

CSE 484 / CSE M 584: Computer Security and Privacy

# Finish Software Security / Intro to Cryptography

Spring 2015

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Thanks to Dan Boneh, Dieter Gollmann, Dan Halperin, Yoshi Kohno, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

# Last Time: Software Security

- Principles:
  - Check inputs, Check all return values, Least privilege, Securely clear memory (passwords, keys, etc.), Failsafe defaults, Defense in depth, NOT security through obscurity, Reduce size of trusted computing base (TCB), Minimize attack surface, Use vetted components, Security by design
  - Simplicity, modularity
    - But: Be careful at interface boundaries!
  - Open design? Open source? Closed source?
    - Different perspectives

# Does Open Source Help?

- Different perspectives...
- Happy example:
  - Linux kernel backdoor attempt thwarted (2003)  
(<http://www.freedom-to-tinker.com/?p=472>)
- Sad example:
  - Heartbleed (2014)
    - Vulnerability in OpenSSL that allowed attackers to read arbitrary memory from vulnerable servers (including private keys)



# Vulnerability Analysis and Disclosure

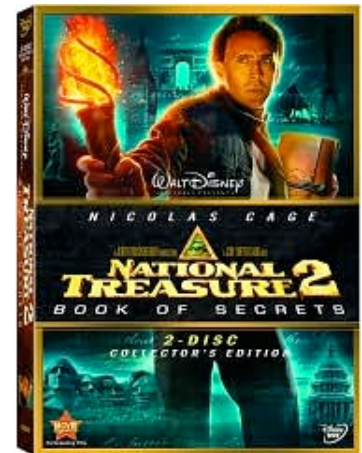
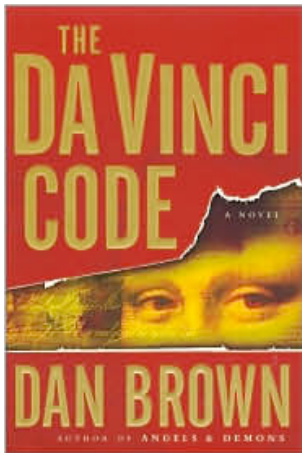
- What do you do if you've found a security problem in a real system?
- Say
  - A commercial website?
  - UW grade database?
  - Boeing 787?
  - TSA procedures?

**Abj sbe fbzr pelcgbtencul!**

**Now for some cryptography!**

# Cryptography and Security

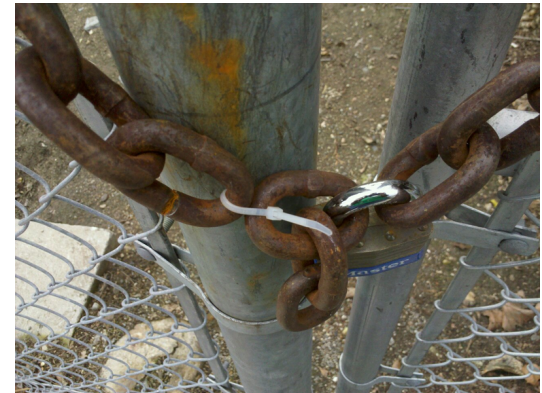
- Art and science of *protecting our information*.
  - Keeping it **private**, if we want privacy.
  - Protecting its **integrity**, if we want to avoid forgeries.



Images from Wikipedia and Barnes and Noble

# Some Thoughts About Cryptography

- Cryptography only one small piece of a larger system
- Must protect entire system
  - Physical security
  - Operating system security
  - Network security
  - Users
  - Cryptography (following slides)
- “Security only as strong as the weakest link”
  - Need to secure weak links
  - But not always clear what the weakest link is (different adversaries and resources, different adversarial goals)
  - Crypto failures may not be (immediately) detected
- Cryptography helps after you’ve identified your threat model and goals
  - Famous quote: “Those who think that cryptography can solve their problems doesn’t understand cryptography and doesn’t understand their problems.”



