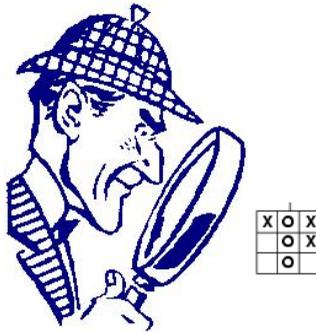


CSE 473

Lecture 7

Playing Games with Alpha-Beta Search



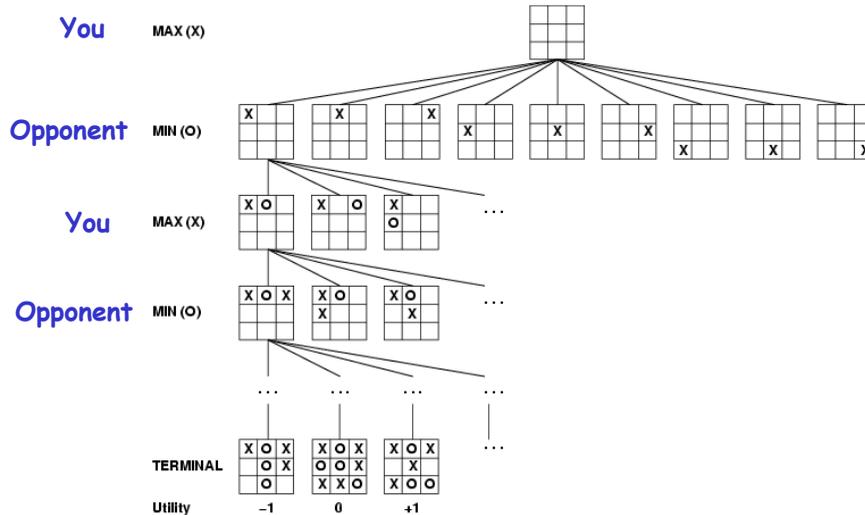
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Today

- Adversarial Search
 - Minimax recap
 - α - β search
 - Evaluation functions
 - State of the art in game playing

Recall: Game Trees

From current position, unwind game into the future



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Recall: Minimax Search

- Find the *best current move* for MAX (you) assuming MIN (opponent) also chooses *its best move*
- Compute for each node n :

MINIMAX-VALUE(n)=

UTILITY(n)

$\max_{s \in \text{succ}(n)} \text{MINIMAX-VALUE}(s)$

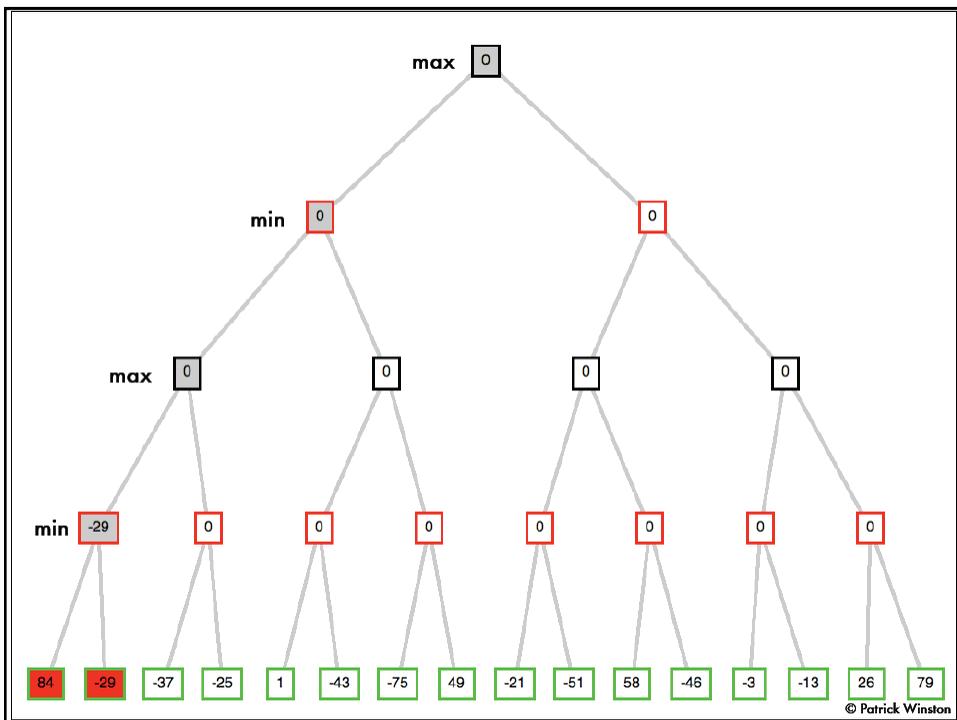
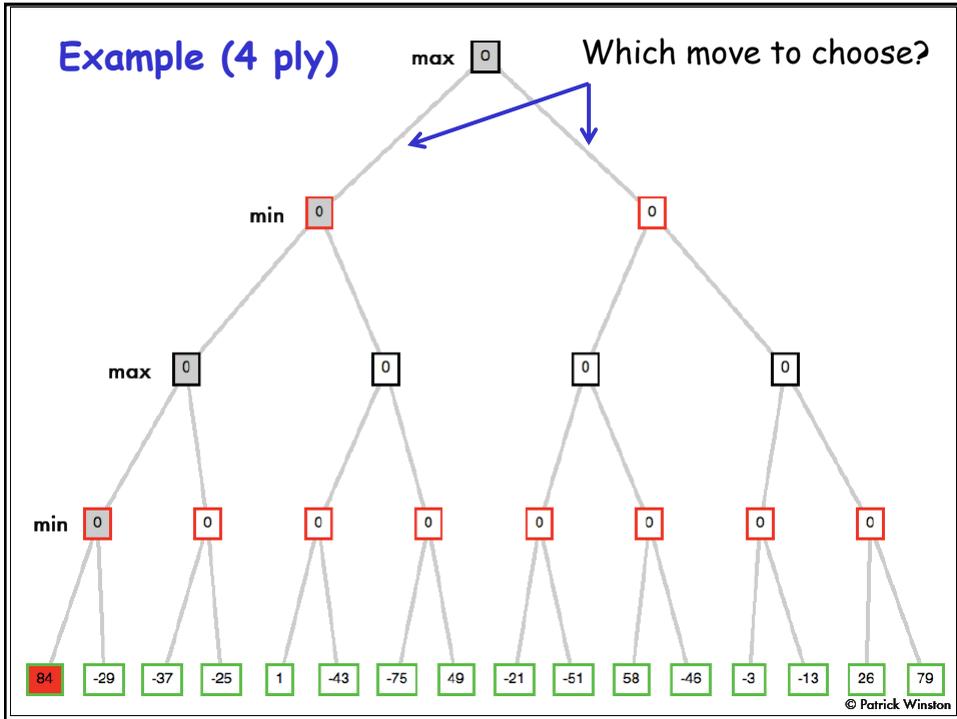
$\min_{s \in \text{succ}(n)} \text{MINIMAX-VALUE}(s)$

