

CSE 473: Artificial Intelligence

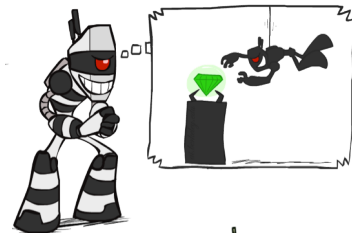
Constraint Satisfaction Problems II Factored (aka Structured) Search



[With many slides by Dan Klein and Pieter Abbeel (UC Berkeley) available at <http://ai.berkeley.edu>.]

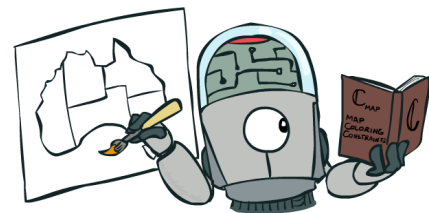
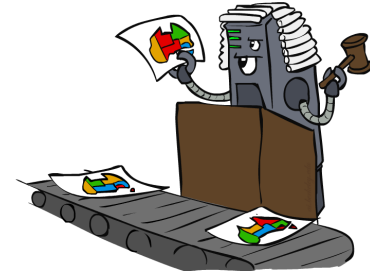
What is Search For?

- **Planning:** sequences of actions
 - The *path to the goal* is the important thing
 - Paths have various costs, depths
 - Assume little about problem structure
- **Identification:** assignments to variables
 - The *goal itself* is important, *not the path*
 - All paths at the same depth (for some formulations)



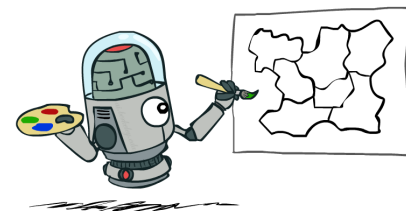
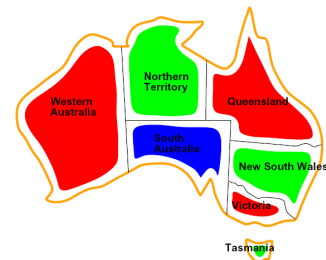
Constraint Satisfaction Problems

- **Standard search problems:**
 - State is a “black box”: arbitrary data structure
 - Goal test can be any function over states
 - Successor function can also be anything
- **Constraint satisfaction problems (CSPs):**
 - A special subset of search problems
 - State is defined by **variables X_i** with values from a **domain D** (sometimes D depends on i)
 - Goal test is a **set of constraints** specifying allowable combinations of values for subsets of variables
- **Making use of CSP formulation allows for optimized algorithms**
 - Typical example of trading generality for utility (in this case, speed)



CSP Example: Map Coloring

- **Variables:** WA, NT, Q, NSW, V, SA, T
- **Domains:** $D = \{\text{red, green, blue}\}$
- **Constraints: adjacent regions must have different colors**
 - Implicit: $WA \neq NT$
 - Explicit: $(WA, NT) \in \{(\text{red, green}), (\text{red, blue}), \dots\}$
- **Solutions are assignments satisfying all constraints, e.g.:**
 - $\{WA=\text{red}, NT=\text{green}, Q=\text{red}, NSW=\text{green}, V=\text{red}, SA=\text{blue}, T=\text{green}\}$



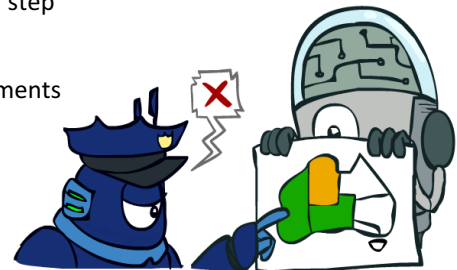
Systematic Search to Solve CSP

- **States** – partial assignments to variables
- **Operators** – assign another variable
- **Initial State** – no variables assigned
- **Goal State** – all vars assigned & constraints satisfied