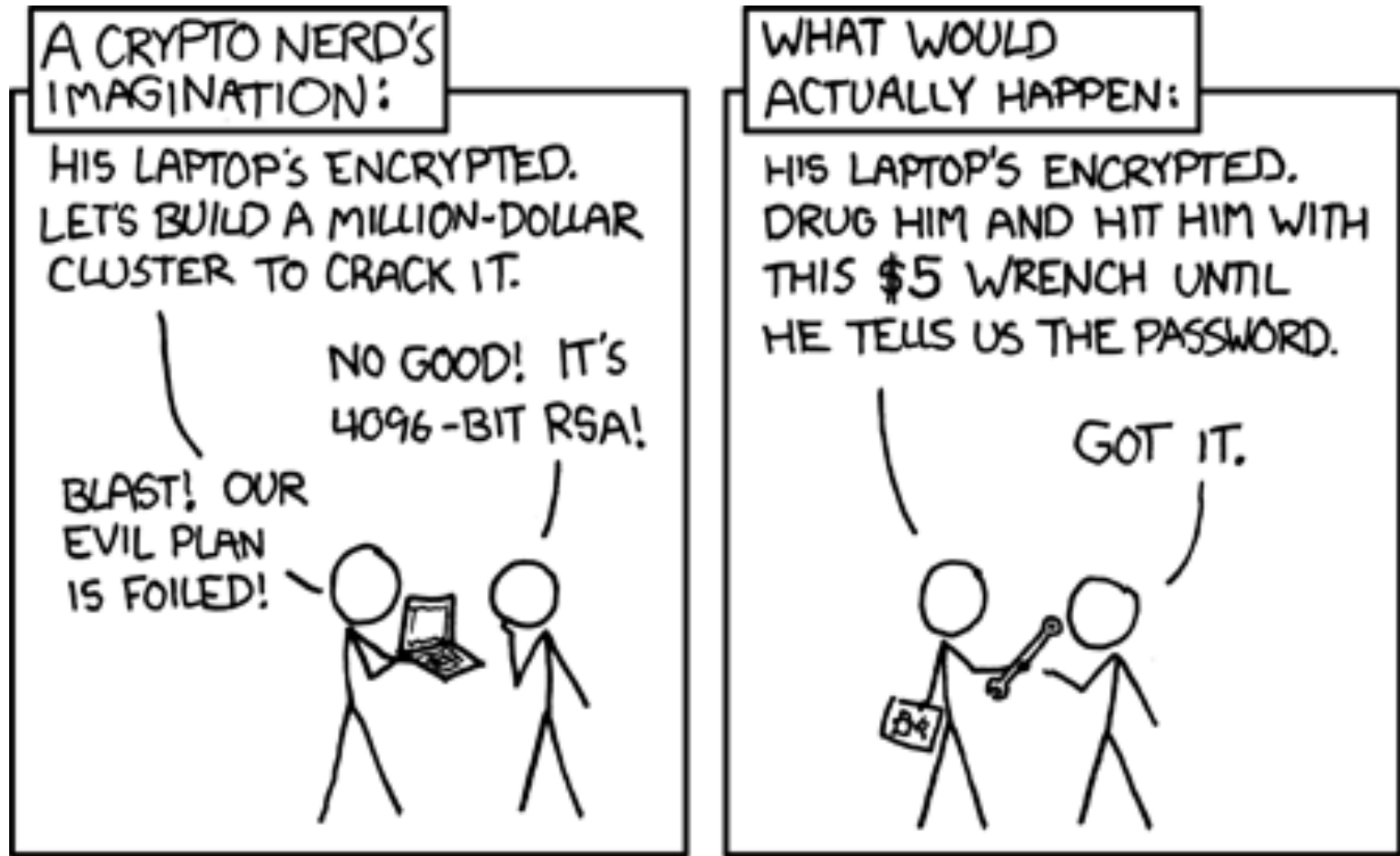
# Improved Security, Increased Risk

- RFIDs in car keys:
  - RFIDs in car keys make it harder to hotwire a car
  - Result: Car jackings increased



- RFIDs in car keys: Biometric car lock defeated by cutting off owner's finger
    - RFIDs in car keys
    - Result: Car jackin
- POSTED BY CORY DOCTOROW, MARCH 31, 2005 7:53 AM | [PERMALINK](#)
- Andrei sez, "'Malaysia car thieves steal finger.' This is what security visionaries Bruce Schneier and Ross Anderson have been warning about for a long time. Protect your \$75,000 Mercedes with biometrics and you risk losing whatever body part is required by the biometric mechanism."
- “ ...[H]aving stripped the car, the thieves became frustrated when they wanted to restart it. They found they again could not bypass the immobiliser, which needs the owner's fingerprint to disarm it.
- They stripped Mr Kumaran naked and left him by the side of the road - but not before cutting off the end of his index finger with a machete.

# XKCD: <http://xkcd.com/538/>



# Key Entry Pad (4-digit PIN)



Image from profmason.com

- This is the key pad on my office safe.
- Inside my safe is a copy of final exam.
- How long would it take you to break in?
- Answer (combinatorics):
  - $10^4$  tries *maximum*
  - $10^4 / 2$  tries *on average*
- Answer (unit conversion):
  - 3 seconds per try --> 4 hours and 10 minutes on average

# Key Entry Pad (4-digit PIN)



Image from profmason.com

- Now assume the safe automatically calls police after 3 failed attempts.
- What is the probability that you will guess the PIN within 3 tries? (Assume no repeat tries.)
- Answer (combinatorics)
  - 10000 choose 3 possible choices for the 3 guesses
  - $1 \cdot (9999 \text{ choose } 2)$  possible choices contain the correct PIN
  - So success probability is  $3 / 10000$

# Key Entry Pad (4-digit PIN)



Image from profmason.com

- Could you do better at guessing the PIN?
- Answer (*chemical* combinatorics):
  - Put different chemical on each key (NaCl, KCl, LiCl, ...)

Idea from <http://eprint.iacr.org/2003/217.ps>



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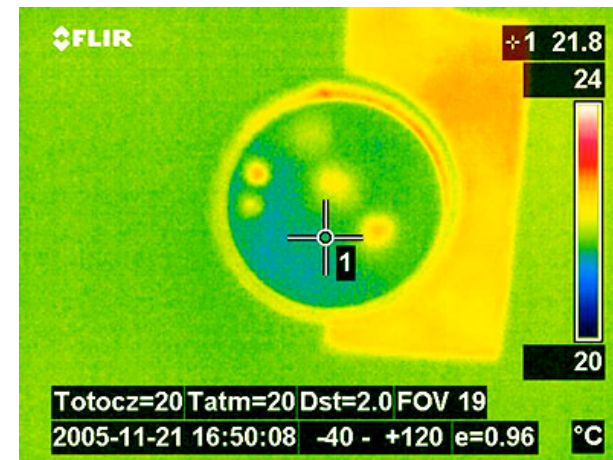
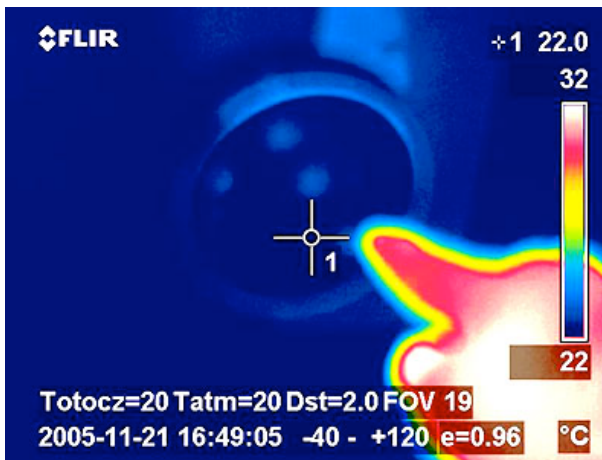
Image from profmason.com

