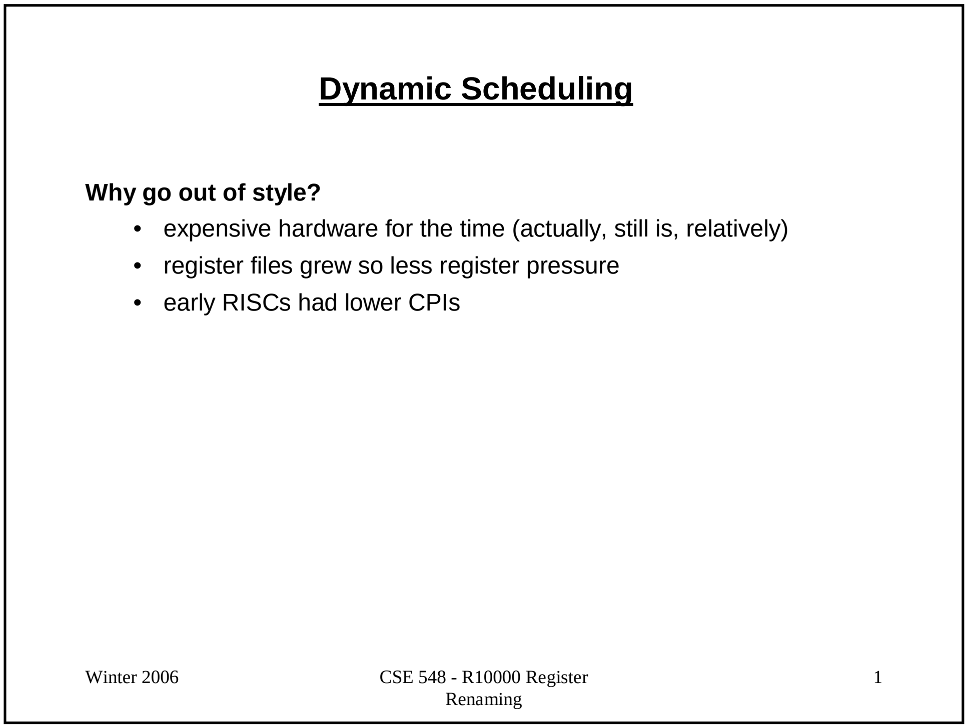
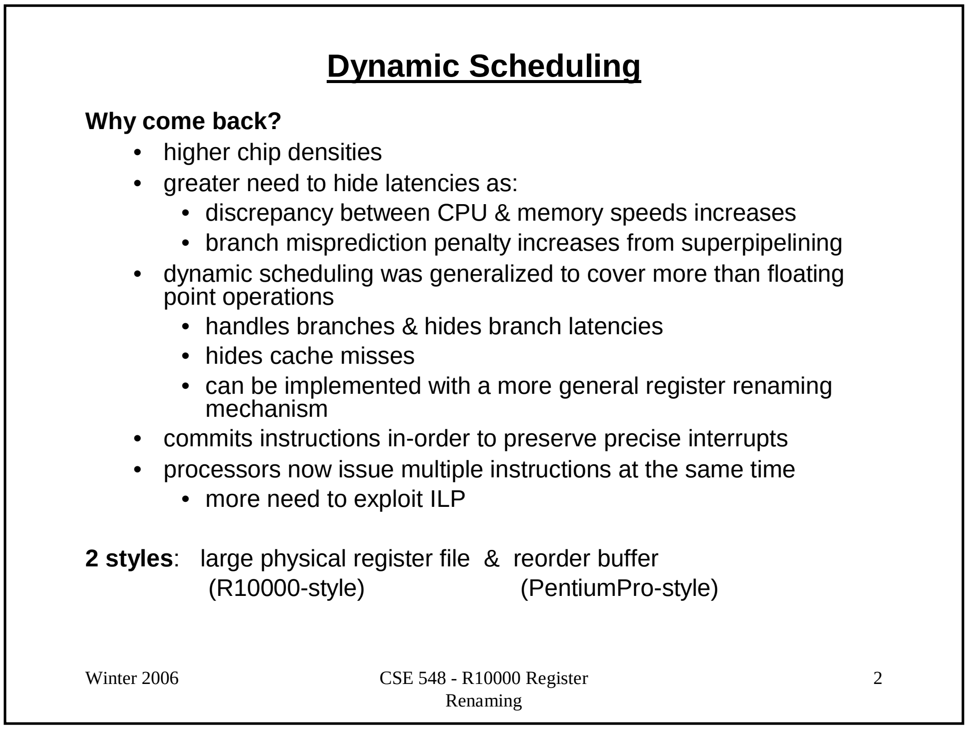


