CSE 484 / CSE M 584: Computer Security and Privacy

Software Security: Buffer Overflow Attacks

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Thanks to Dan Boneh, Dieter Gollmann, Dan Halperin, Yoshi Kohno, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

Announcements

- Remember to sign the ethics form
- Please start forming groups of 3 for lab 1
 - Once formed, post your final group to the forum
 - You can also find group members on the forum

Forum & Homework 1

- Goal: help develop the "security mindset"
- Best way to learn a foreign language: move to that country and immerse yourself in the language.
- Same thing applies to "security thinking"
- Forum: opportunity to think about security on a regular basis -- outside of class
 - Current events
 - New product announcements
- Also in forum:
 - Security in your everyday life
 - Movies / books

– ...

Security: Not Just for PCs



smartphones

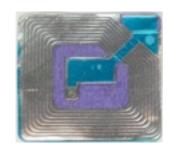


wearables





voting machines



RFID



game platforms



EEG headsets



mobile sensing platforms



medical devices



cars

airplanes

Software Lifecycle (Simplified)

- Requirements
- Design
- Implementation
- Testing
- Use

Software Bug Halts F-22 Flight

Posted by kdawson on Sunday February 25, @06:35PM from the dare-you-to-cross-this-line dept.

mgh02114 writes

"The new US stealth fighter, the F-22 Raptor, was deployed for the first time to Asia earlier this month. On Feb. 11, twelve Raptors flying from Hawaii to Japan were forced to turn back when a software glitch crashed all of the F-22s' on-board computers as they crossed the international date line. The delay in arrival in Japan was previously reported, with rumors of problems with the software. CNN television, however, this morning reported that every fighter completely lost all navigation and communications when they crossed the international date line. They reportedly had to turn around and follow their tankers by visual contact back to Hawaii. According to the CNN story, if they had not been with their tankers, or the weather had been bad, this would have been serious. CNN has not put up anything on their website yet."



1985-1987 -- Therac-25 medical accelerator. A radiation therapy device malfunctions and delivers lethal radiation doses at several medical facilities. Based upon a previous design, the Therac-25 was an "improved" therapy system that could deliver two different kinds of radiation: either a low-power electron beam (beta particles) or X-rays. The Therac-25's X-rays were generated by smashing high-power electrons into a metal target positioned between the electron gun and the patient. A second "improvement" was the replacement of the older Therac-20's electromechanical safety interlocks with software control, a decision made because software was perceived to be more reliable.

What engineers didn't know was that both the 20 and the 25 were built upon an operating system that had been kludged together by a programmer with no formal training. Because of a subtle bug called a "race condition," a quick-fingered typist could accidentally configure the Therac-25 so the electron beam would fire in high-power mode but with the metal X-ray target out of position. At least five patients die; others are seriously injured.

January 15, 1990 -- AT&T Network Outage. A bug in a new release of the software that controls AT&T's #4ESS long distance switches causes these mammoth computers to crash when they receive a specific message from one of their neighboring machines -- a message that the neighbors send out when they recover from a crash.

One day a switch in New York crashes and reboots, causing its neighboring switches to crash, then their neighbors' neighbors, and so on. Soon, 114 switches are crashing and rebooting every six seconds, leaving an estimated 60 thousand people without long distance service for nine hours. The fix: engineers load the previous software release.

- Other serious bugs (many others exist)
 - US Vincennes tracking software
 - MV-22 Osprey



- Medtronic Model 8870 Software Application Card

Adversarial Failures

- Software bugs are bad
 - Consequences can be serious
- Even worse when an intelligent adversary wishes to exploit them!
 - Intelligent adversaries: Force bugs into "worst possible" conditions/states
 - Intelligent adversaries: Pick their targets
- Buffer overflows bugs: Big class of bugs
 - Normal conditions: Can sometimes cause systems to fail
 - Adversarial conditions: Attacker able to violate security of your system (control, obtain private information, ...)