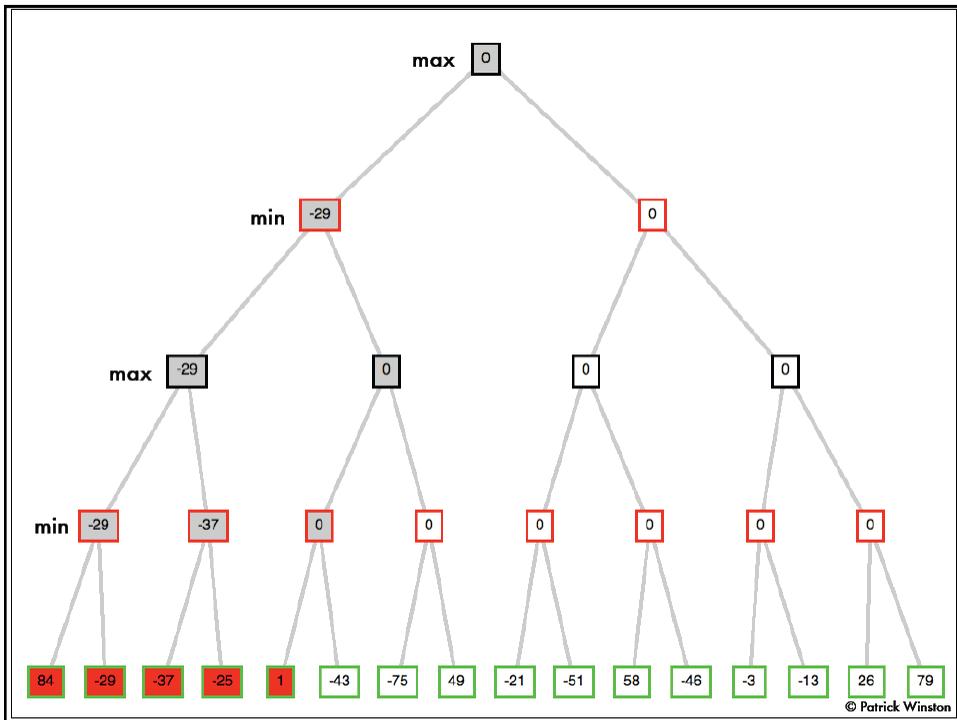
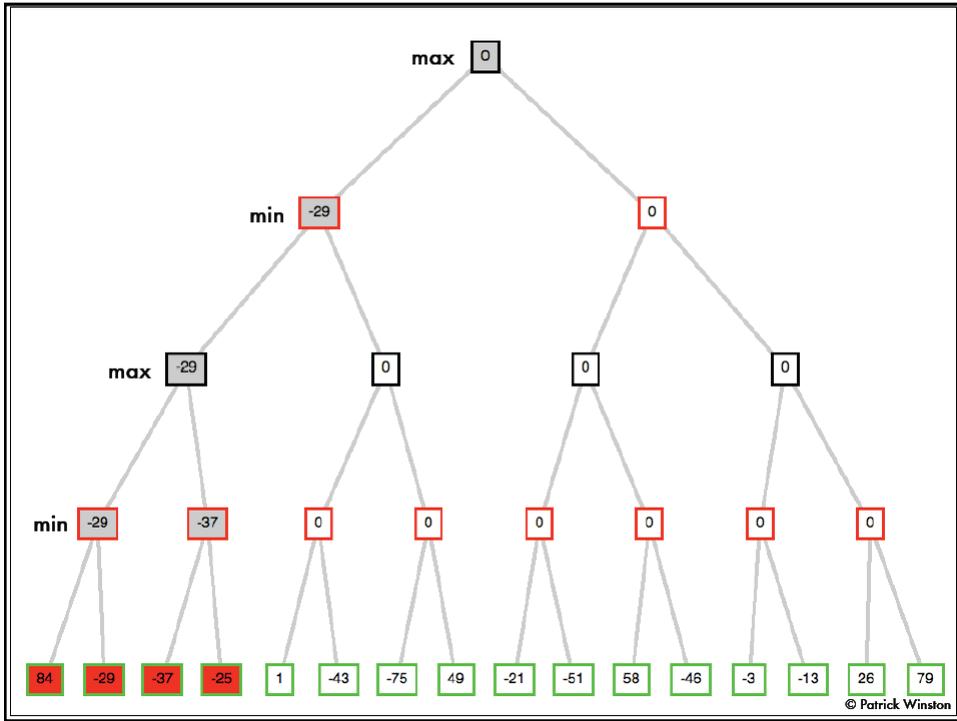
if n is a terminal

