CSE 484 / CSE M 584: Computer Security and Privacy

Web Security
[Overview + Browser Security Model]

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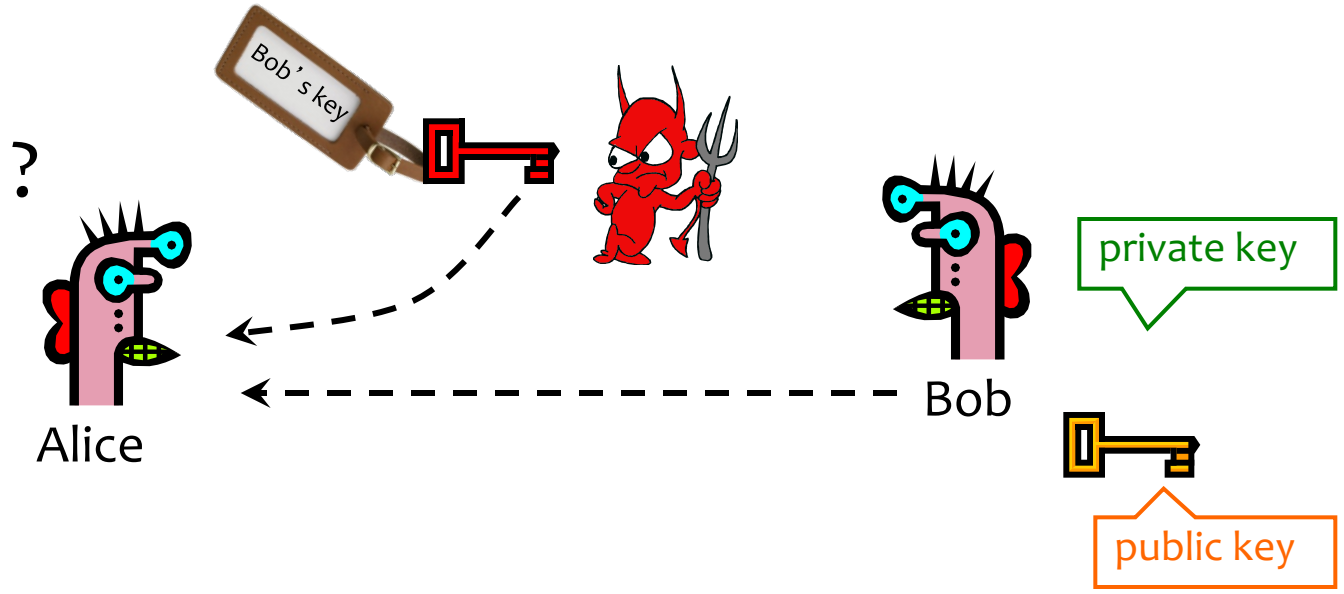
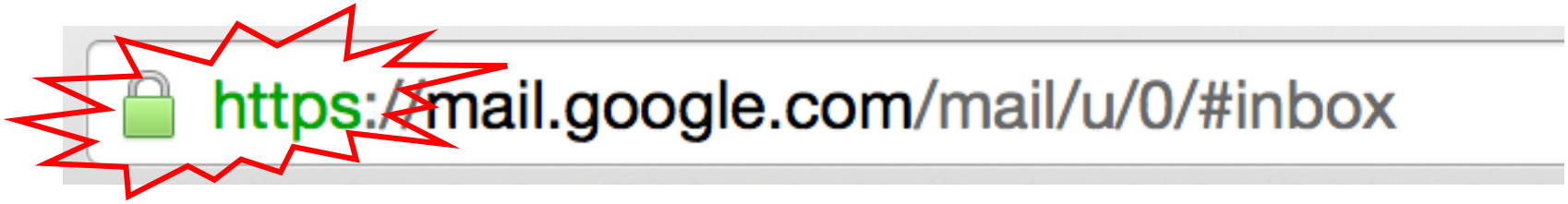
Recall: Authenticity of Public Keys