- Could you do better at guessing the PIN?
- Answer (*chemical* combinatorics):
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- **Lesson:** Consider the complete system, physical security, etc.
- **Lesson:** Think outside the box

Idea from <http://eprint.iacr.org/2003/217.ps>



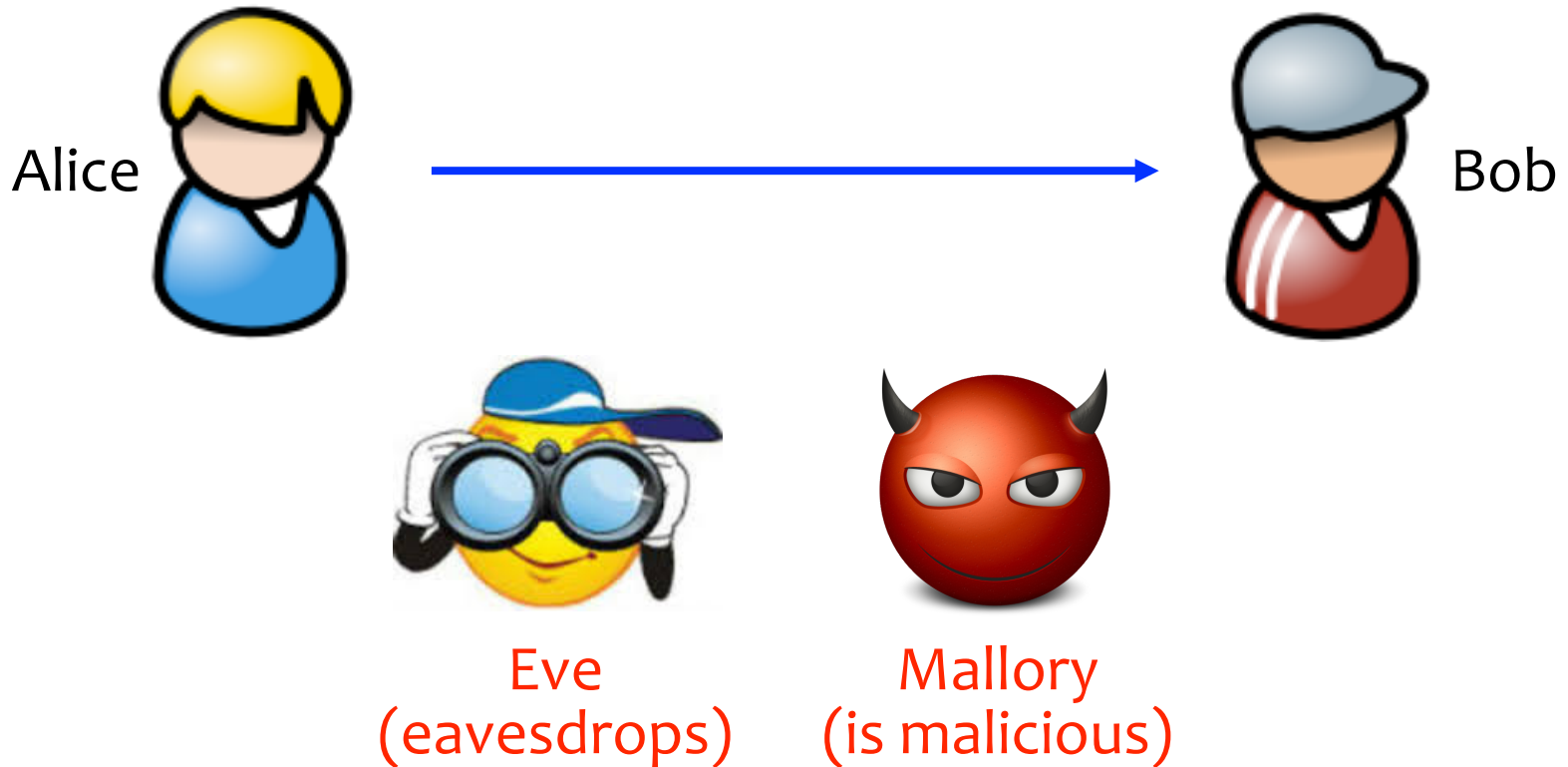
# Thermal Patterns



Images from <http://lcamtuf.coredump.cx/tsafe/>

# Alice and Bob

- Archetypal characters



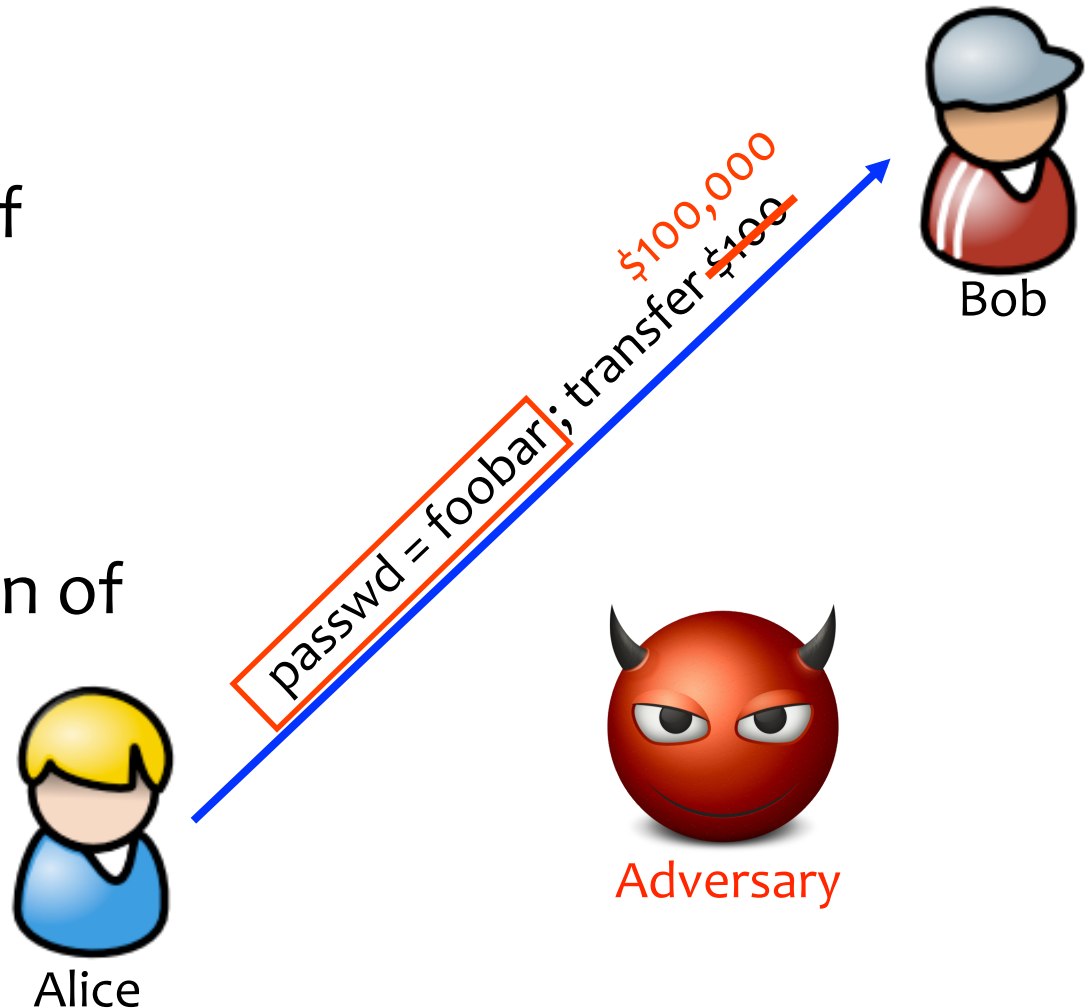
# Common Communication Security Goals

## Privacy of data:

Prevent exposure of information

## Integrity of data:

Prevent modification of information

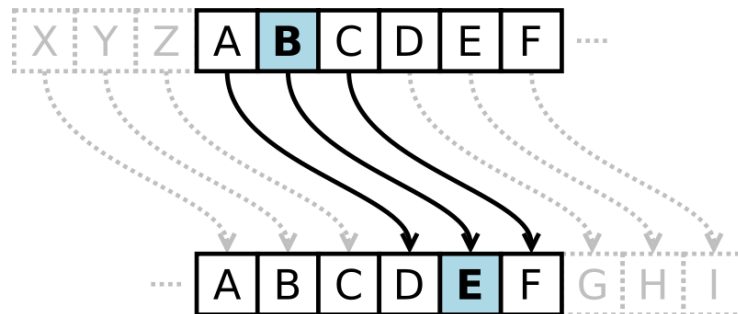


# History

- Substitution Ciphers
  - Caesar Cipher
- Transposition Ciphers
- Codebooks
- Machines
- Recommended Reading: **The Codebreakers** by David Kahn and **The Code Book** by Simon Singh.