Dynamic Scheduling

Why go out of style?

- expensive hardware for the time (actually, still is, relatively)
- register files grew so less register pressure
- early RISCs had lower CPIs



Dynamic Scheduling

Why come back?

- higher chip densities
- greater need to hide latencies as:
  - discrepancy between CPU & memory speeds increases
  - branch misprediction penalty increases from superpipelining
- dynamic scheduling was generalized to cover more than floating point operations
  - handles branches & hides branch latencies
  - hides cache misses
  - can be implemented with a more general register renaming mechanism
- commits instructions in-order to preserve precise interrupts
- processors now issue multiple instructions at the same time
  - more need to exploit ILP

2 styles: large physical register file & reorder buffer
(R10000-style) (PentiumPro-style)

Register Renaming with A Physical Register File

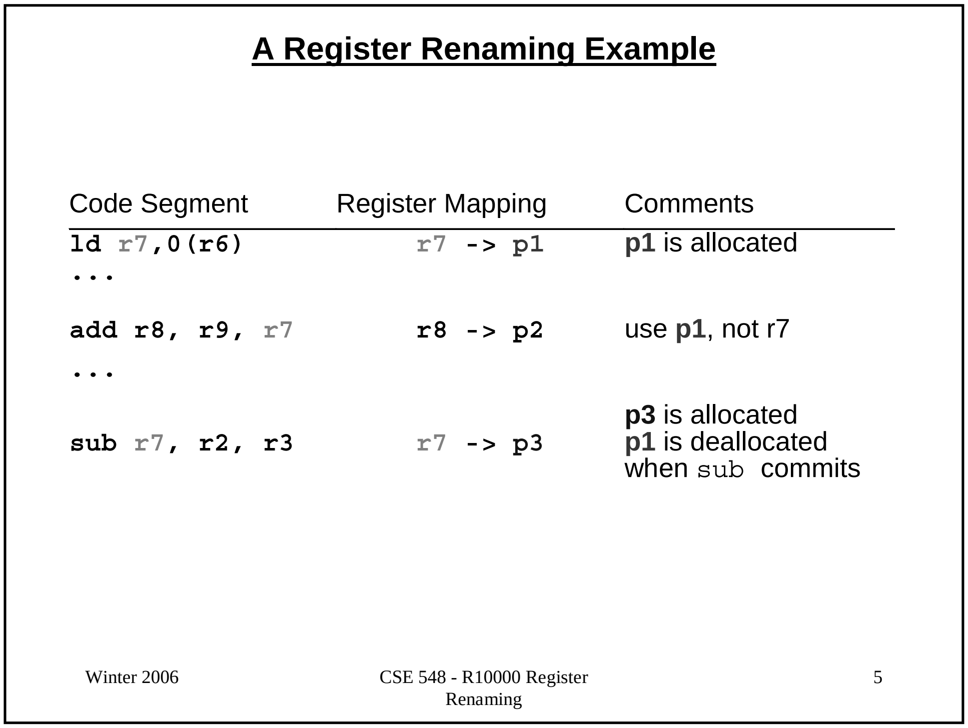
Register renaming provides a **mapping** between 2 register sets

- **architectural registers** defined by the ISA
- **physical registers** implemented in the CPU
  - hold results of the instructions committed so far
  - hold results of subsequent instructions that have not yet committed
  - more of them than architectural registers
    - ~ issue width * # pipeline stages between register renaming & commit

