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Protocol Rollback and Network Security

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Goals for Today

Protocol Rollback Attacks (in SSL)

Network security



What is SSL / TLS?

Transport Layer Security (TLS) protocol, version 1.2

- De facto standard for Internet security
- "The primary goal of the TLS protocol is to provide privacy and data integrity between two communicating applications"
- In practice, used to protect information transmitted between browsers and Web servers (and mail readers and ...)
- https://datatracker.ietf.org/wg/tls/

 Based on Secure Sockets Layers (SSL) protocol, version 3.0

• Same protocol design, different algorithms

Ubiquitously deployed in commercial Web browsers

TLS Basics

TLS consists of two protocols

• Familiar pattern for key exchange protocols

Handshake protocol

 Use public-key cryptography to establish a shared secret key between the client and the server

Record protocol

- Use the secret key established in the handshake protocol to protect communication between the client and the server
- We will focus on the handshake protocol

TLS Handshake Protocol

- Two parties: client and server
- Negotiate version of the protocol and the set of cryptographic algorithms to be used
 - Interoperability between different implementations of the protocol
- Authenticate client and server (optional)
 - Use digital certificates to learn each other's public keys and verify each other's identity
- Use public keys to establish a shared secret

Handshake Protocol Structure



ClientHello



ServerHello



ServerKeyExchange



ClientKeyExchange





Version Rollback Attack



SSL 2.0 Weaknesses (Fixed in 3.0)

- Cipher suite preferences are not authenticated
 - "Cipher suite rollback" attack is possible
- SSL 2.0 uses padding when computing MAC in block cipher modes, but padding length field is not authenticated
 - Attacker can delete bytes from the end of messages
- MAC hash uses only 40 bits in export mode
- No support for certificate chains or non-RSA algorithms, no handshake while session is open

Protocol Rollback Attacks

Why do people release new versions of security protocols? Because the old version got broken!

- New version must be backward-compatible
 - Not everybody upgrades right away
- Attacker can fool someone into using the old, broken version and exploit known vulnerability
 - Similar: fool victim into using weak crypto algorithms
- Defense is hard: must authenticate version in early designs
- Many protocols had "version rollback" attacks
 - SSL, SSH, GSM (cell phones)

Version Check in SSL 3.0 (Approximate)



SSL/TLS Record Protection



Internet Infrastructure



TCP/IP for packet routing and connections
Border Gateway Protocol (BGP) for route discovery
Domain Name System (DNS) for IP address discovery



• Service (can get to Internet)



- Privacy (middle-entities shouldn't know what communicating or with whom)
- Fairness (e.g., get service I paid for)

User

- Integrity (can't impersonate me, modify my data)
- Safety (network shouldn't attack me)



• Service (clients can get to Internet)

• Performance (network works well)

Network Admin

- Identity (know what's on network)
- Safety (no one launching attacks)
- Accountability (can find bad users)

- Service (deliver traffic -> earn \$\$)
- Reliability & Performance (network works well)

Intermediate

ISPs

 Integrity of delivered traffic (can bill customers properly, you're not overcharged by providers)



Server

- Service (deliver traffic -> earn \$\$)
- Reliability & Performance (network works well)
- Analytics (better delivery)
- Accounting (can bill customers properly)
- Safety (not being attacked)



OSI Protocol Stack

email, Web, NFS application presentation RPC session TCP, UDP, ICMP transport IP network Ethernet data link physical

Data Formats



IP (Internet Protocol)

Connectionless

- Unreliable, "best-effort" protocol
- Uses numeric addresses for routing
 - Typically several hops in the route



TCP (Transmission Control Protocol)

Sender: break data into packets

• Sequence number is attached to every packet

Receiver: reassemble packets in correct order

• Acknowledge receipt; lost packets are re-sent

Connection state maintained on both sides



UDP (User Datagram Protocol)

Sender: break data into packets

- Sequence number maybe? If Application wants them
- Receiver: receive packets
 - No acknowledgement
 - Dropped packets are skipped no retransmission



ICMP (Control Message Protocol)

Provides feedback about network operation

- "Out-of-band" messages carried in IP packets
- Error reporting, congestion control, reachability, etc.
- Example messages:
 - Destination unreachable
 - Time exceeded
 - Parameter problem
 - Redirect to better gateway
 - Reachability test (echo / echo reply)
 - Message transit delay (timestamp request / reply)



Detecting attacks



• **Problem:** IP packets contain source IP address

User

Launch undetectable attacks

• Solution: Spoof IP address