# History: Caesar Cipher (Shift Cipher)

- Plaintext letters are replaced with letters a fixed shift away in the alphabet.
- Example:
  - Plaintext: The quick brown fox jumps over the lazy dog
  - Key: Shift 3
  - ABCDEFGHIJKLMNOPQRSTUVWXYZ  
DEFGHIJKLMNOPQRSTUVWXYZABC
  - Ciphertext: WKHTX LFNEU RZQIR AMXPS VRYHU WKHOD CBGRJ



# History: Caesar Cipher (Shift Cipher)

- ROT13: shift 13 (encryption and decryption are symmetric)
- What is the key space?
  - 26 possible shifts.
- How to attack shift ciphers?
  - Brute force.



# History: Substitution Cipher

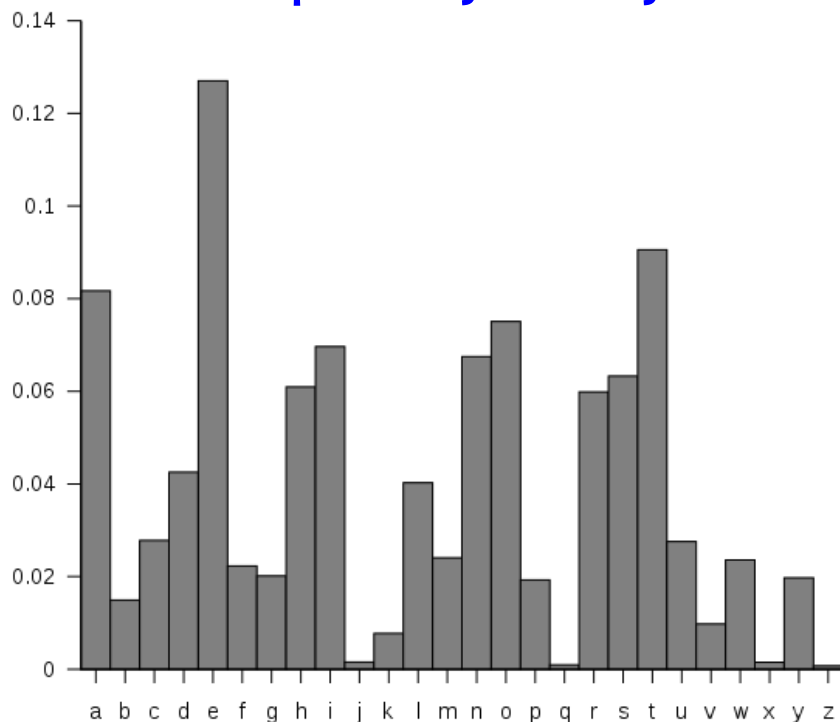
- Superset of shift ciphers: each letter is substituted for another one.
- Add a secret key
- Example:
  - Plaintext: ABCDEFGHIJKLMNOPQRSTUVWXYZ
  - Cipher: ZEBRAS CDEFGHIJKLMNOPQTUVWXY
- “State of the art” for thousands of years