BUFFER OVERFLOWS

A Bit of History: Morris Worm

- Worm was released in 1988 by Robert Morris
 - Graduate student at Cornell, son of NSA chief scientist
 - Convicted under Computer Fraud and Abuse Act, sentenced to 3 years of probation and 400 hours of community service
 - Now an EECS professor at MIT
- Worm was intended to propagate slowly and harmlessly measure the size of the Internet
- Due to a coding error, it created new copies as fast as it could and overloaded infected machines
- \$10-100M worth of damage

Morris Worm and Buffer Overflow

- One of the worm's propagation techniques was a buffer overflow attack against a vulnerable version of fingerd on VAX systems
 - By sending special string to finger daemon, worm caused it to execute code creating a new worm copy
 - Unable to determine remote OS version, worm also attacked fingerd on Suns running BSD, causing them to crash (instead of spawning a new copy)

Famous Internet Worms

- Buffer overflows: very common cause of Internet attacks
 - In 1998, over 50% of advisories published by CERT (computer security incident report team) were caused by buffer overflows
- Morris worm (1988): overflow in fingerd
 - 6,000 machines infected
- CodeRed (2001): overflow in MS-IIS server
 - 300,000 machines infected in 14 hours
- SQL Slammer (2003): overflow in MS-SQL server
 - 75,000 machines infected in 10 minutes (!!)
- Sasser (2005): overflow in Windows LSASS
 - Around 500,000 machines infected

... And More

- Conficker (2008-08): overflow in Windows RPC
 - Around 10 million machines infected (estimates vary)
- Stuxnet (2009-10): several zero-day overflows + same
 Windows RPC overflow as Conficker
 - Windows print spooler service
 - Windows LNK shortcut display
 - Windows task scheduler
- Flame (2010-12): same print spooler and LNK overflows as Stuxnet
 - Targeted cyperespionage virus
- Still ubiquitous, especially in embedded systems

Attacks on Memory Buffers

- Buffer is a pre-defined data storage area inside computer memory (stack or heap)
- Typical situation:
 - A function takes some input that it writes into a preallocated buffer.
 - The developer forgets to check that the size of the input isn't larger than the size of the buffer.
 - Uh oh.
 - "Normal" bad input: crash
 - "Adversarial" bad input: take control of execution

Stack Buffers

buf uh oh!

Suppose Web server contains this function

```
void func(char *str) {
      char buf[126];
      ...
      strcpy(buf,str);
      ...
}
```

- No bounds checking on strcpy()
- If str is longer than 126 bytes
 - Program may crash
 - Attacker may change program behavior

Example: Changing Flags

buf

I (:-)!)

Suppose Web server contains this function

```
void func(char *str) {
     char buf[126];
     ...
     strcpy(buf,str);
     ...
}
```

- Authenticated variable non-zero when user has extra privileges
- Morris worm also overflowed a buffer to overwrite an authenticated flag in fingerd

Memory Layout

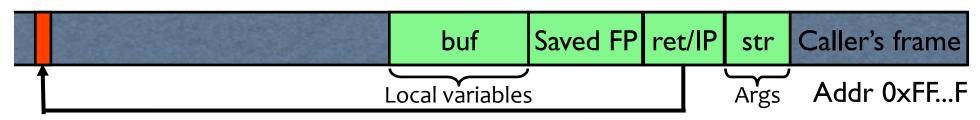
- Text region: Executable code of the program
- Heap: Dynamically allocated data
- Stack: Local variables, function return addresses; grows and shrinks as functions are called and return

		Top	Bottom
Text region	Неар	Stack	
Addr 0x000		Ad	dr 0xFFF

Stack Buffers

Suppose Web server contains this function:

 When this function is invoked, a new frame (activation record) is pushed onto the stack.



Execute code at this address after func() finishes

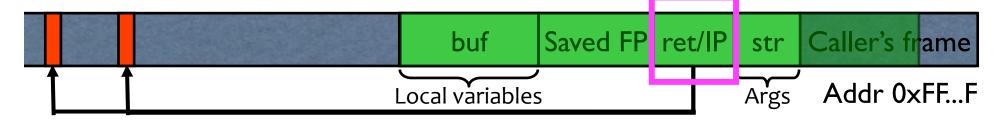
What if Buffer is Overstuffed?

