

## CSE 473 Logic in AI

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(With some slides from Mausam, Stuart Russell, Dieter Fox, Henry Kautz...)

**There is nothing so powerful as truth, and often nothing so strange.**

**- Daniel Webster (1782-1852)**

## Overview

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- Introduction & Agents
- Search, Heuristics & CSPs
- Adversarial Search
- Logical Knowledge Representation
- Planning & MDPs
- Reinforcement Learning
- Uncertainty & Bayesian Networks
- Machine Learning
- NLP & Special Topics

## KR Hypothesis

Any **intelligent process** will have ingredients that

- 1) We as external observers interpret as knowledge
- 2) This knowledge plays a formal, causal & essential role in guiding the behavior

**- Brian Smith (paraphrased)**

## Some KR Languages

- Propositional Logic
- Predicate Calculus
- Frame Systems
- Rules with Certainty Factors
- Bayesian Belief Networks
- Influence Diagrams
- Semantic Networks
- Concept Description Languages
- Non-monotonic Logic

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## Knowledge Representation

- *represent knowledge in a manner that facilitates inferencing (i.e. drawing conclusions) from knowledge.*
- Typically based on
  - Logic
  - Probability
  - Logic and Probability

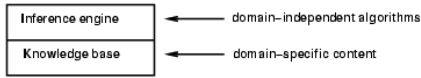
## Basic Idea of Logic

- By starting with true assumptions, you can deduce true conclusions.

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## Knowledge bases

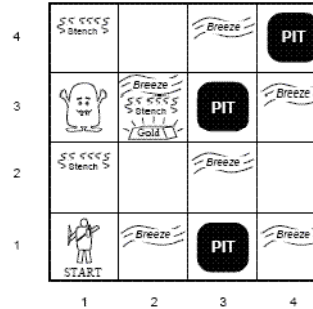


- Knowledge base = set of sentences in a formal language
- Declarative approach to building an agent (or other system):
  - Tell it what it needs to know
- Then it can Ask itself what to do - answers should follow from the KB
- Agents can be viewed at the knowledge level  
i.e., what they know, regardless of how implemented
- Or at the implementation level  
i.e., data structures in KB and algorithms that manipulate them

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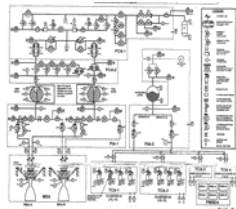
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## When Useful



## Deep Space One

- Autonomous diagnosis & repair “Remote Agent”
- Compiled schematic to 7,000 var SAT problem



## Muddy Children Problem



- Mom to N children “Don’t get dirty”
- While playing,  $K \geq 1$  get mud on forehead
- Father: “Some of you are dirty!”
- Father: “Raise your hand if you are dirty”
  - Noone raises hand
- Father: “Raise your hand if you are dirty”
  - Noone raises hand
- ...
- Father: “Raise your hand if you are dirty”
  - All dirty children raise hand

} K-1 times

## Components of KR

- Syntax: defines the sentences in the language
- Semantics: defines the “meaning” of sentences
- Inference Procedure
  - Algorithm
  - Sound?
  - Complete?
  - Complexity
- Knowledge Base

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## Propositional Logic

- Syntax
  - Atomic sentences: P, Q, ...
  - Connectives:  $\wedge$ ,  $\vee$ ,  $\neg$ ,  $\Rightarrow$
- Semantics
  - Truth Tables
- Inference
  - Modus Ponens
  - Resolution
  - DPLL
  - GSAT
- Complexity

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## Propositional Logic: Syntax

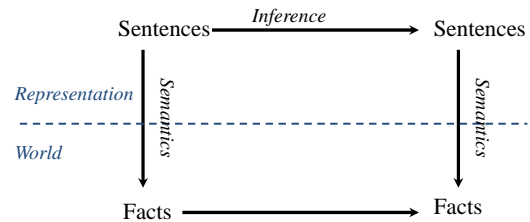
- **Atoms**
  - P, Q, R, ...
- **Literals**
  - P,  $\neg P$
- **Sentences**
  - Any literal is a sentence
  - If S is a sentence
    - Then  $(S \wedge S)$  is a sentence
    - Then  $(S \vee S)$  is a sentence
- **Conveniences**
  - $P \rightarrow Q$  same as  $\neg P \vee Q$

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## Semantics

- **Syntax**: which arrangements of symbols are *legal*
  - (Def “sentences”)
- **Semantics**: what the symbols *mean* in the world
  - (Mapping between symbols and worlds)



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## Propositional Logic: SEMANTICS

- “Interpretation” (or “possible world”)
  - Assignment to each variable either T or F
  - Assignment of T or F to each connective via defns

		Q	
		T	F
P	T	T	F
	F	F	F
		$P \wedge Q$	

		Q	
		T	F
P	T	T	T
	F	F	F
		$P \vee Q$	

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## Satisfiability, Validity, & Entailment

- S is **satisfiable** if it is true in *some* world
- S is **unsatisfiable** if it is false *all* worlds
- S is **valid** if it is true in *all* worlds
- S1 **entails** S2 if *whenever* S1 is true S2 is also true

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## Examples

$$P \rightarrow Q$$

$$X \rightarrow X$$

$$S \wedge (W \wedge \neg S)$$

$$T \vee \neg T$$

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## Notation

$\Rightarrow$   
 $\supset$   
 $\rightarrow$

} **Implication** (syntactic symbol)

