CSE 484 / CSE M 584: Software Security, Buffer Overflows

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Announcements

- Things Due **(submit to Gradescope):**
    - Ethics Form: Due today!
    - Homework #1: Due Friday
    - Research Readings (CSE M 584): Due Thursday (and every Thursday thereafter)
- Lab 1
    - Out later this week **– don't miss section!**
- Office hours
    - See course website for planned schedule **(double check before you go!)**

Bugs, Vulnerabilities, and Exploits

- Bug
  - Not working quite right

- Vulnerability
  - A malfunction that can be used for an adversary's goals

- Exploit
  - The mechanical set of operations to make use of a vulnerability

Adversarial Failures

- Software bugs are bad
    - Consequences can be serious
- Even worse when an intelligent adversary wishes to exploit them!
    - Intelligent adversaries: Force bugs into "worst possible" conditions/states
    - Intelligent adversaries: Pick their targets

A Few Words on Ethics

- Ethics questions can be subtle complex!

- *Ex: If you find a vulnerability in a deployed product, what should you?*

- HW1 will ask you to learn + consider different ethical frameworks:
  - **Consequentialist:** considers **consequences** of different decisions
  - **Deontological:** considers **duties and rights** of different stakeholders

- You might come to the same conclusion with different frameworks, or you may come to different conclusions. **We are not looking for a single "right answer".**

Memory Corruption Bugs

- Buffer overflows bugs: Big class of bugs
  - Normal conditions: Can sometimes cause systems to fail
  - Adversarial conditions: Attacker able to violate security of your system (control, obtain private information, …)
- Stack, Heap both possibilities

BUFFER OVERFLOWS

A Bit of History: Morris Worm

- Worm was released in 1988 by Robert Morris
  - Graduate student at Cornell, son of NSA chief scientist
  - Convicted under Computer Fraud and Abuse Act,
    - 3 years probation and 400 hours of community service
  - Now an EECS professor at MIT
- Worm was intended to propagate slowly and harmlessly measure the size of the Internet
- Due to a coding error, it created new copies as fast as it could and overloaded infected machines
- $10-100M worth of damage (in 1988)

Morris Worm and Buffer Overflow

- One of the worm's propagation techniques was a buffer overflow attack against a vulnerable version of `fingerd` on VAX systems
  - By sending special string to finger daemon, worm caused it to execute code creating a new worm copy

Buffer overflows remain a common source of vulnerabilities and exploits today!

(Especially in embedded systems.)

Attacks on Memory Buffers

- Buffer is a pre-defined data storage area inside computer memory (stack or heap)

- Typical situation:

  – A function takes some input that it writes into a pre-allocated buffer.

  – The developer forgets to check that the size of the input isn't larger than the size of the buffer.

  – Uh oh.

    - "Normal" bad input: crash

    - "Adversarial" bad input : take control of execution

Stack Buffers

	buf	uh oh!	

- Suppose Web server contains this function

```
void func(char *str) {
        char buf[126];
        ...
        strcpy(buf,str);
        ...
    }
```

- No bounds checking on strcpy()
- If str is longer than 126 bytes
  - Program may crash
  - Attacker may change program behavior

Example: Changing Flags

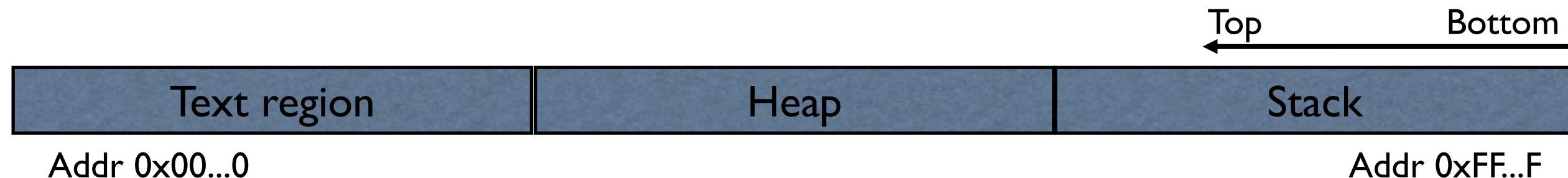|	buf	! (:-) !)	

