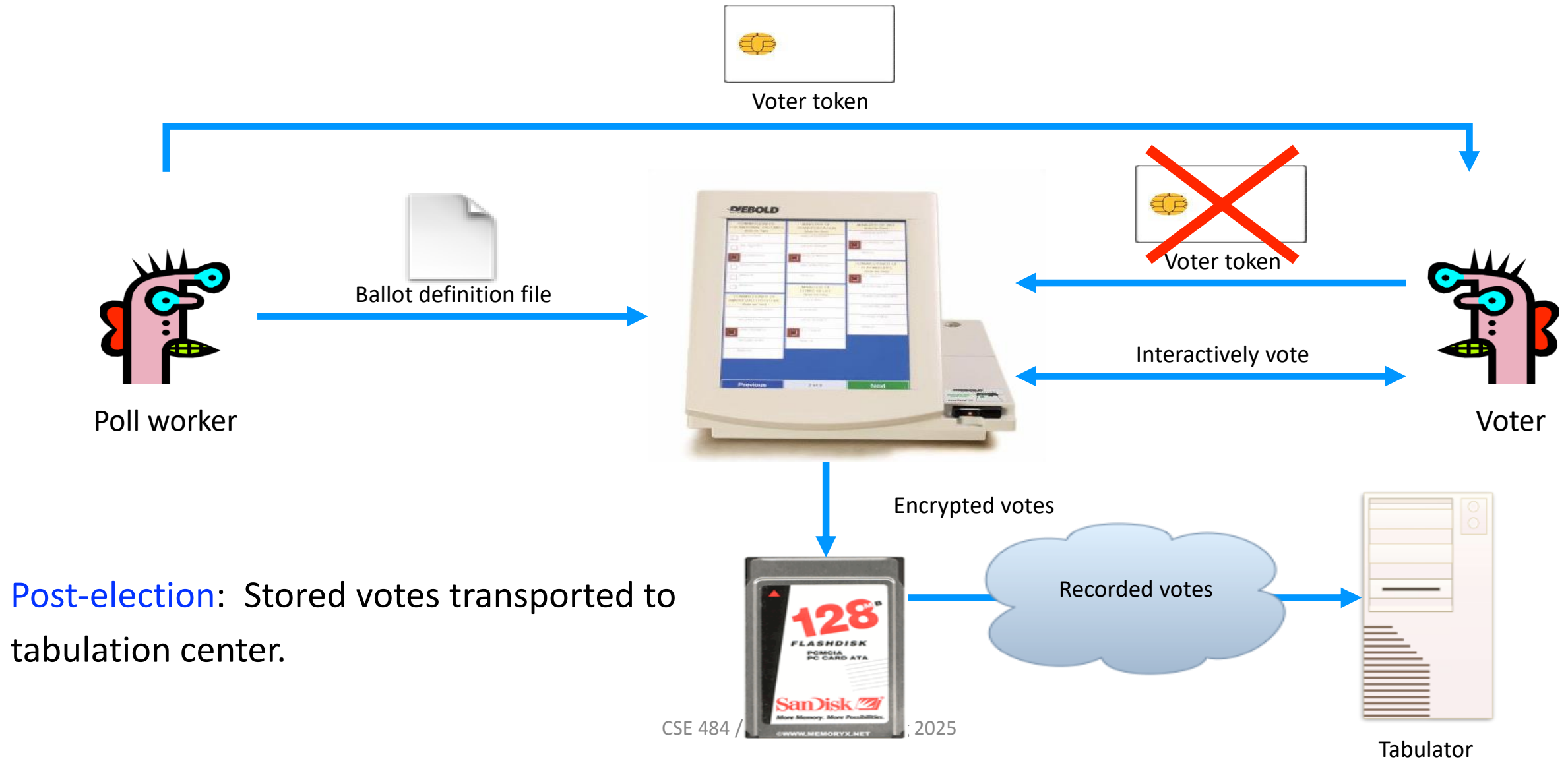
Post-Election



Threat modeling – In detail

- Gradescope
- Fill out the questions while discussing with your neighbors
 - Everyone should submit their own
 - **Polish not required, get down some good ideas!**

Post-Election





KEYS TO THE KINGDOM

Photo taken from Diebold's online store. The keys that open every Diebold touch-screen voting machine. Working copies have been made from the photo.

The x86(_64)

First technical component of the course

- Understanding classic binary vulnerabilities and exploitation techniques
- We'll be doing everything on x86 (32 or 64bit)
- Code will be in C
 - Lab 1: very little C
 - FP: More C

Reminders

- Manual memory management
- Strings are ‘just’ arrays of bytes, *no length field*
 - Strings only end when there is a NUL(NULL) byte.
- Pointers/integers/etc are all just bytes

C strings

```
void main(int argc, char* argv){  
    char string1[32];  
  
    memset(string1, 'a', 32);  
    string1[31] = 0x00; // or, '\0'  
  
    printf("String1: %s\n", string1);  
  
    memset(string1, 'a', 32);  
  
    printf("String1, again: %s\n", string1);  
}
```

Attacks on Memory Buffers

- **Buffer** is a pre-defined data storage area inside computer memory (stack or heap)
- Typical situation:
 - A function takes some input that it writes into a **pre-allocated buffer**.
 - The developer **forgets to check** that the size of the input isn't larger than the size of the buffer.
- **Uh oh.**
 - “Normal” bad input: crash
 - “Adversarial” bad input : take control of execution

Stack Buffers



buf

uh oh!

- Suppose Web server contains this function

```
void func(char *str) {  
    char buf[126];  
    ...  
    strcpy(buf, str);  
    ...  
}
```

- No bounds checking on strcpy()
- If str is longer than 126 bytes
 - Program may crash
 - Attacker may change program behavior

Example: Changing Flags

buf

1 (:-) !)