# History: Substitution Cipher

- What is the key space?  $26! \approx 2^{88}$

- How to attack?

— Frequency analysis.



## Bigrams:

th 1.52%	en 0.55%	ng 0.18%
he 1.28%	ed 0.53%	of 0.16%
in 0.94%	to 0.52%	al 0.09%
er 0.94%	it 0.50%	de 0.09%
an 0.82%	ou 0.50%	se 0.08%
re 0.68%	ea 0.47%	le 0.08%
nd 0.63%	hi 0.46%	sa 0.06%
at 0.59%	is 0.46%	si 0.05%
on 0.57%	or 0.43%	ar 0.04%
nt 0.56%	ti 0.34%	ve 0.04%
ha 0.56%	as 0.33%	ra 0.04%
es 0.56%	te 0.27%	ld 0.02%
st 0.55%	et 0.19%	ur 0.02%

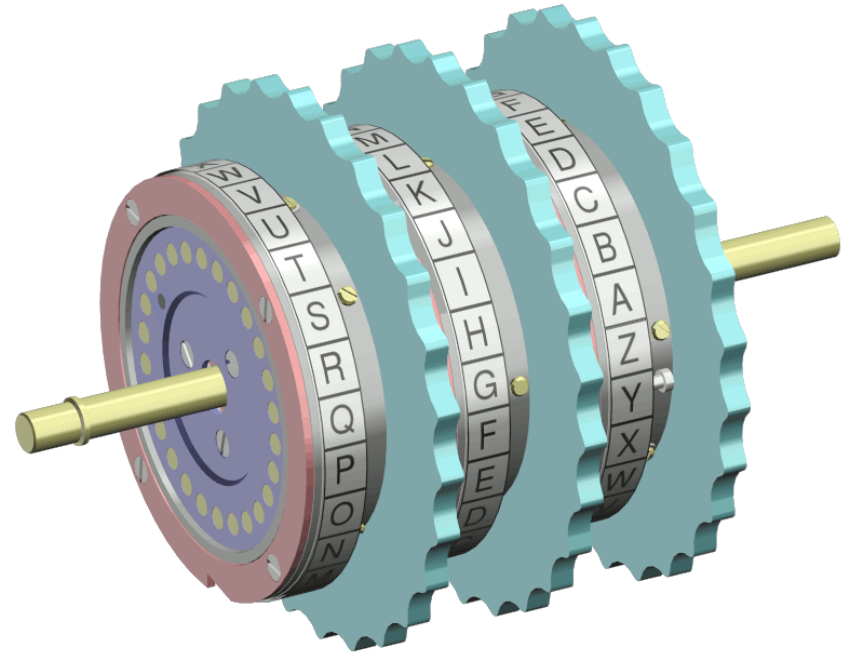
## Trigrams:

1. the	6. ion	11. nce
2. and	7. tio	12. edt
3. tha	8. for	13. tis
4. ent	9. nde	14. oft
5. ing	10. has	15. sth



# History: Enigma Machine

Uses rotors (substitution cipher) that change position after each key.



Key = initial setting of rotors

Key space?

$26^n$  for  $n$  rotors

# Kerckhoff's Principle

- Security of a cryptographic object should depend only on the secrecy of the secret (private) key.
- Security should not depend on the secrecy of the algorithm itself.

