

CSE 484 / CSE M 584: Computer Security and Privacy

Finish Cryptography; Start Web Security

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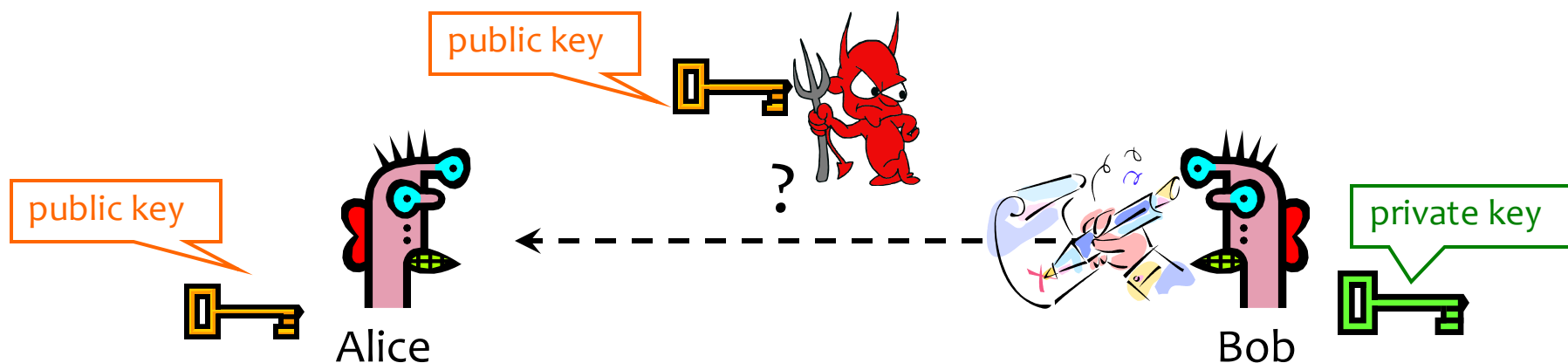
Thanks to Dan Boneh, Dieter Gollmann, Dan Halperin, Yoshi Kohno, Ada Lerner, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

Admin

- Lab 1 due today (10/30)
- Homework 2 due in a week (11/6)
- Final Project checkpoint #1 due in 2 weeks (11/13)

Confidentiality
↳ symmetric (AES)
↳ asymmetric (RSA)
Integrity / Authenticity
↳ hash functions / MAC
(symmetric)

Digital Signatures: Basic Idea



Given: Everybody knows Bob's **public key**
Only Bob knows the corresponding **private key**

Goal: Bob sends a “digitally signed” message

1. To compute a signature, must know the private key
2. To verify a signature, only the public key is needed

RSA Signatures

- Public key is (\underline{n}, e) , private key is (\underline{n}, d)
- To **sign** message m : $s = m^d \bmod n$
 - Signing & decryption are same **underlying** operation in RSA
 - It's infeasible to compute s on m if you don't know d
- To **verify** signature s on message m :
verify that $s^e \bmod n = (m^d)^e \bmod n = m$
 - Just like encryption (for RSA primitive)
 - Anyone who knows n and e (public key) can verify signatures produced with d (private key)
- In practice, also need **padding & hashing**
 - Standard padding/hashing schemes exist for RSA signatures

$$n = p \cdot q$$

$$[msg, sig(msg)]$$

DSS Signatures

- Digital Signature Standard (DSS)
 - U.S. government standard (1991, most recent rev. 2013)
- Public key: $(p, q, g, y=g^x \bmod p)$, private key: x
- Security of DSS requires hardness of discrete log
 - If could solve discrete logarithm problem, would extract x (private key) from $g^x \bmod p$ (public key)
- Again: We've discussed discrete logs modulo integers; significant advantages to using elliptic curve groups instead.