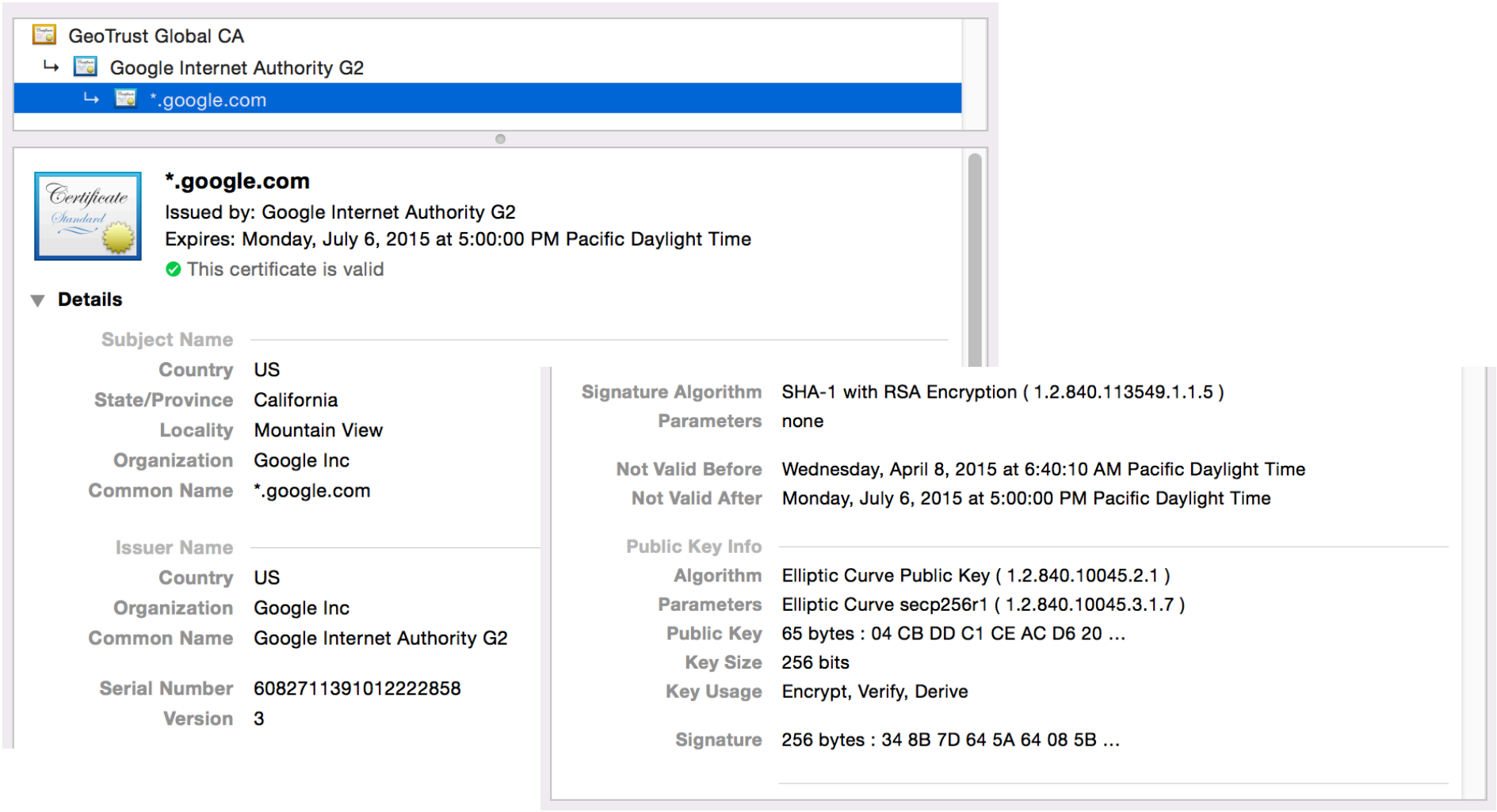
Problem: How does Alice know that the public key she received is really Bob's public key?

You encounter this every day…



SSL/TLS: Encryption & authentication for connections

Example of a Certificate



Many Challenges…

- Hash collisions
- Weak security at CAs
  - Allows attackers to issue rogue certificates
- Users don't notice when attacks happen
  - We'll talk more about this later in the course
- Etc…

Attacking CAs

DigiNotar Hacked by Black.Spook and Iranian Hackers Posted by Mikko @ 09:05 GMT | Comments

DigiNotar is a Dutch Certificate Authority. They sell SSL certificates.

Somehow, somebody managed to get a rogue SSL certificate from them on **July 10th, 2011**. This certificate was issued for domain name **.google.com**.

What can you do with such a certificate? Well, you can impersonate Google — assuming you can first reroute Internet traffic for google.com to you. This is something that can be done by a government or by a rogue ISP. Such a reroute would only affect users within that country or under that ISP.