- Suppose Web server contains this function

```
void func(char *str) {

        byte authenticated = 0;
        char buf[126];
        ...
        strcpy(buf,str);
        ...
    }
```

- Authenticated variable non-zero when user has extra privileges
- Morris worm also overflowed a buffer to overwrite an authenticated flag in fingerd

Refresher: Memory Layout

- Text region: Executable code of the program

- Heap: Dynamically allocated data

- Stack: Local variables, function return addresses; grows and shrinks as functions are called and return

Top ← Bottom

Text region	Heap	Stack

Addr 0x00...0   Addr 0xFF...F
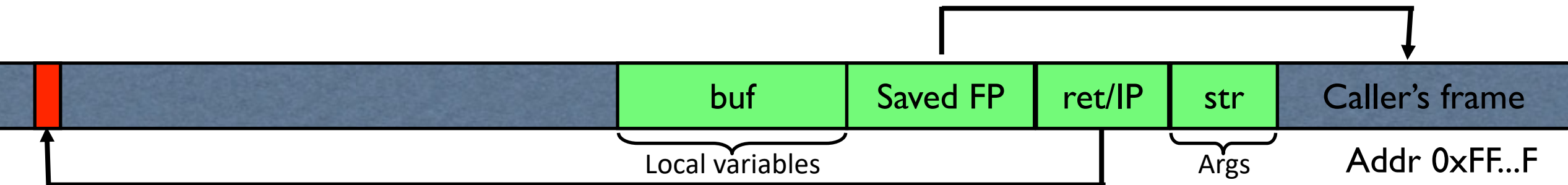
Refresher: Stack Buffers

- Suppose Web server contains this function:

```
void func(char *str) {
    char buf[126];
    strcpy(buf,str);
}
```

Allocate local buffer
(126 bytes reserved on stack)

Copy argument into local buffer

- When this function is invoked, a new frame (activation record) is pushed onto the stack.

| buf | Saved FP | ret/IP | str | Caller's frame |

Local variables

Args

Addr 0xFF...F

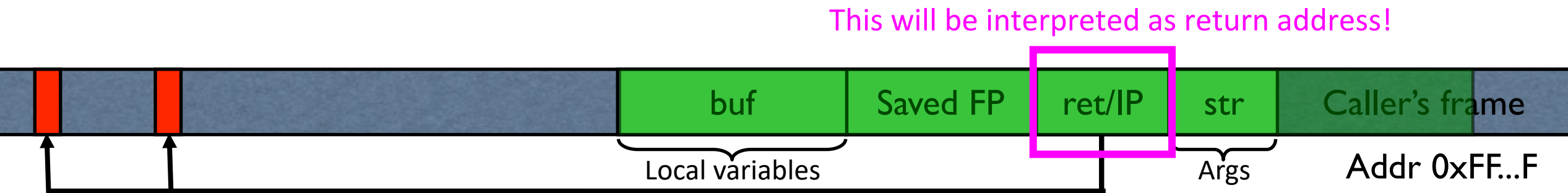Execute code at this address after func() finishes

What if Buffer is Overstuffed?

- Memory pointed to by str is copied onto stack…

```
void func(char *str) {
    char buf[126];
    strcpy(buf,str);
}
```
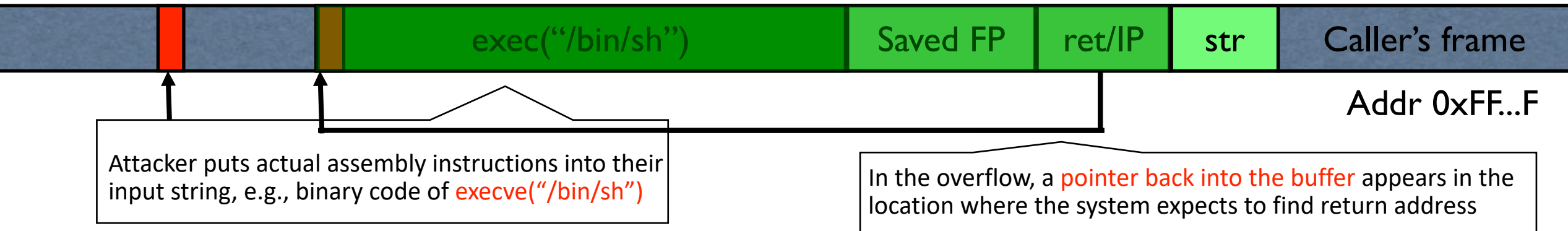
> strcpy does NOT check whether the string at *str contains fewer than 126 characters

- If a string longer than 126 bytes is copied into buffer, it will overwrite adjacent stack locations.

This will be interpreted as return address!

	buf	Saved FP	ret/IP	str	Caller's frame

Local variables

Args

Addr 0xFF...F

Executing Attack Code

- Suppose buffer contains attacker-created string
  - For example, str points to a string received from the network as the URL

| | | | | exec("/bin/sh") | Saved FP | ret/IP | str | Caller's frame |

Addr 0xFF...F

Attacker puts actual assembly instructions into their input string, e.g., binary code of execve("/bin/sh")

In the overflow, a pointer back into the buffer appears in the location where the system expects to find return address

- When function exits, code in the buffer will be executed, giving attacker a shell **("shellcode")**
  - Root shell if the victim program is setuid root

Buffer Overflows Can Be Tricky...