# How Cryptosystems Work Today

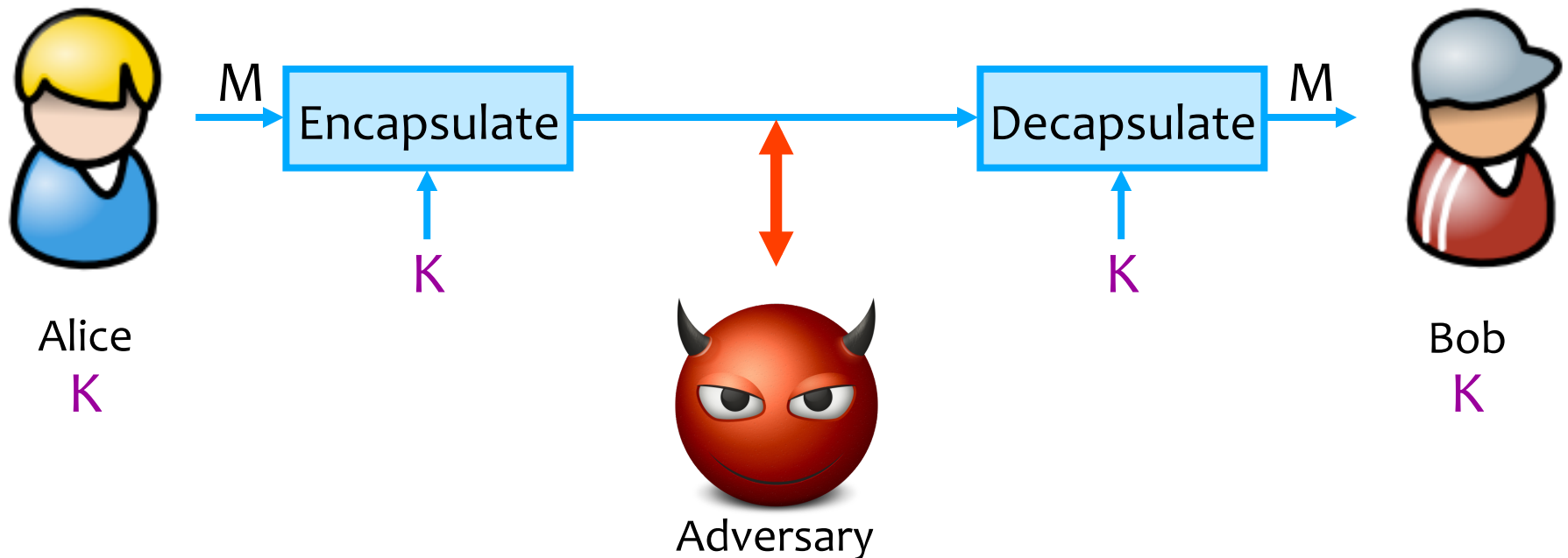
- Layered approach:
  - Cryptographic primitives, like block ciphers, stream ciphers, hash functions, and one-way trapdoor permutations
  - Cryptographic protocols, like CBC mode encryption, CTR mode encryption, HMAC message authentication
- Public algorithms (Kerckhoff's Principle)
- Security proofs based on assumptions (not this course)
- Don't roll your own!

# Flavors of Cryptography

- Symmetric cryptography
  - Both communicating parties have access to a shared random string  $K$ , called the key.
- Asymmetric cryptography
  - Each party creates a public key  $pk$  and a secret key  $sk$ .

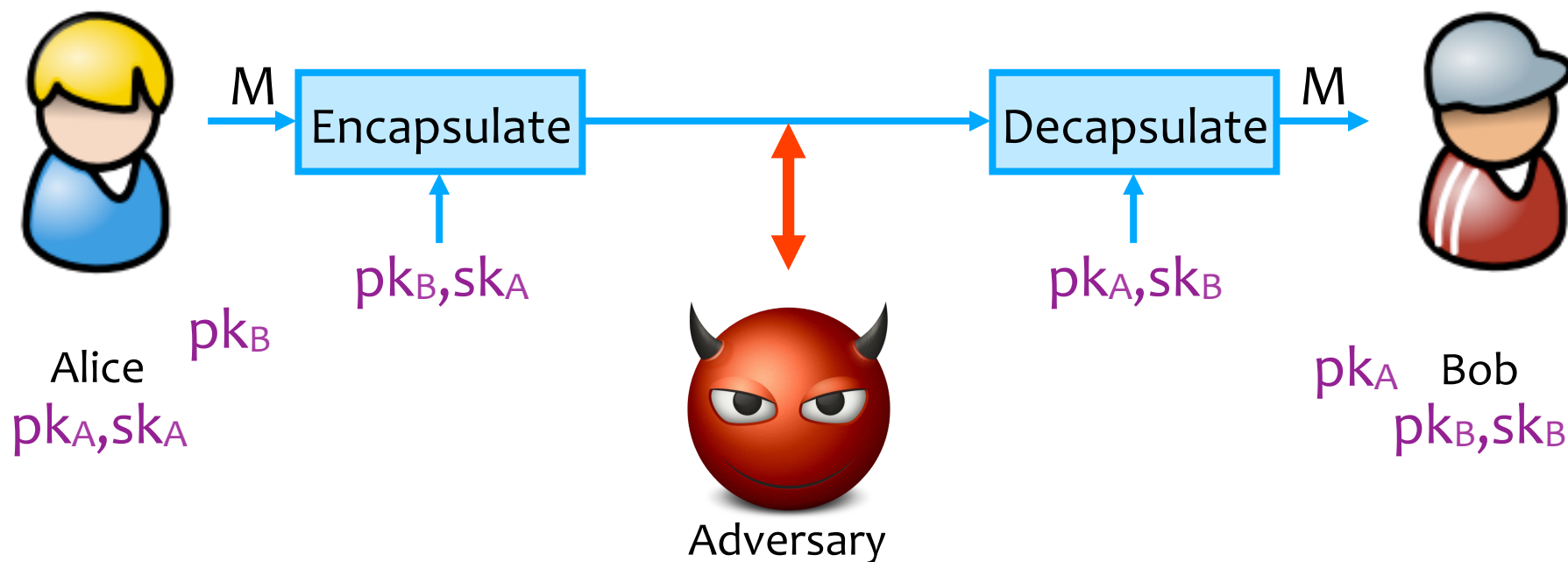
# Symmetric Setting

Both communicating parties have access to a shared random string  $K$ , called the *key*.



