CSE 484: Computer Security and Privacy

# Web Security [Web Application Security]

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. . .

#### Admin

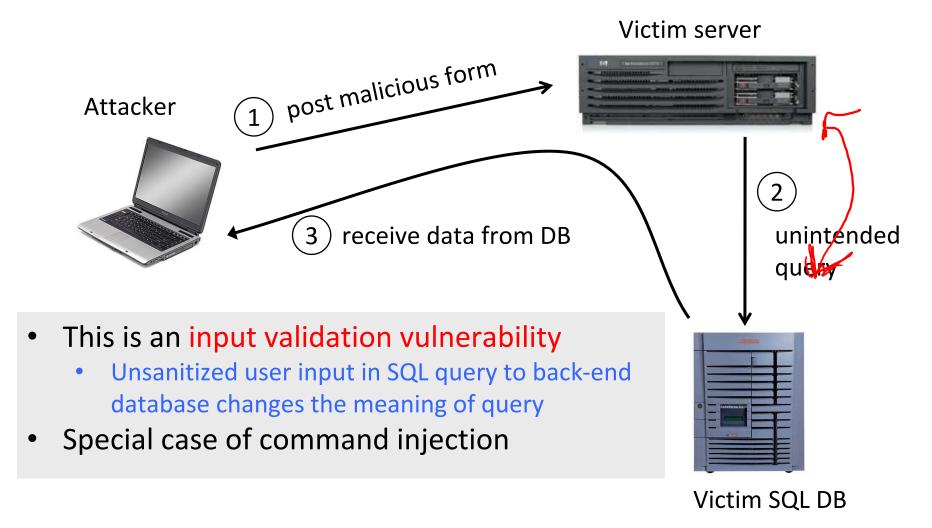
- Lab 2
  - Granting access on a regular basis
  - Please sign up if you haven't already
- Final project

First checkpoint deadline TODAY!

HW2 hit submit

# SQL Injection

### SQL Injection: Basic Idea





#### Authentication with Backend DB

```
set UserFound = execute(
    "SELECT * FROM UserTable WHERE
    username= ' " & form("user") & " ' AND
    password= ' " & form("pwd") & " ' ");
```

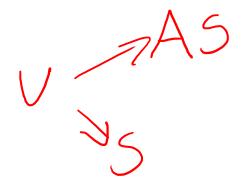


User supplies username and password, this SQL query checks if user/password combination is in the database

If not UserFound.EOF

Authentication correct
else Fail

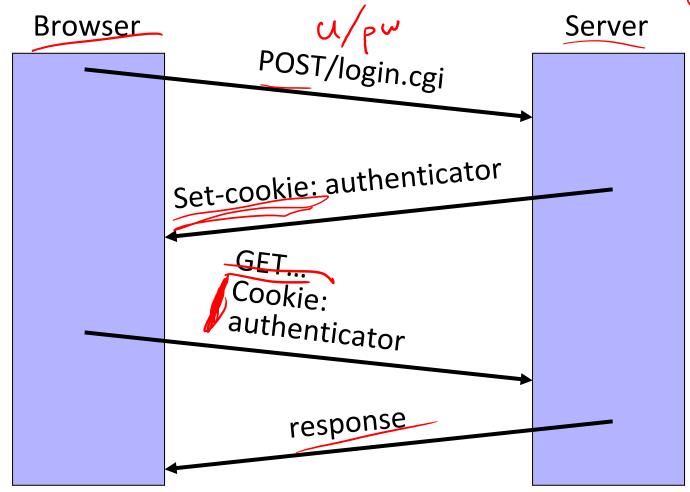
Only true if the result of SQL query is not empty, i.e., user/pwd is in the database



# Cross-Site Request Forgery (CSRF/XSRF)

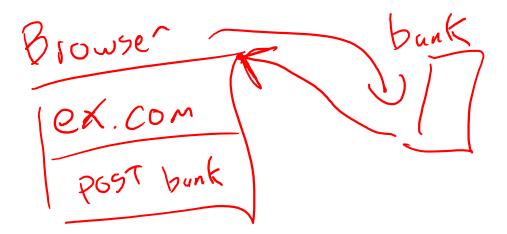
# Cookie-Based Authentication Redux





#### Browser Sandbox Redux

- Based on the same origin policy (SOP)
- Active content (scripts) can send anywhere!
  - For example, can submit a POST request
  - Some ports inaccessible -- e.g., SMTP (email)
- Can only read response from the same origin
  - ... but you can do a lot with just sending!