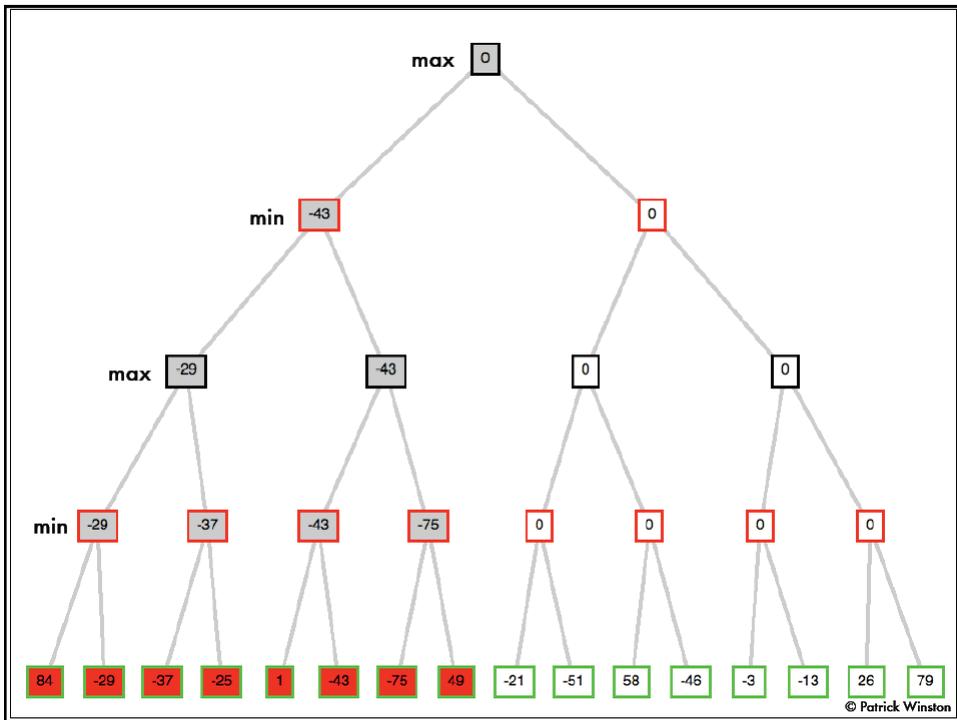
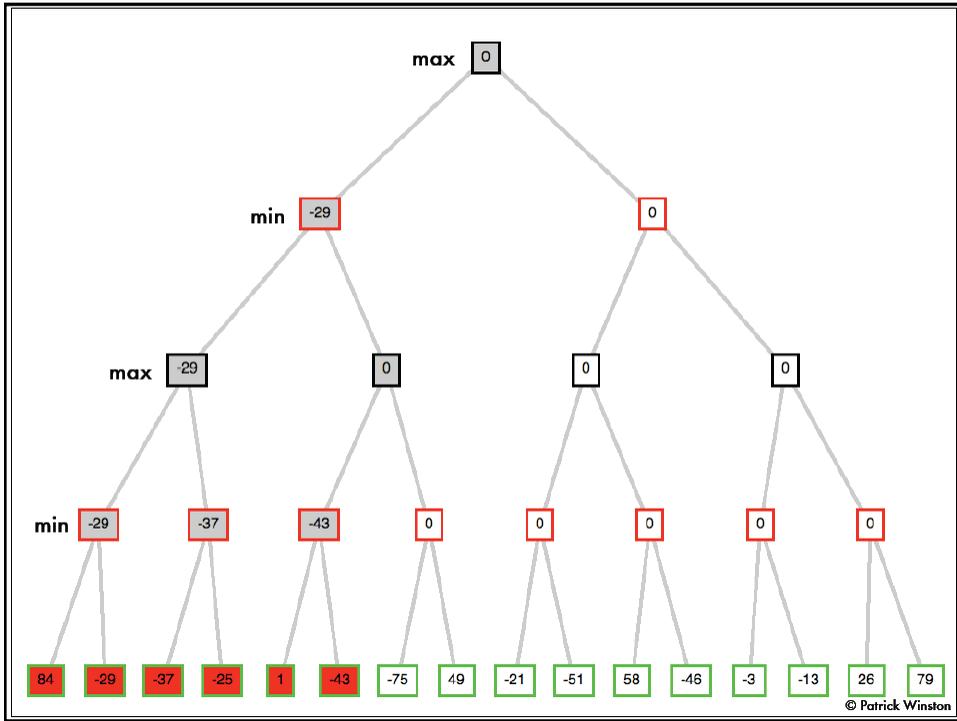
if n is a MAX node

