

# Sampling-Based Motion Planning: From PRMs to RRTs

Instructor: Chris Mavrogiannis

TAs: Kay Ke, Gilwoo Lee, Matt Schmittle

\*Slides based on or adapted from Sanjiban Choudhury, Steve Lavalle, Peter Allen, Pieter Abbeel

# Logistics

- Lab 3
  - Deadline Friday March 6th
  - Demo Thursday March 5th (recitation slots)
  - **Extra Credit important for final project**
- Final Project
  - Out today!
  - Demo Thursday March 12th
  - Short writeup due Monday 16th
- Special Topics

# Probabilistic Roadmap Path Planning

- Explicit Geometry based planners impractical in high dimensional spaces.
- Exact solutions with complex geometries provably exponential
- Sampling-based planners can often create plans in high-dimensional spaces efficiently
- Rather than Compute the C-Space explicitly, we Sample it

# Completeness in Motion Planning

- **Complete Planner:** always answers a path planning query correctly in bounded time (return a path if one exists, otherwise report failure)
- ***Probabilistically* Complete Planner:** Probability of returning a path approaches 1 as more samples are generated.
- ***Resolution-Complete* Planner:** Probability of finding a path approaches 1 as the resolution becomes finer.

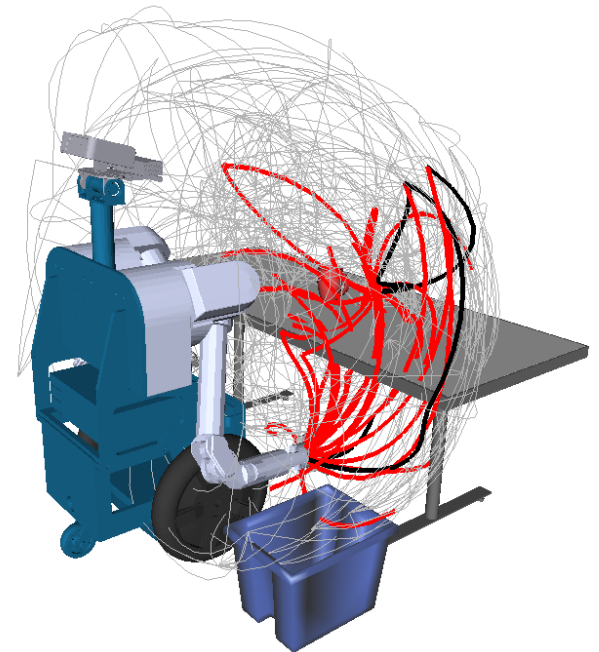
# Sampling-Based Planners

- Do not attempt to explicitly construct the C-Space and its boundaries.
- Simply need to know if a single robot configuration is in collision
- Exploits simple tests for collision with full knowledge of the space
- Collision detection is a separate module- can be tailored to the application
- As collision detection improves, so do these algorithms
- Different approaches for single-query and multi-query requests

# This is the PRM Algorithm!

PRM = Probabilistic Roadmap

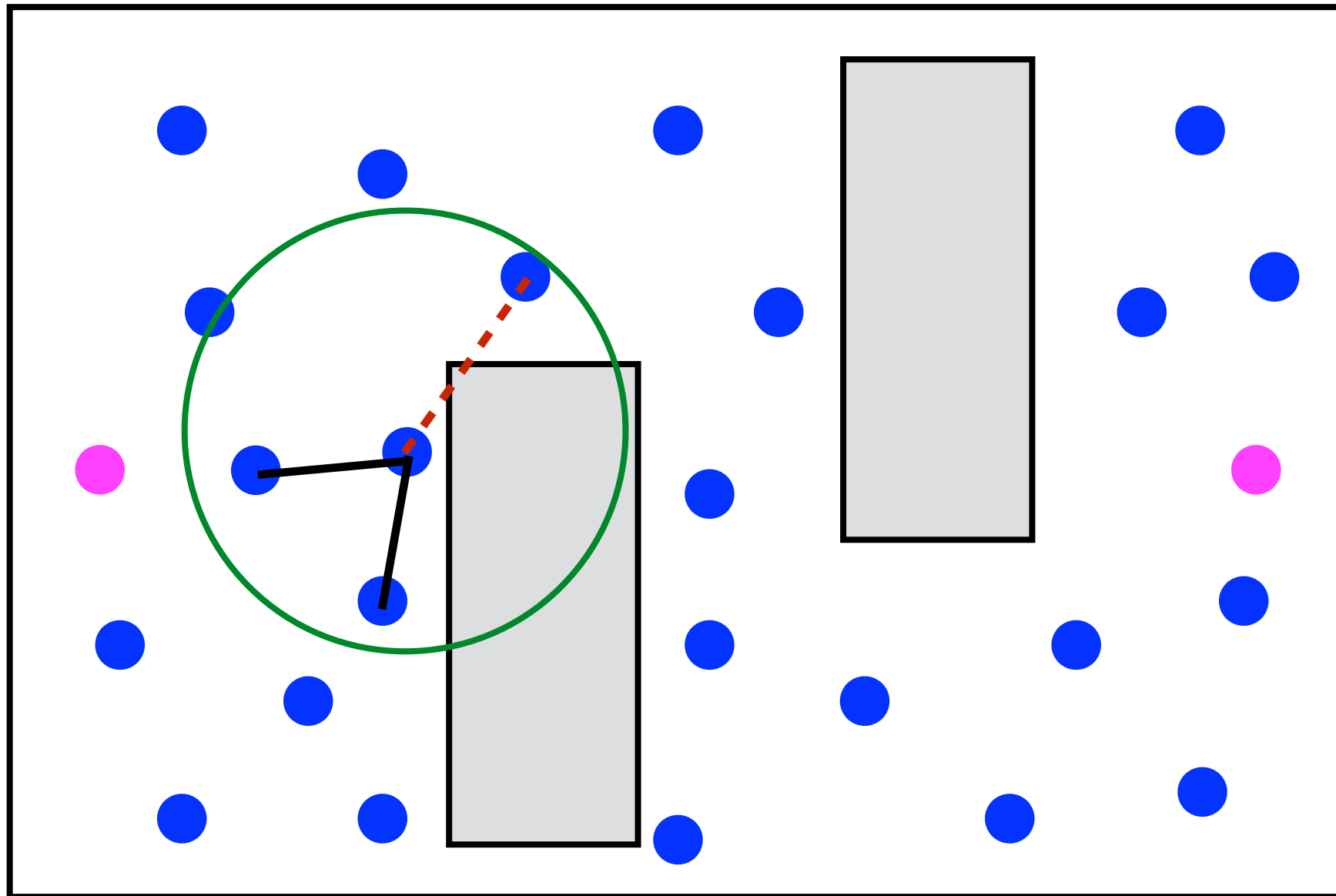
1. Sample vertices randomly and collision check
2. Try to connect vertices within radius (or k NN)
  - collision check potential edges
3. Search graph to find a solution
  - can reuse graph across multiple queries



Probabilistically Complete

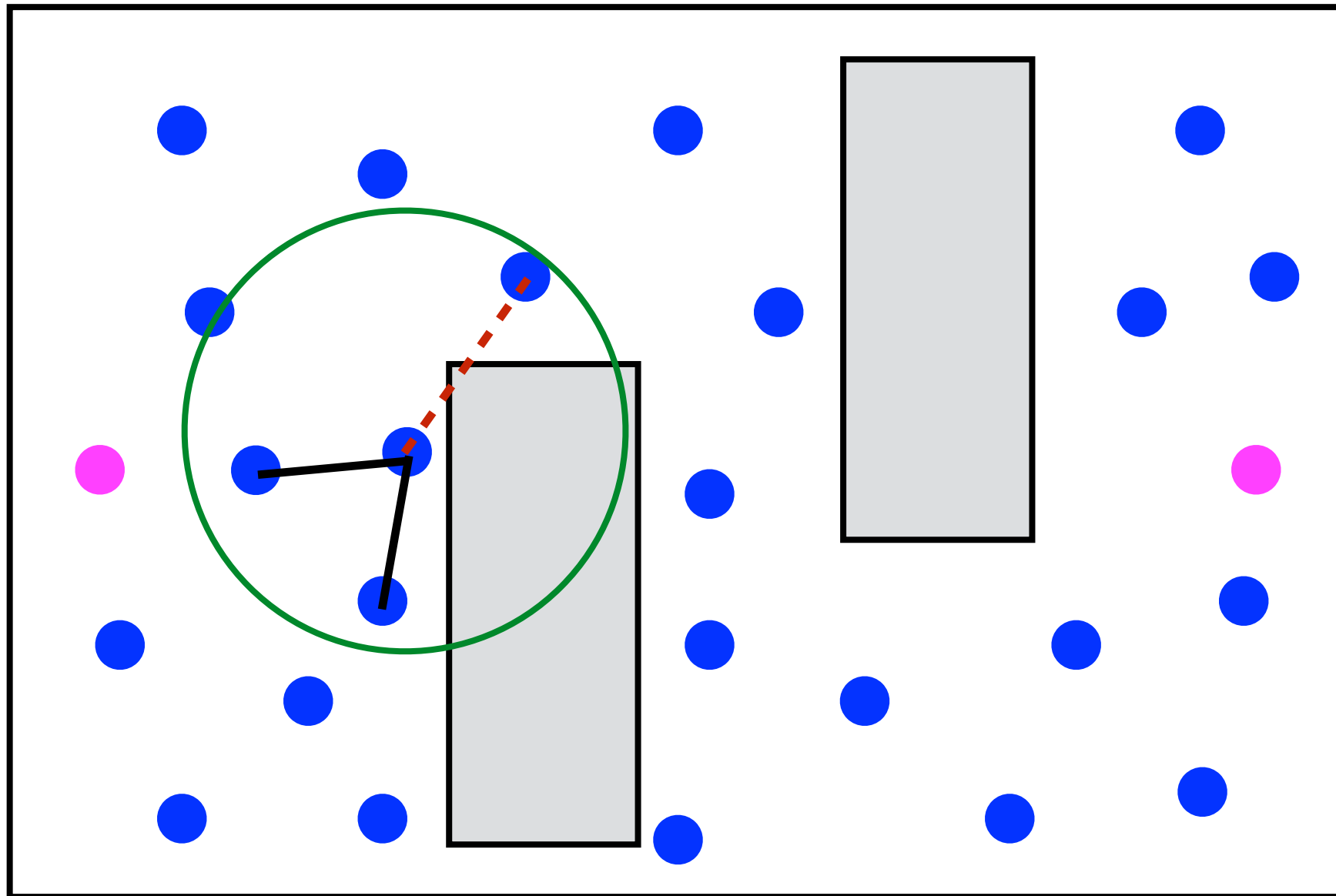
# What is the optimal radius?

What happens if radius too large? too small?