# 508

# Cross-Site Request Forgery

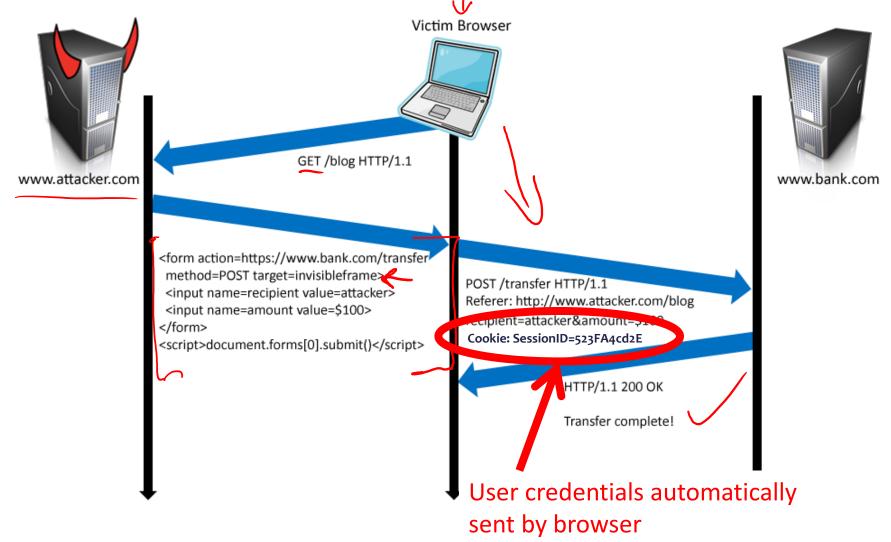
- Users logs into bank.com, forgets to sign off
  - Session cookie remains in browser state
- User then visits a malicious website containing

```
<form name=BillPayForm
action=http://bank.com/BillPay.php>
<input name=recipient value=badguy>...
<script> document.BillPayForm.submit(); </script>
```

- Browser sends cookie, payment request fulfilled!
- <u>Lesson</u>: cookie authentication is not sufficient when side effects can happen

Viewbalance

# Cookies in Forged Requests



### Impact

Hijack any ongoing session (if no protection)

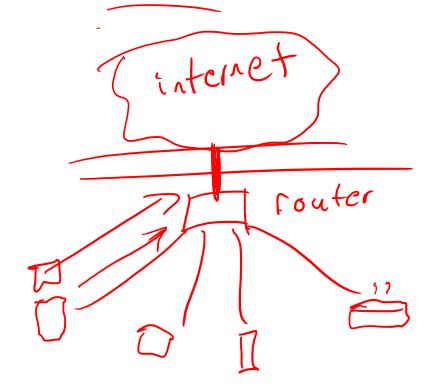
• Netflix: change account settings, Gmail: steal contacts, Amazon: one-click

purchase

Reprogram the user's home router

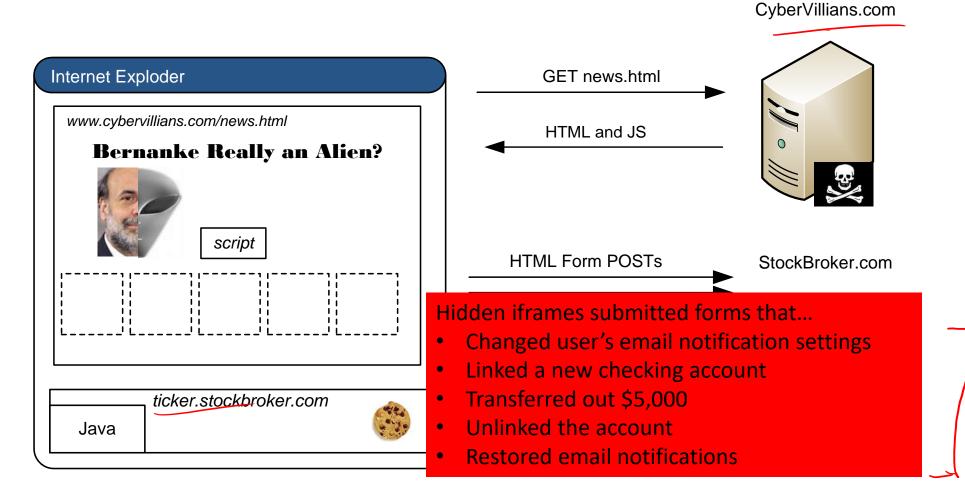
• Login to the *attacker's* account

• Why?

