

## Issues in Multiprocessors

Which **programming model for interprocessor communication**

- shared memory
  - regular loads & stores
  - SPARCcenter, SGI Challenge, Cray T3D, Convex Exemplar, KSR-1&2
- message passing
  - explicit sends & receives
  - TMC CM-5, Intel Paragon, IBM SP-2

Which **execution model**

- control parallel
  - identify & synchronize different asynchronous threads
- data parallel
  - same operation on different parts of the shared data space

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1

## Issues in Multiprocessors

How to **express parallelism**

- language support
  - HPF, ZPL
- runtime library constructs
  - coarse-grain, explicitly parallel C programs
- automatic (compiler) detection
  - implicitly parallel C & Fortran programs, e.g., SUIF & PTRANS compilers

**Application development**

- embarrassingly parallel programs could be easily parallelized
- development of different algorithms for same problem

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## Issues in Multiprocessors

How to get **good parallel performance**

- recognize parallelism
- transform programs to increase parallelism without decreasing processor locality
- decrease sharing costs

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3

## Flynn Classification

**SISD**: single instruction stream, single data stream

- single-context uniprocessors

**SIMD**: single instruction stream, multiple data streams

- exploits data parallelism
- example: Thinking Machines CM

**MISD**: multiple instruction streams, single data stream

- systolic arrays
- example: Intel iWarp, streaming processors

**MIMD**: multiple instruction streams, multiple data streams

- multiprocessors
- multithreaded processors
- parallel programming & multiprogramming
- relies on control parallelism: execute & synchronize different asynchronous threads of control
- example: most processor companies have MP configurations

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4

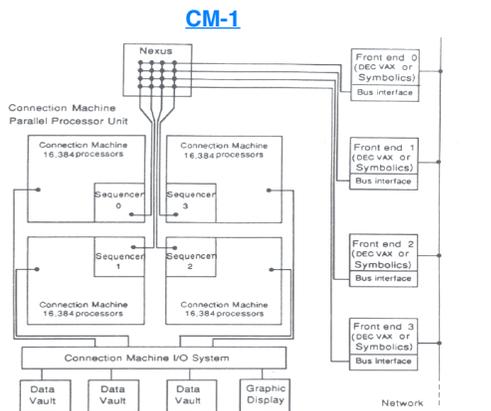
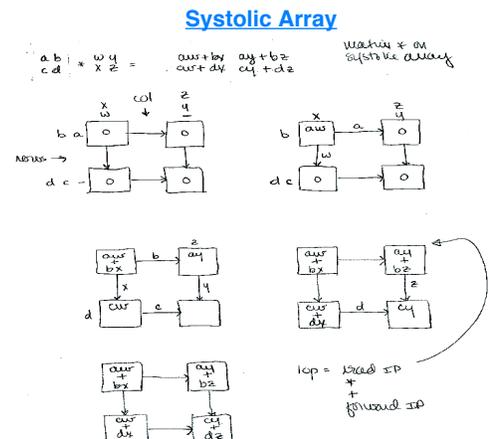


Figure 1. Connection Machine system organization.

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5



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## MIMD

### Low-end

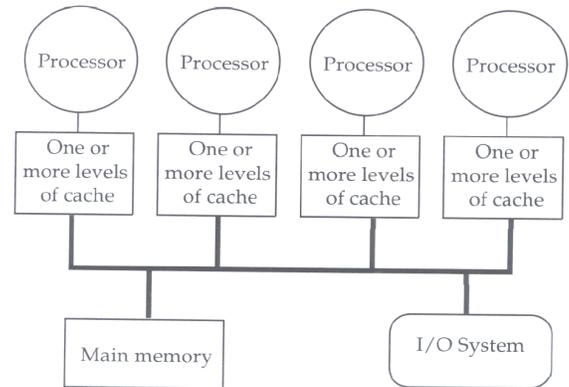
- bus-based
  - simple, but a bottleneck
  - simple cache coherency protocol
- physically centralized memory
- uniform memory access (UMA machine)
- Sequent Symmetry, SPARCcenter, Alpha-, PowerPC- or SPARC-based servers

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7

### Low-end MP



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8

## MIMD

### High-end

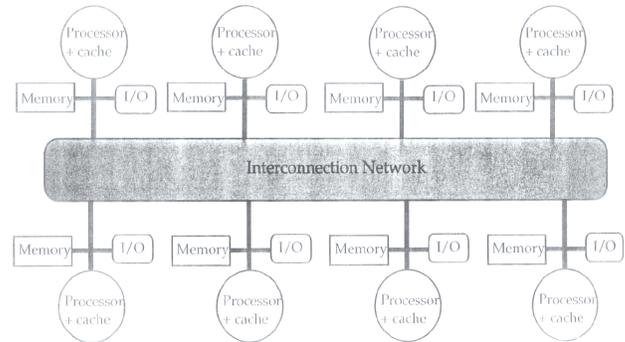
- higher bandwidth, multiple-path interconnect
  - more scalable
  - more complex cache coherency protocol (if shared memory)
  - longer latencies
- physically distributed memory
- non-uniform memory access (NUMA machine)
- could have processor clusters
- SGI Challenge, Convex Exemplar, Cray T3D, IBM SP-2, Intel Paragon

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9

### High-end MP

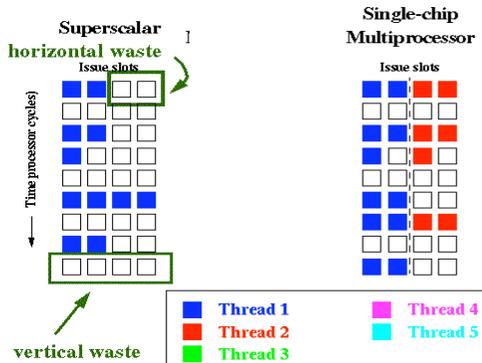


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## Comparison of Issue Capabilities



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## Shared Memory vs. Message Passing

### Shared memory

- + simple parallel programming model
  - global shared address space
  - not worry about data locality *but*
    - get better performance when program for data placement
    - lower latency when data is local
      - but* can do data placement if it is crucial, but don't have to
  - hardware maintains data coherence
    - synchronize to order processor's accesses to shared data
  - like uniprocessor code so parallelizing by programmer or compiler is easier
- ⇒ can focus on program semantics, not interprocessor communication

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## Shared Memory vs. Message Passing

### Shared memory

- + low latency (no message passing software) **but**  
*overlap of communication & computation*  
*latency-hiding techniques can be applied to message passing machines*
- + higher bandwidth for small transfers **but**  
*usually the only choice*

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## Shared Memory vs. Message Passing

### Message passing

- + abstraction in the programming model encapsulates the communication costs **but**  
*more complex programming model*  
*additional language constructs*  
*need to program for nearest neighbor communication*
- + no coherency hardware
- + good throughput on large transfers **but**  
*what about small transfers?*
- + more scalable (memory latency doesn't scale with the number of processors) **but**  
*large-scale SM has distributed memory also*
  - **hah!** so you're going to adopt the message-passing model?

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14

## Shared Memory vs. Message Passing

Why there was a debate

- little experimental data
- not separate implementation from programming model
- can emulate one paradigm with the other
  - MP on SM machine  
message buffers in local (to each processor) memory  
copy messages by ld/st between buffers
  - SM on MP machine  
ld/st becomes a message copy  
*slooooooooow*

Who won?

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15