

CSE 484 / CSE M 584 (Autumn 2011)

# Security and Networks

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Thanks to Dan Boneh, Dieter Gollmann, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

# Class updates

- Homework 3 due today
- My office hours this week:
  - **CSE 210:** W,Th,F in the after-class slot
  - Other times by appointment.
  - Come pick up graded Homework #2

# Lab 3

- **Posted on website and on Catalyst.**

- <https://catalyst.uw.edu/collectit/assignment/dhalperi/17513/72548>
- Hack my privacy!

- ***This lab is optional***

- Can only help your grade.
- Lots of opportunity for extra credit.
- I really think this lab is fun, and encourage you to do it, but we're not going to require it.

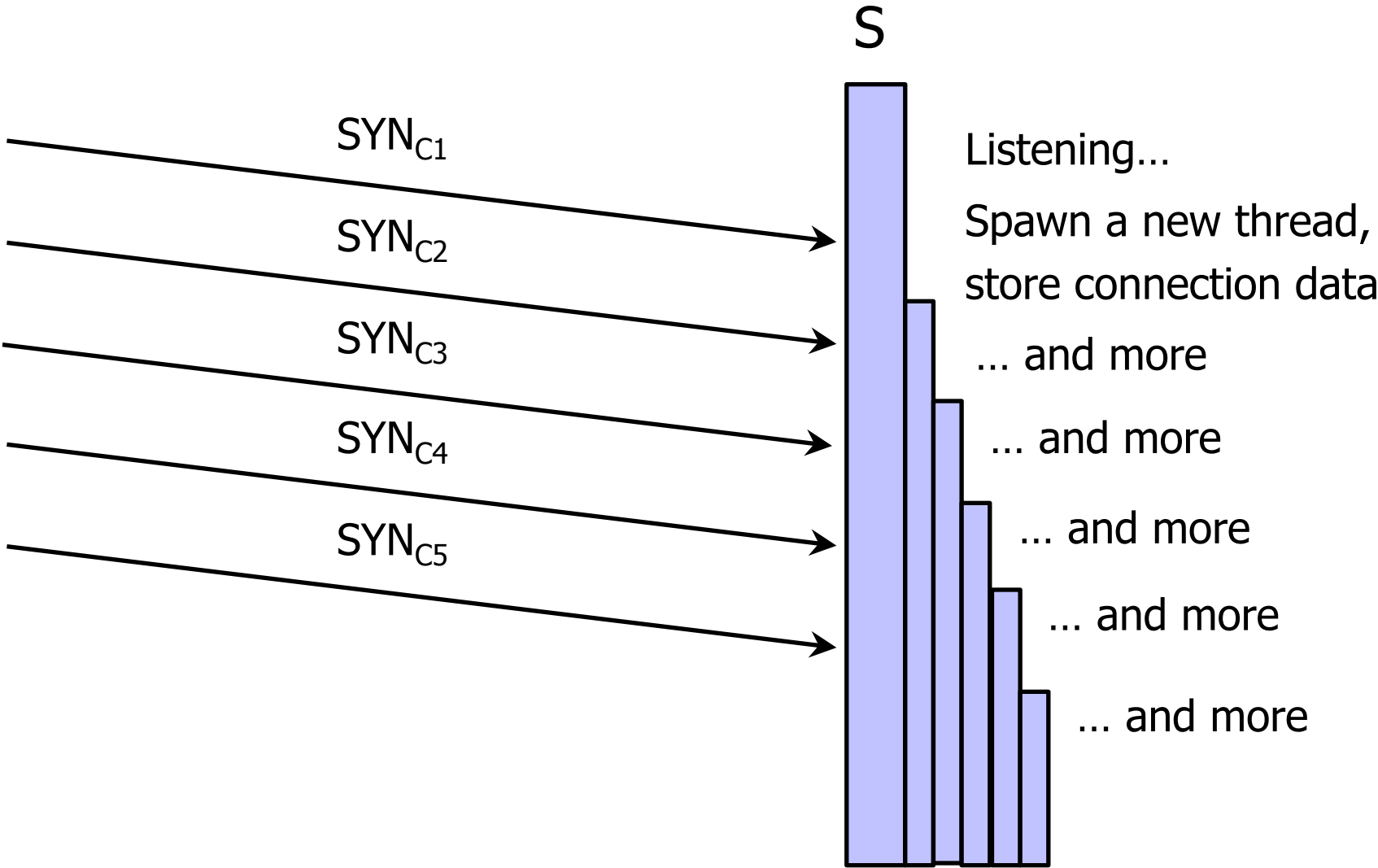
# This week

- **Today:** Finish networks, Final, & Course Evals
- **Friday:** Any questions you have
  - Submit to my email, cse484-tas
  - Submit anonymously via the feedback form on the website