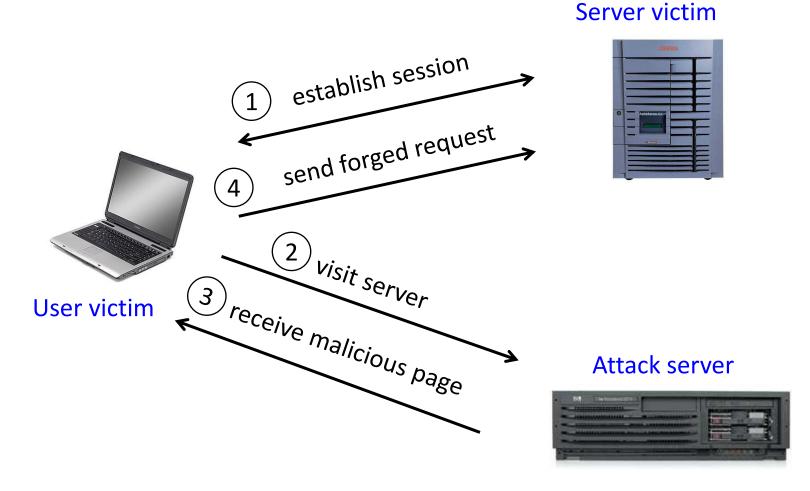


### XSRF True Story





# XSRF (aka CSRF): Summary



Q: how long do you stay logged on to Gmail? Financial sites?

#### Broader View of XSRF



- Abuse of cross-site data export
  - SOP does not control data export
  - Malicious webpage can initiates requests from the user's browser to an honest server
  - Server thinks requests are part of the established session between the browser and the server (automatically sends cookies)

#### XSRF Defenses

Secret validation token



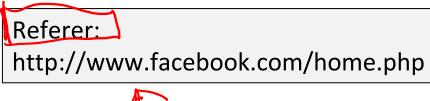


<input type=hidden value=23a3af01b>



Referer validation







#### Add Secret Token to Forms

<input type=hidden value=23a3af01b>

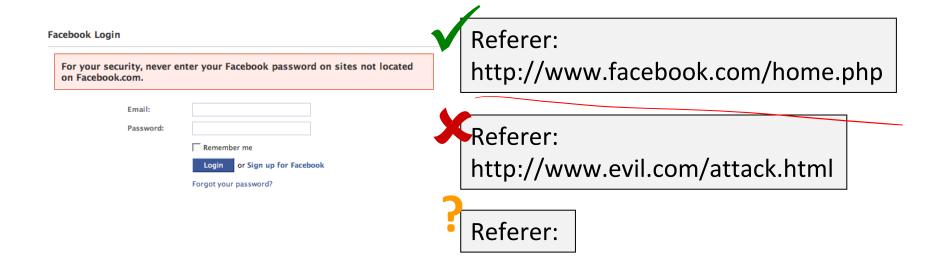
"Synchronizer Token Pattern"



- Include a secret challenge token as a hidden input in forms
  - Token often based on user's session ID
  - Server must verify correctness of token before executing sensitive operations
- Why does this work?
  - Same-origin policy: attacker can't read token out of legitimate forms loaded in user's browser, so can't create fake forms with correct token



#### Referer Validation



- Lenient referer checking header is optional
- Strict referer checking header is required

Why Not Always Strict Checking? NSAwiki/Project for Cellphones/

• Why might the referer header be suppressed?

Stripped by the organization's network filter

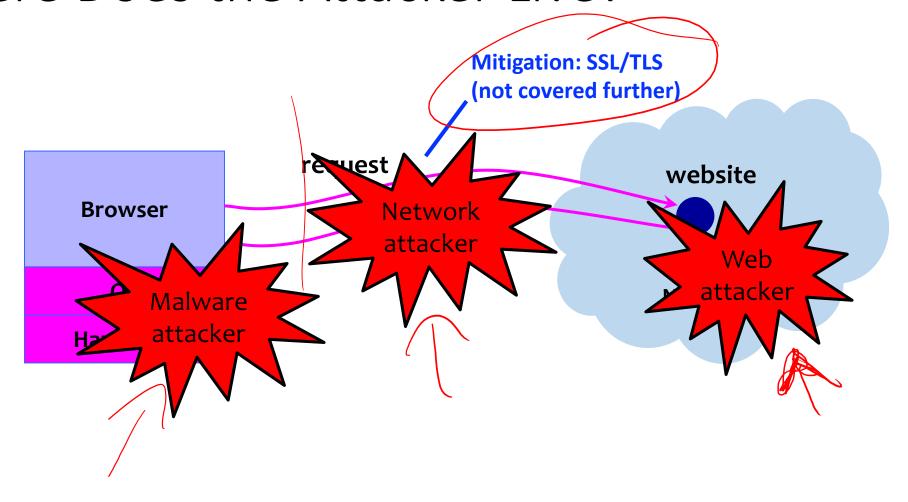
- Stripped by the local machine
- Stripped by the browser for HTTPS → HTTP transitions
- User preference in browser
- Buggy browser
- Web applications can't afford to block these users
- Many web application frameworks include CSRF defenses today

tokens!