Register Renaming with A Physical Register File

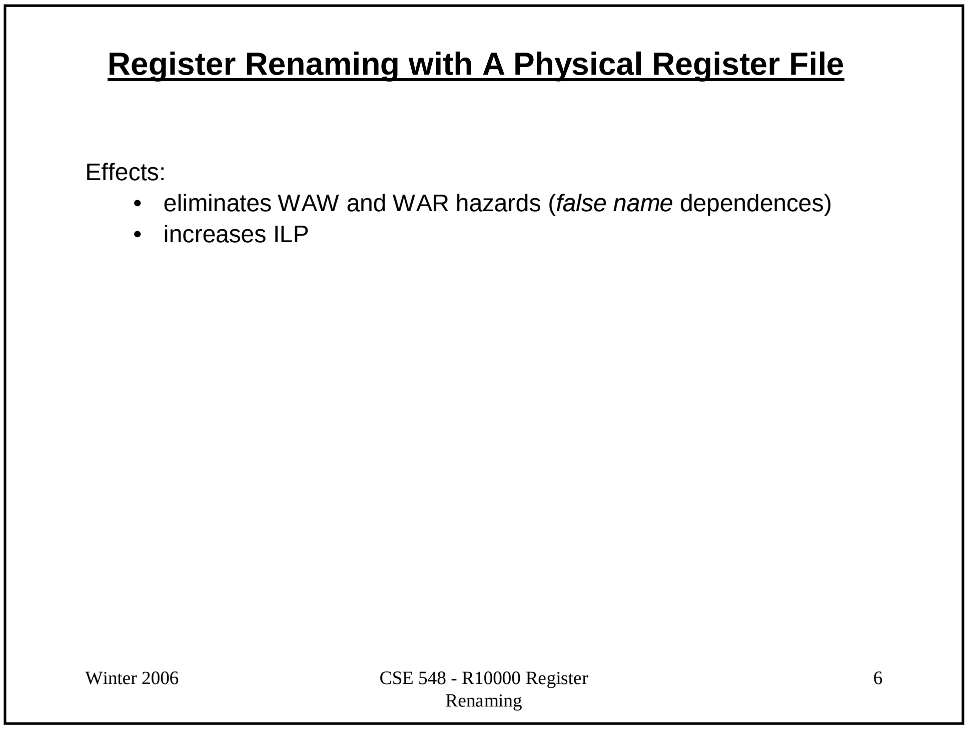
How does it work?:

- an architectural register is mapped to a physical register during a register renaming stage in the pipeline
  - destination registers create mappings
  - source registers use mappings
- operands thereafter are called by their physical register number
  - hazards determined by comparing physical register numbers, not architectural register numbers



A Register Renaming Example

Code Segment	Register Mapping	Comments
`ld r7,0(r6)`	`r7 -> p1`	**p1** is allocated
`...`		
`add r8, r9, r7`	`r8 -> p2`	use **p1**, not r7
`...`		
`sub r7, r2, r3`	`r7 -> p3`	**p3** is allocated **p1** is deallocated when `sub` commits



Register Renaming with A Physical Register File

Effects:

- eliminates WAW and WAR hazards (*false name* dependences)
- increases ILP

An Implementation (R10000)

Modular design with regular hardware data structures

Structures for register renaming

- 64 **physical registers** (each, for integer & FP)
- **map tables** for the ***current*** architectural-to-physical register mapping (separate, for integer & FP)
  - accessed with an architectural register number
  - produces a physical register number
- source operands refer to the latest defined destination register, i.e., the current mappings
- a destination register is assigned a new physical register number from a **free register list** (separate, for integer & FP)

An Implementation (R10000)

Instruction “queues” (integer, FP & data transfer)