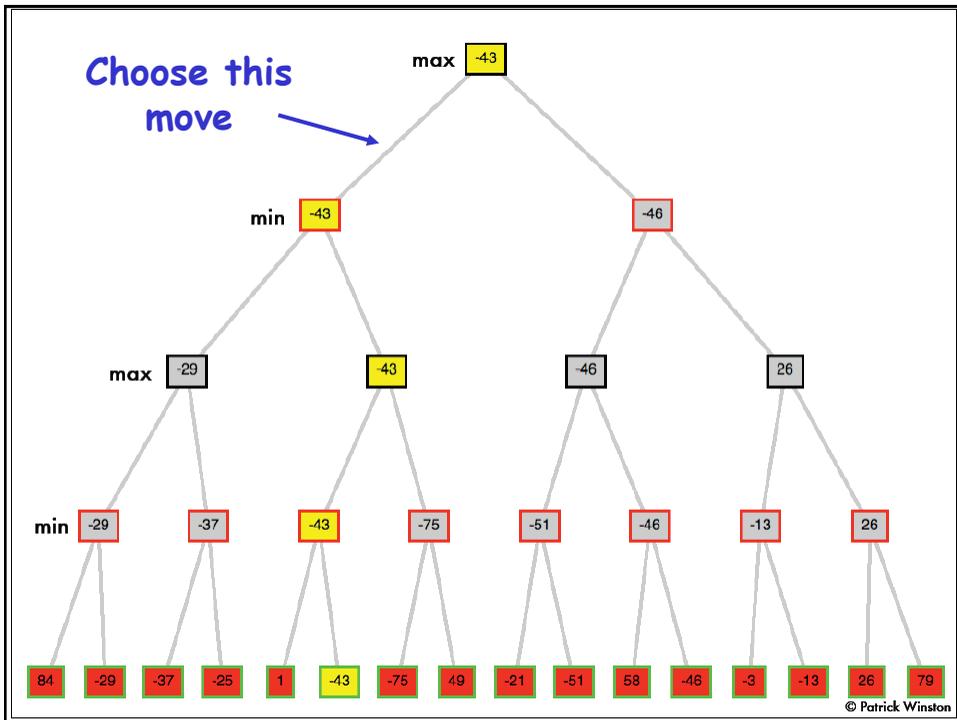
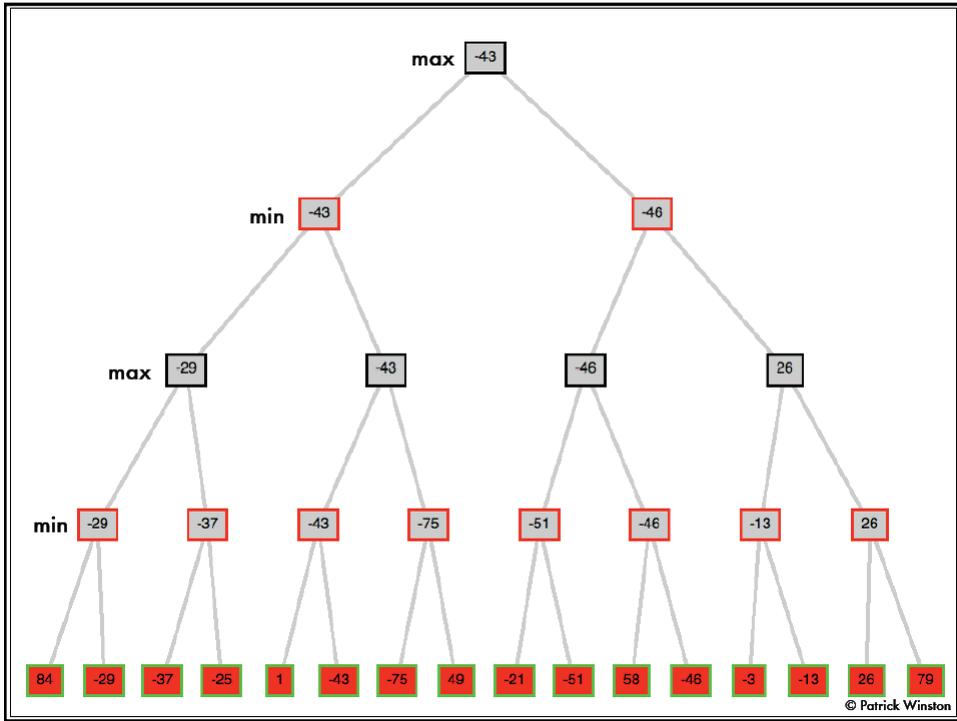
if n is a MIN node

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Minimax Algorithm

```
function MINIMAX-DECISION(state) returns an action
```

```
   $v \leftarrow \text{MAX-VALUE}(\textit{state})$ 
```

```
  return the action in SUCCESSORS(state) with value  $v$ 
```

```
function MAX-VALUE(state) returns a utility value
```

```
  if TERMINAL-TEST(state) then return UTILITY(state)
```

```
   $v \leftarrow -\infty$ 
```

```
  for  $a, s$  in SUCCESSORS(state) do
```

```
     $v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s))$ 
```

```
  return  $v$ 
```

```
function MIN-VALUE(state) returns a utility value
```

```
  if TERMINAL-TEST(state) then return UTILITY(state)
```

```
   $v \leftarrow \infty$ 
```

```
  for  $a, s$  in SUCCESSORS(state) do
```

```
     $v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s))$ 
```

```
  return  $v$ 
```

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Properties of minimax

- Complete? Yes (if tree is finite)
- Optimal? Yes (against an optimal opponent)
- Time complexity? $O(b^m)$
- Space complexity? $O(bm)$ (depth-first exploration)

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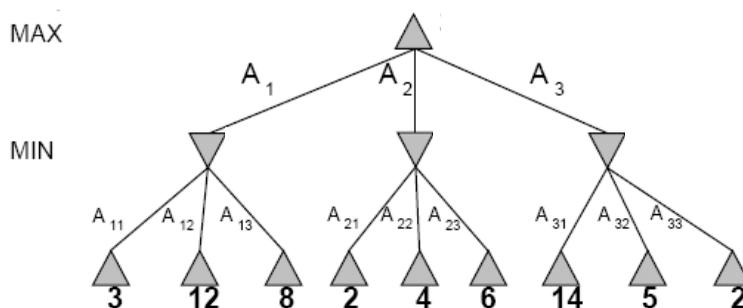
Good enough?