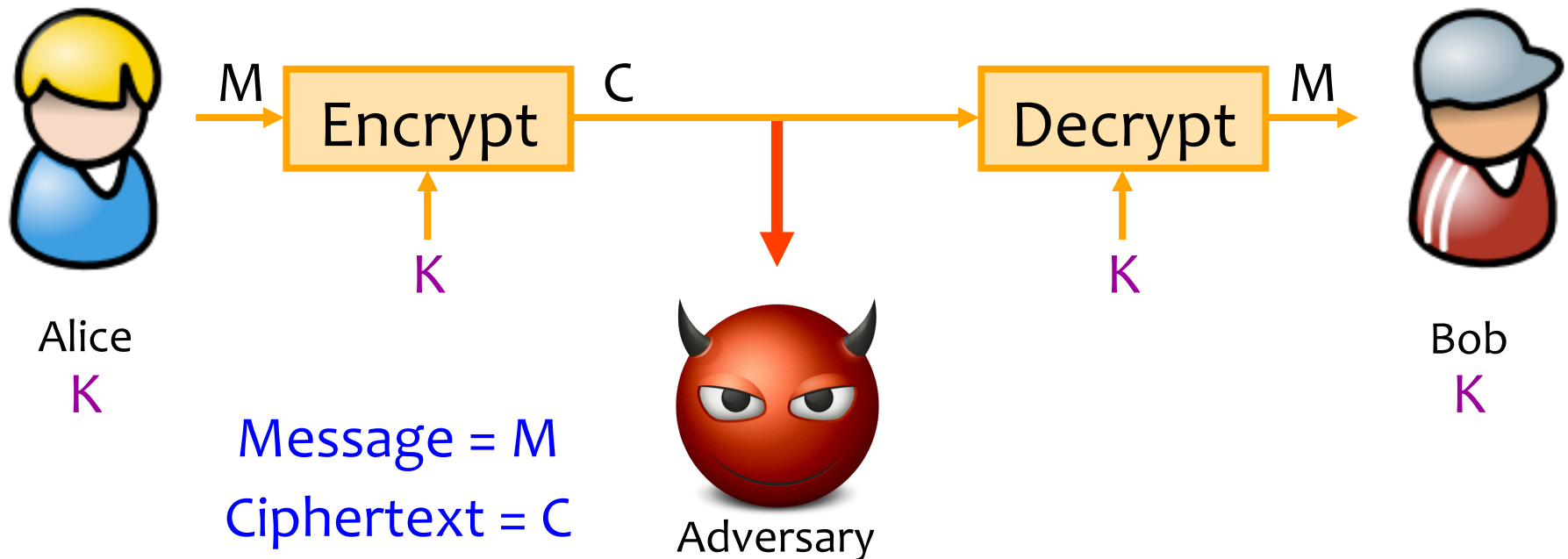
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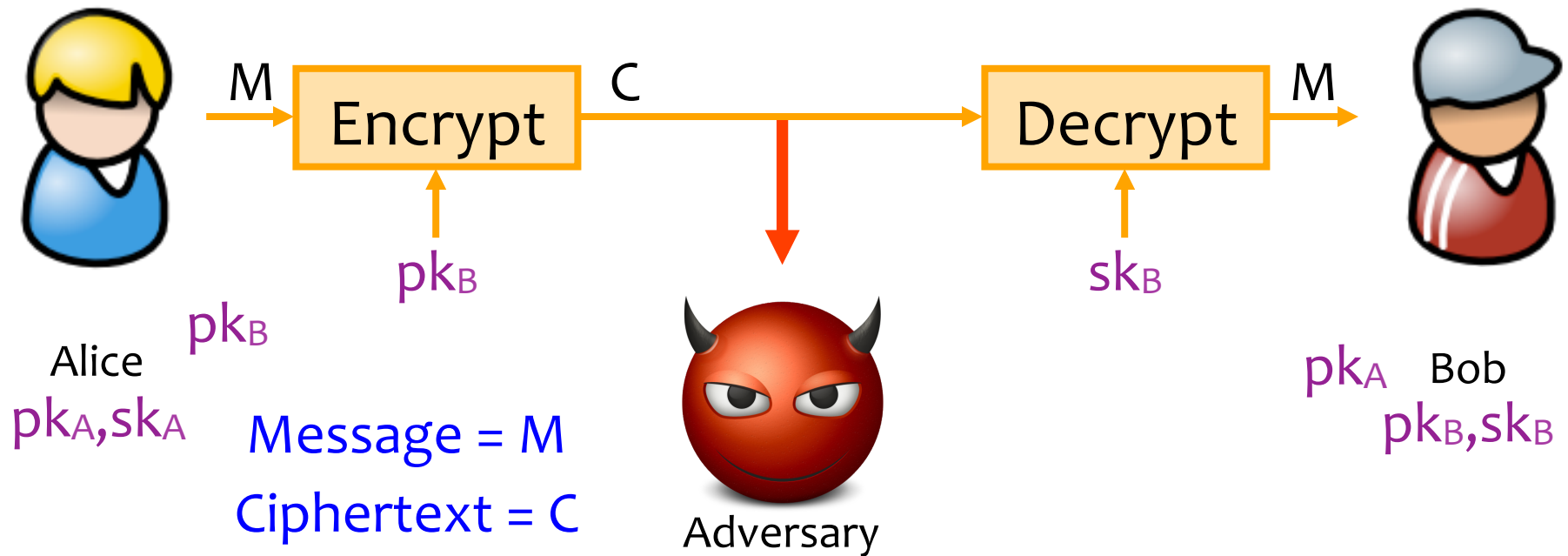
# Achieving Privacy (Symmetric)

Encryption schemes: A tool for protecting **privacy**.