- contains decoded & mapped instructions with the current physical register mappings
  - instructions entered into free locations in the IQ
  - sit there until they are dispatched to functional units
  - somewhat analogous to Tomasulo reservation stations without value fields or valid bits
- used to determine when operands are available
  - compare each source operand of instructions in the IQ to destination values just computed
- determines when an appropriate functional unit is available
- dispatches instructions to functional units

An Implementation (R10000)

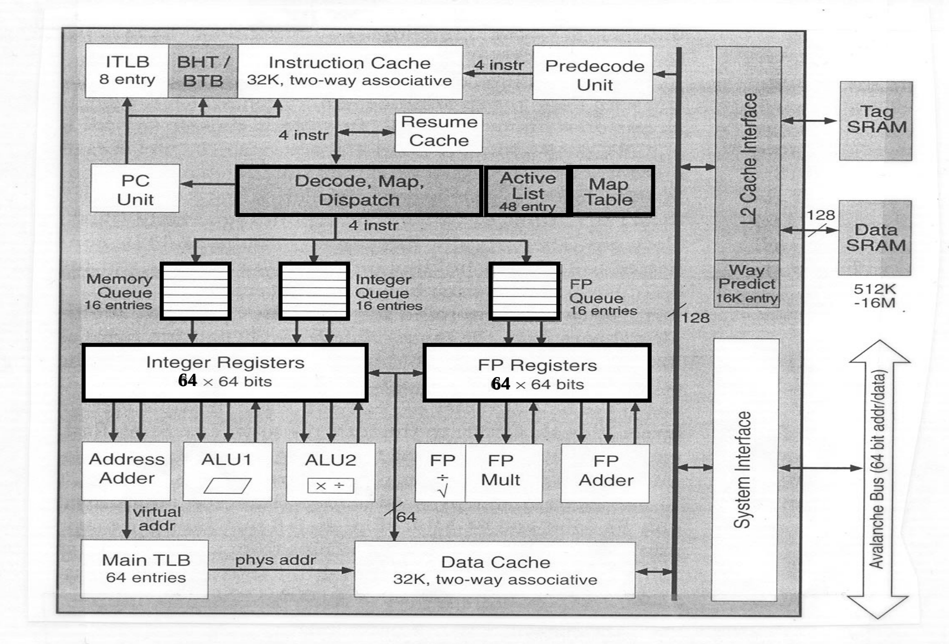
active list for all uncommitted instructions

- the mechanism for maintaining precise interrupts
  - instructions entered in program-generated order
  - allows instructions to complete in program-generated order
- instructions removed from the active list when:
  - an instruction commits:
    - the instruction has completed execution
    - all instructions ahead of it have also completed
  - branch is mispredicted
  - an exception occurs
- contains the ***previous*** architectural-to-physical destination register mapping
  - used to recreate the map table for instruction restart after an exception
- instructions in the other hardware structures & the functional units are identified by their active list location

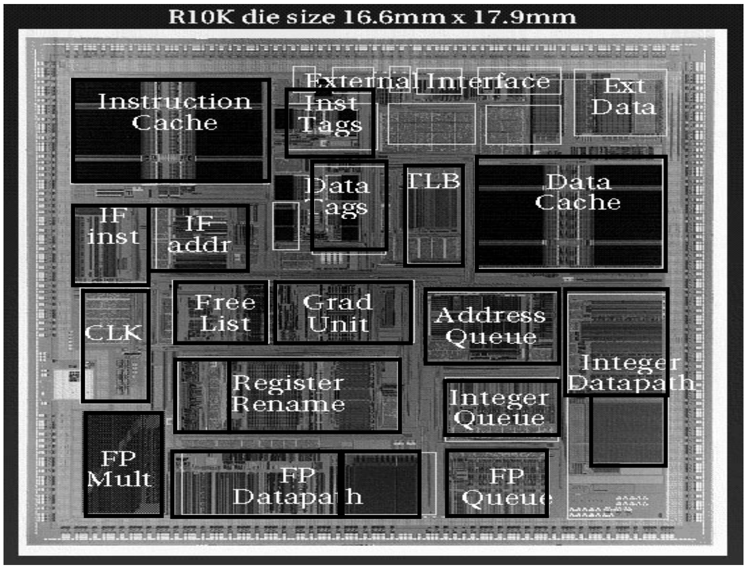
An Implementation (R10000)

busy-register table (integer & FP):