- Chess:
 - branching factor $b \approx 35$
 - game length $m \approx 100$
 - search space $b^m \approx 35^{100} \approx 10^{154}$
- The Universe:
 - number of atoms $\approx 10^{78}$
 - age $\approx 10^{21}$ milliseconds

Can we search more efficiently?

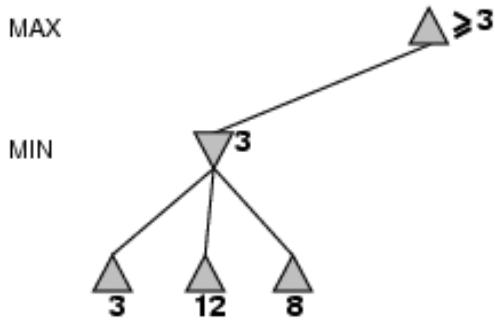
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Two-Ply Game Tree



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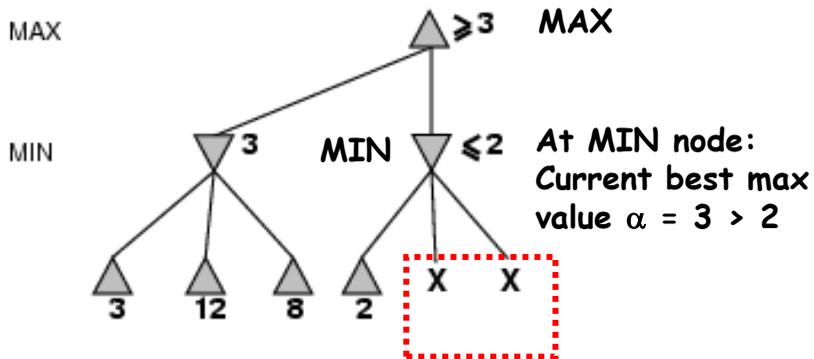
Pruning trees



Minimax algorithm explores depth-first

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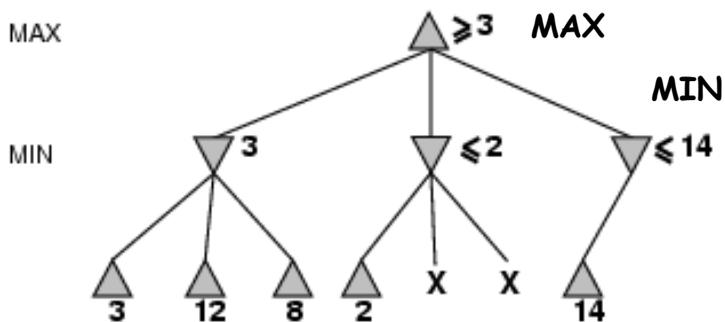
Pruning trees



No need to look at these nodes!! (these nodes can only decrease MIN value from 2)

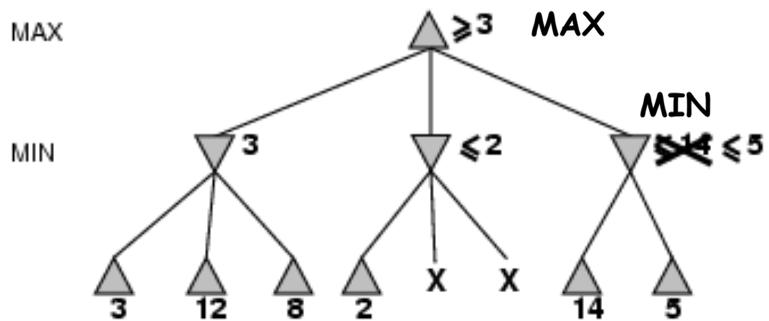
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Pruning trees



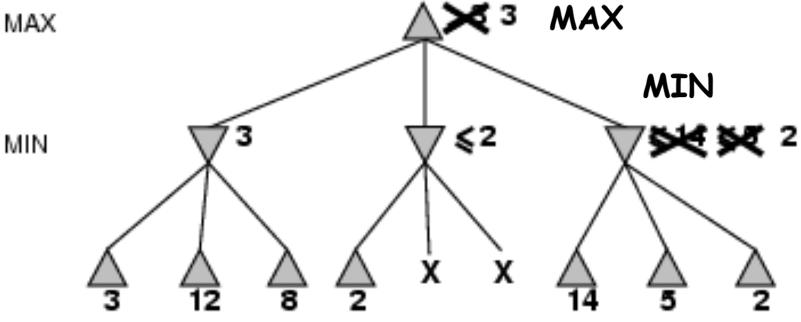
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Pruning trees