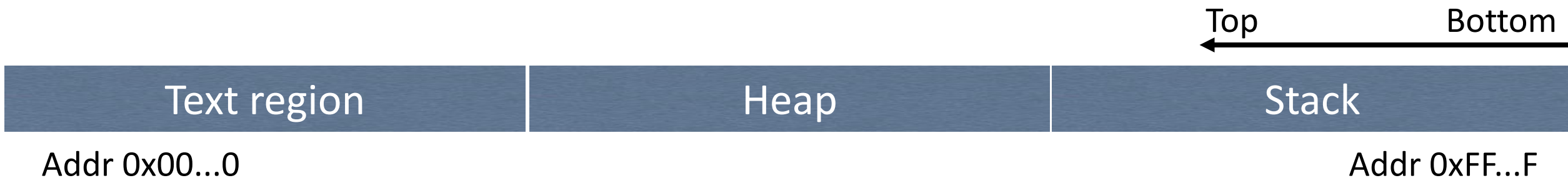
- Suppose Web server contains this function

```
void func(char *str) {  
    byte auth = 0;  
    char buf[126];  
    ...  
    strcpy(buf, str);  
    ...  
}
```

- **Authenticated** variable non-zero when user has extra privileges
- Morris worm also overflowed a buffer to overwrite an authenticated flag in fingerd

Memory Layout

- **Text region:** Executable code of the program
- **Heap:** Dynamically allocated data
- **Stack:** Local variables, function return addresses; grows and shrinks as functions are called and return



Stack Buffers

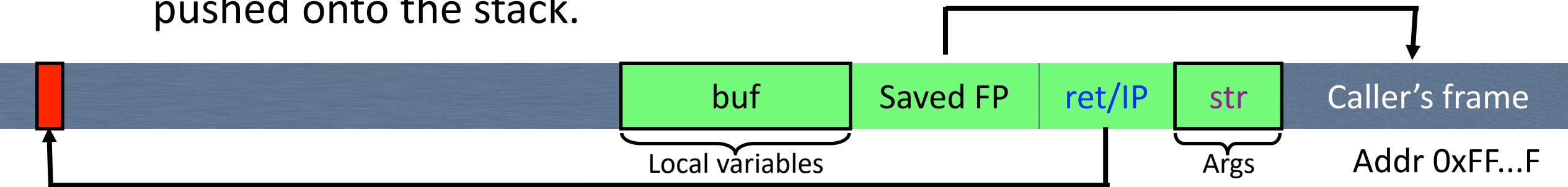
- Suppose Web server contains this function:

```
void func(char *str) {  
    char buf[126];  
    strcpy(buf, str);  
}
```

Allocate local buffer
(126 bytes reserved on stack)

Copy argument into local buffer

- When this function is invoked, a new **frame** (activation record) is pushed onto the stack.



Execute code at this address after func() finishes

What if Buffer is Overstuffed?

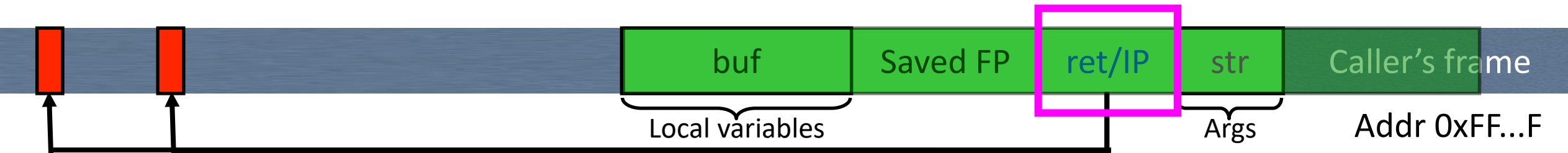
- Memory pointed to by str is copied onto stack...

```
void func(char *str) {  
    char buf[126];  
    strcpy(buf, str);  
}
```

strcpy does NOT check whether the string at *str contains fewer than 126 characters

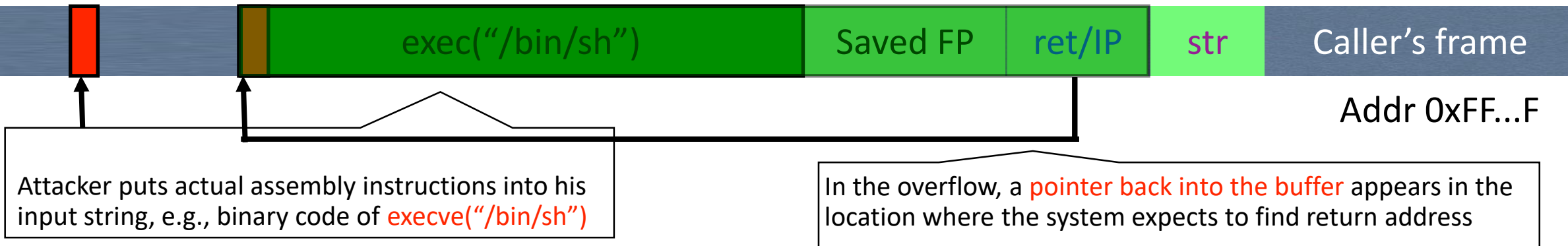
- If a string longer than 126 bytes is copied into buffer, it will overwrite adjacent stack locations.

This will be interpreted as return address!



Executing Attack Code

- Suppose buffer contains attacker-created string
 - For example, `str` points to a string received from the network as the URL



- When function exits, code in the buffer will be executed, giving attacker a shell (**"shellcode"**)
 - **Root shell** if the victim program is setuid root