Security of DigiNotar servers:

- All core certificate servers controlled by a single admin password (Prod@dm1n)
- Software on public-facing servers out of date, unpatched
- No anti-virus (could have detected attack)

Consequences

- Attacker needs to first divert users to an attacker-controlled site instead of Google, Yahoo, Skype, but then…
  - For example, use DNS to poison the mapping of mail.yahoo.com to an IP address
- … "authenticate" as the real site
- … decrypt all data sent by users
  - Email, phone conversations, Web browsing

Attempt to Fix CA Problems:

Certificate Transparency

- **Problem:** browsers will think nothing is wrong with a rogue certificate until revoked
- **Goal:** make it impossible for a CA to issue a bad certificate for a domain *without the owner of that domain knowing*
  - (Then what?)
- **Approach:** auditable certificate logs

www.certificate-transparency.org

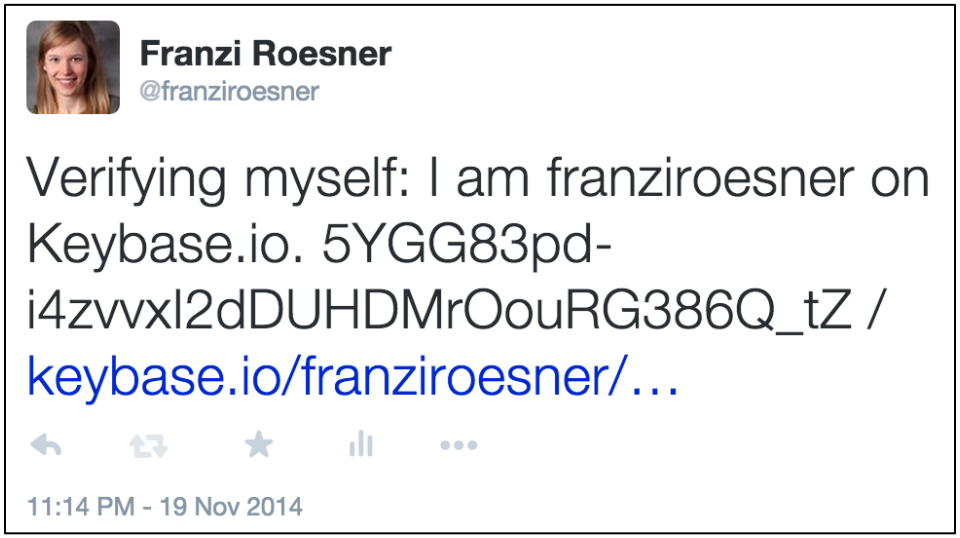
Attempt to Fix CA Problems:

Certificate Pinning

- Trust on first access: tells browser how to act on subsequent connections
- HPKP – HTTP Public Key Pinning
  - Use these keys!
  - HTTP response header field "`Public-Key-Pins`"
- HSTS – HTTP Strict Transport Security
  - Only access server via HTTPS
  - HTTP response header field "`Strict-Transport-Security`"

Keys for People: Keybase

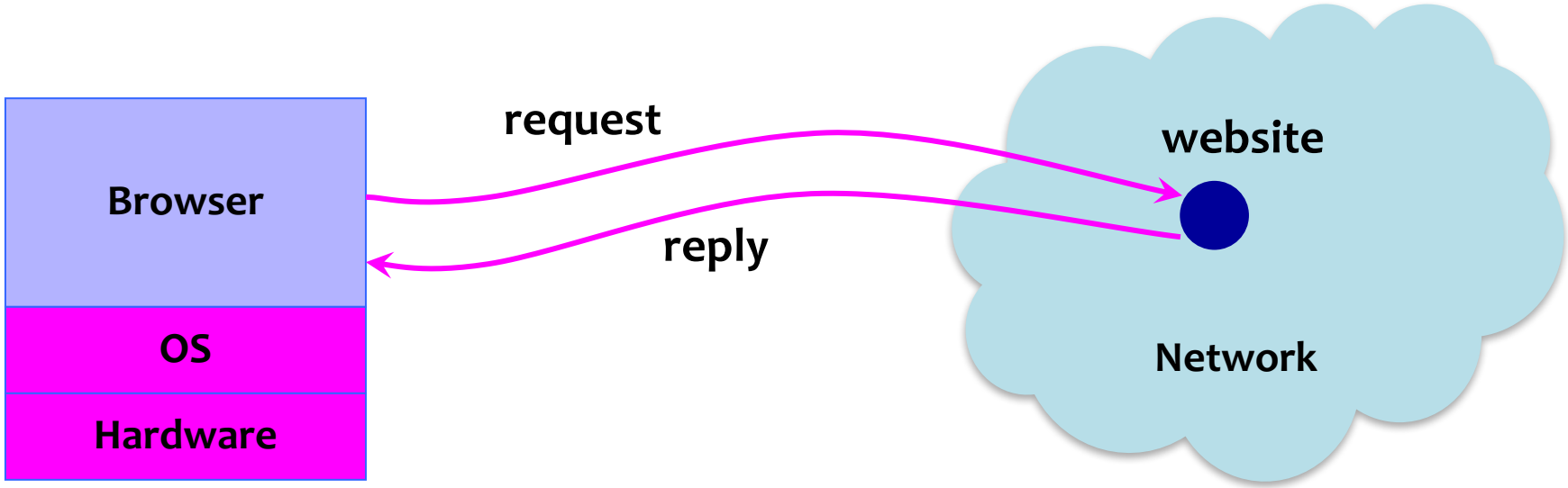
- Basic idea:
  - Rely on existing trust of a person's ownership of other accounts (e.g., Twitter, GitHub, website)
  - Each user publishes signed proofs to their linked account