# What is the optimal radius?

Set the radius to  $r = \gamma \left( \frac{\log |V|}{|V|} \right)^{1/d}$  where magic constant!  
 $\gamma \geq 2(1 + 1/d)^{1/d} \frac{\mu(\mathcal{C}_{free})}{\zeta_d}$



Also known as a Random Geometric Graph (RGG)



# This is the PRM\* Algorithm!

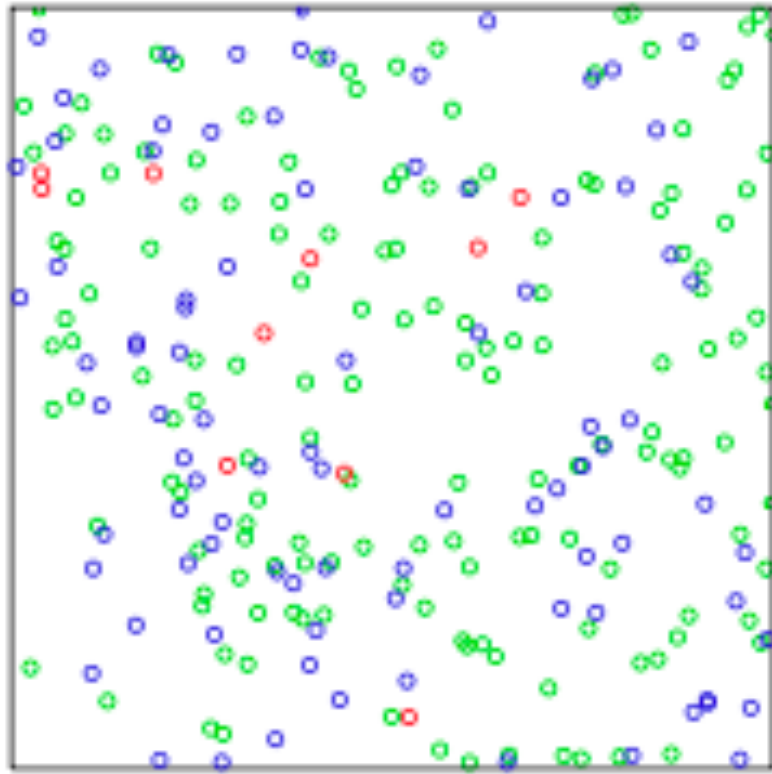
1. Sample vertices randomly
2. Use **optimal radius formula** to connect vertices
3. Search graph to find a solution

**Theorem:** Probabilistically complete AND Asymptotically optimal

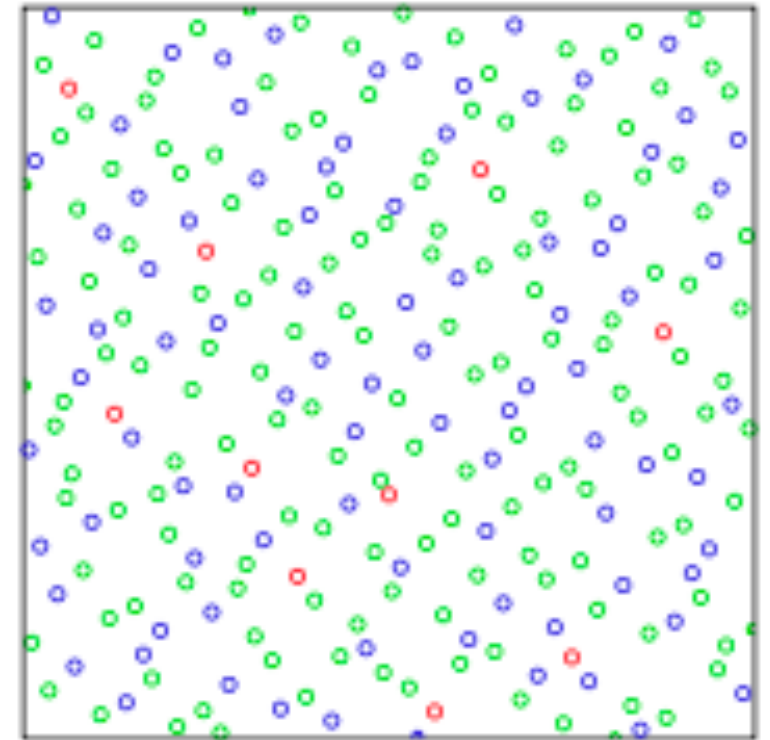
“Sampling-based Algorithms for Optimal Motion Planning”

Sertac Karaman and Emilio Frazzoli, IJRR 2011

# Can we do better than random?



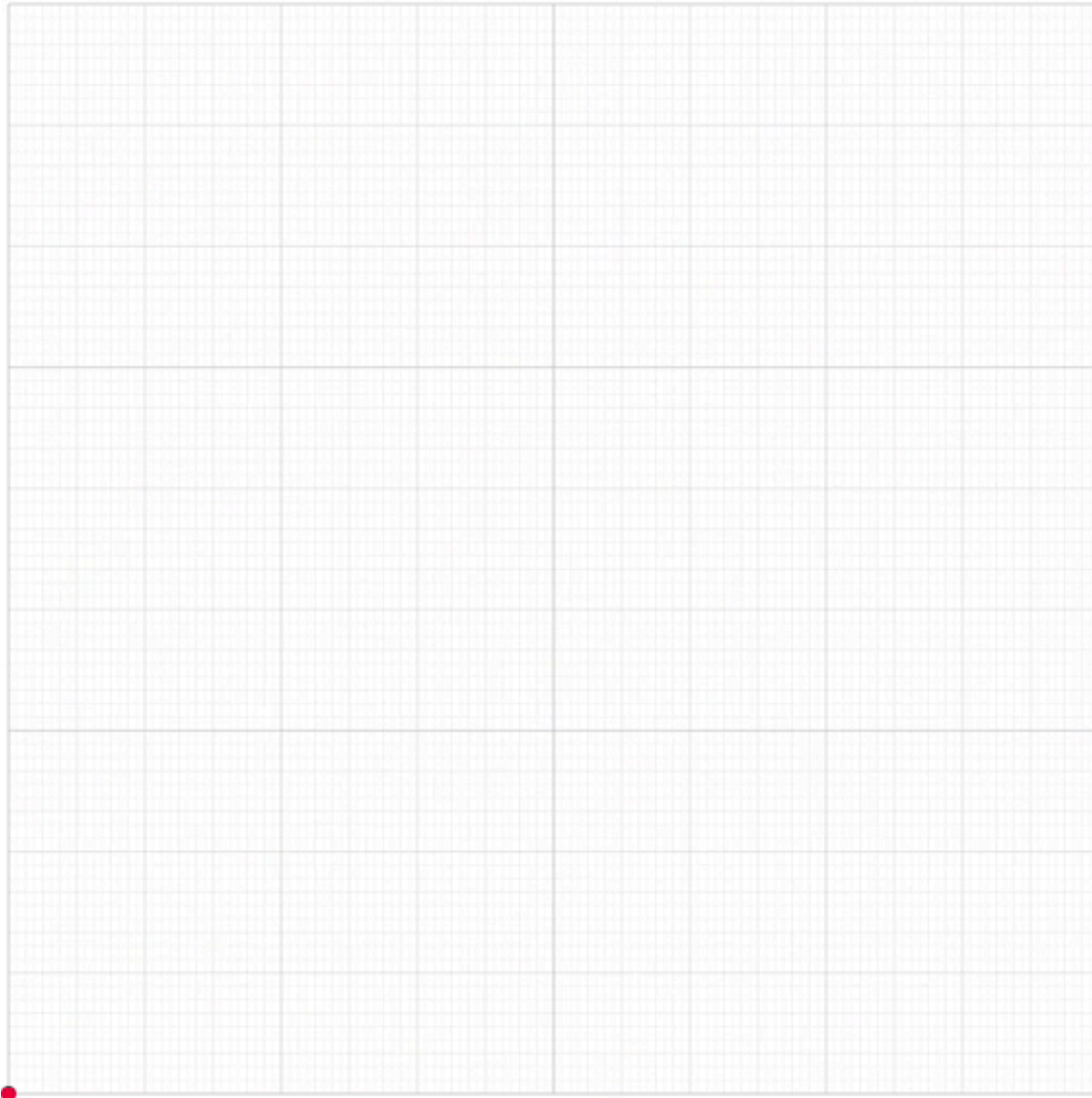
Uniform random  
sampling tends to  
clump



Ideally we would  
want points to be  
spread out evenly

**Question:** How do we do this without discretization?

# Halton Sequence

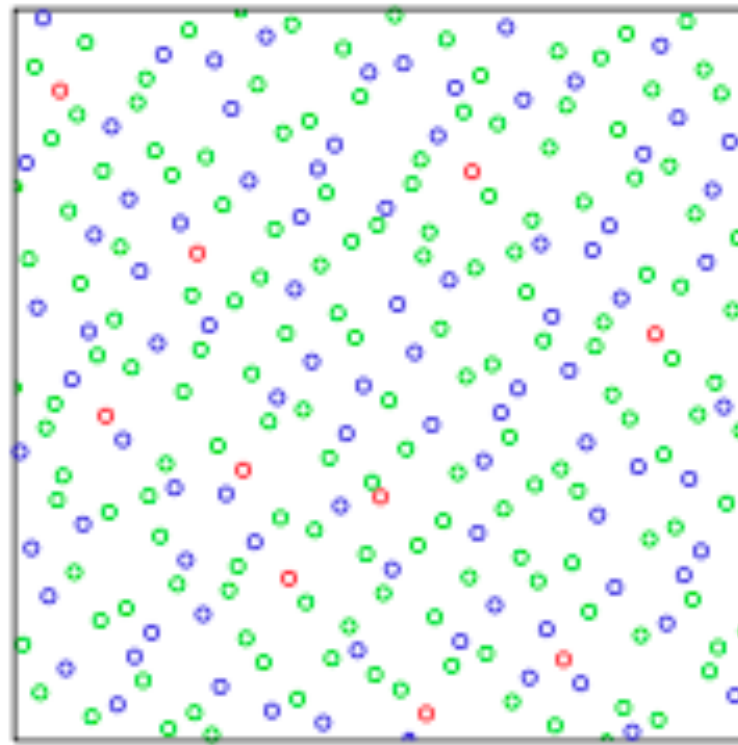


Generalization of  
Van de Coruput Sequence

**Intuition:** Create a sequence  
using prime numbers that  
uniformly densify space

Link for exact algorithm:  
<https://observablehq.com/@jrus/halton>

# How do we connect vertices?



Halton sequences have much better coverage  
(i.e. they are low dispersion)

Connect vertices that are within a radius of

$$r = \gamma \left( \frac{1}{|V|} \right)^{1/d} \quad \text{(as opposed to:)} \quad r = \gamma \left( \frac{\log |V|}{|V|} \right)^{1/d}$$

# This is the gPRM Algorithm!

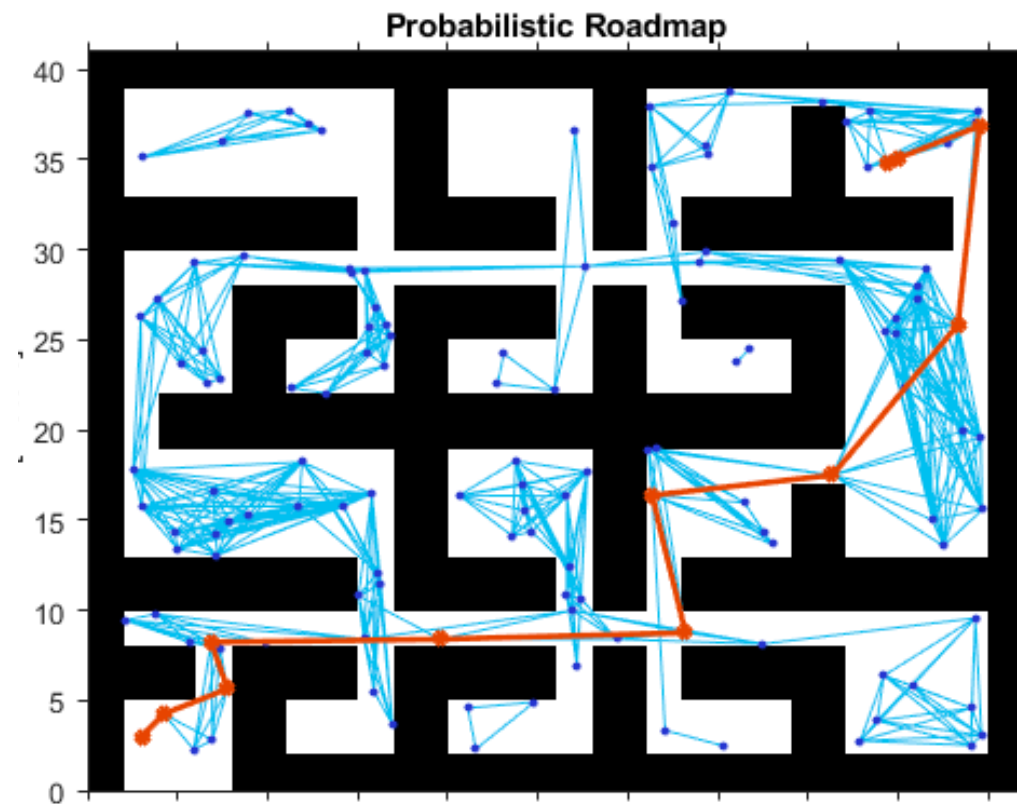
1. Sample vertices randomly
2. Use **optimal radius formula** to connect vertices
3. Search graph to find a solution

**Theorem:** Probabilistically complete AND Asymptotically optimal  
AND Asymptotic rate of convergence

“Deterministic Sampling-Based Motion Planning: Optimality, Complexity, and Performance”  
Lucas Janson, Brian Ichter, Marco Pavone, IJRR 2017