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# Bonus topic: Consider the network

#### Where Does the Attacker Live?



#### Network attacker

Lives between you and your destination server

Person-in-the-middle

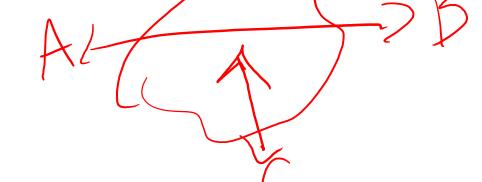






Passive/active

Physical/remote



TREVOR PAGLEN

185.jpg

NSA-Tapped Undersea Cables, North Pacific Ocean, 2016

# What might they be interested in?

• Eavesdropping  $\angle$ 

Making us talk to the wrong server

Denial-of-service

Corrupting our conversation with a real server

Domain Name System

Background: DNS

**HTTP Start!** 

Who is www.google.com?

172.217.14.228

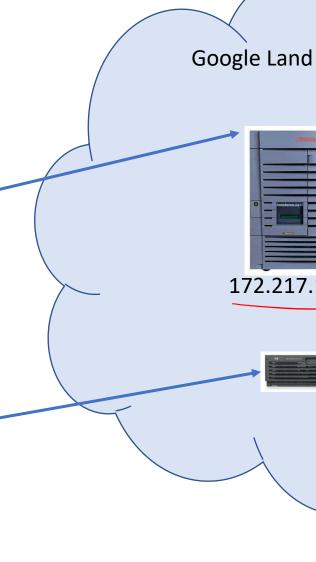
Who is

www.google.com?

Who is .com?







172.217.14.228

US TCP

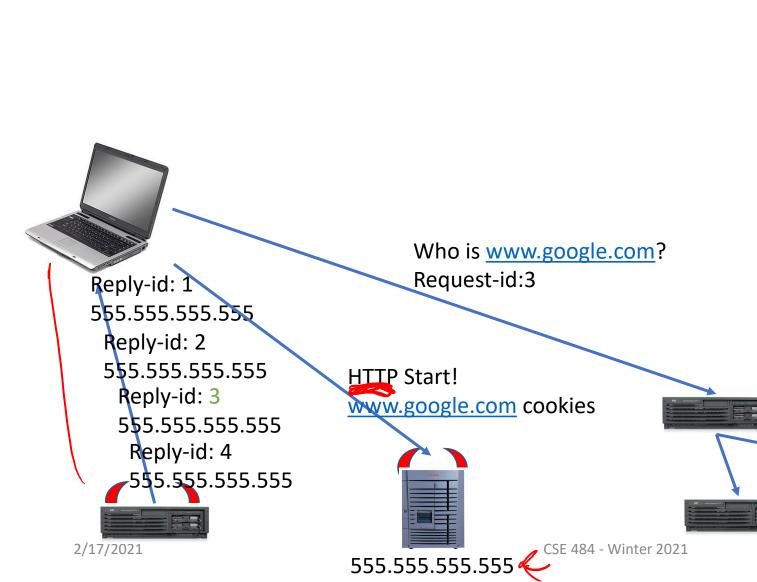
# DNS is unauthenticated and over UDP

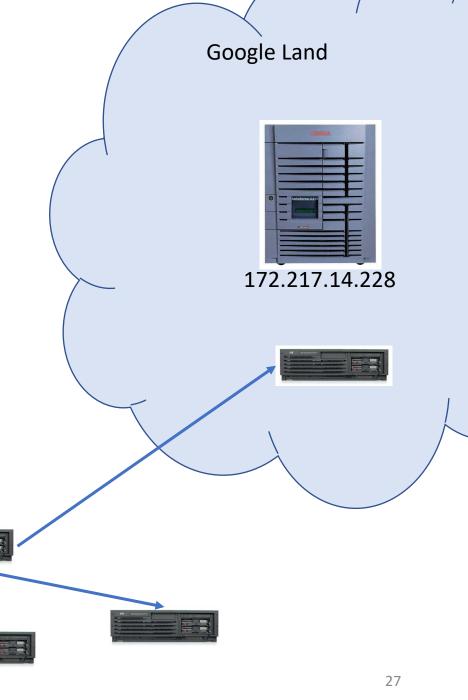
- 16-bit 'request ID'
  - Used to be sequential
  - Now random
- Reply is cleartext and 'simple'