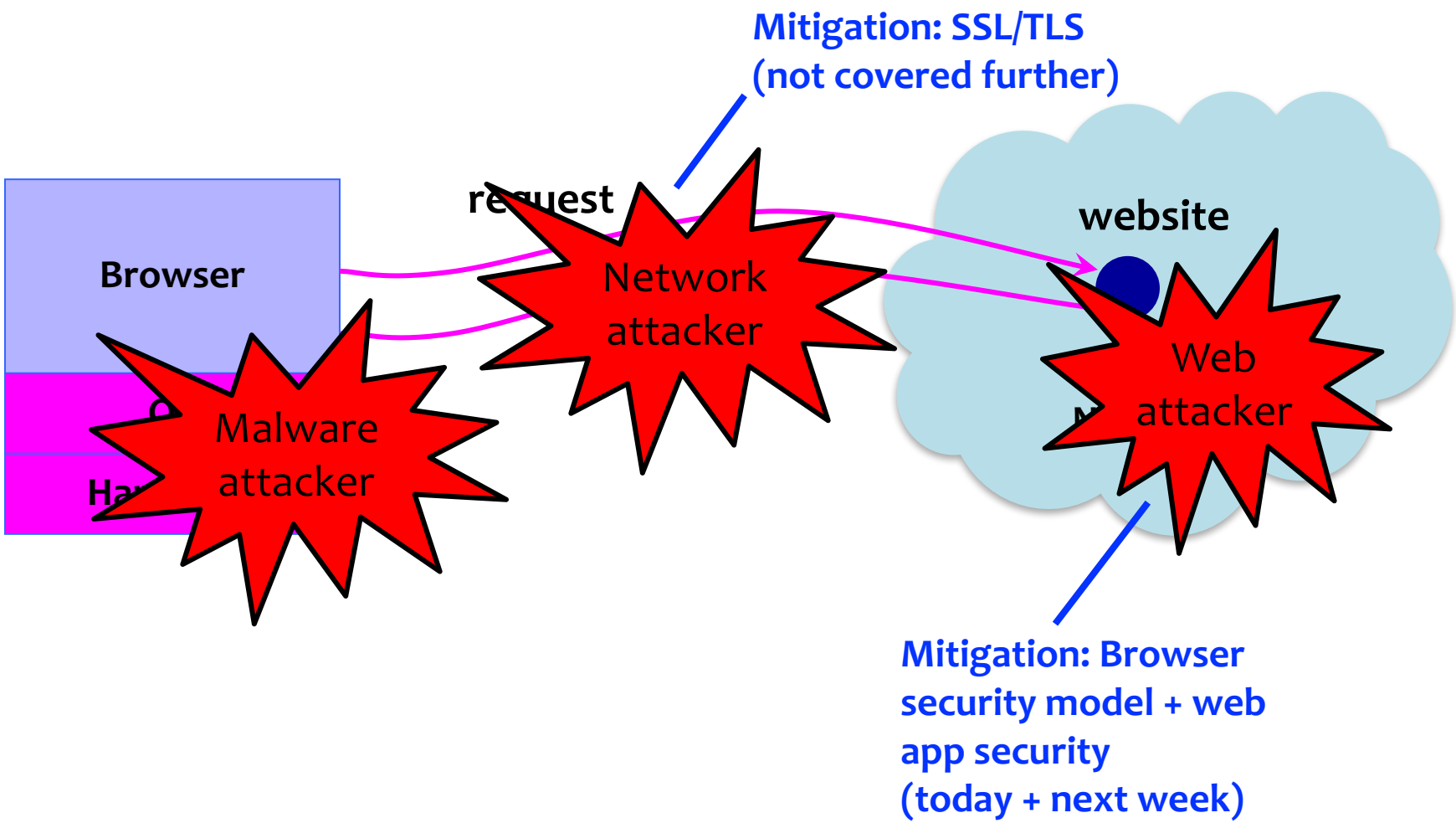
https://keybase.io/

Web+Browser Security

Big Picture: Browser and Network



Where Does the Attacker Live?



