Contiki – a Lightweight and Flexible Operating System for Tiny Networked Sensors

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Sensor OS trade-offs:

static vs dynamic

event-driven vs multi-threaded



What we have done

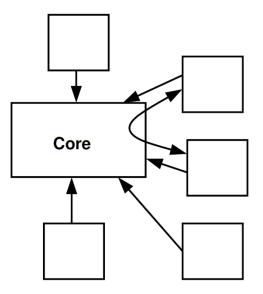
- Contiki an OS for sensor network nodes
- Ported Contiki to a number of platforms
 - MSP430, AVR, HC12, Z80, 6502, x86, ...
 - Simulation environment for BSD/Linux/Windows
- Built a few applications for experimental network deployments

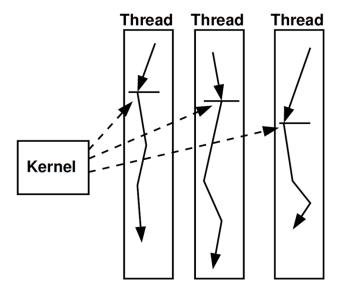


Contributions

Dynamic loading of programs

- Selective reprogramming
- Static vs dynamic linking
- Concurrency management
 mechanisms
 - Events vs threads
 - Trade-offs: preemption, size







Contiki design target

- "Mote"-class device
 - 10-100 kilobytes of code ROM
 - 1-10 kilobytes of RAM
 - Communication (radio)
- ESB from FU Berlin
 - MSP430, 2k RAM, 60k ROM









Contiki size (bytes)

Module	Code MSP430	Code AVR	RAM
Kernel	810	1044	10 + e + p
Program loader	658	-	8
Multi-threading library	582	678	8+s
Timer library	60	90	0
Memory manager	170	226	0
Event log replicator	1656	1934	200
µIP TCP/IP stack	4146	5218	18 + b



Run-time reprogramming and loadable programs



Reprogramming sensor nodes

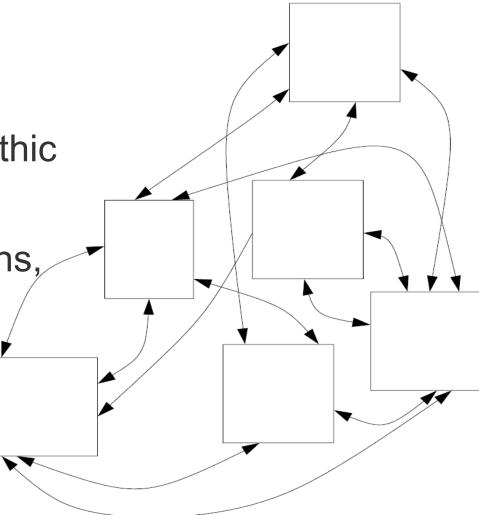
- . Software development for sensor nets
 - Need to reprogram many nodes quite often
- Utilize radio for reprogramming
 - Radio inherently broadcast
- Reprogram many nodes at once
 - Much faster than firmware download via cable or programming adapter
- Reprogram deployed networks



Traditional systems: entire system a monolithic binary

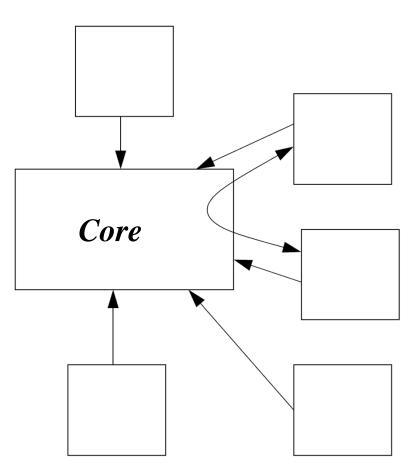
- Most systems statically linked at compile-time
 - Entire system is a monolithic binary
 - Compile-time optimizations, analysis possible
 - Makes code smaller
- But: hard to change
 - . Must download entire





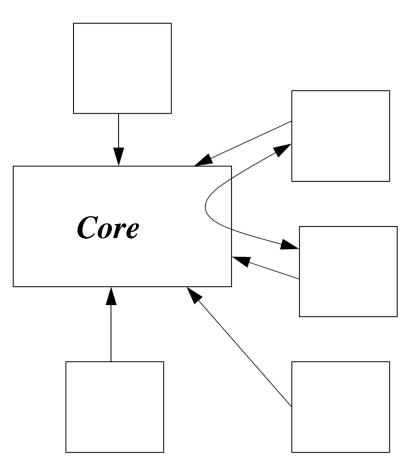
Contiki: loadable programs

- Contiki: one-way
 dependencies
 - Core resident in memory
 - Language run-time, communication
 - Programs "know" the core
 - Statically linked against core
- Individual programs can be
 Individual programs
 Con be



Loadable programs

- Programs can be loaded from anywhere
 - Radio (multi-hop, singlehop), EEPROM, etc
- During software development, usually change only one module





How well does it work?