- indicates whether a physical register contains a value
- somewhat analogous to Tomasulo’s register status
- used to determine operand availability
  - bit is set when a register is mapped & leaves the free list (not available yet)
  - cleared when a FU writes the register (now there’s a value)

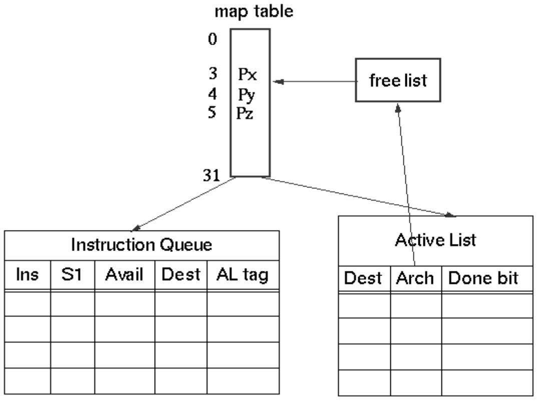


R10000 Die Photo



The R10000 in Action 1

→	`ld A3, #(reg)`	arch register A3 defined potential multi-cycle
	`add A4, A3, reg`	arch register A3 used
	`sub A3, reg, reg`	arch register A3 redefined name dependence
	`or A5, A3, reg`	arch register A3 used



Instruction Queue

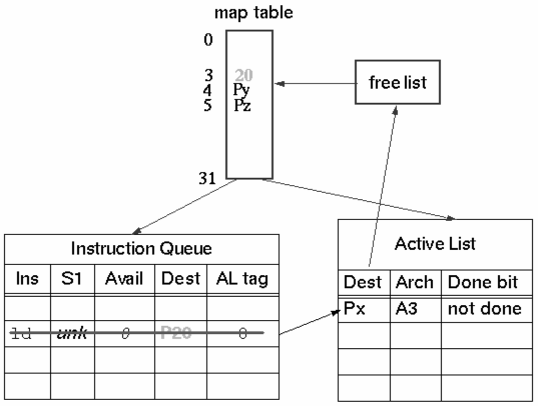
Ins	S1	Avail	Dest	AL tag

Active List

Dest	Arch	Done bit

The R10000 in Action 2

→	`ld A3, #(reg)`	arch register A3 defined potential multi-cycle
	`add A4, A3, reg`	
	`sub A3, reg, reg`	
	`or A5, A3, reg`	