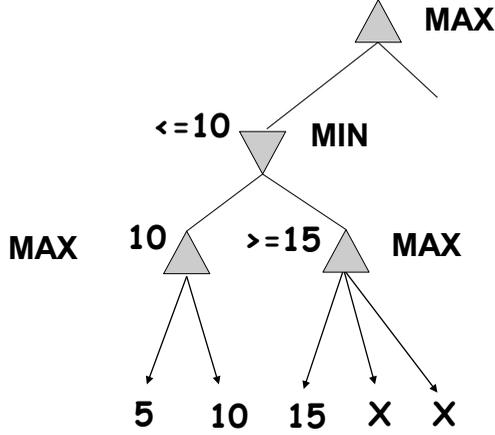


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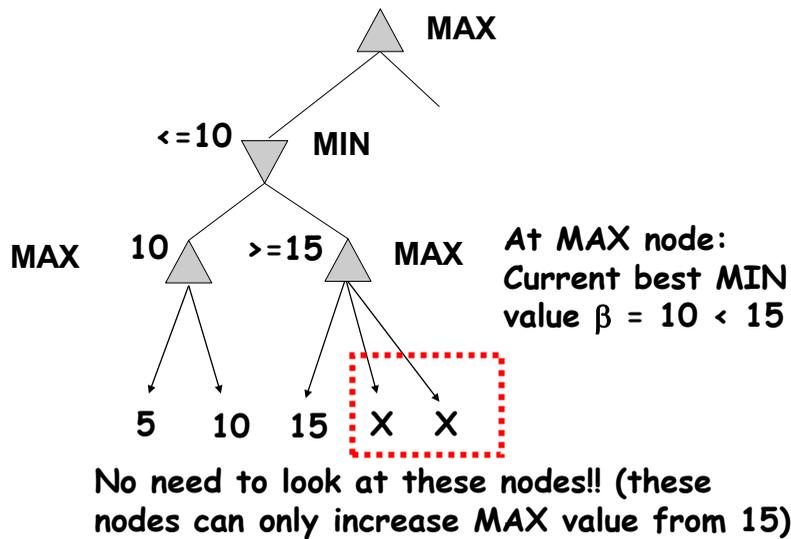
Pruning trees



One more example



One more example



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This form of tree pruning is
known as alpha-beta pruning

alpha = highest (best) value for MAX along current path
from root

beta = lowest (best) value for MIN along current path
from root

The α - β algorithm

(minimax with four lines of added code)

function ALPHA-BETA-SEARCH(*state*) returns an action

inputs: *state*, current state in game

$v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$

return the action in SUCCESSORS(*state*) with value *v*

function MAX-VALUE(*state*, α , β) returns a utility value

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) then return UTILITY(*state*)

$v \leftarrow -\infty$

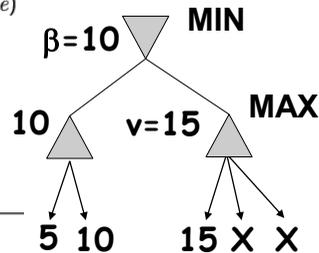
for *a, s* in SUCCESSORS(*state*) do

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))$

New { if $v \geq \beta$ then return *v* → Pruning

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return *v*



The α - β algorithm (cont.)

function MIN-VALUE(*state*, α , β) returns a utility value

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) then return UTILITY(*state*)

$v \leftarrow +\infty$

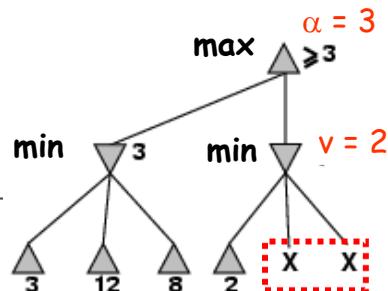
for *a, s* in SUCCESSORS(*state*) do

$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s, \alpha, \beta))$

{ if $v \leq \alpha$ then return *v* → Pruning

$\beta \leftarrow \text{MIN}(\beta, v)$

return *v*



Properties of α - β

- Pruning **does not** affect final result
- *Good move ordering* improves effectiveness of pruning
- With "perfect ordering," time complexity = $O(b^{m/2})$
 - allows us to search deeper - **doubles** depth of search
- α - β search is a simple example of the value of reasoning about which computations are relevant (a form of **metareasoning**)

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Good enough?

- **Chess:**
 - branching factor $b \approx 35$
 - game length $m \approx 100$
 - α - β search space $b^{m/2} \approx 35^{50} \approx 10^{77}$
- **The Universe:**
 - number of atoms $\approx 10^{78}$
 - age $\approx 10^{21}$ milliseconds