- Works well
 - Program typically much smaller than entire system image (1-10%)
 - Much quicker to transfer over the radio
 - Reprogramming takes seconds
- Static linking can be a problem
 - Small differences in core means module cannot be run
 - We are implementing a dynamic linker



Concurrency in Contiki



Concurrency is tricky!

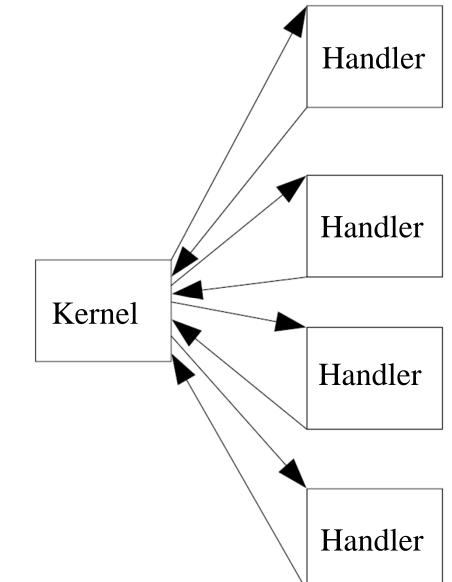
- Event-driven vs multi-threaded
- Event-driven (TinyOS)
 - Compact, low context switching overhead, fits well for reactive systems
 - Not suitable for e.g. long running computations
 - Public/private key cryptography
- Multi-threading
 - Suitable for long running computations
 - Requires more resources





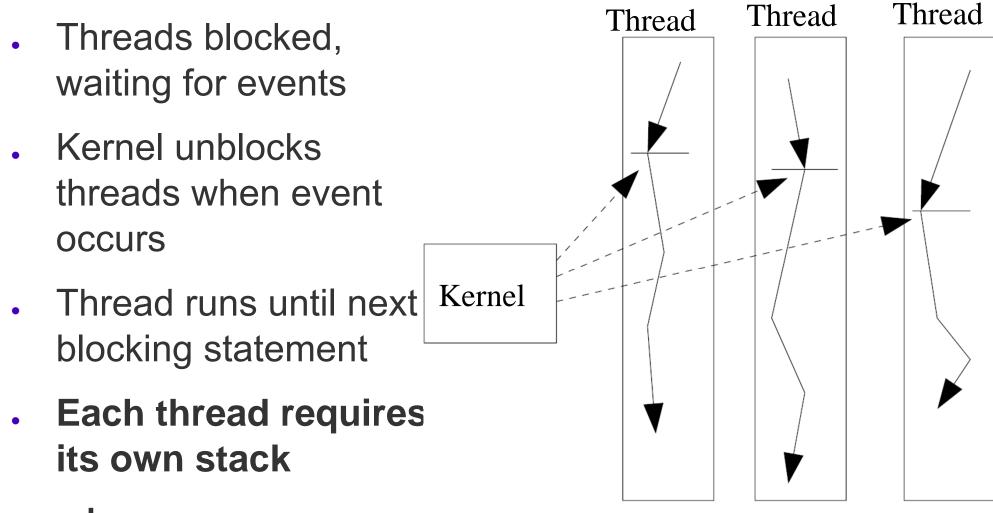
Event-driven

- Event-driven (TinyOS)
 - Processes do not run without events
 - Event occurs: kernel
 invokes event handler
 - Event handler runs to completion (explicit return;)





Multi-threaded



. Larger memory usage



Event-driven vs multi-threaded

Event-driven

- No wait () statements
- No preemption
- State machines
- + Compact code
- + Locking less of a problem
- + Memory efficient

Multi-threaded

- + wait() statements
- + Preemption possible
- + Sequential code flow
- Larger code overhead
- Locking problematic
- Larger memory requirements

Why don't we try to combine them?