Memory pointed to by str is copied onto stack…

```
void func(char *str) {
    char buf[126];
    strcpy does NOT check whether the string
    at *str contains fewer than 126 characters
}
```

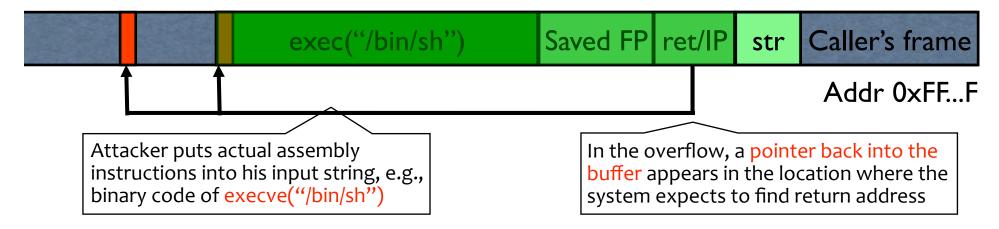
• If a string longer than 126 bytes is copied into buffer, it will overwrite adjacent stack locations.

This will be interpreted as return address!



Executing Attack Code

- Suppose buffer contains attacker-created string
 - For example, str points to a string received from the network as the URL



- When function exits, code in the buffer will be executed, giving attacker a shell
 - Root shell if the victim program is setuid root

Buffer Overflow Issues

- Executable attack code is stored on stack, inside the buffer containing attacker's string
 - Stack memory is supposed to contain only data, but...
- Overflow portion of the buffer must contain correct address of attack code in the RET position
 - The value in the RET position must point to the beginning of attack assembly code in the buffer
 - Otherwise application will (probably) crash with segmentation violation
 - Attacker must correctly guess in which stack position his/her buffer will be when the function is called

Problem: No Bounds Checking

- strcpy does <u>not</u> check input size
 - strcpy(buf, str) simply copies memory contents into buf starting from *str until "\o" is encountered, ignoring the size of area allocated to buf
- Many C library functions are unsafe
 - strcpy(char *dest, const char *src)
 - strcat(char *dest, const char *src)
 - gets(char *s)
 - scanf(const char *format, ...)
 - printf(const char *format, ...)

Does Bounds Checking Help?

- strncpy(char *dest, const char *src, size_t n)
 - If strncpy is used instead of strcpy, no more than n characters will be copied from *src to *dest
 - Programmer has to supply the right value of n
- Potential overflow in htpasswd.c (Apache 1.3):

```
strcpy(record, user);
strcat(record, ":");
strcat(record, cpw);
Copies username ("user") into buffer
("record"), then appends ":" and
hashed password ("cpw")
```

Published "fix":

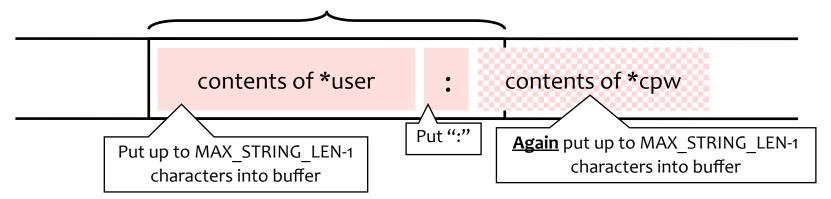
```
strncpy(record,user,MAX_STRING_LEN-1);
strcat(record,":")
strncat(record,cpw,MAX_STRING_LEN-1);
```

Misuse of strncpy in htpasswd "Fix"

Published "fix" for Apache htpasswd overflow:

```
strncpy(record, user, MAX_STRING_LEN-1);
strcat(record, ":")
strncat(record, cpw, MAX_STRING_LEN-1);
```

MAX STRING LEN bytes allocated for record buffer



What About This?