Mail. google.com

reply id of one con 12 juny 1

# DNS Hijacking





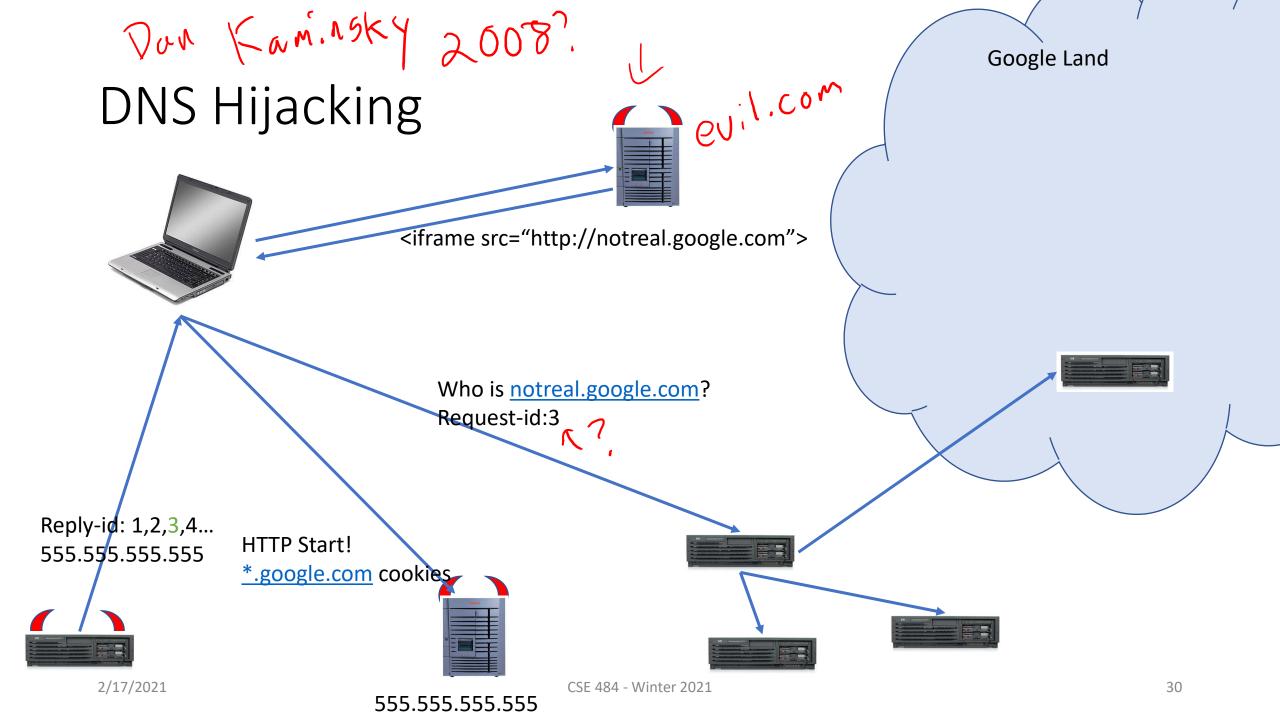
# Throwback: Birthday Paradox

- Are there two people in the first 1/8 of this class that have the same birthday?
  - 365 days in a year (366 some years)
    - Pick one person. To find another person with same birthday would take on the order of 365/2 = 182.5 people
    - Expect birthday "collision" with a room of only 23 people.
    - For simplicity, approximate when we expect a collision as sqrt(365).
- Why is this important for cryptography?
  - 2<sup>128</sup> different 128-bit values
    - Pick one value at random. To exhaustively search for this value requires trying on average 2<sup>127</sup> values.
      - Expect "collision" after selecting approximate y 2<sup>64</sup> random values.
      - 64 bits of security against collision attacks, not 128 bits.

# DNS Hijacking Continued

• 16-bit ID: 2^8 for collision (256!)

- How do we get the victim to as for <a href="www.google.com">www.google.com</a>?
  - How about "notreal.google.com" instead?



#### The state of DNS

- Randomize:
  - Request ID \
  - Port number \\( \( \)

• ... hope!



# Network security

• All our protocols weren't built for security 😊

- DNS
- BGP
- DHCP
- · ... ARP

DNSSE