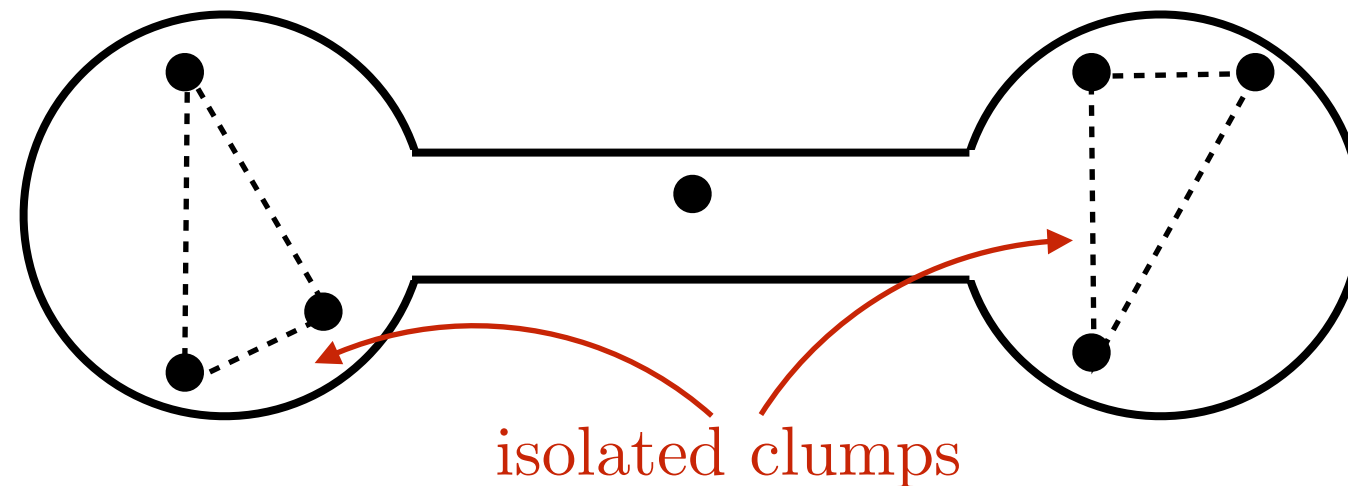
# What makes a good graph?

1. A good graph must be **sparse** (both in vertices and edges)
2. A good graph must have **good coverage**



3. A good graph must have the **same connectivity** of free space

# The Narrow Passage: Planning's boogie man!



Why is narrow passage mathematically hard to plan in?

Mathematical Question:

How many samples do we need to connect the space  
(with high probability)?

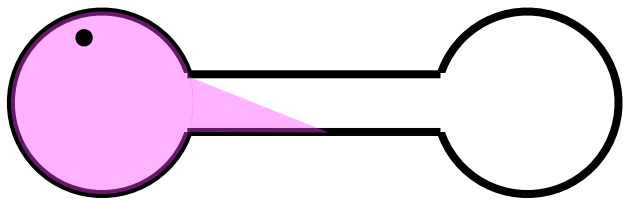
# How many samples do we need?

**Theorem** [Hsu et al., 1999] Let  $2n$  vertices be sampled from  $X_{\text{free}}$ . Then the roadmap is connected with probability at least  $1 - \gamma$  if:

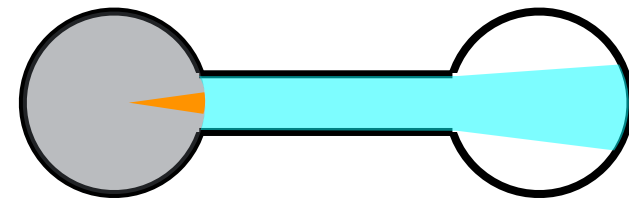
$$n \geq \left\lceil 8 \frac{\log(\frac{8}{\epsilon \alpha \gamma})}{\epsilon \alpha} + \frac{3}{\beta} + 2 \right\rceil$$

The shape of free C-space is dictated by  $\alpha$ ,  $\beta$ ,  $\epsilon \in [0, 1]$

Visibility of free space ( $\epsilon$ )



Expansion of visibility ( $\alpha$ ,  $\beta$ )



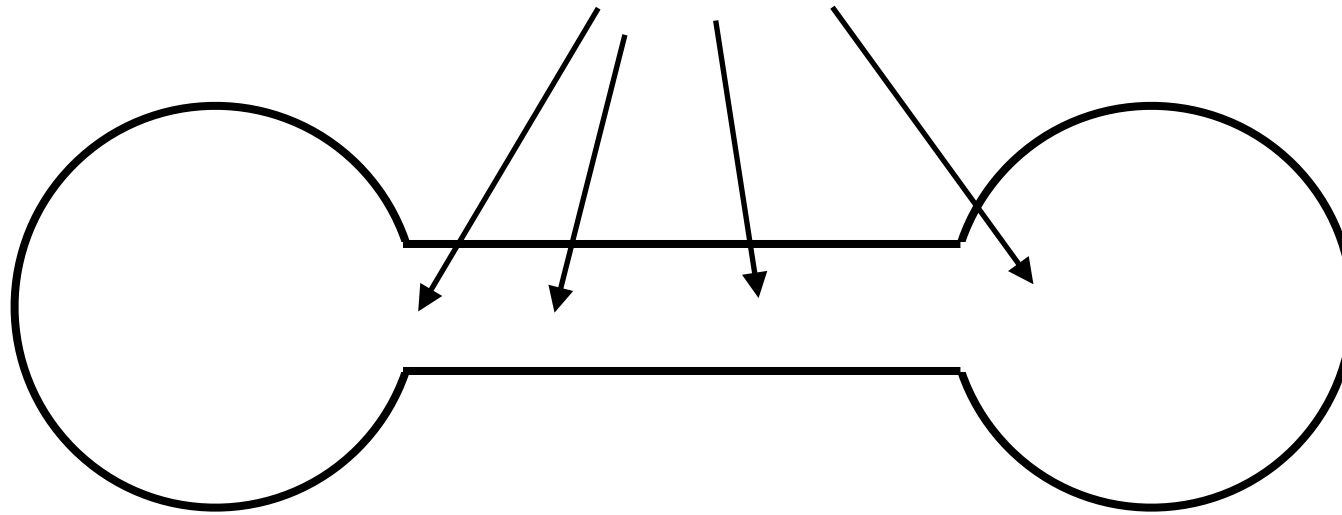
Narrow passage has small values of  $\alpha$ ,  $\beta$ ,  $\epsilon$

Hence, **needs more samples** to find a path



# How do we bias sampling?

We somehow need more samples here



## 1. Sample near obstacle surface?

V. Boor, M. H. Overmars, and A. F. van der Stappen. The Gaussian sampling strategy for probabilistic roadmap planners. 1999

## 2. Add samples that are in between two obstacles?

D. Hsu, T. Jiang, J. Reif, and Z. Sun. The bridge test for sampling narrow passages with probabilistic roadmap planners. 2003.

## 3. Train a learner to detect the narrow passages?

B. Ichter, J. Harrison, M. Pavone. Learning Sampling Distributions for Robot Motion Planning, 2018

# Summary of ways to create graphs

Algorithms	How to sample vertices?	How to connect vertices?
Lattice	Discretize	connectivity rule
PRM	Uniform random	r-disc, k-nn
PRM*	Uniform random	optimal r-disc, k-nn
gPRM	Halton sequence	optimal r-disc, k-nn
Bridge	Sample with bridge test	any visible points
Gaussian	Sample near obstacles	r-disc, k-nn
MAPRM	Sample along medial axis	r-disc, k-nn
Approx. Visibility Graph	Sample on surface of obstacles	any visible points
Learnt Sampler	Use CVAE to approximate free space	optimal r-disc, k-nn

# What graph should I use?

Low dim (2-3): Discretize evenly

Higher dim ( $\geq 4$ ): Halton sequence

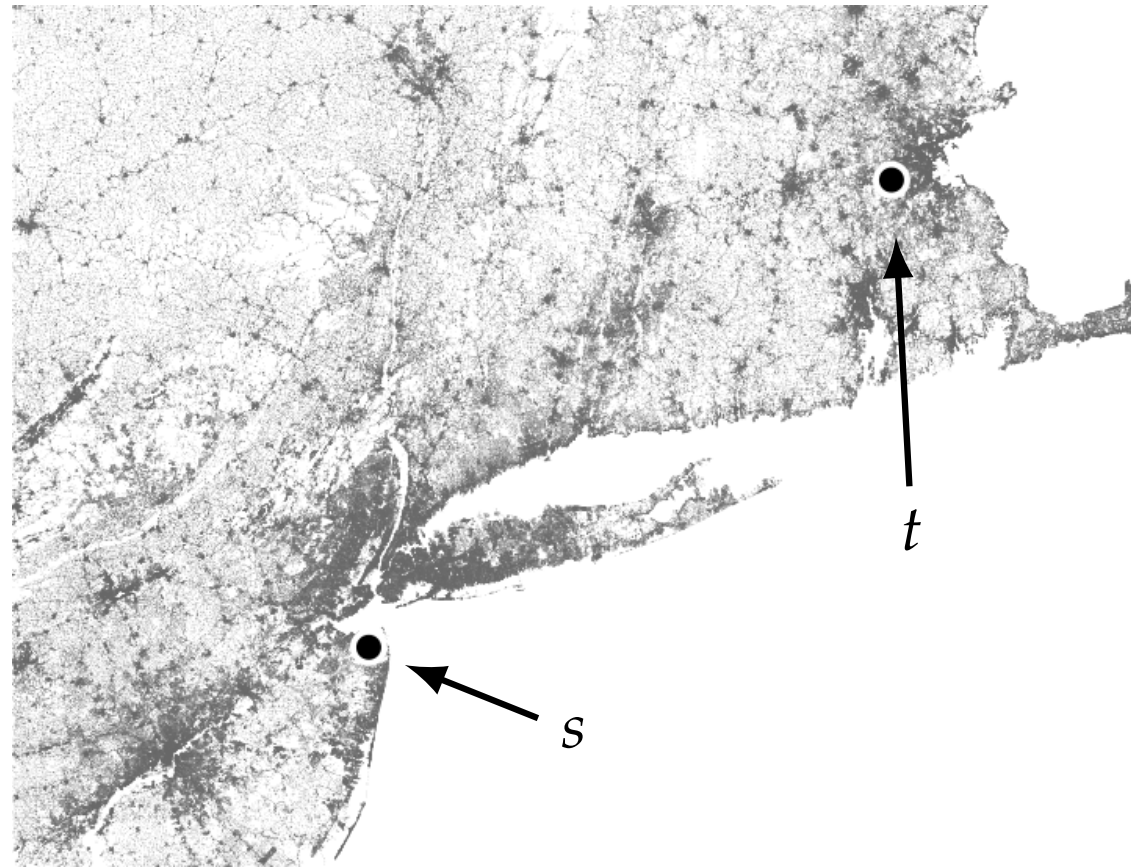
Narrow passage: Bias sampling

# So far we have looked at an Explicit Evaluated Graph

Is it a good idea to evaluate **every edge** that we discover?

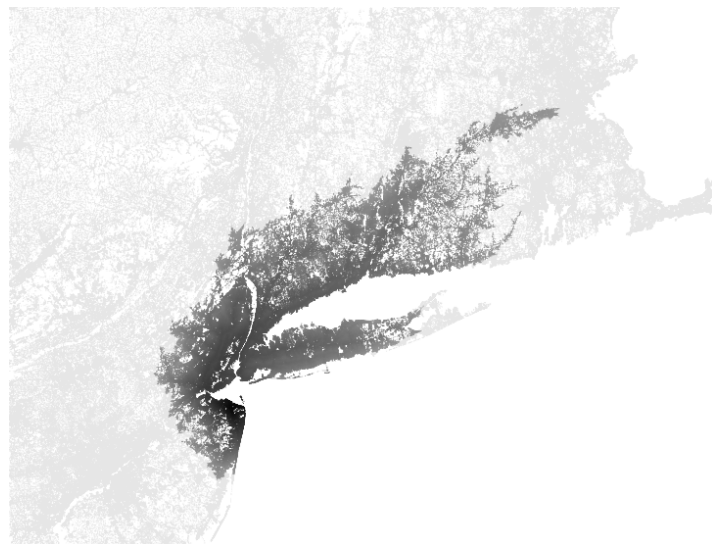
Is it a good idea to **explicitly store** the entire graph in memory?