Home-brewed range-checking string copy

```
void mycopy(char *input) {
    char buffer[512]; int i;

for (i=0; i<=512; i++)
    buffer[i] = input[i];
}
void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}</pre>
```

Off-By-One Overflow

Home-brewed range-checking string copy

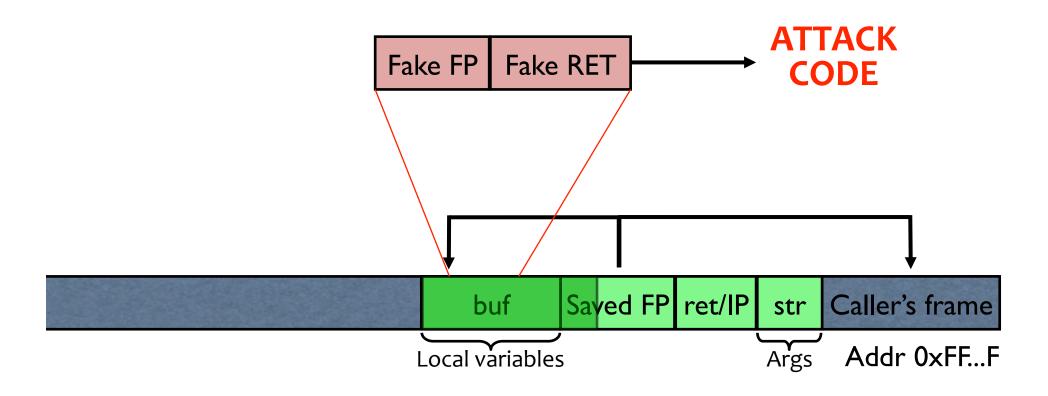
```
void mycopy(char *input) {
    char buffer[512]; int i;

    for (i=0; i<=)12; i++)
        buffer[i] = input[i];
}

void main(int argc, char *argv[]) {
    if (argc==2)
        mycopy(argv[1]);
}</pre>
This will copy 513
characters into
buffer. Oops!
```

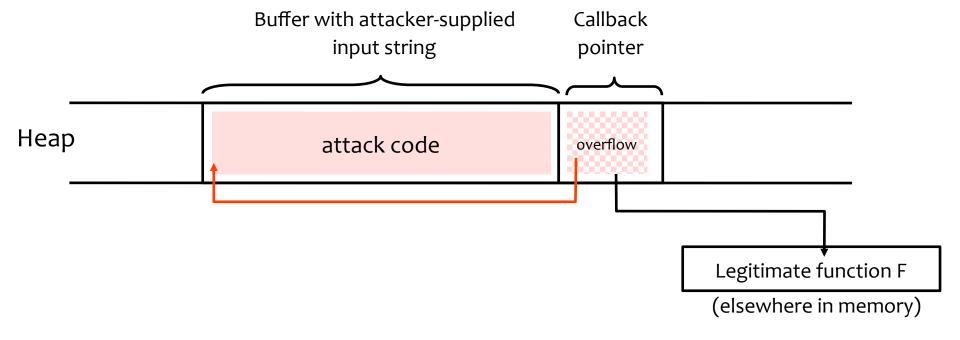
- 1-byte overflow: can't change RET, but can change pointer to previous stack frame
 - On little-endian architecture, make it point into buffer
 - RET for previous function will be read from buffer!

Frame Pointer Overflow



Another Variant: Function Pointer Overflow

• C uses function pointers for callbacks: if pointer to F is stored in memory location P, then another function G can call F as (*P)(...)



Other Overflow Targets

- Format strings in C
 - More details next time
- Heap management structures used by malloc()
 - More details in section

 These are all attacks you can look forward to in Lab #1 ⁽²⁾

Looking Forward

- Ethics form due April 8
- Homework #1 due April 17
- Next week
 - Monday: more buffer overflows
 - Wednesday: start crypto
 - Friday: guest lecture by David Aucsmith
- Section next week about Lab 1