

Clearer Definition

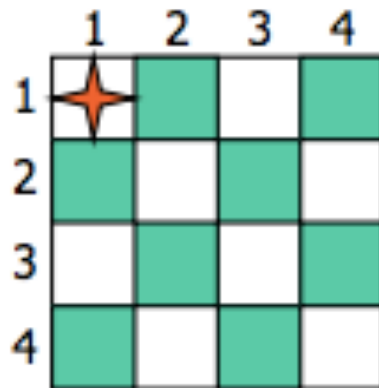
Definition: Arc consistency

- A constraint C_{xy} is said to *be arc consistent* wrt x if for each value v of x there is an allowed value of y
- Similarly, we define that C_{xy} is arc consistent wrt y
- A binary CSP is arc consistent iff every constraint C_{xy} is arc consistent wrt x as well as y
- When a CSP is not arc consistent, we can make it arc consistent, e.g. by using AC3
 - This is also called “enforcing arc consistency”

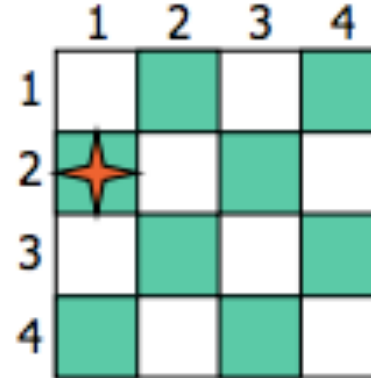
x is the tail

Chess as a CSP

Let's define the 4-queens problem as a CSP with the variable X_i denoting the position (row) of the queen on column i .



$X_1 = 1$



$X_1 = 2$

Remember the constraints: two queens attack each other when they are in the same row, the same column or on the same diagonal. We want to place $n=4$ queens on the board so no queen is attacking another.

CSP Challenge Question 1

Suppose we set $X_1 = 1$

Show the effect of forward checking on the domains of the remaining variables

(I suggest crossing off values in the lists below:)

	1	2	3	4
1	★			
2				
3				
4				

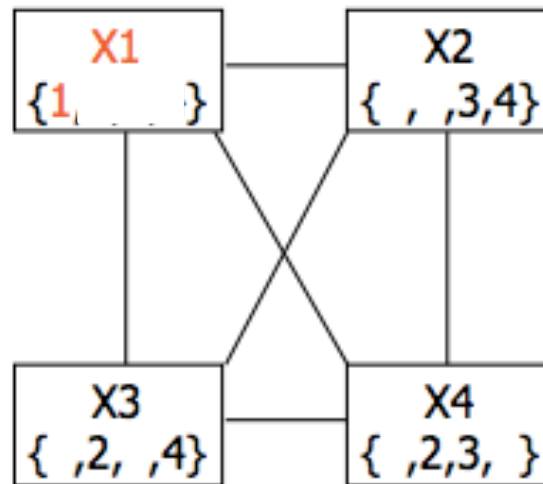
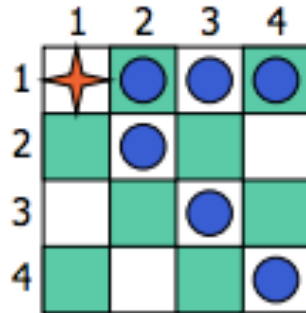
$X_1 = 1$