# Do we need to store this whole graph?



DIMACS  
dataset

1.5 million vertices,  
3.67 million edges



A\* search touches a small fraction



Bidirectional A\*

Key Idea:

Use **implicit unevaluated** graphs

# Implicit graph

1. Initialize with at least one seed vertex

E.g. The start state, the goal state, both (for bi-directional search)

2. Provide a **generative** function for **producing successors**

$$\text{succ}(u) = \{(u, v) \mid (u, v) \in E\}$$

function returns both

- new vertices to add to  $V$
- new edges to add to  $E$

# Searching with implicit graph

1. Start the search with initial vertex (start, goal, or both)
2. Whenever search chooses to “expand a vertex”
  - call `succ( $u$ )`
  - get successor vertices
  - evaluate the edges
  - add these vertices to the search queue



# Explicit vs Implicit Graphs

- Explicit Graph:

- All **vertices** are identified individually and represented separately.
- All **edges** are identified individually and represented separately.

- Implicit Graph:

- Only a subset, possibly only one, of the vertices is given an explicit representation. (The others are implied.)
- Only a subset, and possibly zero, of the edges is given an explicit representation.
- A set of “operators” is provided that can be used to construct “new” edges and vertices.

# Planning with differential constraints



# Differential constraints

So far we assumed **only kinematic constraints**

$$q \notin \mathcal{C}_{obs}$$

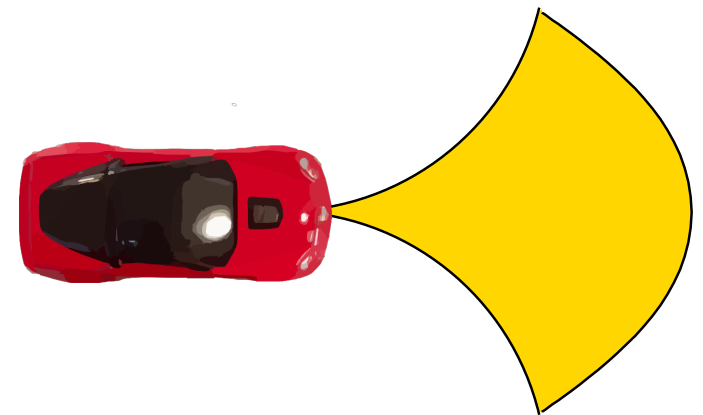
When is this assumption true?

- when controller can track any path
- when robots move very slowly (stop and turn)

# Differential constraints

We now introduce differential constraints

$$\dot{q} = f(q, u)$$



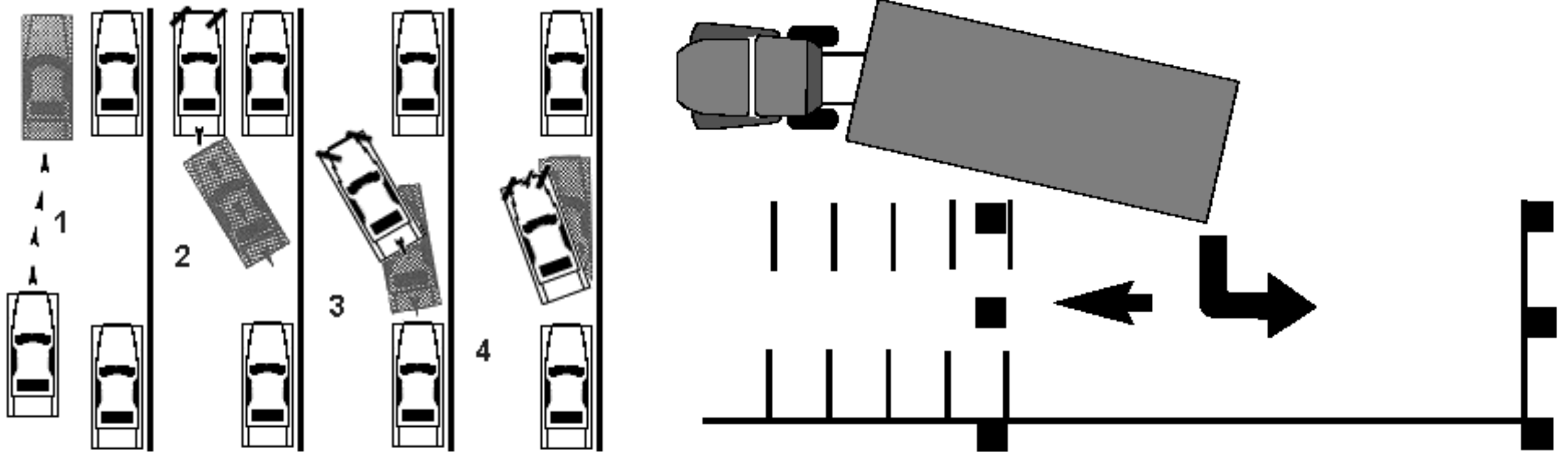
Two new terms:

1. Introduction of control space
2. Introduction of an equality constraint

# Differential constraints make things **even harder**

**Holonomic constraints:** Constraints that can be integrated, i.e.

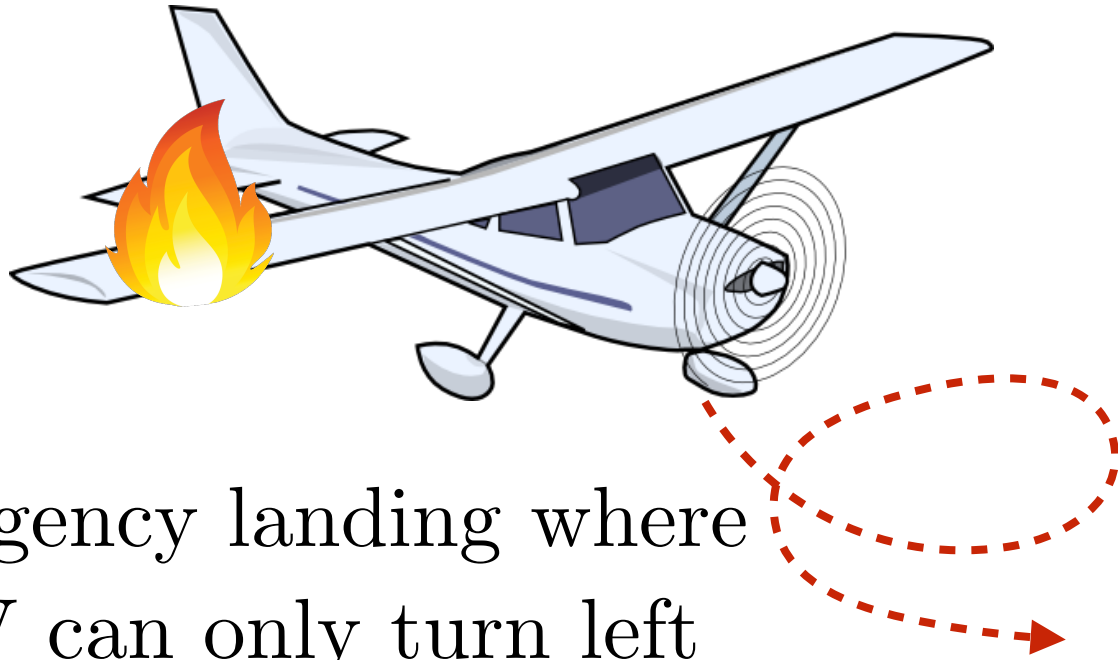
$$\dot{q} = f(q, u) \longrightarrow g(q, u) = 0$$



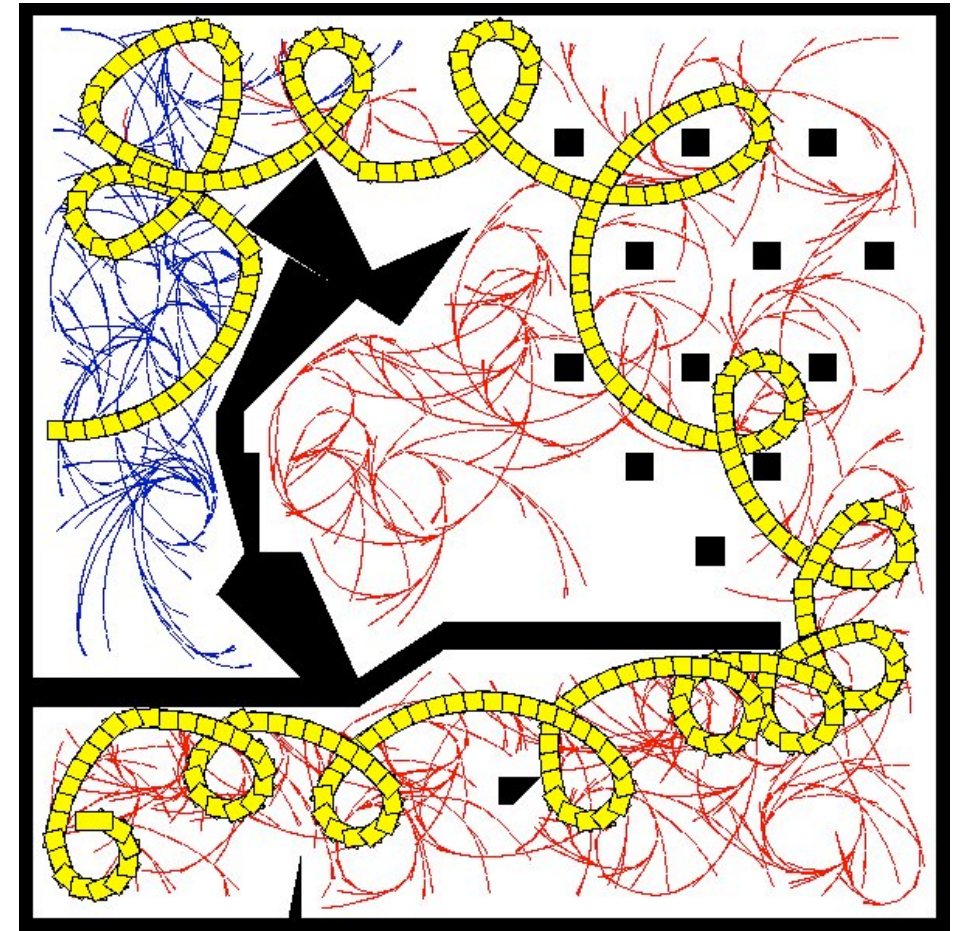
**Non-holonomic constraints:** Constraints that can't be integrated, i.e.  
i.e. the system is trapped in some sub-manifold of the config space



# Differential constraints make things **even harder**



Emergency landing where  
UAV can only turn left



“Left-turning-car”

**Non-holonomic constraints:** Constraints that can't be integrated, i.e.

i.e. the system is trapped in some sub-manifold of the config space

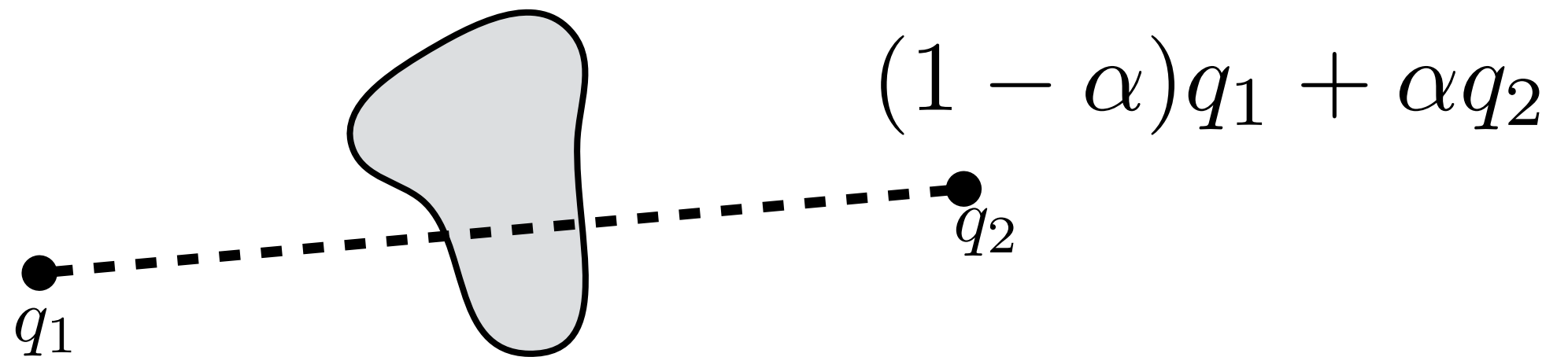
How do we incorporate differential constraints in our framework?