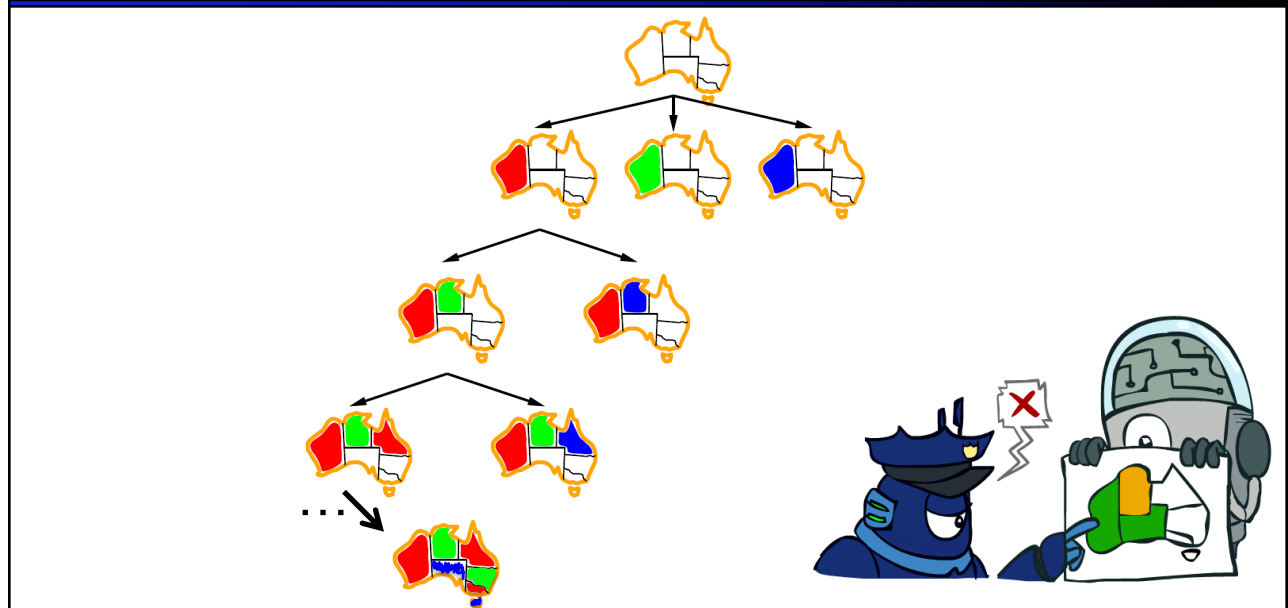
We'll improve this basic method to exploit structure

Backtracking Search

- Backtracking search is the basic uninformed algorithm for solving CSPs
- Start with Depth First Search
 - “backtracking search” **IS a Kind of** depth first search with these 2 details:
- **Idea 1: One variable at a time**
 - Variable assignments are commutative, so **fix ordering**
 - I.e., [WA = red then NT = green] same as [NT = green then WA = red]
 - Only need to consider assignments to a single variable at each step
- **Idea 2: Check constraints as you go**
 - I.e. consider only values which do not conflict previous assignments
 - Might have to do some computation to check the constraints
 - “Incremental goal test”
- Can solve n-queens for $n \approx 25$



Backtracking Example



Backtracking Search

```

function BACKTRACKING-SEARCH(csp) returns solution/failure
  return RECURSIVE-BACKTRACKING({ }, csp)

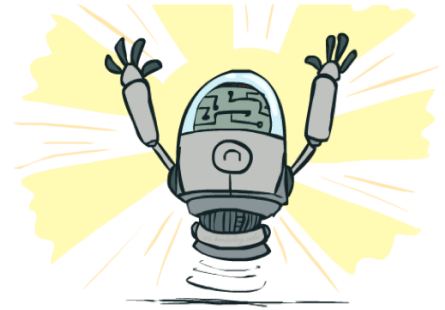
function RECURSIVE-BACKTRACKING(assignment, csp) returns soln/failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment given CONSTRAINTS[csp] then
      add {var = value} to assignment
      result ← RECURSIVE-BACKTRACKING(assignment, csp)
      if result ≠ failure then return result
      remove {var = value} from assignment
  return failure
    
```

- What are the choice points?

