

# **CSE 484 / CSE M 584:** **Software Security, Buffer Overflows**

Fall 2023

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# Announcements

- Things Due:
  - Ethics Form: Due today!
  - Homework #1: Due Friday
  - Research Readings (CSE M 584): Due Thursday (and every Thursday thereafter)
- Lab 1
  - Out later this week
  - Start forming groups!
- Office hours
  - See course website for **tentative** schedule (check before you go!)

# Bugs, Vulnerabilities, and Exploits

- Bug
  - Not working quite right
- Vulnerability
  - A malfunction that can be used for an adversary's goals
- Exploit
  - The mechanical set of operations to make use of a vulnerability

# Adversarial Failures

- Software bugs are bad
  - Consequences can be serious
- Even worse when an **intelligent adversary** wishes to **exploit** them!
  - Intelligent adversaries: Force bugs into “**worst possible**” conditions/states
  - Intelligent adversaries: Pick their targets

# Memory Corruption Bugs

- **Buffer overflows bugs:** Big class of bugs
  - Normal conditions: Can sometimes cause systems to fail
  - Adversarial conditions: Attacker able to violate security of your system (control, obtain private information, ...)
- Stack, Heap both possibilities

# **BUFFER OVERFLOWS**

# A Bit of History: Morris Worm

- Worm was released in 1988 by Robert Morris
  - Graduate student at Cornell, son of NSA chief scientist
  - Convicted under Computer Fraud and Abuse Act,
    - 3 years probation and 400 hours of community service
- Worm was intended to propagate slowly and harmlessly measure the size of the Internet
- Due to a coding error, it created new copies as fast as it could and overloaded infected machines
- \$10-100M worth of damage (in 1988)

# Morris Worm and Buffer Overflow

- One of the worm's propagation techniques was a **buffer overflow attack** against a vulnerable version of `fingerd` on VAX systems
  - By sending special string to `finger` daemon, worm caused it to execute code creating a new worm copy

Buffer overflows remain a common source of vulnerabilities and exploits today!  
(Especially in embedded systems.)



# Attacks on Memory Buffers

- **Buffer** is a pre-defined data storage area inside computer memory (stack or heap)
- Typical situation:
  - A function takes some input that it writes into a **pre-allocated buffer**.
  - The developer **forgets to check** that the size of the input isn't larger than the size of the buffer.
  - **Uh oh.**
    - “Normal” bad input: crash
    - “Adversarial” bad input : take control of execution

# Stack Buffers



buf

uh oh!

- Suppose Web server contains this function

```
void func(char *str) {  
    char buf[126];  
    ...  
    strcpy(buf, str);  
    ...  
}
```

- No bounds checking on `strcpy()`
- If `str` is longer than 126 bytes
  - Program may crash
  - Attacker may change program behavior

# Example: Changing Flags

buf

! (:-) !)

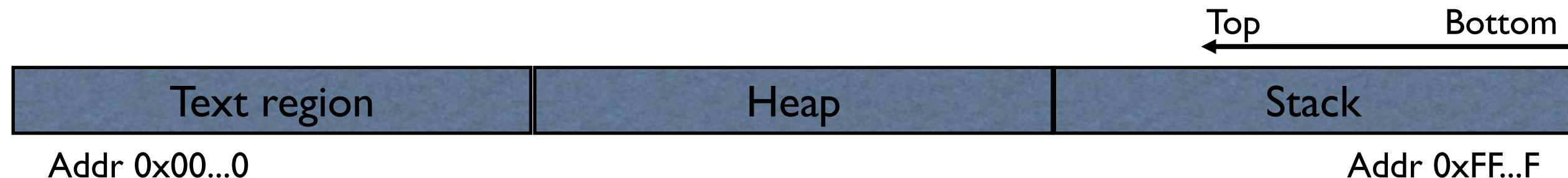
- Suppose Web server contains this function

```
void func(char *str) {  
    byte auth = 0;  
    char buf[126];  
    ...  
    strcpy(buf, str);  
    ...  
}
```

- **Authenticated** variable non-zero when user has extra privileges
- Morris worm also overflowed a buffer to overwrite an authenticated flag in fingerd

# Memory Layout

- **Text region:** Executable code of the program
- **Heap:** Dynamically allocated data
- **Stack:** Local variables, function return addresses; grows and shrinks as functions are called and return