Only change the STEER function!

# Recap: STEER function for geometric

$$\text{steer}(q_1, q_2)$$

A steer function tries to join two configurations  
with a feasible path



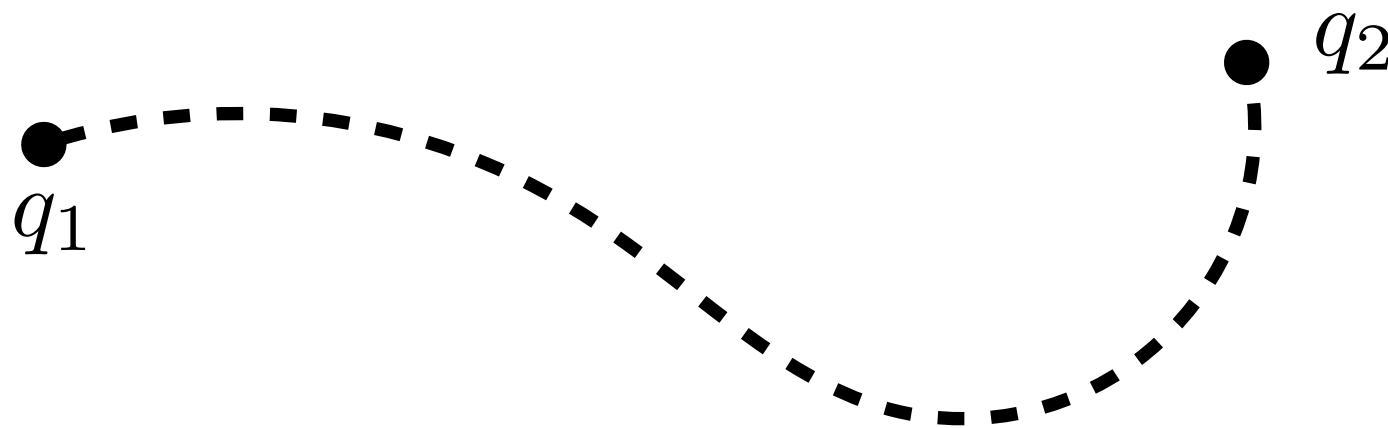
Example: Connect them with a straight line and check for feasibility



# STEER function incorporate dynamics!

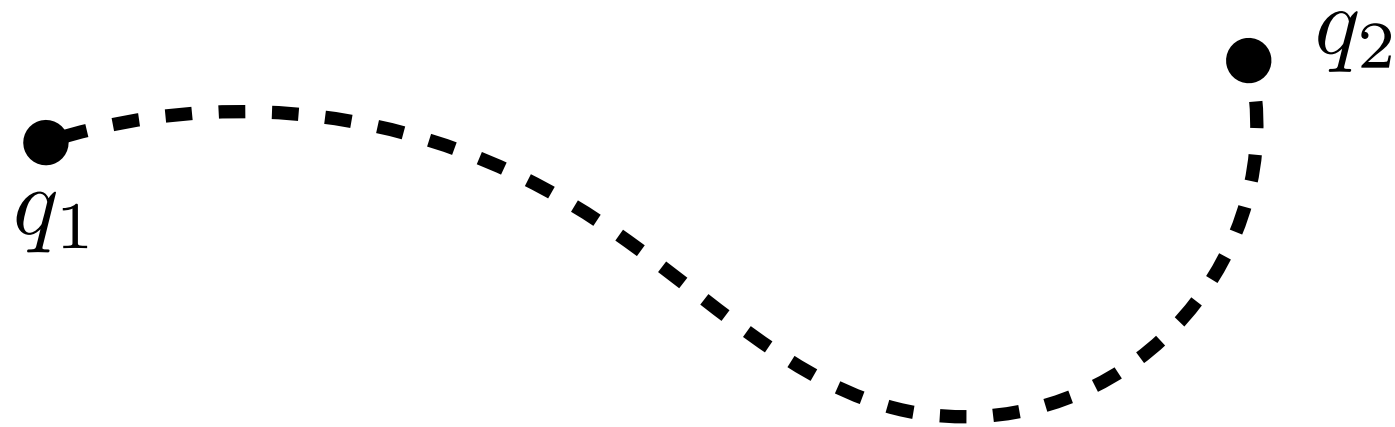
$$\text{steer}(q_1, q_2)$$

A steer function tries to join two configurations  
with a **dynamically** feasible path



# STEER function incorporate dynamics!

$\text{steer}(q_1, q_2)$



Formally called the **boundary value problem (BVP)**

Find a control trajectory  $u(t) \in U$

Such that  $q(0) = q_1$  ,  $q(t_f) = q_2$

$$\dot{q}(t) = f(q(t), u(t))$$

# How do we come up with STEER function?

There are three possible cases

Case 1: We can analytically solve the BVP :)



Case 2: We need to numerically solve the BVP.

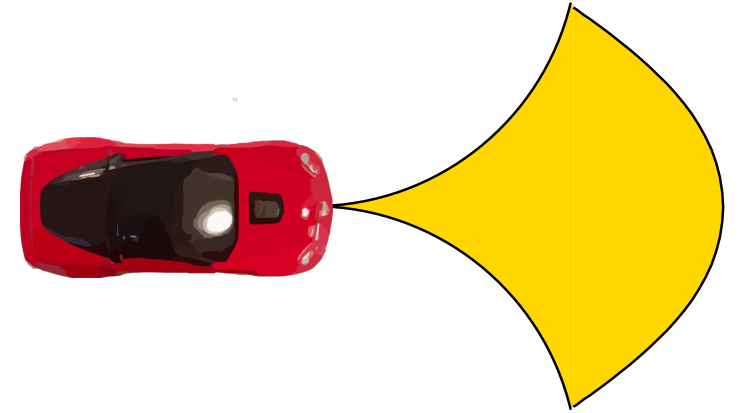


Case 3: We can't even solve the BVP!



# Case 1: We can analytically solve the BVP

Consider the dynamics of your racecar



$$\begin{bmatrix} \dot{x} \\ \dot{y} \\ \dot{\theta} \end{bmatrix} = \begin{bmatrix} V \cos \theta \\ V \sin \theta \\ C \tan \delta \end{bmatrix}$$

$$q_1 = (x_1, y_1, \theta_1)$$

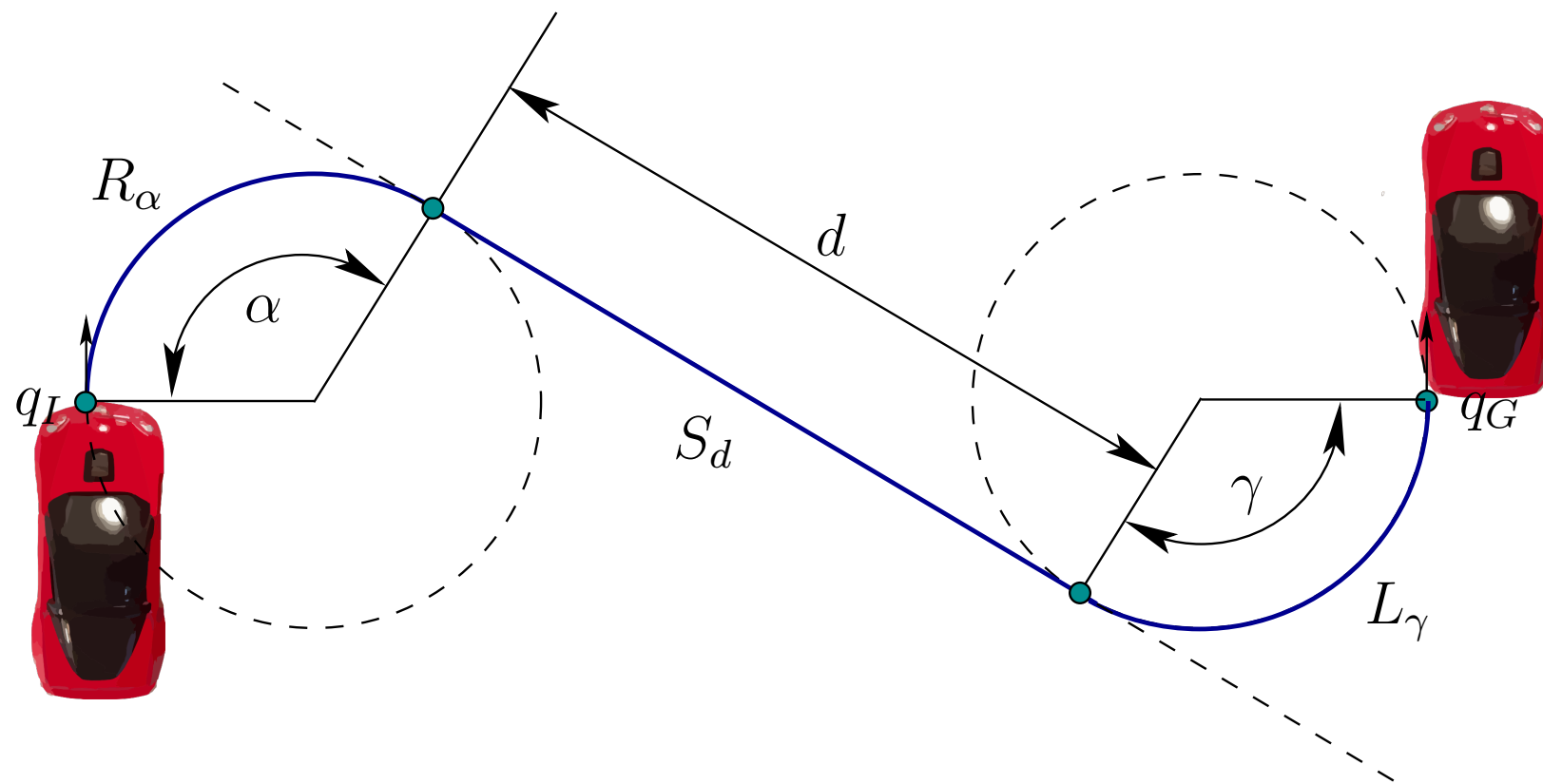
$$q_2 = (x_2, y_2, \theta_2)$$

$$|\delta| \leq \delta_{max}$$

Can we solve this **analytically**?