Cryptography Summary

- Goal: Privacy

- Symmetric keys:

- One-time pad, Stream ciphers
 - Block ciphers (e.g., DES, AES) → modes: EBC, CBC, CTR

- Public key crypto (e.g., Diffie-Hellman, RSA)

Key exchange
to get shared key

- Goal: Integrity

- MACs, often using hash functions (e.g, SHA-256)

- Goal: Privacy and Integrity

- Encrypt-then-MAC (not Encrypt and - MAC)

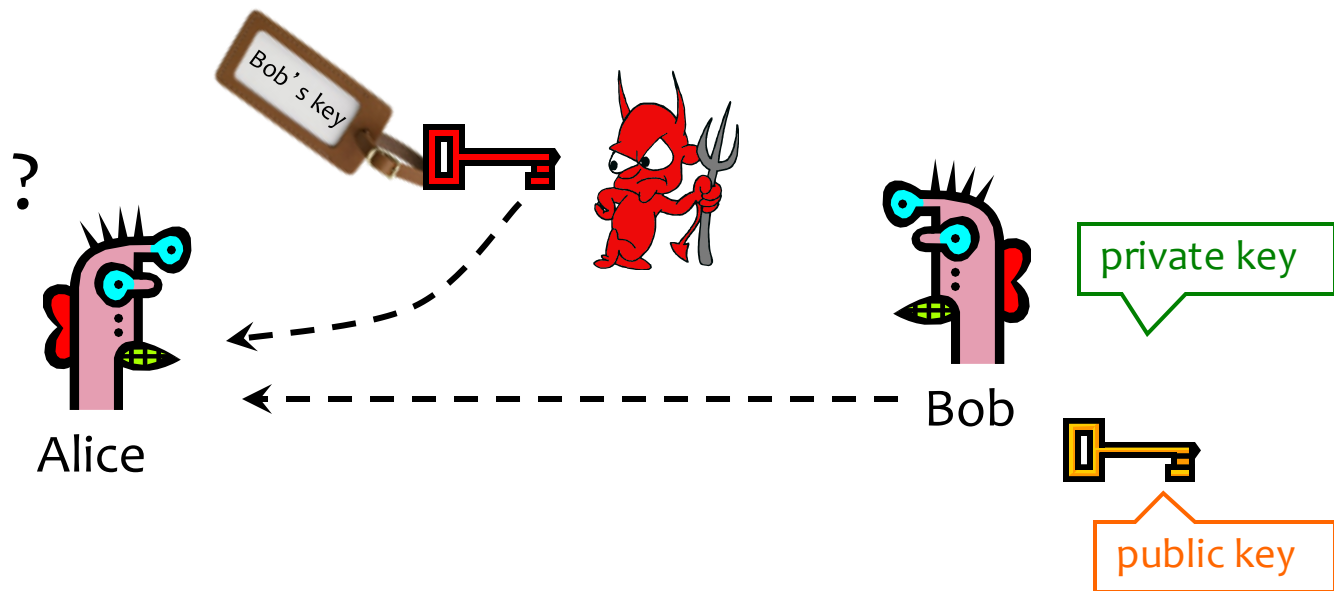
- Goal: Authenticity (and Integrity)

- Digital signatures (e.g., RSA, DSS)

Want More Crypto?

- Some suggestions:
 - CSE 490C (Rachel Lin):
<https://courses.cs.washington.edu/courses/cse490c/20au/>
 - Stanford Coursera (Dan Boneh):
<https://www.coursera.org/learn/crypto>

Authenticity of Public Keys



Problem: How does Alice know that the public key she received is really Bob's public key?