Domain $X_2 = \{1, 2, 3, 4\}$

Domain $X_3 = \{1, 2, 3, 4\}$

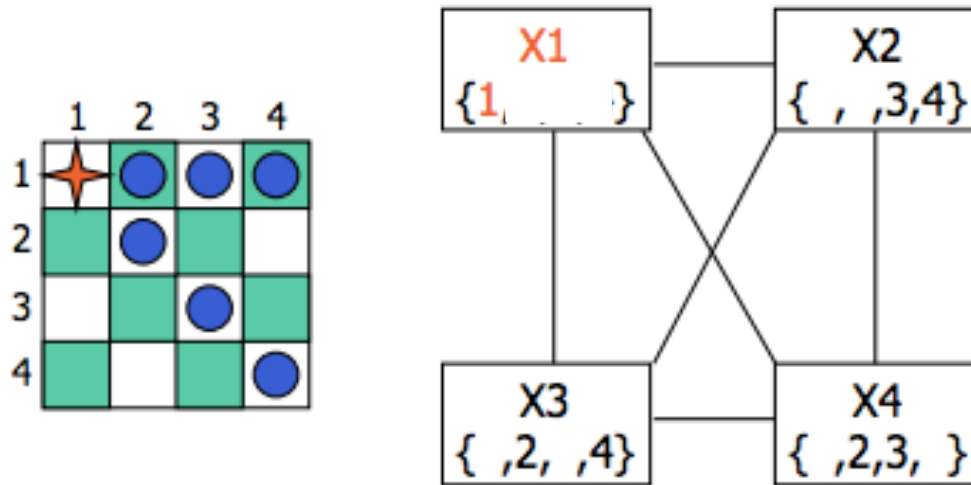
Domain $X_4 = \{1, 2, 3, 4\}$

Answer 1



Forward checking will delete values from the domains of all other variables, as shown

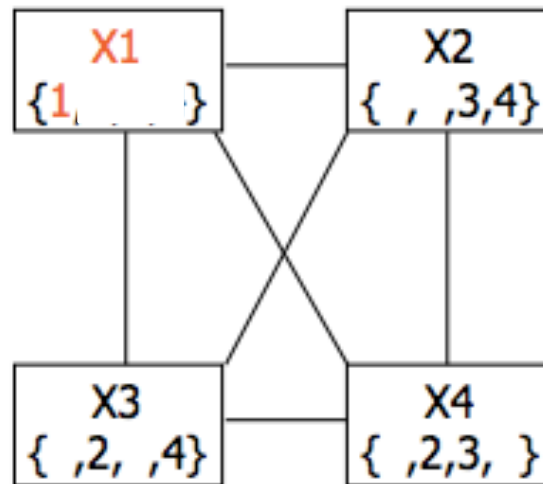
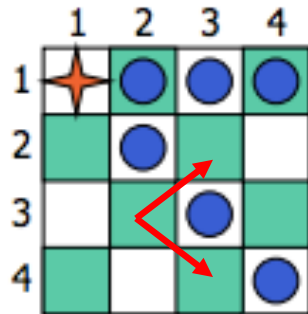
Question 2



Is this CSP now arc consistent?

(for the purposes of this question – assume that there is one constraint between each pair of queens that rules out all attacks)

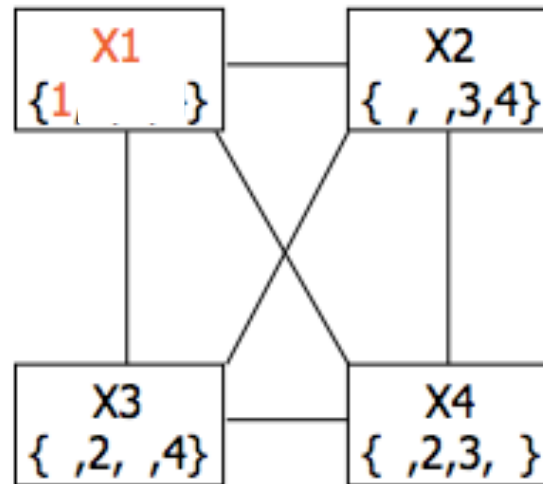
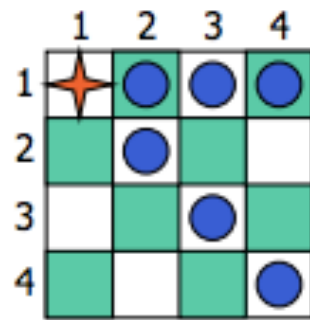
Answer 2



No, the constraint between X2 and X3 is not consistent with respect to X2. There exists a value in the domain of X2 (specifically X2=3) such that NO value for X3 will work.