Web Attacker

- Controls a malicious website (**attacker.com**)
  - Can even obtain SSL/TLS certificate for site 

- User visits attacker.com – why?
  - Phishing email, enticing content, search results, placed by an ad network, blind luck …
- Attacker has no other access to user machine!
- Variation: good site **honest.com**, but:
  - An iframe with malicious content included
  - Website has been compromised

Two Sides of Web Security

(1) Web browser

- Responsible for securely confining content presented by visited websites

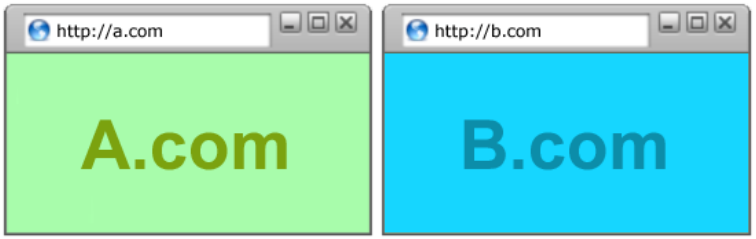
(2) Web applications

- Online merchants, banks, blogs, Google Apps …
- Mix of server-side and client-side code
  - Server-side code written in PHP, Ruby, ASP, JSP
  - Client-side code written in JavaScript
- Many potential bugs: XSS, XSRF, SQL injection

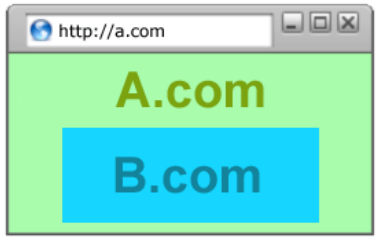
All of These Should Be Safe

- Safe to visit an evil website

- Safe to visit two pages at the same time



- Safe delegation



Browser Security Model

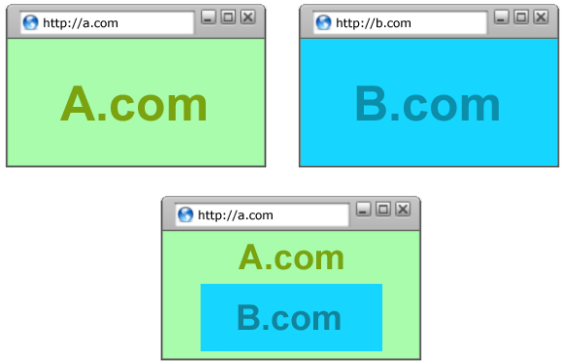
Goal 1: Protect local system from web attacker

→ Browser Sandbox

Goal 2: Protect/isolate web content from other web content

→ Same Origin Policy

(plus sandbox)



Browser Sandbox

Goals: Protect local system from web attacker; protect websites from each other

- E.g., safely execute JavaScript provided by a website
- No direct file access, limited access to OS, network, browser data, content from other websites
- Tabs **(new: also iframes!)** in their own processes
- Implementation is browser and OS specific*

*For example, see: https://chromium.googlesource.com/chromium/src/+/master/docs/design/sandbox.md

	High-quality report with functional exploit [1]
Sandbox Escape [5]	$15,000

From Chrome Bug Bounty Program

Same Origin Policy

Goal: Protect/isolate web content from other web content

Website origin = (scheme, domain, port)

Compared URL	Outcome	Reason
http://www.example.com/dir/page.html	Success	Same protocol and host
http://www.example.com/dir2/other.html	Success	Same protocol and host
http://www.example.com:**81**/dir/other.html	Failure	Same protocol and host but different port
https://www.example.com/dir/other.html	Failure	Different protocol
http://**en.example.com**/dir/other.html	Failure	Different host
http://**example.com**/dir/other.html	Failure	Different host (exact match required)
http://**v2.www.example.com**/dir/other.html	Failure	Different host (exact match required)

[Example from Wikipedia]

Same Origin Policy is Subtle!

- Some examples of how messy it gets in practice…
- Browsers don't (or didn't) always get it right...

- We'll talk about:
  - DOM / HTML Elements
  - Navigation
  - Cookie Reading
  - Cookie Writing
  - Iframes vs. Scripts