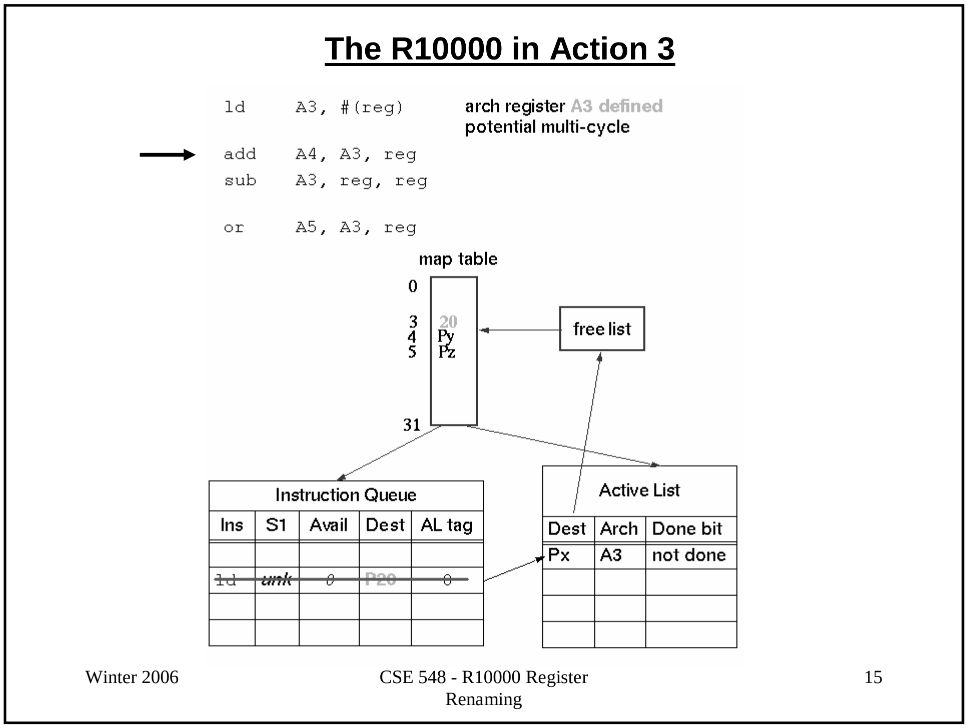


Instruction Queue

Ins	S1	Avail	Dest	AL tag
~~ld~~	~~*unk*~~	~~*0*~~	~~P20~~	~~0~~

Active List

Dest	Arch	Done bit
Px	A3	not done



The R10000 in Action 3

```
    ld      A3, #(reg)
→   add     A4, A3, reg
    sub     A3, reg, reg
    or      A5, A3, reg
```

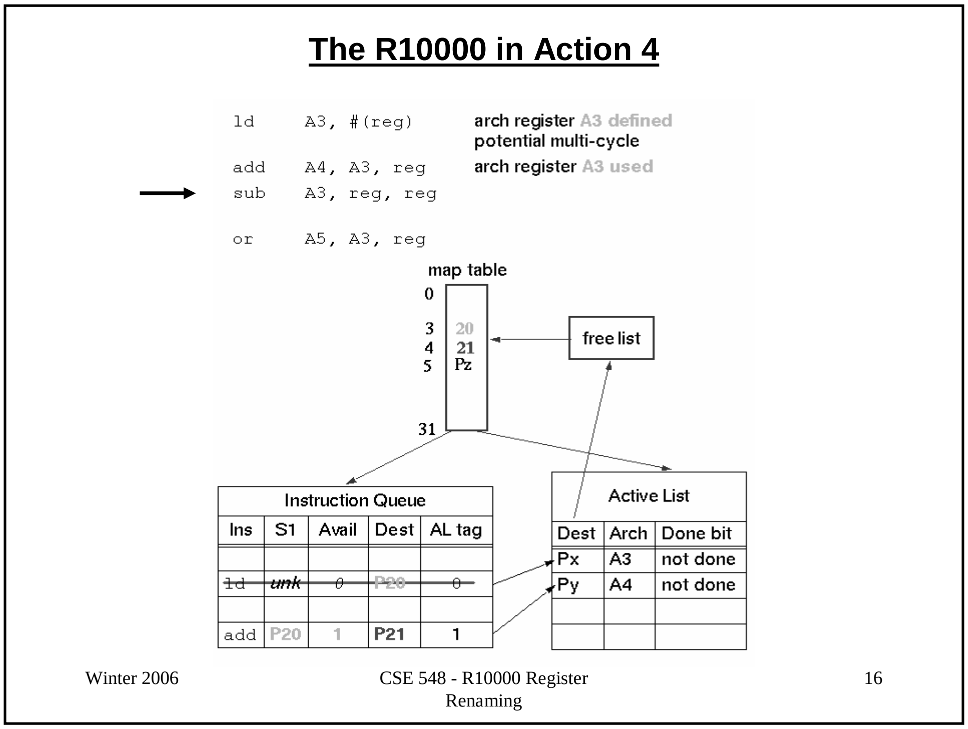
arch register A3 defined
potential multi-cycle