# Grading?

**Final?**

# SYN Flooding Attack



# SYN Flooding Explained

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- ◆ Attacker sends many connection requests with spoofed source addresses
- ◆ Victim allocates resources for each request
  - Connection state maintained until timeout
  - Fixed bound on half-open connections
- ◆ Once resources exhausted, requests from legitimate clients are denied
- ◆ This is a classic **denial of service (DoS)** attack
  - Common pattern: it costs nothing to TCP initiator to send a connection request, but TCP responder must allocate state for each request (**asymmetry!**)



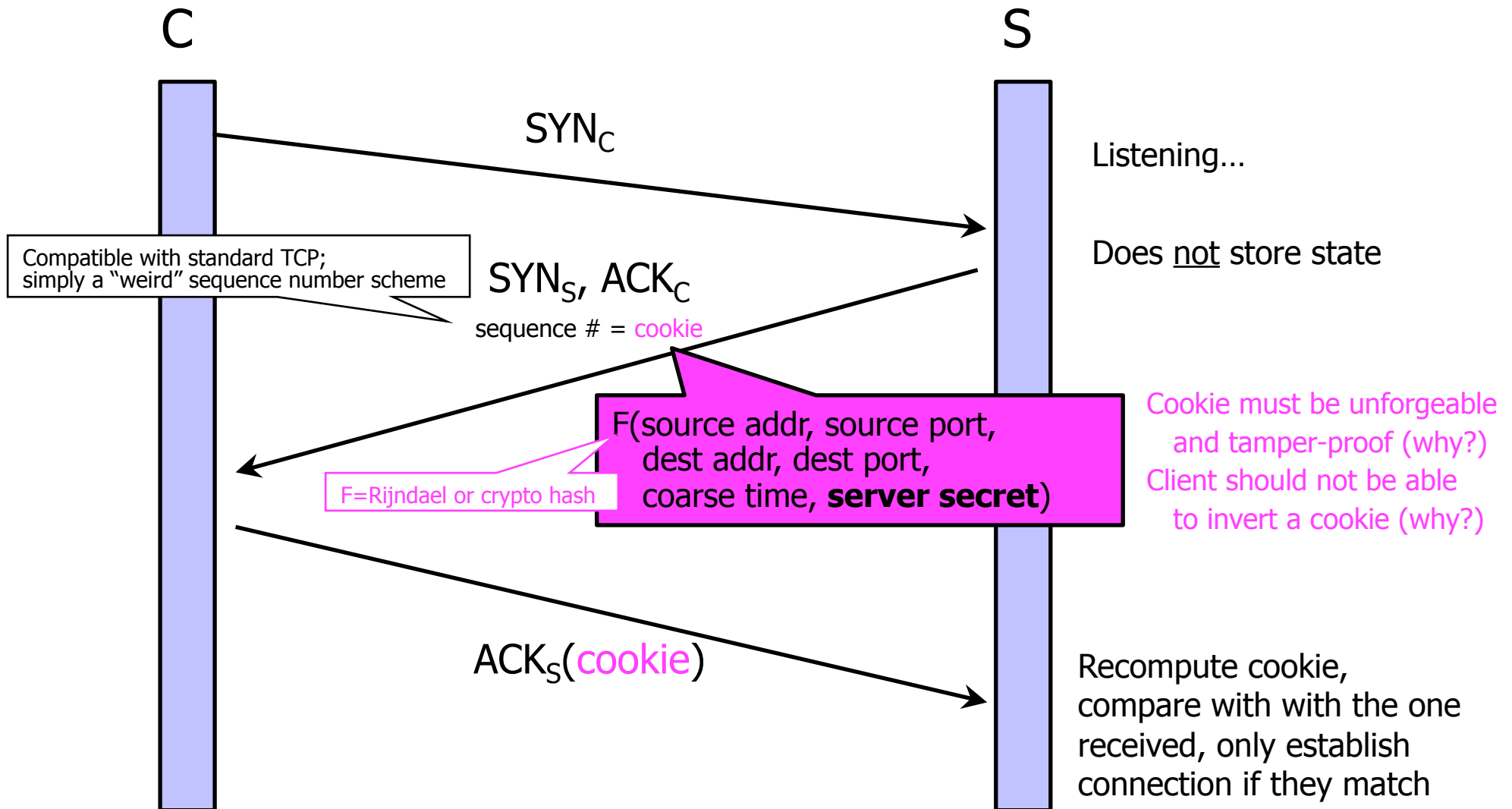
# Preventing Denial of Service

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- ◆ DoS is caused by asymmetric state allocation
  - If responder opens a state for each connection attempt, attacker can initiate thousands of connections from bogus or forged IP addresses