$\sim 10^{15} \text{ m}^3$

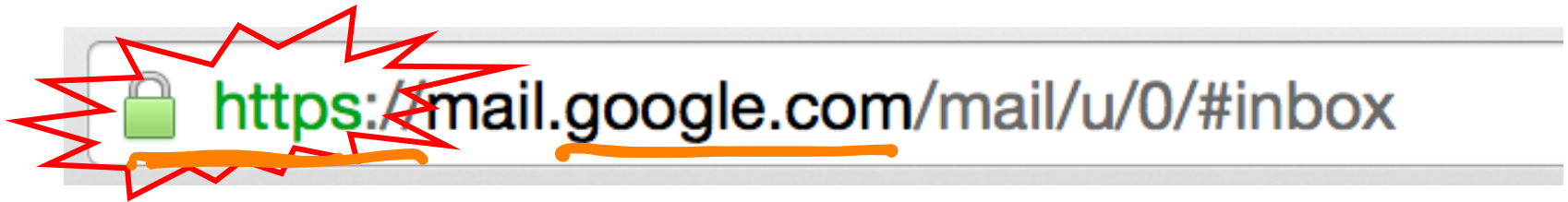


Distribution of Public Keys

- Public announcement or public directory
 - Risks: forgery and tampering
- Public-key certificate *CA = certificate authority*
 - Signed statement specifying the key and identity
 - $\text{sig}_{\text{CA}}(\text{"Bob"}, \text{PK}_B)$
- Common approach: certificate authority (CA) *– verisign*
 - Single agency responsible for certifying public keys
 - After generating a private/public key pair, user proves his identity and knowledge of the private key to obtain CA's certificate for the public key (offline)
 - Every computer is pre-configured with CA's public key

After School's
key

You encounter this every day...



SSL/TLS: Encryption & authentication for connections

SSL/TLS High Level

- SSL/TLS consists of **two** protocols
 - Familiar pattern for key exchange protocols
- Handshake protocol
 - Use **public-key cryptography** to establish a shared secret key between the client and the server
- Record protocol
 - Use the **secret symmetric key** established in the handshake protocol to protect communication between the client and the server

Example of a Certificate

GeoTrust Global CA
Google Internet Authority G2
*.google.com

***.google.com**
Issued by: Google Internet Authority G2
Expires: Monday, July 6, 2015 at 5:00:00 PM Pacific Daylight Time
✓ This certificate is valid

▼ **Details**

Subject Name	
Country	US
State/Province	California
Locality	Mountain View
Organization	Google Inc
Common Name	*.google.com
Issuer Name	
Country	US
Organization	Google Inc
Common Name	Google Internet Authority G2
Serial Number	6082711391012222858
Version	3

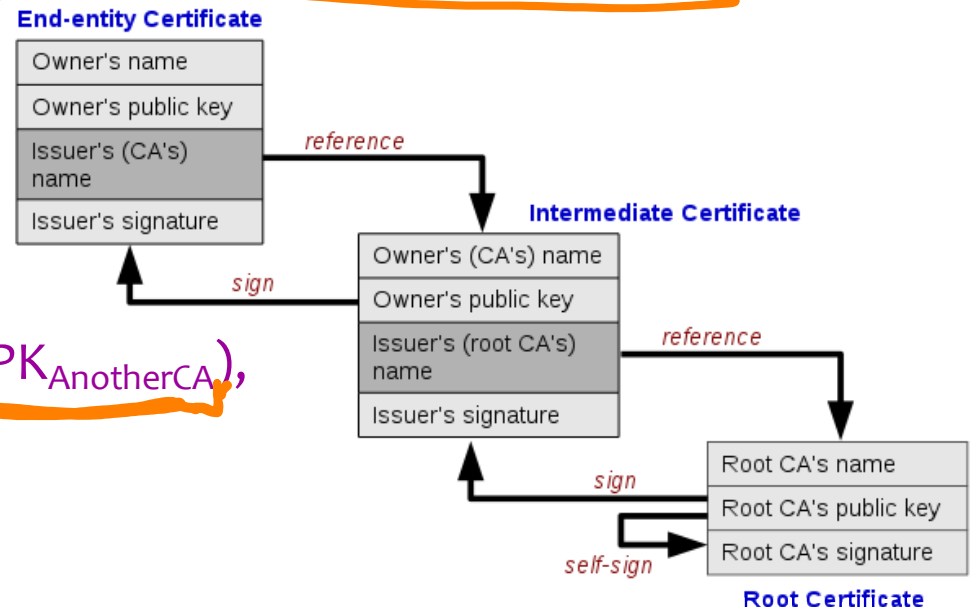
Signature Algorithm	SHA-1 with RSA Encryption (1.2.840.113549.1.1.5)
Parameters	none
Not Valid Before	Wednesday, April 8, 2015 at 6:40:10 AM Pacific Daylight Time
Not Valid After	Monday, July 6, 2015 at 5:00:00 PM Pacific Daylight Time
Public Key Info	
Algorithm	Elliptic Curve Public Key (1.2.840.10045.2.1)
Parameters	Elliptic Curve secp256r1 (1.2.840.10045.3.1.7)
Public Key	65 bytes : 04 CB DD C1 CE AC D6 20 ...
Key Size	256 bits •
Key Usage	Encrypt, Verify, Derive
Signature	256 bytes : 34 8B 7D 64 5A 64 08 5B ...