⊢ **Proves**:  $S1 \vdash S2$  if inference algo, i, says ‘S2’ from S1

⊨ **Entails**:  $S1 \models S2$  if whenever S1 is true S2 is also true

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## Resolution

*If the unicorn is mythical, then it is immortal, but if it is not mythical, it is a reptile. If the unicorn is either immortal or a reptile, then it is horned.*

**M = mythical**  
**I = immortal**  
**R = reptile**  
**H = horned**

$$\begin{array}{ll} (\neg R \vee H) & (\neg I \vee H) \\ (M \vee R) & (\neg M \vee I) \end{array}$$

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## Prop. Logic: Knowledge Engr

- 1) One of the women is a biology major
- 2) Lisa is not next to Dave in the ranking
- 3) Dave is immediately ahead of Jim
- 4) Jim is immediately ahead of a bio major
- 5) Mary or Lisa is ranked first

### 1. Choose Vocabulary

Universe: Lisa, Dave, Jim, Mary  
 LD = "Lisa is immediately ahead of Dave"  
 D = "Dave is a Bio Major"

### 2. Choose initial sentences (wffs)

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## Reasoning Tasks

### • Model finding

KB = background knowledge  
 S = description of problem  
 Show  $(KB \wedge S)$  is satisfiable  
 A kind of **constraint satisfaction**

### • Deduction

S = question  
 Prove that  $KB \models S$   
 Two approaches:

- Rules to derive new formulas from old (inference)
- Show  $(KB \wedge \neg S)$  is unsatisfiable

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## Special Syntactic Forms

### • General Form:

$$((q \wedge \neg r) \rightarrow s) \wedge \neg (s \wedge t)$$

### • Conjunction Normal Form (CNF)

$$(\neg q \vee r \vee s) \wedge (\neg s \vee \neg t)$$

Set notation:  $\{(\neg q, r, s), (\neg s, \neg t)\}$

empty clause  $() = \text{false}$

### • Binary clauses: 1 or 2 literals per clause

$$(\neg q \vee r) \quad (\neg s \vee \neg t)$$

### • Horn clauses: 0 or 1 positive literal per clause

$$(\neg q \vee \neg r \vee s) \quad (\neg s \vee \neg t)$$

$$(q \wedge r) \rightarrow s \quad (s \wedge t) \rightarrow \text{false}$$

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## Propositional Logic: Inference

A **mechanical** process for computing new sentences

1. Backward & Forward Chaining
2. Resolution (Proof by Contradiction)
3. GSAT
4. Davis Putnam

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## Inference 1: Forward Chaining

### Forward Chaining

Based on rule of **modus ponens**

If know  $P_1, \dots, P_n$  & know  $(P_1 \wedge \dots \wedge P_n) \rightarrow Q$

Then can conclude Q

### Backward Chaining: search

start from the query and go backwards

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## Analysis

- Sound?
- Complete?
  - Can you prove  $\{ \} \models Q \vee \neg Q$
- If KB has only Horn clauses & query is a single literal
  - Forward Chaining is complete
  - Runs linear in the size of the KB

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## Propositional Logic: Inference

A *mechanical* process for computing new sentences

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## Conversion to CNF

$$B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$$

1. Eliminate  $\Leftrightarrow$ , replacing  $\alpha \Leftrightarrow \beta$  with  $(\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)$ .

$$(B_{1,1} \Rightarrow (P_{1,2} \vee P_{2,1})) \wedge ((P_{1,2} \vee P_{2,1}) \Rightarrow B_{1,1})$$

2. Eliminate  $\Rightarrow$ , replacing  $\alpha \Rightarrow \beta$  with  $\neg \alpha \vee \beta$ .

$$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}) \wedge (\neg(P_{1,2} \vee P_{2,1}) \vee B_{1,1})$$

3. Move  $\neg$  inwards using de Morgan's rules and double-negation:

$$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}) \wedge ((\neg P_{1,2} \wedge \neg P_{2,1}) \vee B_{1,1})$$

4. Apply distributivity law ( $\vee$  over  $\wedge$ ) and flatten:

$$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}) \wedge (\neg P_{1,2} \vee B_{1,1}) \wedge (\neg P_{2,1} \vee B_{1,1})$$

## Inference 2: Resolution

[Robinson 1965]

$$\{ (p \vee \alpha), (\neg p \vee \beta \vee \gamma) \} \vdash_{\text{R}} (\alpha \vee \beta \vee \gamma)$$

Correctness

If  $S1 \vdash_{\text{R}} S2$  then  $S1 \models S2$

Refutation Completeness:

If  $S$  is unsatisfiable then  $S \vdash_{\text{R}} ()$

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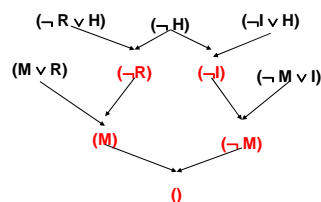
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## Resolution

*If the unicorn is mythical, then it is immortal, but if it is not mythical, it is a reptile. If the unicorn is either immortal or a reptile, then it is horned.*

**Prove: the unicorn is horned.**

**M** = mythical  
**I** = immortal  
**R** = reptile  
**H** = horned



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## Resolution as Search

- States?
- Operators

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