Instruction Queue

Ins	S1	Avail	Dest	AL tag
~~ld~~	~~*unk*~~	~~*0*~~	~~P20~~	~~0~~

Active List

Dest	Arch	Done bit
Px	A3	not done



The R10000 in Action 4

```
    ld      A3, #(reg)
    add     A4, A3, reg
→   sub     A3, reg, reg
    or      A5, A3, reg
```

arch register A3 defined
potential multi-cycle
arch register A3 used

Instruction Queue

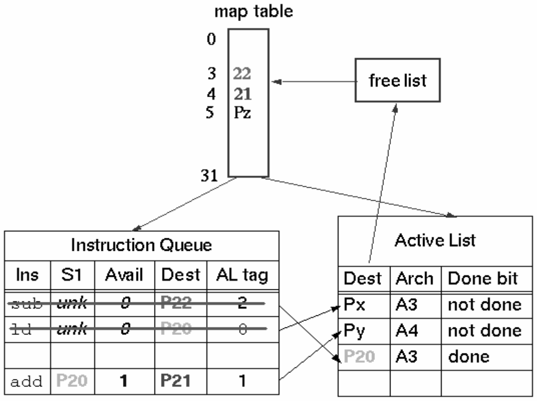
Ins	S1	Avail	Dest	AL tag
~~ld~~	~~*unk*~~	~~*0*~~	~~P20~~	~~0~~
add	P20	1	P21	1

Active List

Dest	Arch	Done bit
Px	A3	not done
Py	A4	not done

The R10000 in Action 5

```
    ld     A3, #(reg)       arch register A3 defined
                            potential multi-cycle
    add    A4, A3, reg      arch register A3 used
    sub    A3, reg, reg     arch register A3 redefined
                            name dependence
→   or     A5, A3, reg
```



Instruction Queue

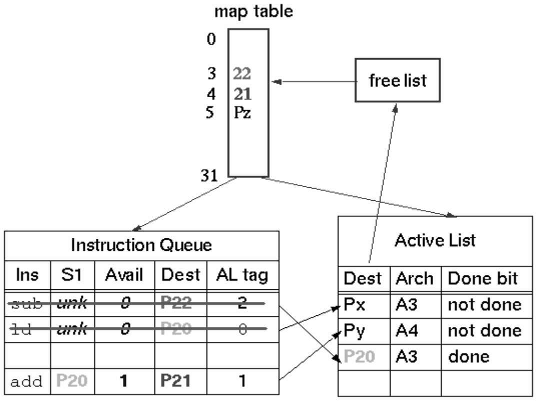
Ins	S1	Avail	Dest	AL tag
~~sub~~	~~unk~~	~~0~~	~~P22~~	~~2~~
~~ld~~	~~unk~~	~~0~~	~~P20~~	~~0~~
add	P20	1	P21	1

Active List

Dest	Arch	Done bit
Px	A3	not done
Py	A4	not done
P20	A3	done

The R10000 in Action 5 : Interrupts 1

```
    ld     A3, #(reg)       arch register A3 defined
                            potential multi-cycle
    add    A4, A3, reg      arch register A3 used
    sub    A3, reg, reg     arch register A3 redefined
                            name dependence
→   or     A5, A3, reg
```



Instruction Queue

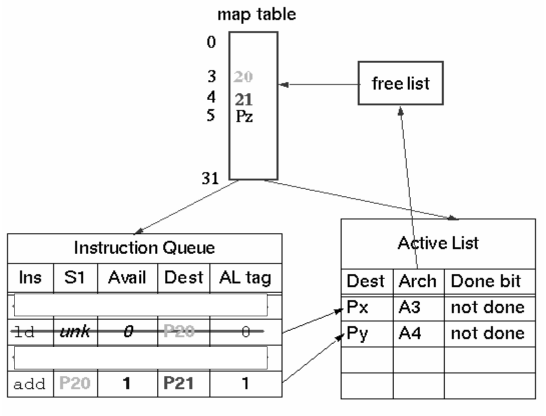
Ins	S1	Avail	Dest	AL tag
~~sub~~	~~unk~~	~~0~~	~~P22~~	~~2~~
~~ld~~	~~unk~~	~~0~~	~~P20~~	~~0~~
add	P20	1	P21	1