Hierarchical Approach

- Single CA certifying every public key is impractical
- Instead, use a trusted **root authority** (e.g., Verisign)

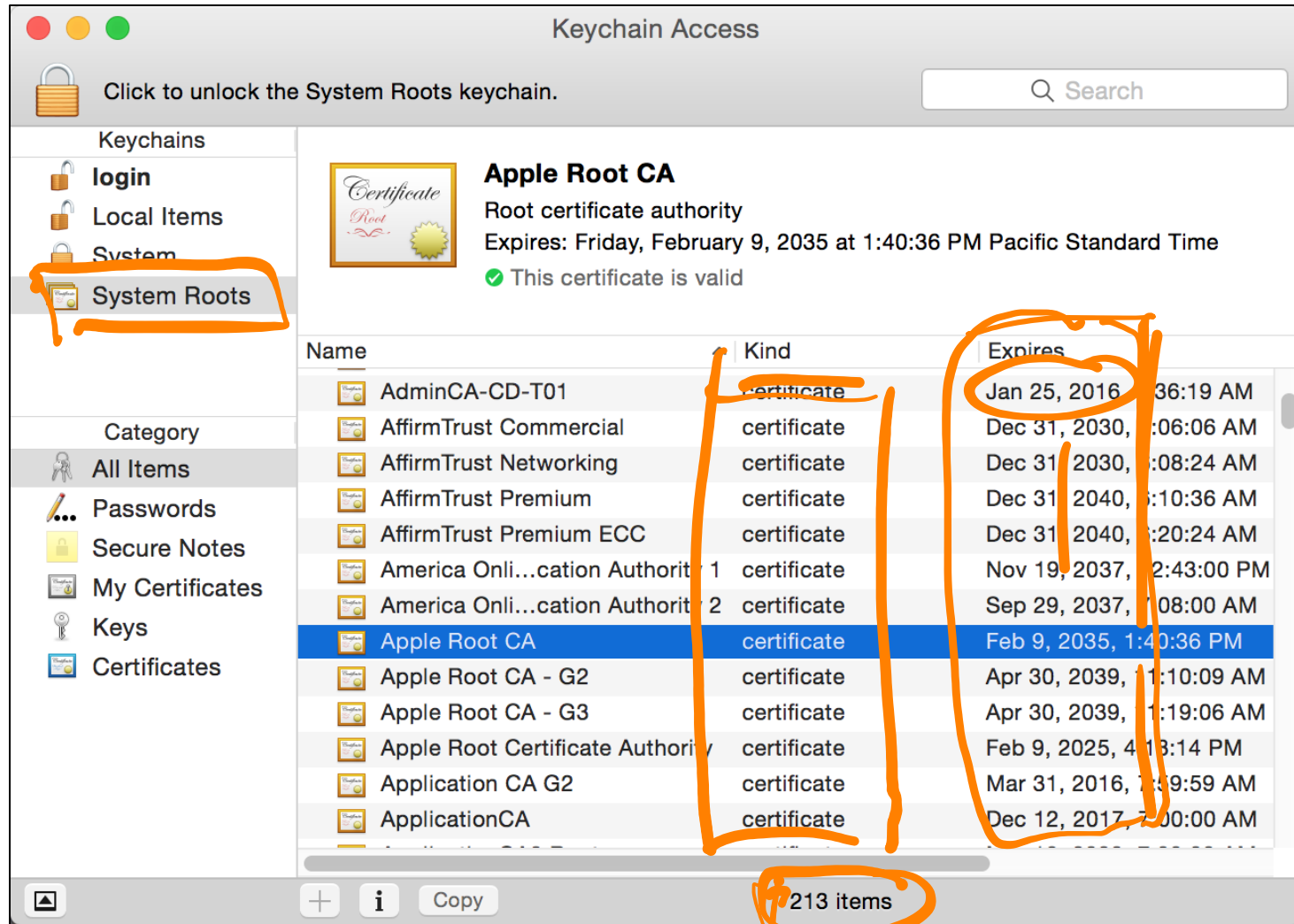
- Everybody must know the root's public key
- Instead of single cert, use a **certificate chain**