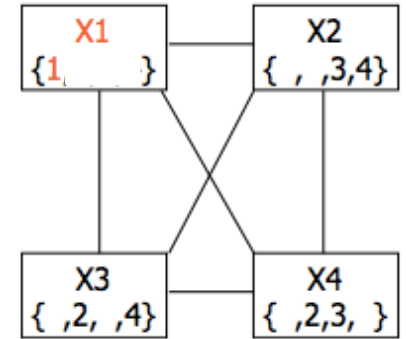
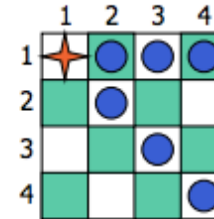
Furthermore, the constraint between X4 and X3 is not consistent with respect to X4, because X4=3 also leaves X3 with no legal values.

Question 3



Simulate the behavior of AC3 to make the CSP arc consistent
First subquestion, what goes on the queue?

Answer 4



For each pair of variables, you need to put a directed constraint.

I'll write $X2 \rightarrow X3$ to mean the constraint wrt $X2$ (ie $X2$ is the tail)

For this example, let's ignore constraints with $X1$ because those constraints are consistent (as a result of forward checking) and can't become inconsistent because we've chose a single value for $X1$.

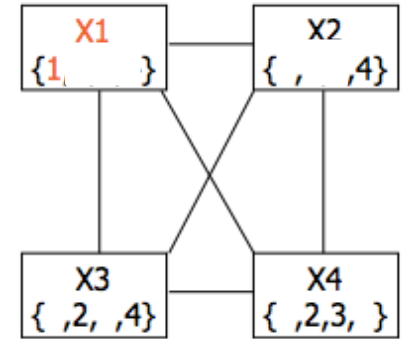
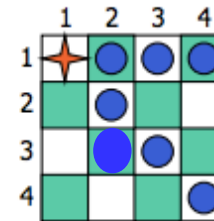
So the queue might be

$\langle X2 \rightarrow X3, X2 \rightarrow X4, X3 \rightarrow X2, X3 \rightarrow X4, X4 \rightarrow X2, X4 \rightarrow X3 \rangle$

We've already established that $X2 \rightarrow X3$ is inconsistent.

What does AC-3 do to fix this?

Answer 5



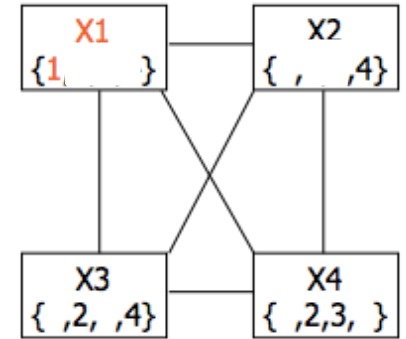
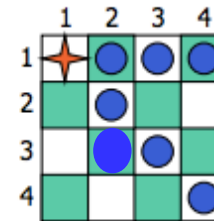
AC-3 deletes from the domain of ... X2..

So now $\text{Domain}(X2) = \{4\}$

AC-3 also adds some more constraints onto the queue, $X3 \rightarrow X2$ and $X4 \rightarrow X2$, but since they are already there there is no change. So the queue is $\langle X2 \rightarrow X4, X3 \rightarrow X2, X3 \rightarrow X4, X4 \rightarrow X2, X4 \rightarrow X3 \rangle$

Is $X2 \rightarrow X4$ consistent?

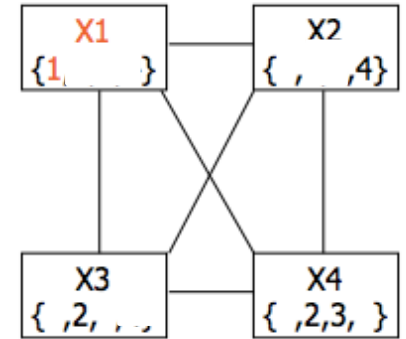
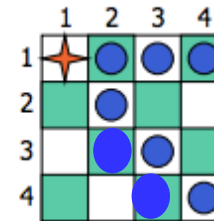
Answer 6



Yep. Now the queue is
<X3→X2, X3→X4, X4→X2, X4→X3>

Is X3→X2 consistent?