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Transposition Tables

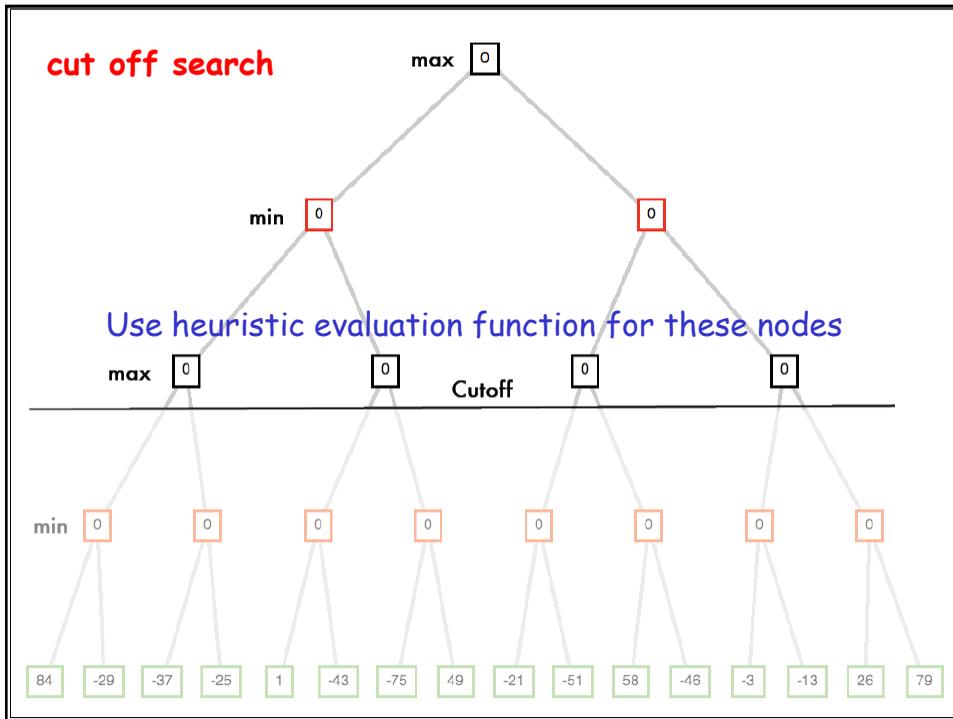
- Game trees contain **repeated states**
- In chess, e.g., the game tree may have 35^{100} nodes, but there are only 10^{40} different board positions
- Similar to **explored set** in graph-search, maintain a **transposition table**
 - Got its name from the fact that the same state is reached by a transposition of moves.
- 10^{40} is still huge!

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Can we do better?

- **Strategies:**
 - search to a fixed depth (**cut off search**)
 - iterative deepening* (most common)

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Evaluation Function

- When **search space is too large**, create game tree up to a certain depth only.
 - Art is to **estimate utilities of positions that are not terminal states**.
 - Example of simple **evaluation criteria in chess**:
 - Material worth: pawn=1, knight =3, rook=5, queen=9.
 - Other: king safety, good pawn structure
- eval(s) =**

$$\begin{aligned}
 &w1 * \text{material}(s) + \\
 &w2 * \text{mobility}(s) + \\
 &w3 * \text{king safety}(s) + \\
 &w4 * \text{center control}(s) + \dots
 \end{aligned}$$

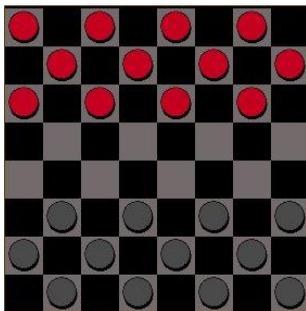
Cutting off search

- Does it work in practice?
Suppose $b^m = 10^6$ and $b=35 \Rightarrow m=4$
- 4-ply lookahead is a **hopeless chess player!**
 - 4-ply \approx human novice
 - 8-ply \approx typical PC, human master
 - 12-ply \approx Deep Blue, Kasparov

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Game Playing State-of-the-Art

- **Checkers:** Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used an endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 443,748,401,247 positions. Checkers is now solved!