• $\text{sig}_{\text{Verisign}}(\text{"AnotherCA"}, \text{PK}_{\text{AnotherCA}}),$
 $\text{sig}_{\text{AnotherCA}}(\text{"Alice"}, \text{PK}_A)$



- What happens if root authority is ever compromised?

Trusted(?) Certificate Authorities



Turtles All The Way Down...



The saying holds that the world is supported by a chain of increasingly large turtles. Beneath each turtle is yet another: it is "turtles all the way down".

[Image from Wikipedia]

Many Challenges...

- Hash collisions
- Weak security at CAs
 - Allows attackers to issue rogue certificates
- Users don't notice when attacks happen
 - We'll talk more about this later in the course
- How do you revoke certificates? .)

DigiNotar is a Dutch Certificate Authority. They sell SSL certificates.



Attacking CAs

Security of DigiNotar servers:

- All core certificate servers controlled by a single admin password (Prod@dm1n)
- Software on public-facing servers out of date, unpatched
- No anti-virus (could have detected attack)

Somehow, somebody managed to get a rogue SSL certificate from them on **July 10th, 2011**. This certificate was issued for domain name .google.com.


What can you do with such a certificate? Well, you can impersonate Google — assuming you can first reroute Internet traffic for google.com to you. This is something that can be done by a government or by a rogue ISP. Such a reroute would only affect users within that country or under that ISP.

Consequences

- Attacker needs to first divert users to an attacker-controlled site instead of Google, Yahoo, Skype, but then...
 - For example, use DNS to poison the mapping of mail.yahoo.com to an IP address
- ... “authenticate” as the real site
- ... decrypt all data sent by users
 - Email, phone conversations, Web browsing

Attempt to Fix CA Problems:

Certificate Transparency

- **Problem:** browsers will think nothing is wrong with a rogue certificate until revoked
- **Goal:** make it impossible for a CA to issue a bad certificate for a domain *without the owner of that domain knowing* 
 - (Then what?)
- **Approach:** auditable certificate logs

www.certificate-transparency.org

Attempt to Fix CA Problems:

Certificate Pinning

- **Trust on first access:** tells browser how to act on subsequent connections
- HPKP – HTTP Public Key Pinning
 - Use these keys!
 - HTTP response header field `Public-Key-Pins`
- HSTS – HTTP Strict Transport Security
 - Only access server via HTTPS
 - HTTP response header field `Strict-Transport-Security`

Next Major Topic!
Web+Browser Security