- ◆ **Cookies** ensure that the responder is stateless until initiator produced at least 2 messages
  - Responder's state (IP addresses and ports of the connection) is stored in a cookie and sent to initiator
  - After initiator responds, cookie is regenerated and compared with the cookie returned by the initiator

# SYN Cookies

[Bernstein and Schenk]



More info: <http://cr.yp.to/syncookies.html>

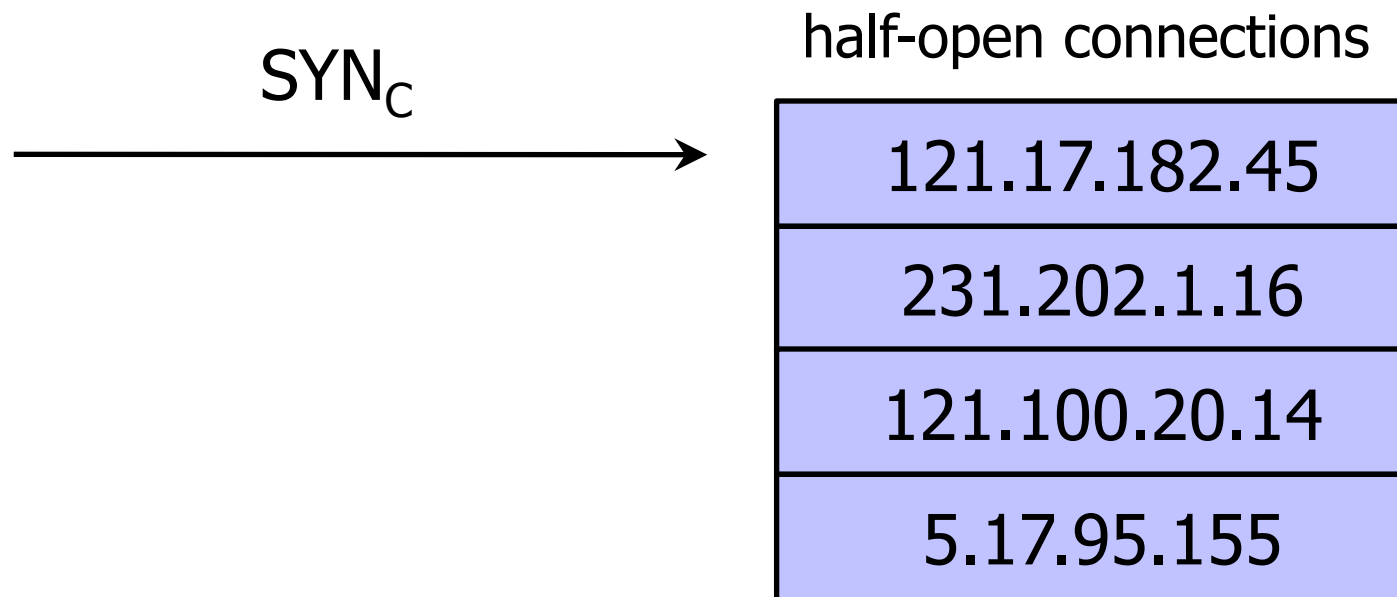
# Anti-Spoofing Cookies: Basic Pattern

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- ◆ Client sends request (message #1) to server
- ◆ Typical protocol:
  - Server sets up connection, responds with message #2
  - Client may complete session or not (potential DoS)
- ◆ Cookie version:
  - Server responds with hashed connection data instead of message #2
  - Client confirms by returning hashed data
    - If source IP address is bogus, attacker can't confirm
  - Need an extra step to send postponed message #2, except in TCP (SYN-ACK already there)