# Case 1: We can analytically solve the BVP

Yes! The solution is called the Dubins path



$$R_\alpha S_d L_\gamma$$

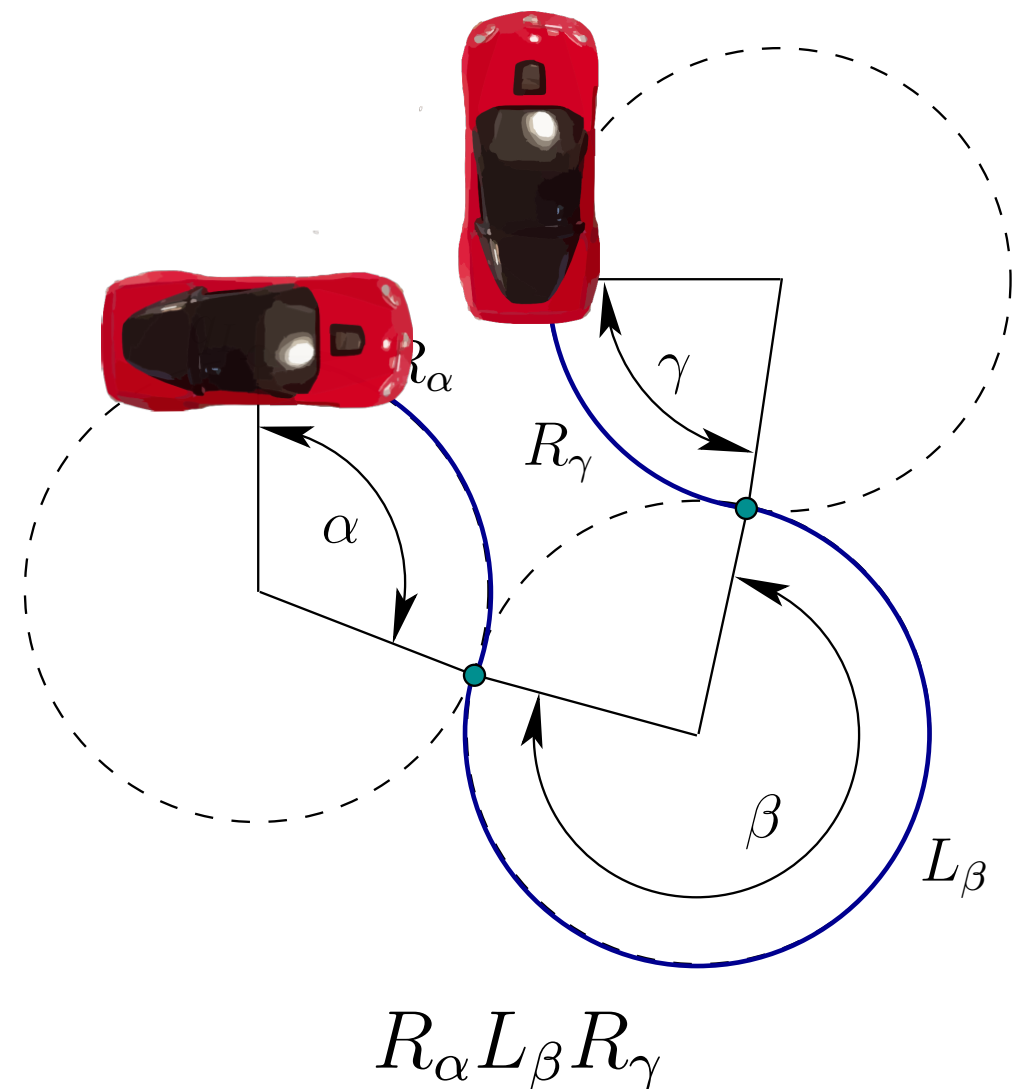
Right Straight Left

# Case 1: We can analytically solve the BVP

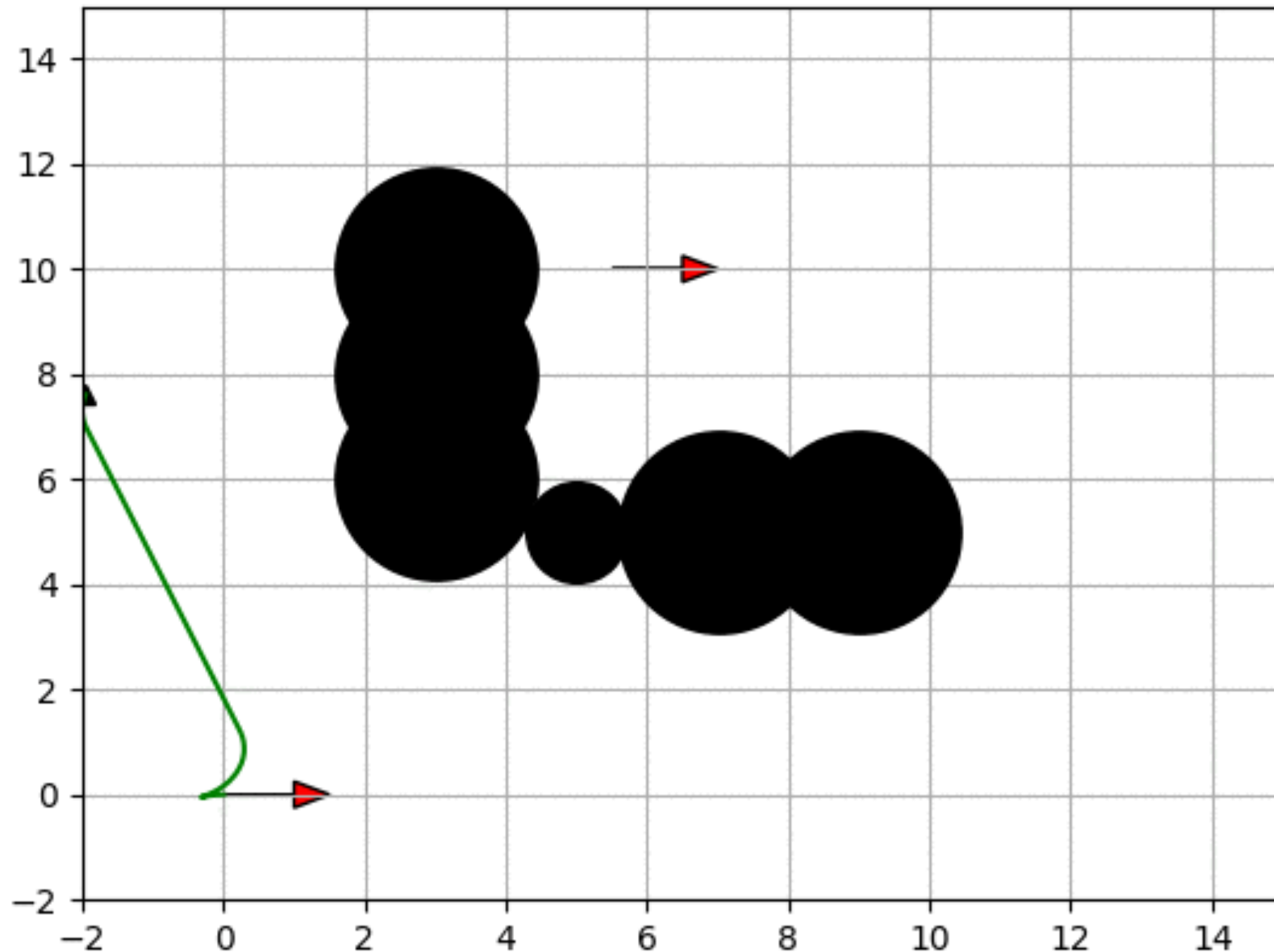
Dubins showed that ALL solutions had to be one of **6 classes**

$$\{LRL, RLR, LSL, LSR, RSL, RSR\}.$$

Hence, given a query,  
evaluate ALL 6 options,  
and pick the shortest one!



# Random sampling with Dubins steering



## Case 2: We need to numerically solve the BVP

What if you were changing the **steering rate**?

$$\begin{bmatrix} \dot{x} \\ \dot{y} \\ \dot{\theta} \\ \ddot{\theta} \end{bmatrix} = \begin{bmatrix} V \cos \theta \\ V \sin \theta \\ \dot{\theta} \\ u \end{bmatrix}$$

$$|\dot{\theta}| \leq C_1 \quad |u| \leq C_2$$

$$q_1 = (x_1, y_1, \theta_1, \dot{\theta}_1)$$

$$q_2 = (x_2, y_2, \theta_2, \dot{\theta}_2)$$

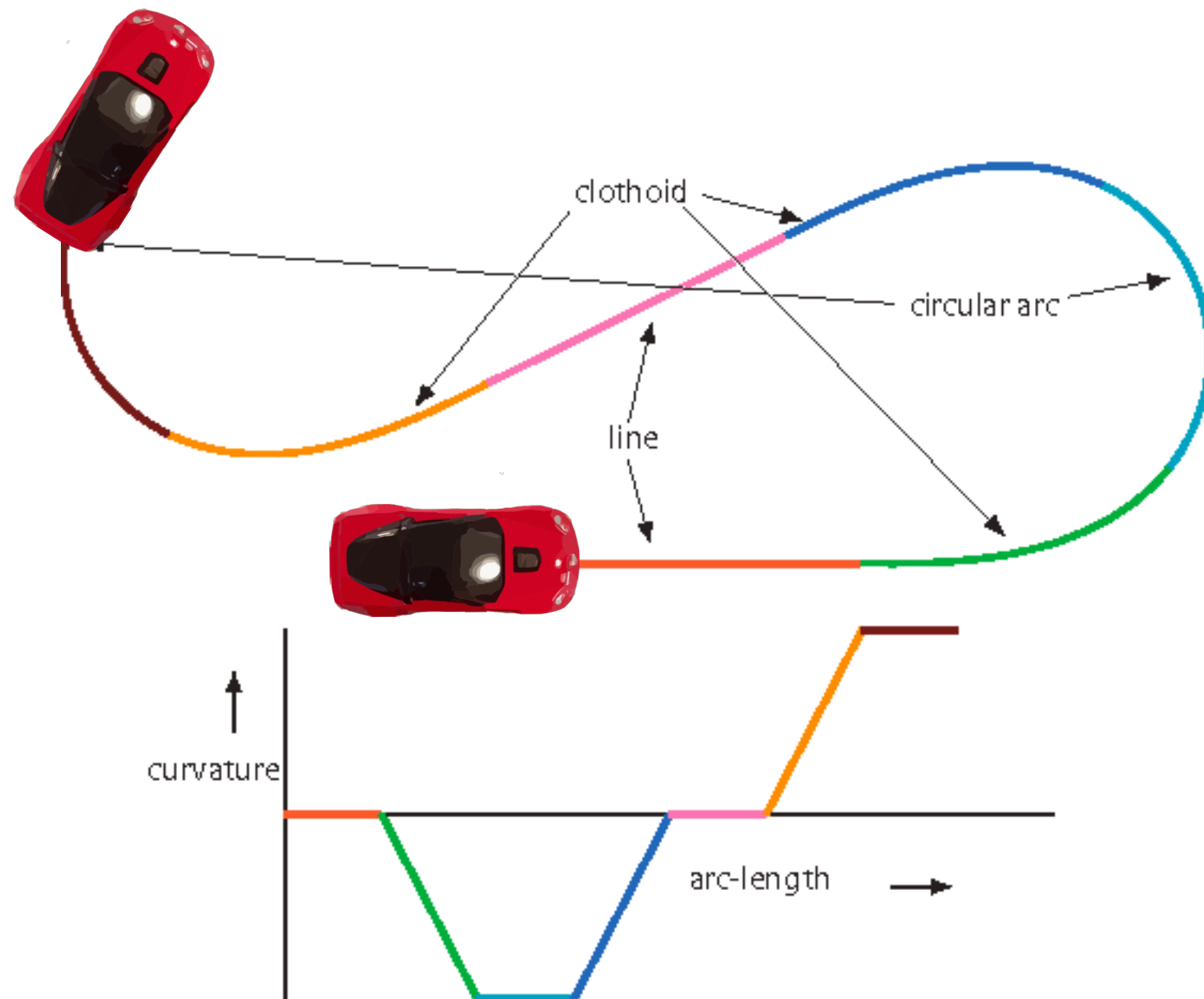
No longer called Dubins path, called **Clothoid path**

Can't solve analytically ... need numerical solutions





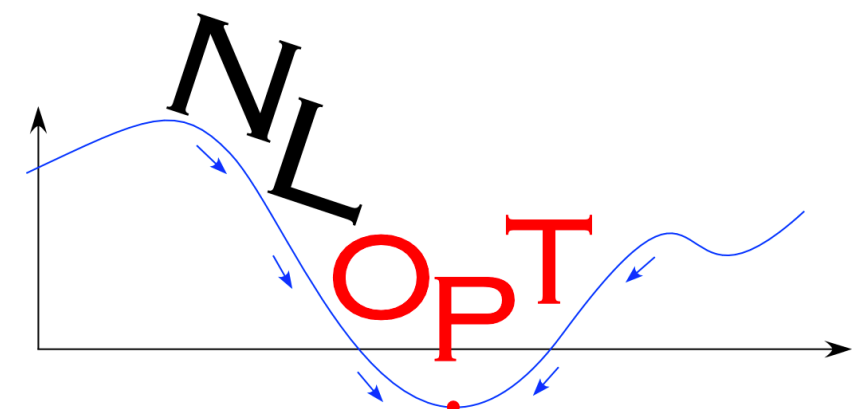
# Case 2: We need to numerically solve the BVP