# Achieving Privacy (Asymmetric)

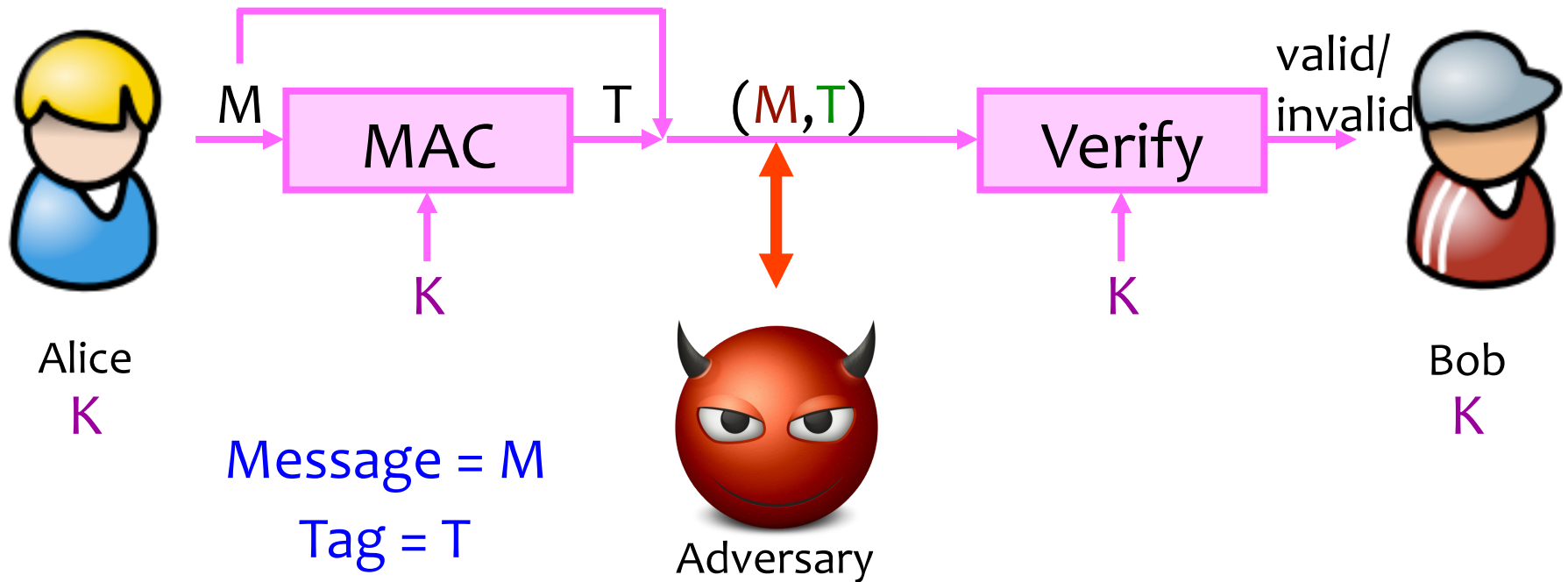
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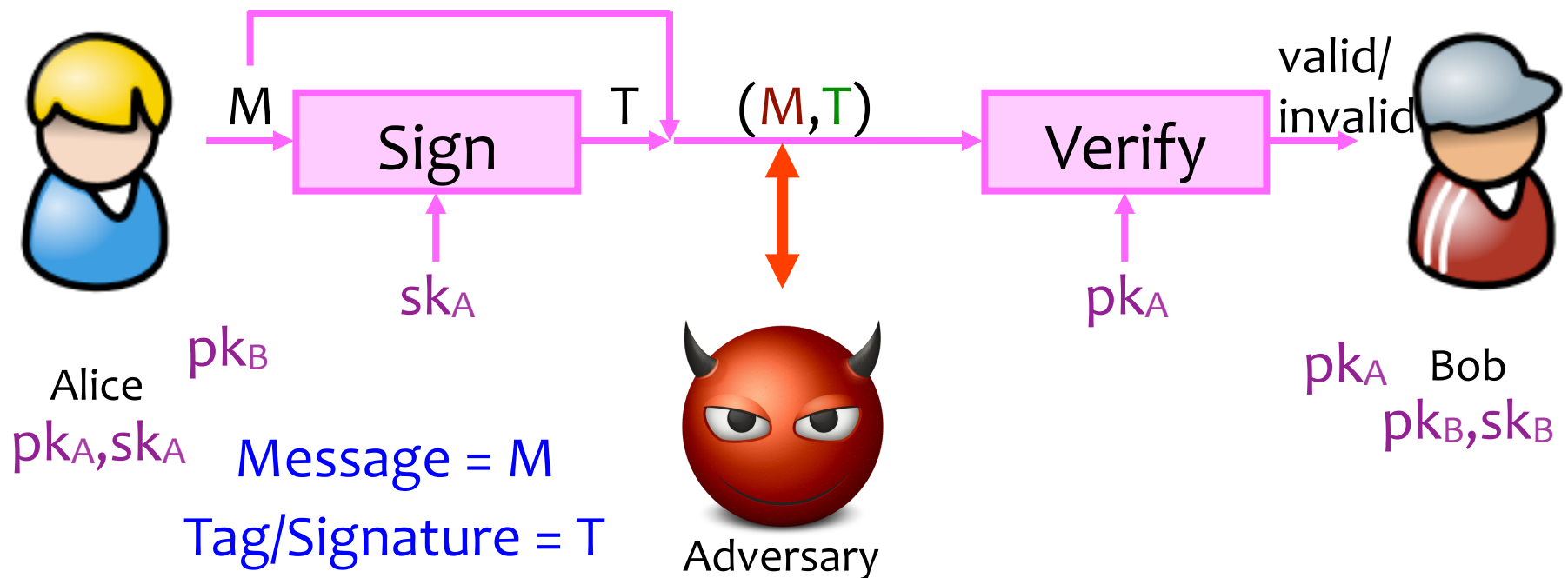
# Achieving Integrity (Symmetric)

Message authentication schemes: A tool for protecting integrity.  
(Also called message authentication codes or MACs.)



# Achieving Integrity (Asymmetric)

Digital signature schemes: A tool for protecting integrity and authenticity.



# Flavors of Cryptography

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# Flavors of Cryptography

- Symmetric cryptography
  - Both communicating parties have access to a shared random string  $K$ , called the key.
  - Challenge: How do you privately share a key?
- Asymmetric cryptography
  - Each party creates a public key  $pk$  and a secret key  $sk$ .
  - Challenge: How do you validate a public key?