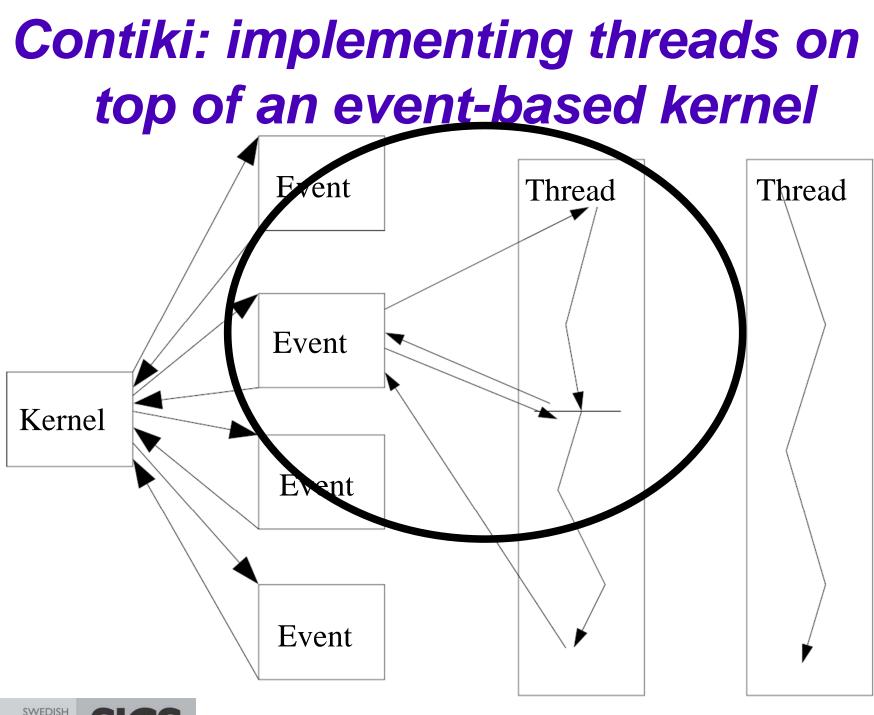


Contiki: event-based kernel with threads

- Contiki: kernel is event-based
 - Most programs run directly on top of the kernel
- Multi-threading implemented as a library
- Threads only used if **explicitly** needed
 - Long running computations, ...
- . Preemption possible
 - Responsive system with running computations













Conclusions



Conclusions

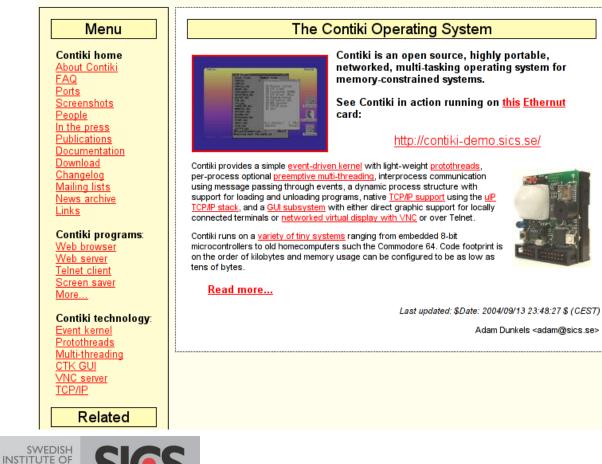
- Contiki OS for "mote"-class sensor nodes
- Contiki explores trade-offs in
 - static vs dynamic
 - event-driven vs multi-threaded
- Loadable programs, works well
 - Static linking can be a problem
- Threads on an event-driven kernel
 - Multi-threading suitable for certain applications





Thank you!

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COMPUTER

Latest news

2004-09-19

Version 1.2-devel1 is released! Go to the <u>download page</u> to get it or to the <u>changelog page</u> to check out what is new.

2004-09-15

A number of small bug devestating bugs had sneaked into the C64 1.2-devel0 release and a bugfixed version has been uploaded as 1.2-devel0-2. Get it on the <u>download page</u>.

2004-09-14

Contiki version 1.2-devel0 is released! Go to the <u>download page</u> to get it or to the <u>changelog page</u> to check out what is new.

2004-09-04

Automated daily development snapshots are now available. See the <u>download</u> <u>page</u> for details.

2004-08-27

Added Contiki - a Lightweight and Flexible Operating System for Tiny Networked Sensors to the publications page. The paper is to be presented at the First IEEE Workshop on Embedded Networked Sensors (EmNetS-I) in Tampa, Florida, USA on the 16th of November 2004.

News archive

Backup slides



Memory management

- . Memory allocated when module is loaded
 - Both ROM and RAM
 - Fixed block memory allocator
- Code relocation made by module loader
 - Exercises flash ROM evenly



Protothreads: light-weight stackless threads

- Protothreads: mixture between event-driven and threaded
 - A third concurrency mechanism
- Allows blocked waiting
- Requires per-thread **no stack**
- Each protothread runs inside a single C function
- 2 bytes of per-protothread state