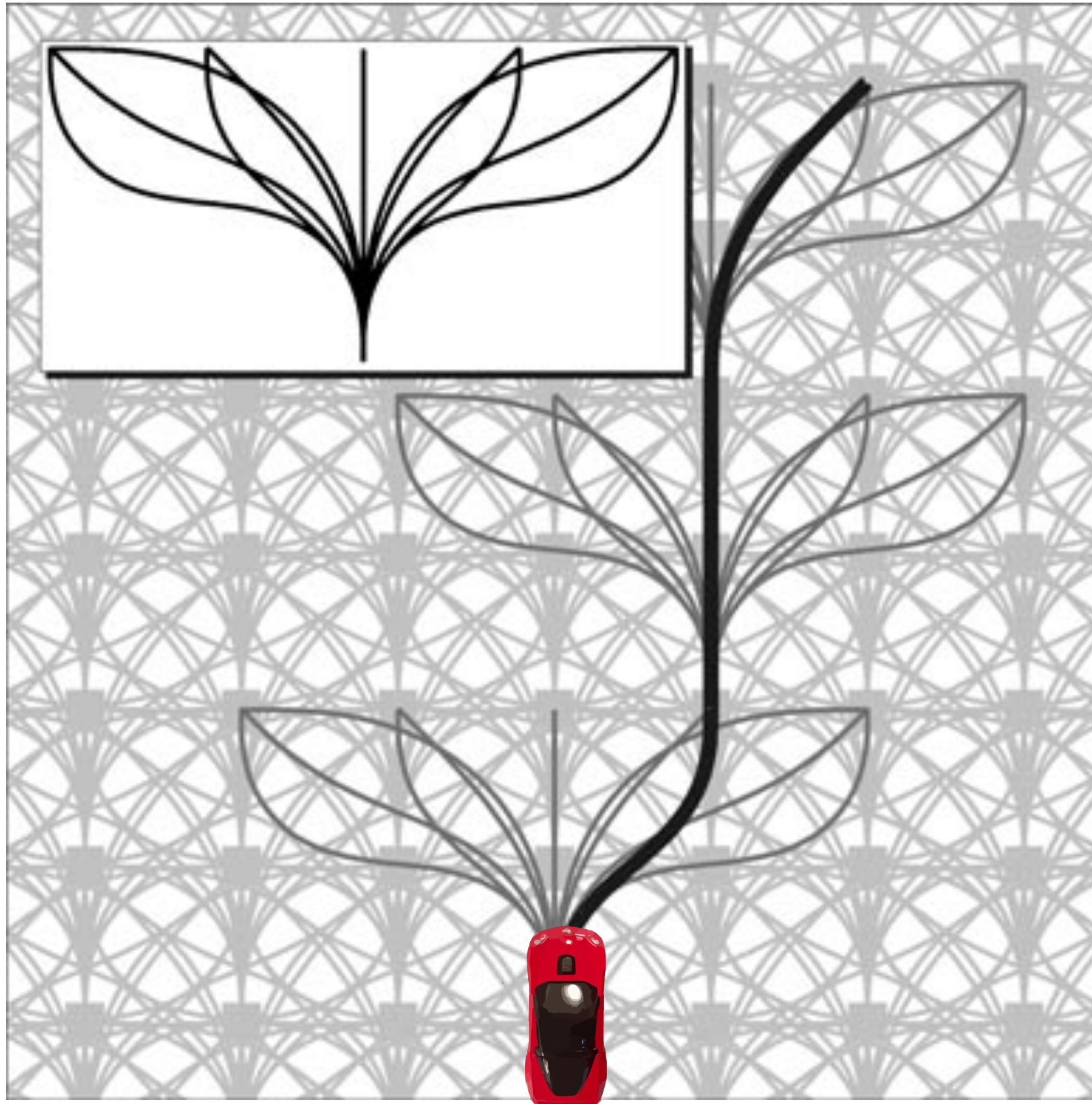
Create a non-linear optimization problem

Represent the path as a curvature polynomial

Enforce end point constraints



# Example of a state lattice



Pivtoraiko  
et al.  
2007

# Case 3: We can't even solve the BVP

Typically rules out roadmap methods  
because we cannot exactly connect 2 states

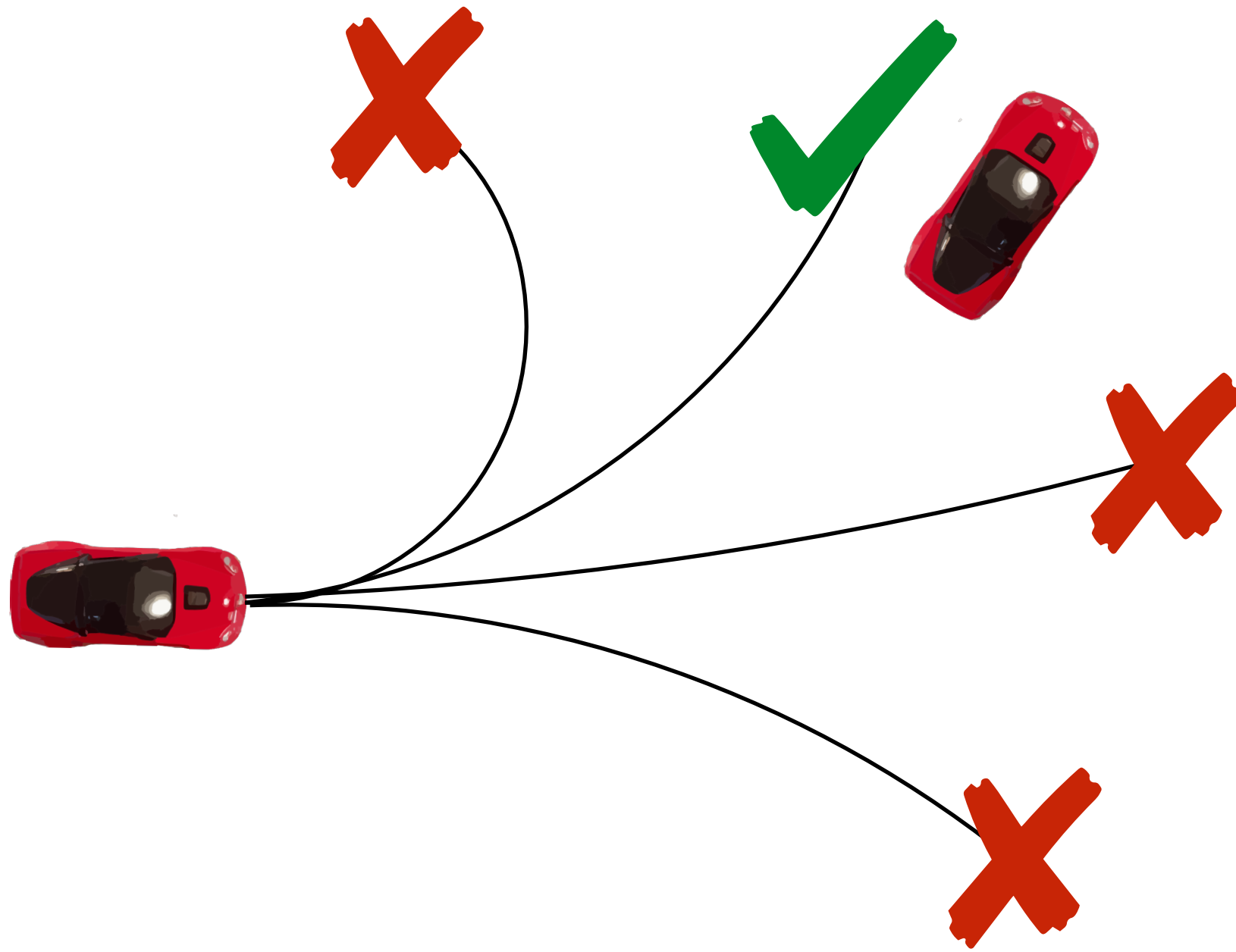


Why?

1. Dynamics too complicated - nonlinear optimization doesn't converge
2. Even if it did, too expensive to run online

# Case 3: We can't even solve the BVP

If we are working with **implicit graphs**, we don't need to exactly connect two states.



Pretend  
this was  
the state  
you wanted  
all along!

# Incremental Sampling-Based Planners

- Multiple-query methods such as PRM valuable in highly structured environments.
- Online planning does not require multiple queries
- Computing roadmap a priori computationally challenging/infeasible.
- Incremental sampling-based planning algorithms single-query counterpart to PRMs.
- Key idea:
  - Avoid necessity to set number of samples *a priori*;
  - return solution as soon as set of trajectories is rich enough.

Karaman, S., & Frazzoli, E. (2011), “Sampling-based algorithms for optimal motion planning”, The International Journal of Robotics Research, 30(7), 846–894. <https://doi.org/10.1177/0278364911406761>



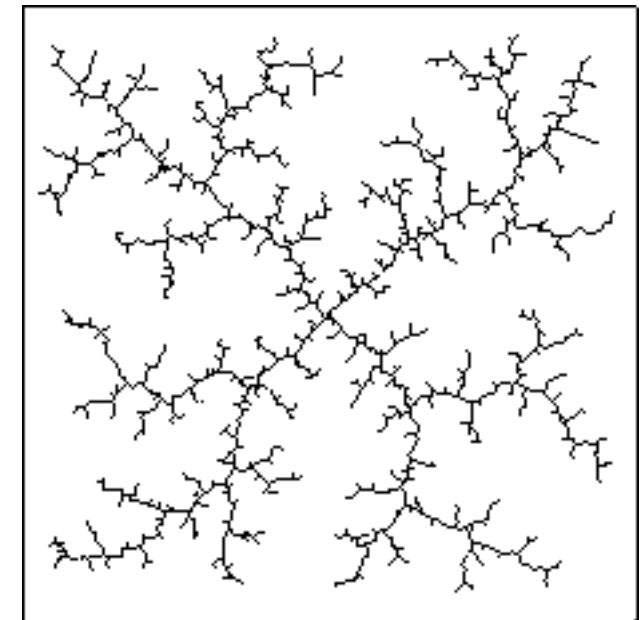
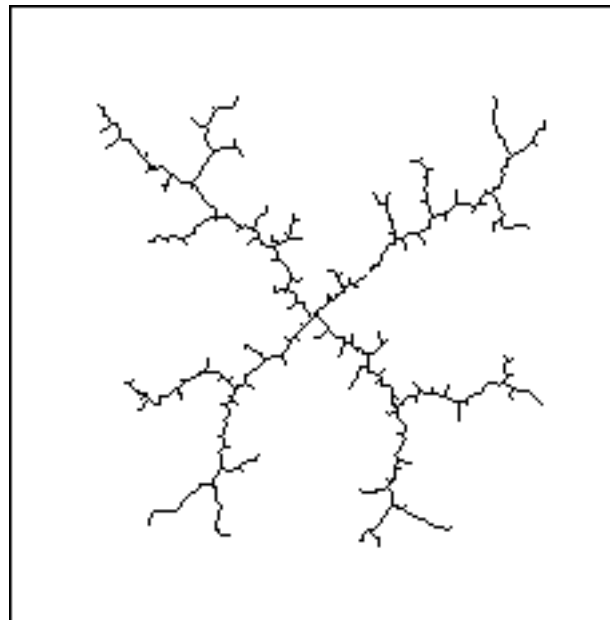
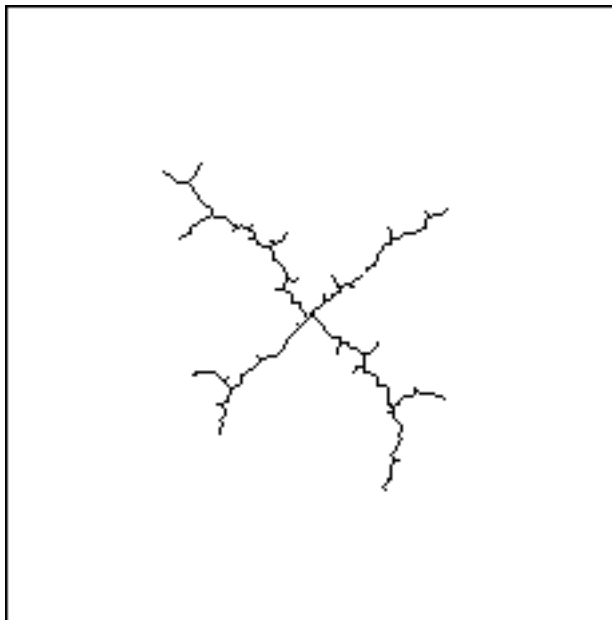
# Rapidly Exploring Random Trees (RRTs)



<http://msl.cs.uiuc.edu/rrt>

# RRTs

- Build a tree by iteratively connecting *next* states via the execution of random controls
- Carefully sample controls to ensure good coverage