Answer 7

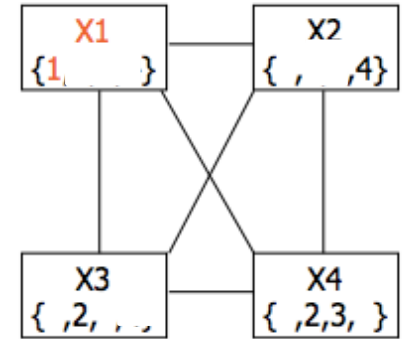
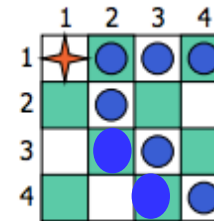


Nope. We need to delete $X3=4$.

That means we need to add some stuff to the queue.

So what's the queue become?

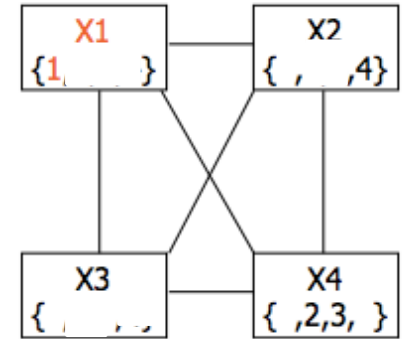
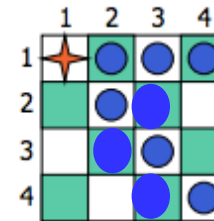
Answer 8



We add $X2 \rightarrow X3$ and $X4 \rightarrow X3$ but the latter was already there so we get $\langle X3 \rightarrow X4, X4 \rightarrow X2, X4 \rightarrow X3, X2 \rightarrow X3 \rangle$

Now what happens when we process the next constraint?

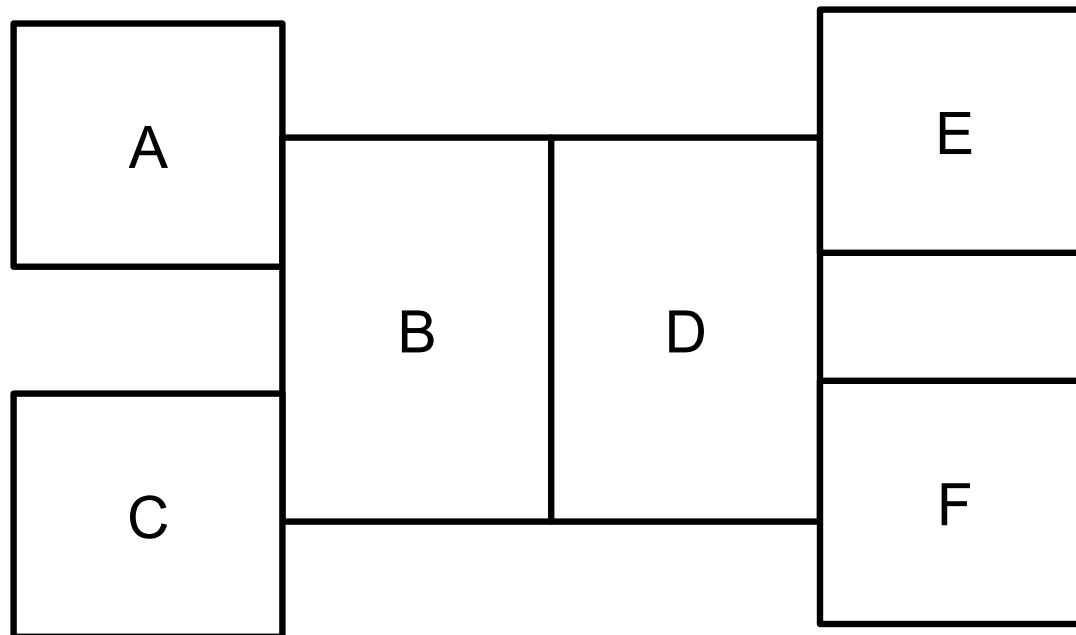
Answer 9



$X3 \rightarrow X4$ is inconsistent so we need to remove $X3=2$, but now $X3$'s domain is empty, which means that the CSP is unsolvable. So the very first decision to Assign $X1=1$ was a mistake.

In fact, following the pseudocode, AC3 will keep running and remove some more stuff – a bit pointless. But I'll stop here.