# Stack Buffers

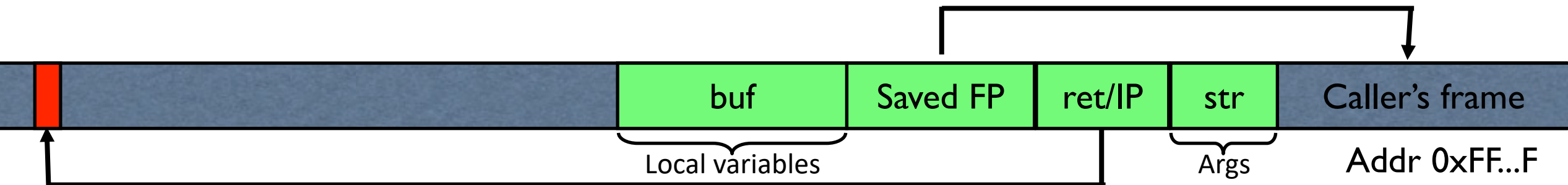
- Suppose Web server contains this function:

```
void func(char *str) {  
    char buf[126];  
    strcpy(buf, str);  
}
```

Allocate local buffer  
(126 bytes reserved on stack)

Copy argument into local buffer

- When this function is invoked, a new **frame** (activation record) is pushed onto the stack.



Execute code at this address after func() finishes

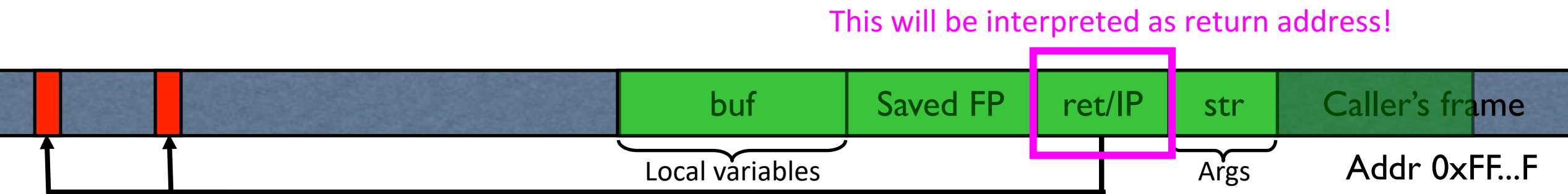
# What if Buffer is Overstuffed?

- Memory pointed to by str is copied onto stack...

```
void func(char *str) {  
    char buf[126];  
    strcpy(buf, str);  
}
```

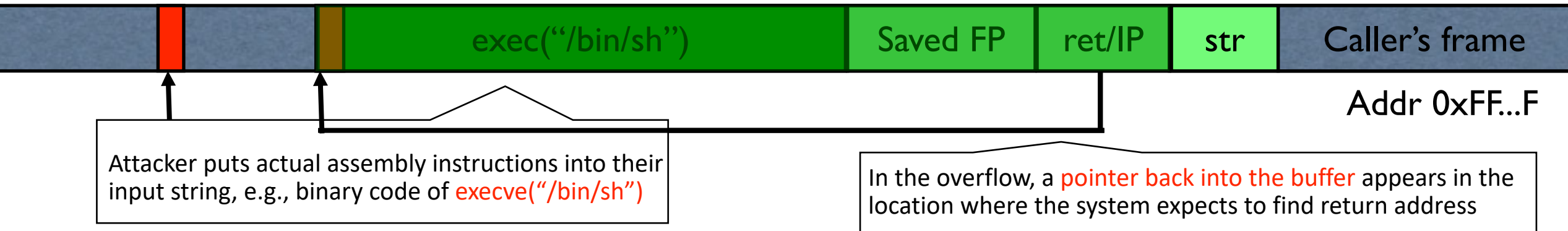
strcpy does NOT check whether the string at \*str contains fewer than 126 characters

- If a string longer than 126 bytes is copied into buffer, it will overwrite adjacent stack locations.



# Executing Attack Code

- Suppose buffer contains attacker-created string
  - For example, `str` points to a string received from the network as the URL



- When function exits, code in the buffer will be executed, giving attacker a shell (**“shellcode”**)
  - **Root shell** if the victim program is `setuid root`

# Buffer Overflows Can Be Tricky...

- Overflow portion of the buffer must contain **correct address of attack code** in the RET position
  - The value in the RET position must point to the beginning of attack assembly code in the buffer
    - Otherwise application will (probably) crash with segfault
  - **Attacker must correctly guess in which stack position his/her buffer will be when the function is called**



# Problem: No Bounds Checking

- strcpy does not check input size
  - strcpy(buf, str) simply copies memory contents into buf starting from \*str until “\0” is encountered, ignoring the size of area allocated to buf
- Many C library functions are unsafe
  - strcpy(char \*dest, const char \*src)
  - strcat(char \*dest, const char \*src)
  - gets(char \*s)
  - scanf(const char \*format, ...)
  - printf(const char \*format, ...)