[Demo: coloring -- backtracking]

Improving Backtracking

- General-purpose ideas give huge gains in speed
- Ordering:
 - Which variable should be assigned next?
 - In what order should its values be tried?
- Filtering: Can we detect inevitable failure early?
- Structure: Can we exploit the problem structure?



Filtering



Filtering: Forward Checking

- Filtering: Keep track of domains for unassigned variables and cross off bad options
- Forward checking: Cross off values that violate a constraint when added to the existing assignment



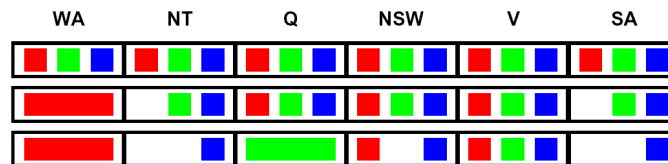
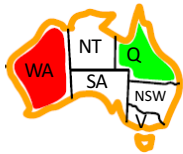
[Demo: coloring -- forward checking]

Video of Demo Coloring – Backtracking with Forward Checking



Filtering: Constraint Propagation

- Forward checking only propagates information from assigned to unassigned
- It doesn't catch when *two unassigned variables* have no consistent assignment:



- NT and SA cannot both be blue!
- Why didn't we detect this yet?
- Constraint propagation*: reason from constraint to constraint

Consistency of a Single Arc

- An arc $X \rightarrow Y$ is **consistent** iff for every x in the tail there is *some* y in the head which could be assigned without violating a constraint

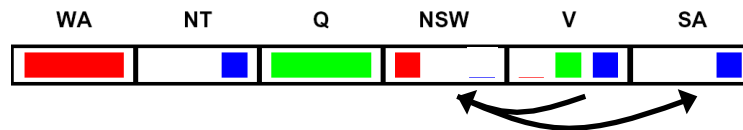
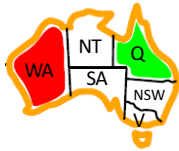


Delete from the tail!

- Forward checking: Enforcing consistency *of arcs pointing to each new assignment*

Arc Consistency of an Entire CSP

- A simple form of propagation makes sure **all** arcs are consistent:



- Important: If X loses a value, neighbors of X need to be rechecked!
- Arc consistency detects failure **earlier** than forward checking
- Can be run as a preprocessor **or** after each assignment
- What's the **downside** of enforcing arc consistency?

Remember:
Delete from the
tail!

AC-3 algorithm for Arc Consistency

```

function AC-3(csp) returns the CSP, possibly with reduced domains
inputs: csp, a binary CSP with variables  $\{X_1, X_2, \dots, X_n\}$ 
local variables: queue, a queue of arcs, initially all the arcs in csp

while queue is not empty do
     $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$ 
    if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) then
        for each  $X_k$  in NEIGHBORS[ $X_i$ ] do
            add  $(X_k, X_i)$  to queue



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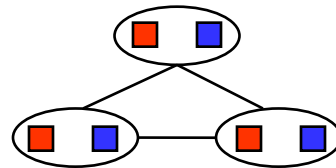
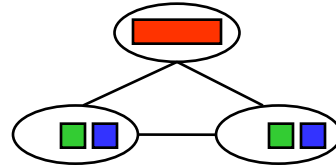

function REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) returns true iff succeeds
    removed  $\leftarrow$  false
    for each  $x$  in DOMAIN[ $X_i$ ] do
        if no value  $y$  in DOMAIN[ $X_j$ ] allows  $(x, y)$  to satisfy the constraint  $X_i \leftrightarrow X_j$ 
            then delete  $x$  from DOMAIN[ $X_i$ ]; removed  $\leftarrow$  true
    return removed
    
```

- Runtime: $O(n^2d^3)$, can be reduced to $O(n^2d^2)$
- ... but detecting **all** possible future problems is NP-hard – why?

[Demo: CSP applet (made available by ainspace.org) -- n-queens]

Limitations of Arc Consistency

- After enforcing arc consistency:
 - Can have one solution left
 - Can have multiple solutions left
 - Can have no solutions left (and not know it)
- Arc consistency still runs inside a backtracking search!



What went wrong here?

[Demo: coloring -- forward checking]
[Demo: coloring -- arc consistency]