# The RRT Algorithm

GENERATE\_RRT( $x_{init}, K, \Delta t$ )

```
1   $\mathcal{T}.\text{init}(x_{init});$   
2  for  $k = 1$  to  $K$  do  
3       $x_{rand} \leftarrow \text{RANDOM\_STATE}();$   
4       $x_{near} \leftarrow \text{NEAREST\_NEIGHBOR}(x_{rand}, \mathcal{T});$   
5       $u \leftarrow \text{SELECT\_INPUT}(x_{rand}, x_{near});$   
6       $x_{new} \leftarrow \text{NEW\_STATE}(x_{near}, u, \Delta t);$   
7       $\mathcal{T}.\text{add\_vertex}(x_{new});$   
8       $\mathcal{T}.\text{add\_edge}(x_{near}, x_{new}, u);$   
9  Return  $\mathcal{T}$ 
```

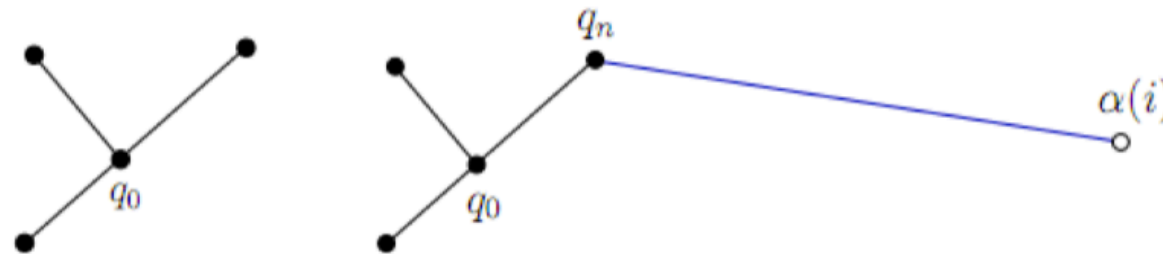
Probabilistically Complete

Exponential rate of decay for the probability of failure

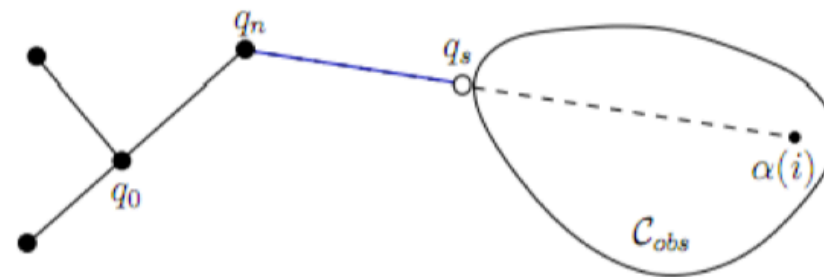


# Extend Function

- No obstacles, holonomic:



- With obstacles, holonomic:

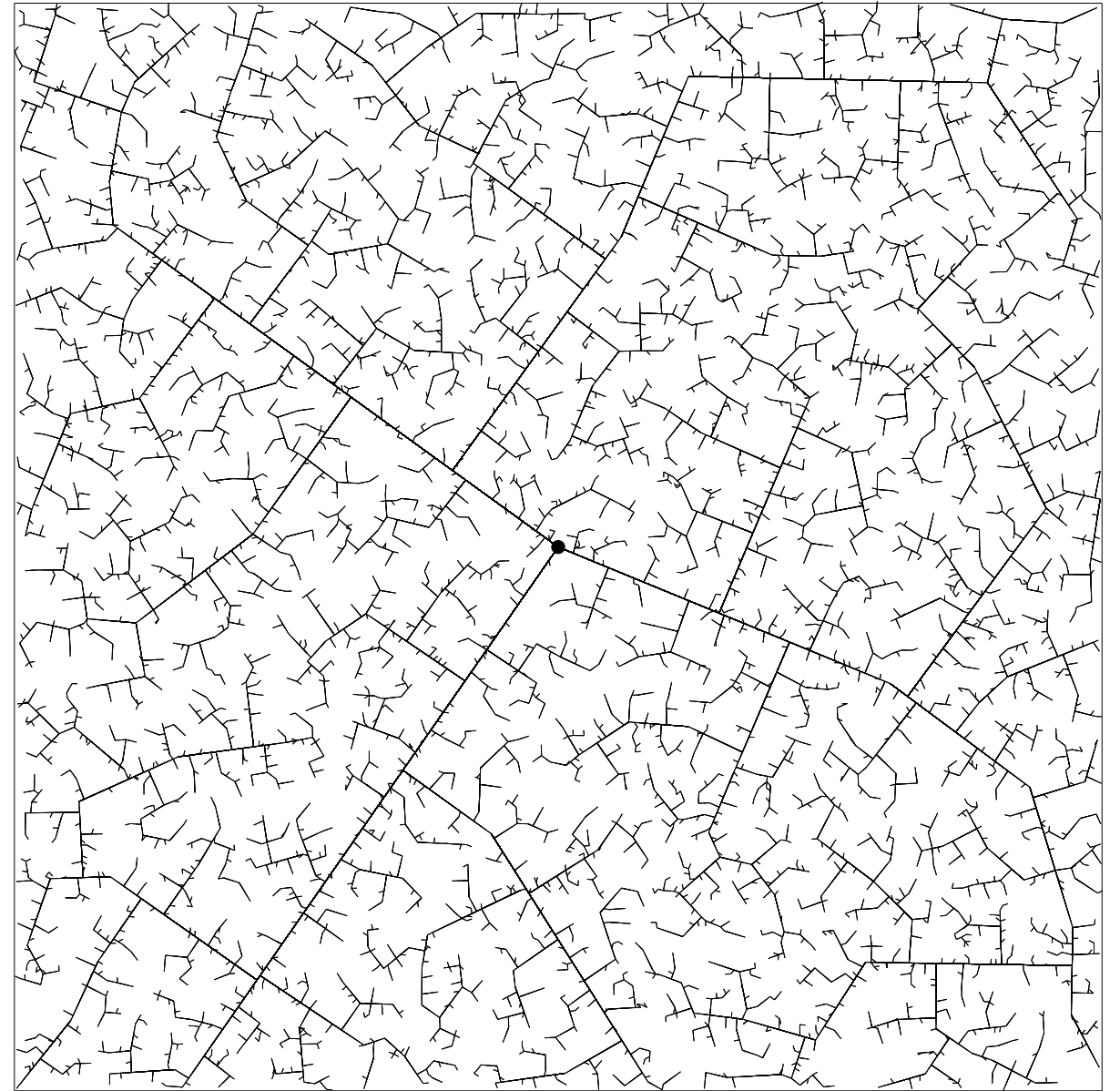


- Non-holonomic: approximately (sometimes as approximate as picking best of a few random control sequences) solve two-point boundary value problem

# RRT Expansion



45 iterations



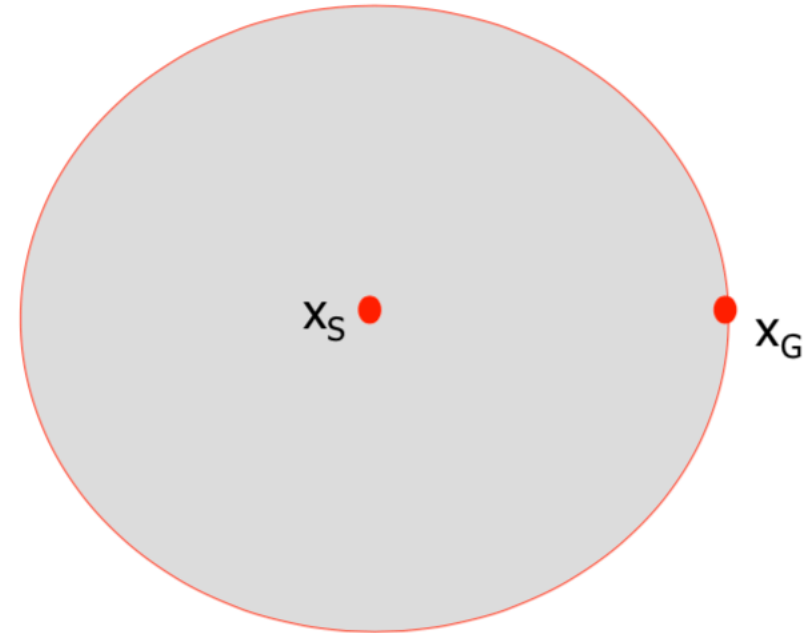
2345 iterations

# Notable variation: Bidirectional RRT

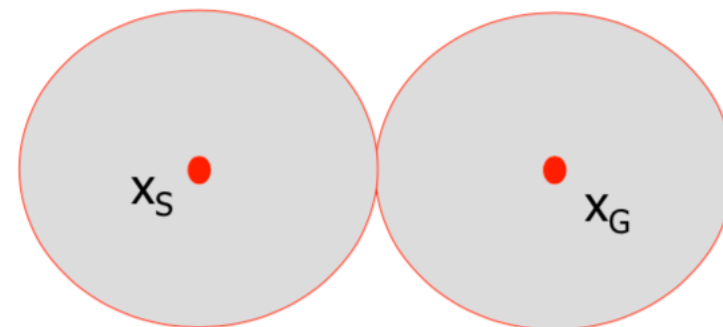
Key idea:

Grow 2 trees, 1 starting from start, 1 from goal, connect them, done.

- Volume swept out by unidirectional RRT:



- Volume swept out by bi-directional RRT:



- Difference becomes even more pronounced in higher dimensions

# RRT\*

- Asymptotically optimal
- Probabilistically complete
- Computationally efficient

## Algorithm 6: RRT\*

```
1  $V \leftarrow \{x_{\text{init}}\}; E \leftarrow \emptyset;$ 
2 for  $i = 1, \dots, n$  do
3    $x_{\text{rand}} \leftarrow \text{SampleFree}_i;$ 
4    $x_{\text{nearest}} \leftarrow \text{Nearest}(G = (V, E), x_{\text{rand}});$ 
5    $x_{\text{new}} \leftarrow \text{Steer}(x_{\text{nearest}}, x_{\text{rand}});$ 
6   if  $\text{ObstacleFree}(x_{\text{nearest}}, x_{\text{new}})$  then
7      $X_{\text{near}} \leftarrow \text{Near}(G = (V, E), x_{\text{new}}, \min\{\gamma_{\text{RRT}^*}(\log(\text{card}(V))/\text{card}(V))^{1/d}, \eta\});$ 
8      $V \leftarrow V \cup \{x_{\text{new}}\};$ 
9      $x_{\text{min}} \leftarrow x_{\text{nearest}}; c_{\text{min}} \leftarrow \text{Cost}(x_{\text{nearest}}) + c(\text{Line}(x_{\text{nearest}}, x_{\text{new}}));$ 
10    foreach  $x_{\text{near}} \in X_{\text{near}}$  do // Connect along a minimum-cost path
11      if  $\text{CollisionFree}(x_{\text{near}}, x_{\text{new}}) \wedge \text{Cost}(x_{\text{near}}) + c(\text{Line}(x_{\text{near}}, x_{\text{new}})) < c_{\text{min}}$  then
12         $x_{\text{min}} \leftarrow x_{\text{near}}; c_{\text{min}} \leftarrow \text{Cost}(x_{\text{near}}) + c(\text{Line}(x_{\text{near}}, x_{\text{new}}))$ 
13     $E \leftarrow E \cup \{(x_{\text{min}}, x_{\text{new}})\};$ 
14    foreach  $x_{\text{near}} \in X_{\text{near}}$  do // Rewire the tree
15      if  $\text{CollisionFree}(x_{\text{new}}, x_{\text{near}}) \wedge \text{Cost}(x_{\text{new}}) + c(\text{Line}(x_{\text{new}}, x_{\text{near}})) < \text{Cost}(x_{\text{near}})$ 
16        then  $x_{\text{parent}} \leftarrow \text{Parent}(x_{\text{near}});$ 
17         $E \leftarrow (E \setminus \{(x_{\text{parent}}, x_{\text{near}})\}) \cup \{(x_{\text{new}}, x_{\text{near}})\}$ 
17 return  $G = (V, E);$ 
```

Karaman, S., & Frazzoli, E. (2011), “Sampling-based algorithms for optimal motion planning”, The International Journal of Robotics Research, 30(7), 846–894. <https://doi.org/10.1177/0278364911406761>

# Smoothing

Paths extracted from sampling-based motion planners tend to be jerky.

Remedies:

1. Shortcutting

Along the found path, pick two vertices  $x_{t1}$ ,  $x_{t2}$  and try to connect them directly (skipping over all intermediate vertices)

2. Nonlinear optimization for optimal control

Allows to specify an objective function that includes smoothness in state, control, small control inputs, etc.