# Does Bounds Checking Help?

- `strncpy(char *dest, const char *src, size_t n)`
  - For `strncpy` (unlike `strcpy`), no more than `n` characters will be copied from `*src` to `*dest`
    - Programmer has to supply the right value of `n`
- Potential overflow in `htpasswd.c` (Apache 1.3):

```
strcpy(record, user);  
strcat(record, ":");  
strcat(record, cpw);
```

Copies username (“user”) into buffer (“record”), then appends “:” and hashed password (“cpw”)

- Published fix:

```
strncpy(record, user, MAX_STRING_LEN-1);  
strcat(record, ":");  
strncat(record, cpw, MAX_STRING_LEN-1);
```

# In-Class Activity

Canvas -> Quizzes -> Oct 2

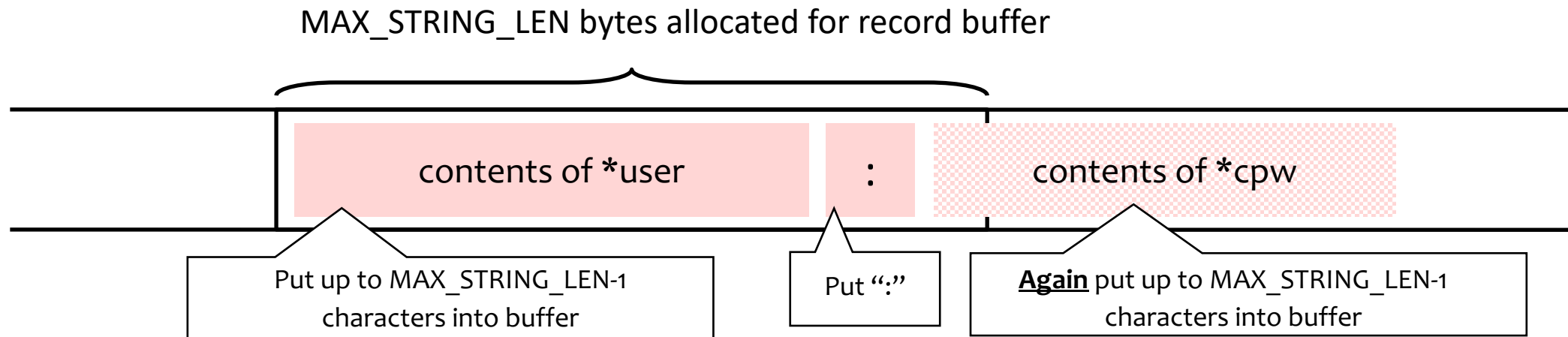
*(This is the first one that will be graded.*

*Reminder that you can submit up to half of them a week late.)*

# Misuse of strncpy in httpasswd “Fix”

- Published “fix” for Apache httpasswd overflow:

```
strncpy(record,user,MAX_STRING_LEN-1);  
strcat(record,":")  
strncat(record,cpw,MAX_STRING_LEN-1);
```



# What About This?

- Home-brewed range-checking string copy

```
void mycopy(char *input) {
    char buffer[512]; int i;

    for (i=0; i<=512; i++)
        buffer[i] = input[i];
}

void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}
```

# In-Class Activity

Canvas -> Quizzes -> Oct 2

# Off-by-One Overflow

- Home-brewed range-checking string copy

```
void mycopy(char *input) {
    char buffer[512]; int i;

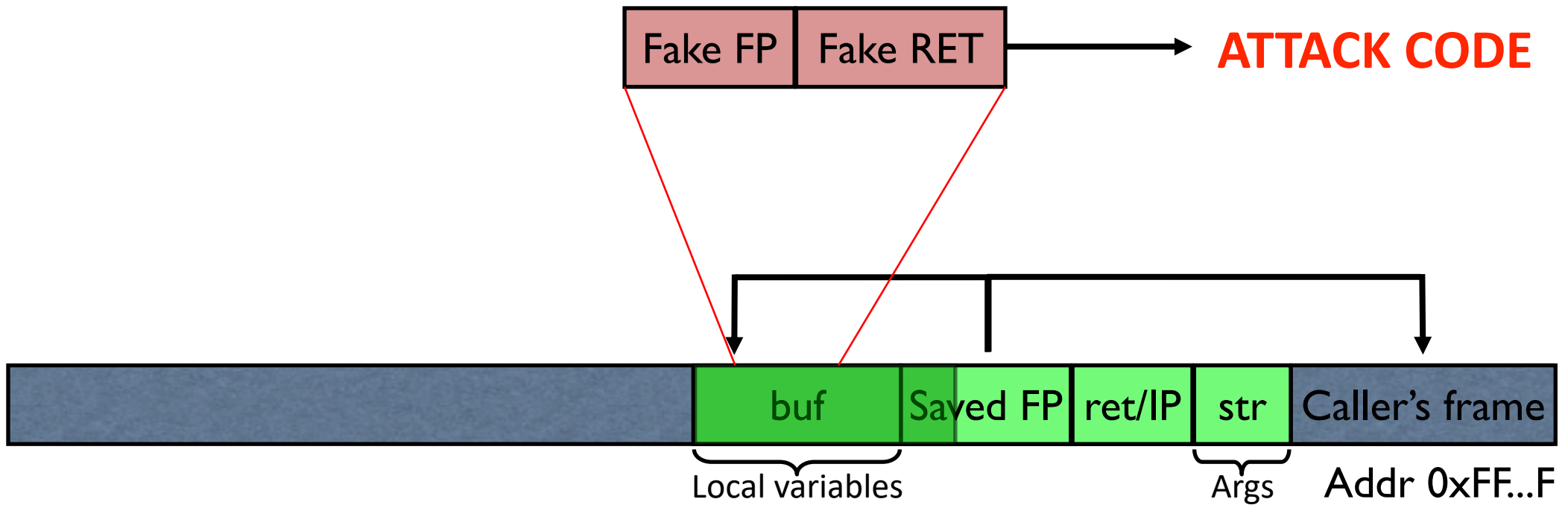
    for (i=0; i<=512; i++)
        buffer[i] = input[i];
}

void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}
```