- Overflow portion of the buffer must contain correct address of attack code in the RET position

  - The value in the RET position must point to the beginning of attack assembly code in the buffer

    - Otherwise application will (probably) crash with segfault

  - **Attacker must correctly guess in which stack position his/her buffer will be when the function is called**

Problem: No Bounds Checking

- strcpy does <u>not</u> check input size
  - strcpy(buf, str) simply copies memory contents into buf starting from *str until "\0" is encountered, ignoring the size of area allocated to buf
- Many C library functions are unsafe
  - strcpy(char *dest, const char *src)
  - strcat(char *dest, const char *src)
  - gets(char *s)
  - scanf(const char *format, …)
  - printf(const char *format, …)

Does Bounds Checking Help?

- strncpy(char *dest, const char *src, size_t n)
  - For strncpy (unlike strcpy), no more than n characters will be copied from *src to *dest

- Potential overflow in htpasswd.c (Apache 1.3):

```
strcpy(record,user);
strcat(record,":");
strcat(record,cpw);
```

Copies username ("user") into buffer ("record"), then appends ":" and hashed password ("cpw")

- Published fix:

```
strncpy(record,user,MAX_STRING_LEN-1);
strcat(record,":")
strncat(record,cpw,MAX_STRING_LEN-1);
```

In-Class Activity

In-Class Activity Gradescope – 9/30 Activity
https://www.gradescope.com/courses/881751/assignments/5072027

(This is the first one that will be graded.

Reminder that you can submit up to half of them a week late.)

Does Bounds Checking Help?

- strncpy(char *dest, const char *src, size_t n)
    - For strncpy (unlike strcpy), no more than n characters will be copied from *src to *dest
        - Programmer has to supply the right value of n

- Potential overflow in htpasswd.c (Apache 1.3):

```
strcpy(record,user);
strcat(record,":");
strcat(record,cpw);
```

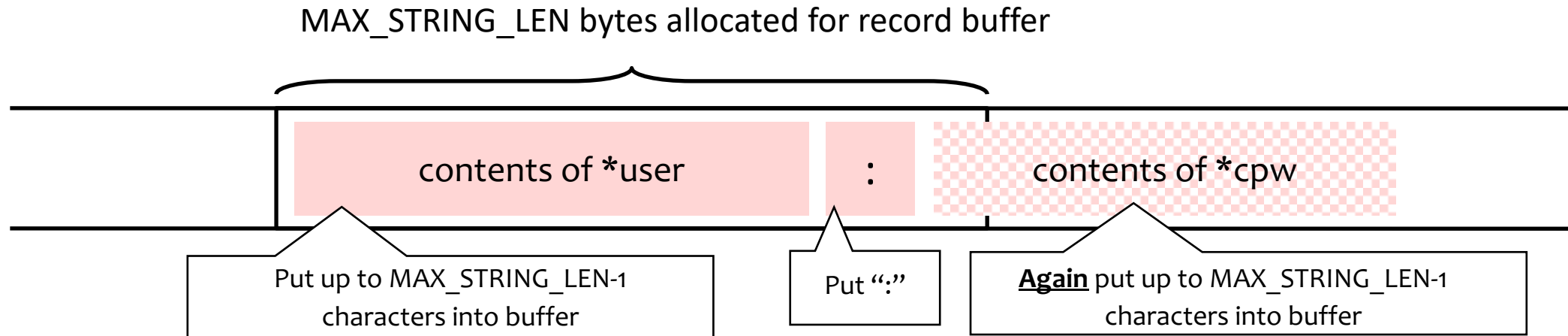Copies username ("user") into buffer ("record"), then appends ":" and hashed password ("cpw")

- Published fix:

```
strncpy(record,user,MAX_STRING_LEN-1);
strcat(record,":")
strncat(record,cpw,MAX_STRING_LEN-1);
```

Misuse of strncpy in htpasswd "Fix"

- Published "fix" for Apache htpasswd overflow:

```
strncpy(record,user,MAX_STRING_LEN-1);
strcat(record,":")
strncat(record,cpw,MAX_STRING_LEN-1);
```

MAX_STRING_LEN bytes allocated for record buffer

| contents of *user | : | contents of *cpw |

Put up to MAX_STRING_LEN-1 characters into buffer

Put ":"

Again put up to MAX_STRING_LEN-1 characters into buffer

What About This?

- Home-brewed range-checking string copy

```
void mycopy(char *input) {
    char buffer[512]; int i;

    for (i=0; i<=512; i++)
        buffer[i] = input[i];
}

void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}
```

In-Class Activity

In-Class Activity Gradescope → 9/30

Off-by-One Overflow

- Home-brewed range-checking string copy

```
void mycopy(char *input) {
    char buffer[512]; int i;

    for (i=0; i<=512; i++)
        buffer[i] = input[i];
}

void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}
```

This will copy **513** characters into buffer. Oops!

- 1-byte overflow: can't change RET, but can change pointer to previous stack frame…