# Another Defense: Random Deletion

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- ◆ If SYN queue is full, delete random entry
  - Legitimate connections have a chance to complete
  - Fake addresses will be eventually deleted
- ◆ Easy to implement

# “Ping of Death”

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- ◆ If an old Windows machine received an ICMP packet with a payload longer than 64K, machine would crash or reboot
  - Programming error in older versions of Windows
  - Packets of this length are illegal, so programmers of Windows code did not account for them
- ◆ Recall “security theme” of this course - every line of code might be the target of an adversary

Solution: patch OS, filter out ICMP packets

# Intrusion Detection Systems

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- ◆ Advantage: can recognize new attacks and new versions of old attacks
- ◆ Disadvantages
  - High false positive rate
  - Must be trained on known good data
    - Training is hard because network traffic is very diverse
  - Definition of “normal” constantly evolves
    - What’s the difference between a **flash crowd** and a **denial of service** attack?

# Intrusion Detection Problems

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- ◆ Lack of training data with real attacks
  - But lots of “normal” network traffic, system call data
- ◆ Data drift
  - Statistical methods detect changes in behavior
  - Attacker can attack gradually and incrementally
- ◆ Main characteristics not well understood
  - By many measures, attack may be within bounds of “normal” range of activities
- ◆ False identifications are very costly
  - Sysadm will spend many hours examining evidence

# Intrusion Detection Errors

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- ◆ **False negatives:** attack is not detected
  - Big problem in signature-based misuse detection
- ◆ **False positives:** harmless behavior is classified as an attack
  - Big problem in statistical anomaly detection
- ◆ Both types of IDS suffer from both error types
- ◆ Which is a bigger problem?
  - Attacks are fairly rare events



# Base-Rate Fallacy

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- ◆ 1% of traffic is SYN floods; IDS accuracy is 90%
  - IDS classifies a SYN flood as attack with prob. 90%, classifies a valid connection as attack with prob. 10%
- ◆ What is the probability that a connection flagged by IDS as a SYN flood is actually valid traffic?

# Conditional Probability

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- ◆ Suppose two events A and B occur with probability  $\Pr(A)$  and  $\Pr(B)$ , respectively
- ◆ Let  $\Pr(AB)$  be probability that both A and B occur
- ◆ What is the **conditional probability** that A occurs assuming B has occurred?

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$$\Pr(A \mid B) = \frac{\Pr(AB)}{\Pr(B)}$$

# Bayes' Theorem

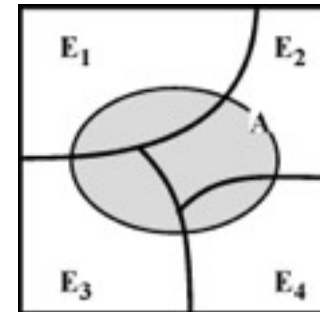
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- ◆ Suppose mutually exclusive events  $E_1, \dots, E_n$  together cover the entire set of possibilities
- ◆ Then probability of any event  $A$  occurring is

$$\Pr(A) = \sum_{1 \leq i \leq n} \Pr(A | E_i) \cdot \Pr(E_i)$$

– Intuition: since  $E_1, \dots, E_n$  cover entire

probability space, whenever  $A$  occurs,  
some event  $E_i$  must have occurred



- ◆ Can rewrite this formula as

$$\Pr(E_i | A) = \frac{\Pr(A | E_i) \cdot \Pr(E_i)}{\Pr(A)}$$

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$$\Pr(\text{valid} \mid \text{alarm}) = \frac{\Pr(\text{alarm} \mid \text{valid}) \cdot \Pr(\text{valid})}{\Pr(\text{alarm})}$$

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