This will copy 513 characters into buffer. Oops!

- 1-byte overflow: can't change RET, but can change pointer to previous stack frame...

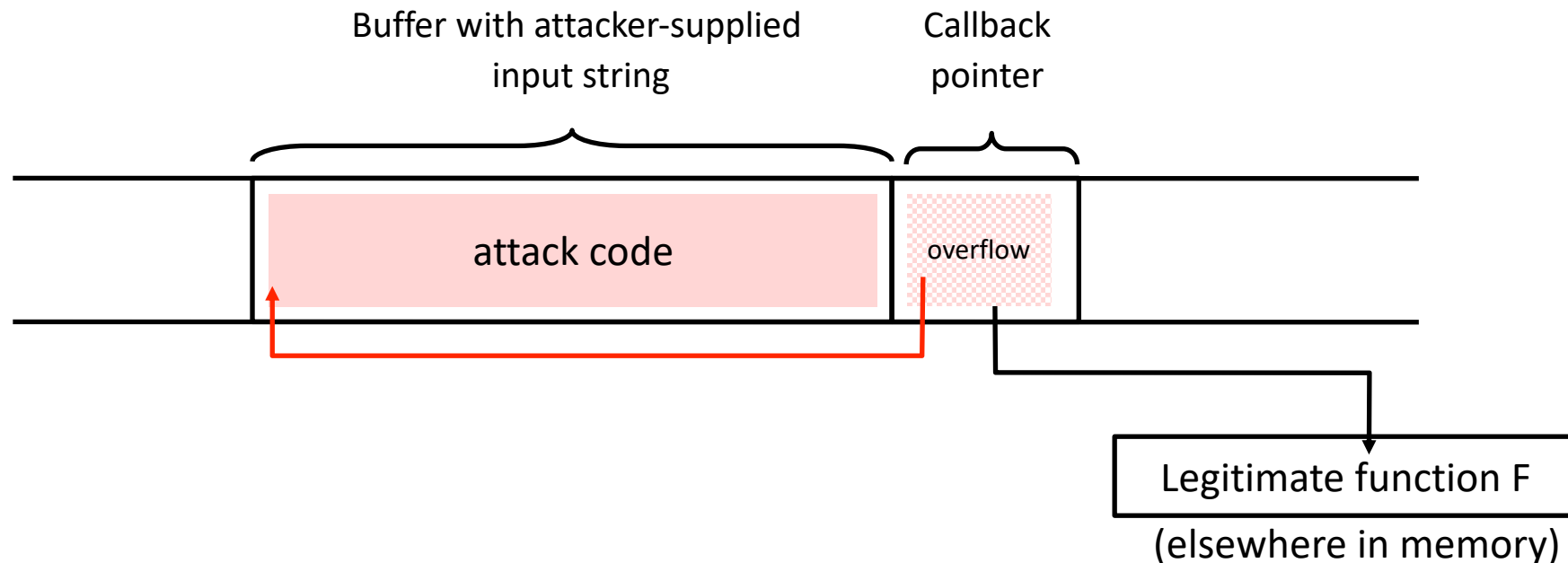
# Frame Pointer Overflow





# Another Variant: Function Pointer Overflow

- C uses **function pointers** for callbacks: if pointer to F is stored in memory location P, then one can call F as  $(*P)(\dots)$



# Other Overflow Targets

- Format strings in C
  - We'll walk through this later
- Heap management structures used by malloc()
  - More details in section
  - Techniques have changed wildly over time
- These are all attacks you can look forward to in Lab #1 😊