Active List

Dest	Arch	Done bit
Px	A3	not done
Py	A4	not done
P20	A3	done

The R10000 in Action: Interrupts 2

	`ld`	`A3, #(reg)`	arch register A3 defined potential multi-cycle
	`add`	`A4, A3, reg`	arch register A3 used
→	`sub`	`A3, reg, reg`	arch register A3 redefined name dependence
	`or`	`A5, A3, reg`	arch register A3 used



Instruction Queue

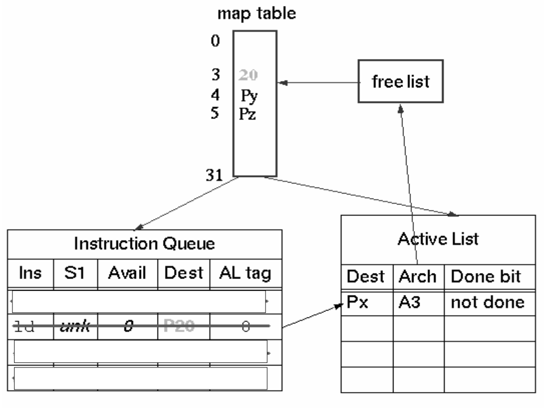
Ins	S1	Avail	Dest	AL tag
~~ld~~	~~*unk*~~	~~*0*~~	~~P20~~	~~0~~
add	P20	1	P21	1

Active List

Dest	Arch	Done bit
Px	A3	not done
Py	A4	not done

The R10000 in Action: Interrupts 3

	`ld`	`A3, #(reg)`	arch register A3 defined potential multi-cycle
→	`add`	`A4, A3, reg`	arch register A3 used
	`sub`	`A3, reg, reg`	arch register A3 redefined name dependence
	`or`	`A5, A3, reg`	arch register A3 used



Instruction Queue

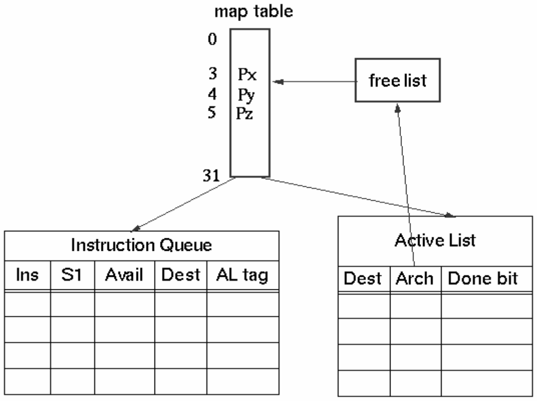
Ins	S1	Avail	Dest	AL tag
~~ld~~	~~*unk*~~	~~*0*~~	~~P20~~	~~0~~

Active List

Dest	Arch	Done bit
Px	A3	not done

The R10000 in Action: Interrupts 4

```
→ ld    A3, #(reg)      arch register A3 defined
                        potential multi-cycle
  add   A4, A3, reg     arch register A3 used
  sub   A3, reg, reg    arch register A3 redefined
                        name dependence
  or    A5, A3, reg     arch register A3 used
```



Instruction Queue

Ins	S1	Avail	Dest	AL tag

Active List

Dest	Arch	Done bit

R10000 Execution

In-order issue (have already fetched instructions)
- rename architectural registers to physical registers via a map table
- detect structural hazards for instruction queues (integer, memory & FP) & active list
- issue up to 4 instructions to the instruction queues

Out-of-order execution (to increase ILP)
- reservation-station-like instruction queues that indicate when an operand has been calculated
  - each instruction monitors the setting of the busy-register table
- set busy-register table entry for the destination register
- detect functional unit structural & RAW hazards
- dispatch instructions to functional units

In-order commit (to preserve precise interrupts)
- this & previous program-generated instructions have completed
- physical register in previous mapping returned to free list
- rollback on interrupts