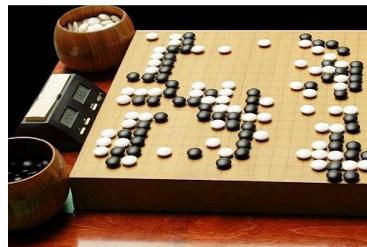
Game Playing State-of-the-Art

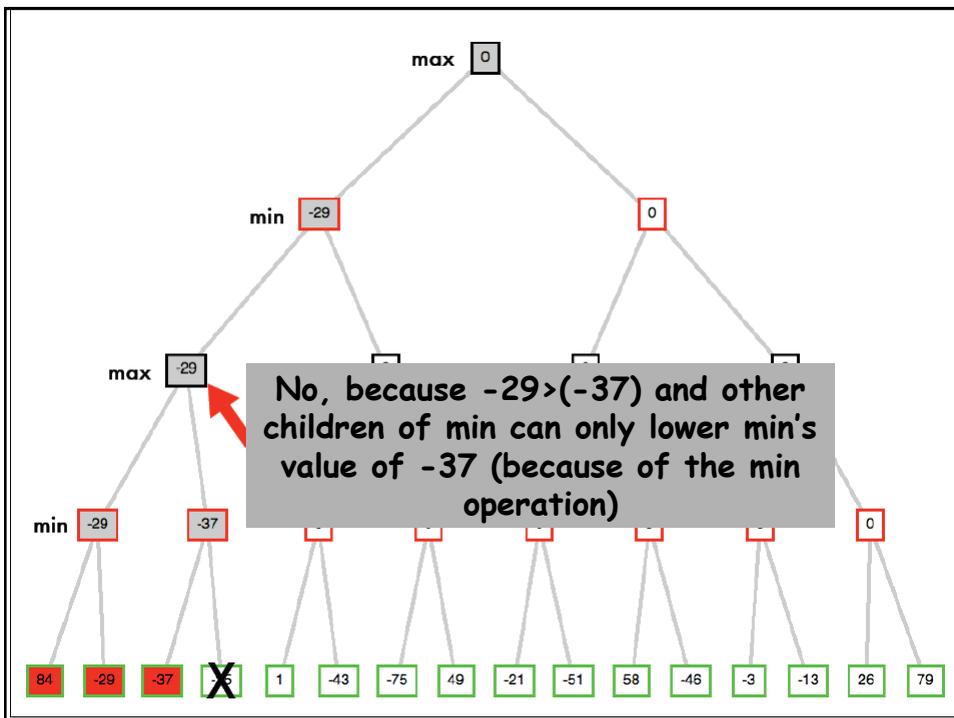
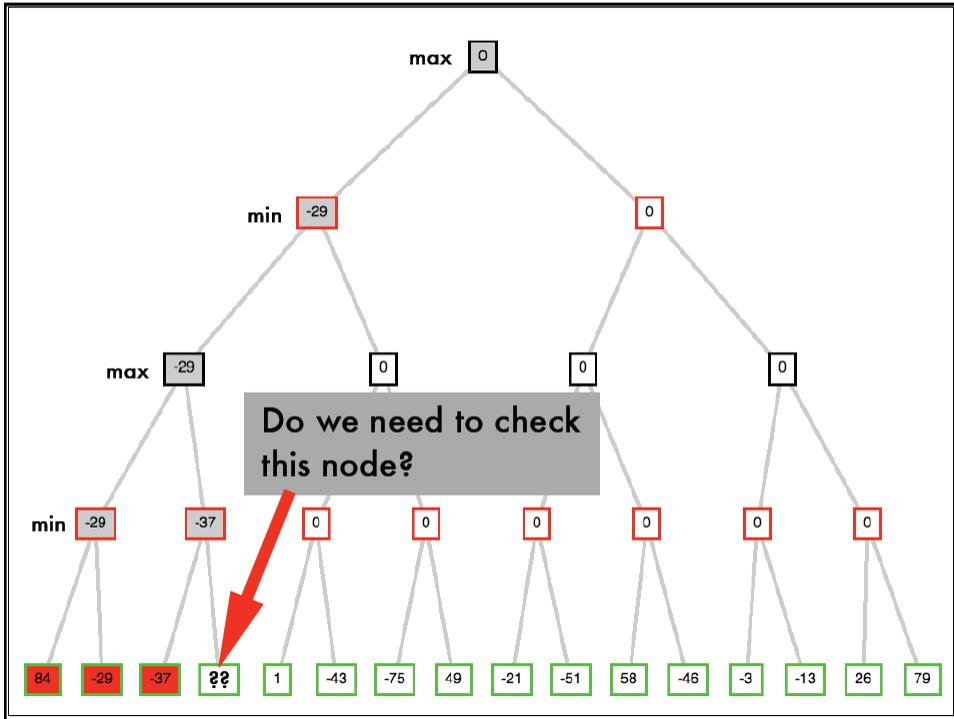
- **Chess:** Deep Blue defeated human world champion Gary Kasparov in a six-game match in 1997. Deep Blue examined 200 million positions per second, used very sophisticated evaluation function and undisclosed methods for extending some lines of search up to 40 ply. Current programs are even better, if less historic.

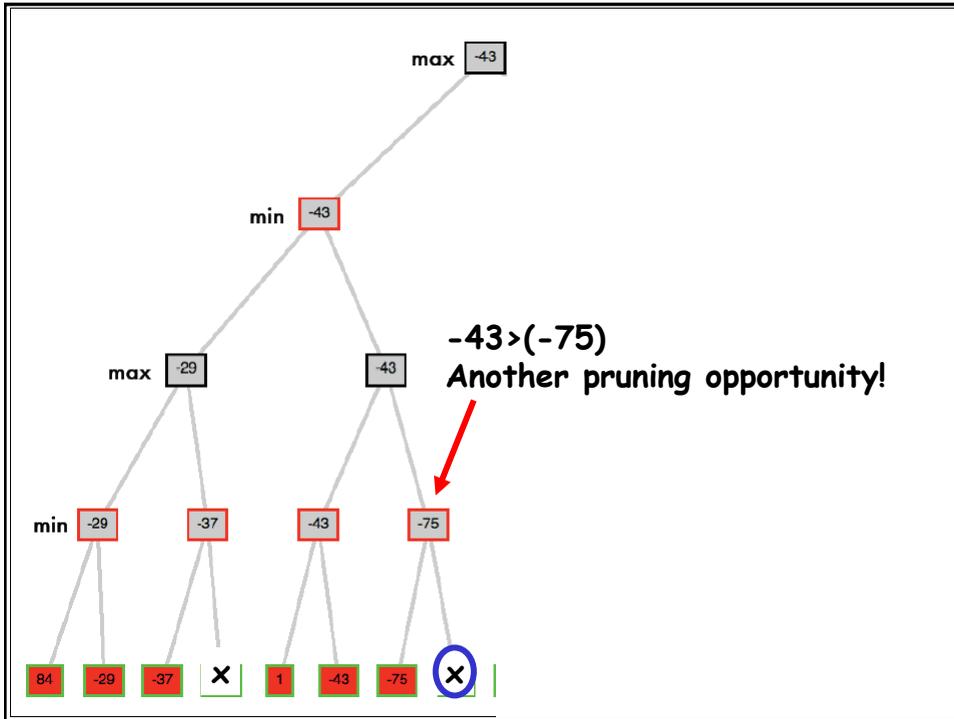


Game Playing State-of-the-Art

- **Othello:** Human champions refuse to play against computers because **software is too good**
- **Go:** Human champions refuse to play against computers because **software is too bad**.
 - In Go, $b > 300$, so need pattern databases and Monte Carlo search (UCT)
 - Human champions are now beginning to be challenged by machines.
- **Pacman:** The reigning champion is [your CSE 473 program here](#)







Pruning can eliminate entire subtrees!