Part II – Tree structured CSPs



Let's color this!

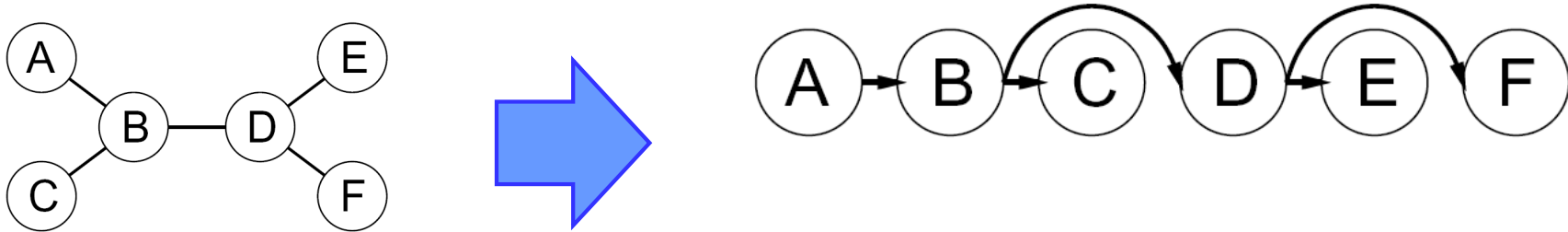
Tree-Structured CSPs

- Algorithm for tree-structured CSPs:
 1. Order: Choose a root variable, order variables so that parents precede children
 2. Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$
 3. Assign forward: For $i = 1 : n$, assign X_i consistently with $\text{Parent}(X_i)$

Tree-Structured CSPs

- Algorithm for tree-structured CSPs:

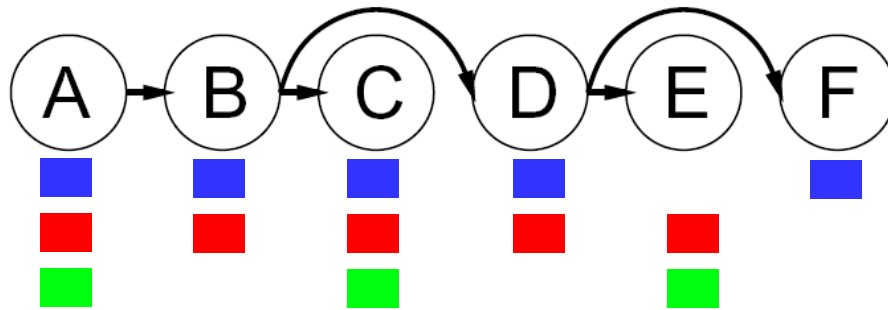
1. Order: Choose a root variable, order variables so that parents precede children



My choice to start with A as the root is arbitrary – could have started with anything else. It also doesn't matter if B comes before C in the ordering etc.

Question 10

- Algorithm for tree-structured CSPs:
 - Order: Choose a root variable, order variables so that parents precede children
 - Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$



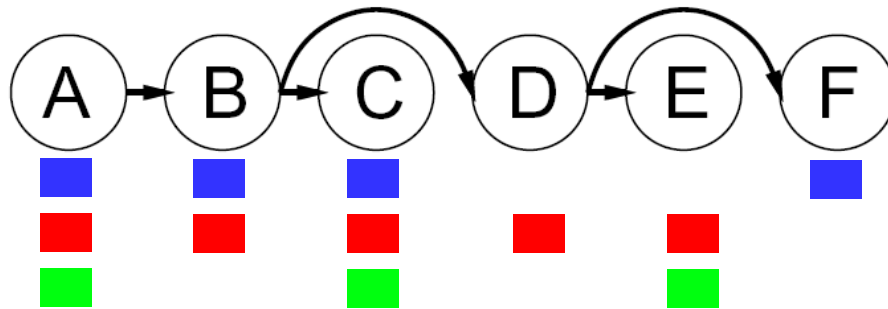
Suppose that the initial legal colors are as I show above

Simulate step 2 of the algorithm (I suggest cross off colors in the diagram above)

Answer 10

- Algorithm for tree-structured CSPs:

1. Order: Choose a root variable, order variables so that parents precede children
2. Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$



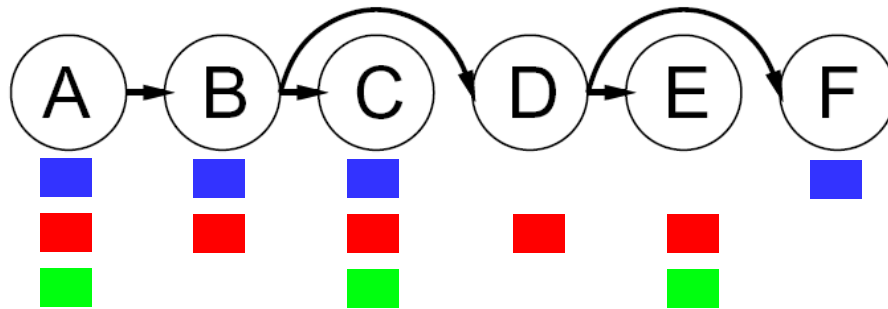
When processing $D \rightarrow F$, we need to remove blue from the domain of D

What about when we process $D \rightarrow E$?

Answer 11

- Algorithm for tree-structured CSPs:

1. Order: Choose a root variable, order variables so that parents precede children
2. Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$



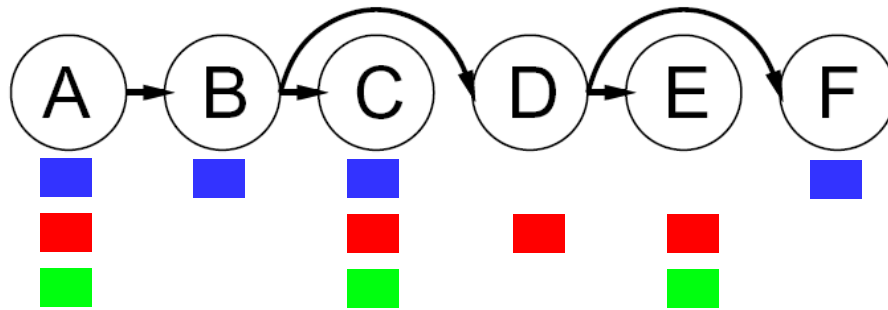
When processing $D \rightarrow E$, we don't do anything.

We would only remove something from the parent, D, but red is consistent, because we can make E green. So we just leave it as is.

What about $B \rightarrow D$?

Answer 12

- Algorithm for tree-structured CSPs:
 - Order: Choose a root variable, order variables so that parents precede children
 - Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$



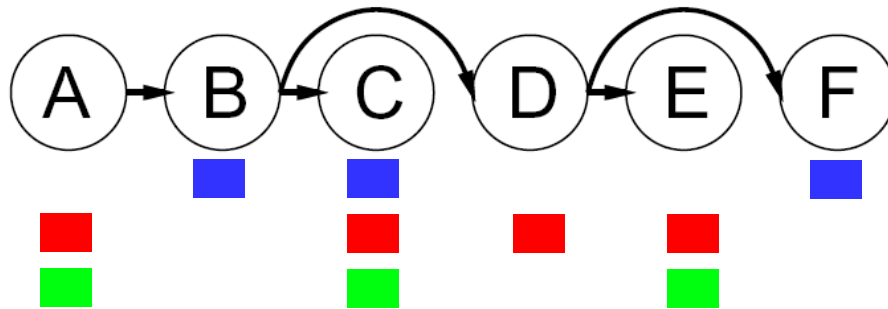
When processing $B \rightarrow C$, we don't do anything.

What about $A \rightarrow B$?

Answer 13

- Algorithm for tree-structured CSPs:

1. Order: Choose a root variable, order variables so that parents precede children
2. Remove backward: For $i = n : 2$, apply $\text{RemoveInconsistent}(\text{Parent}(X_i), X_i)$
3. Assign forward: For $i = 1 : n$, assign X_i consistently with $\text{Parent}(X_i)$



Right, we delete blue from A.

Now simulate step 3.

Any choice for A is ok. B will be blue. C can be red or green. D is red, E will be green... It all works!!