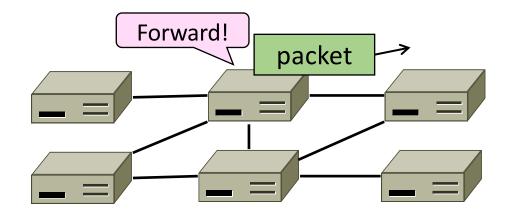
# Routing and forwarding

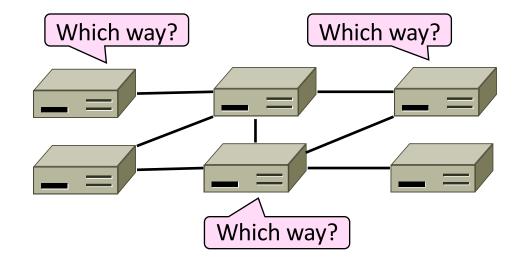
CSE 461 Ratul Mahajan

#### Recap: Routing versus Forwarding

 Forwarding is the process of sending a packet on its way



• <u>Routing</u> is the process of deciding in which direction to send traffic



## Overview of Internet Routing and Forwarding

- Hosts on same network have IPs in the same IP prefix
- Hosts send off-network traffic to the gateway router

- Routers discover routes to different prefixes (routing)
- Routers use <u>longest prefix matching</u> to send packets to the right next hop (forwarding)

## Longest Prefix Matching

 Prefixes in the forwarding table can overlap

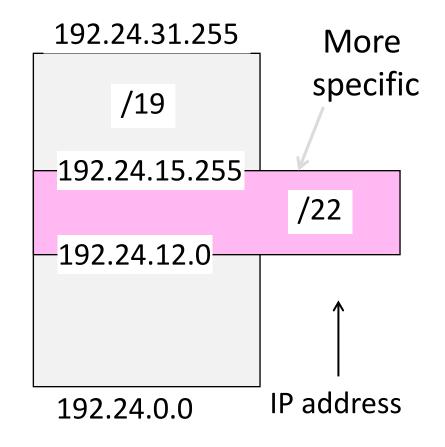
Prefix	Next Hop
0.0.0/0	А
192.24.0.0/19	В
192.24.12.0/22	С

- Longest prefix matching forwarding rule:
  - For each packet, find the longest prefix that contains the destination address, i.e., the most specific entry
  - Forward the packet to the next hop router for that prefix

# Longest Prefix Matching (2)

Prefix	Next Hop
192.24.0.0/19	D
192.24.12.0/22	В

192.24.6.0 $\rightarrow$  ?192.24.14.32 $\rightarrow$  ?192.24.54.0 $\rightarrow$  ?



# Flexibility of Longest Prefix Matching

- Can provide default behavior, with less specifics
  - Send traffic going outside an organization to a border router (gateway)
- Can special case behavior, with more specifics
  - For performance, economics, security, ...

# Performance of Longest Prefix Matching

- Uses hierarchy for a compact table
  Relies on use of large prefixes
- Lookup more complex than table
  - Used to be a concern for fast routers
  - Not an issue in practice these days

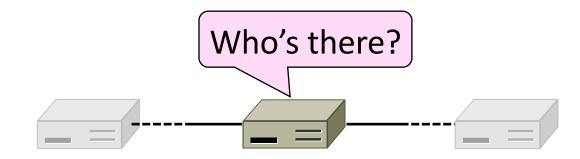
# Goals of Routing Algorithms

• We want several properties of any routing scheme:

Property	Meaning	
Correctness	Finds paths that work	
Efficient paths	Uses network bandwidth well	
Fair paths	Doesn't starve any nodes	
Fast convergence	Recovers quickly after changes	
Scalability	Works well as network grows large	

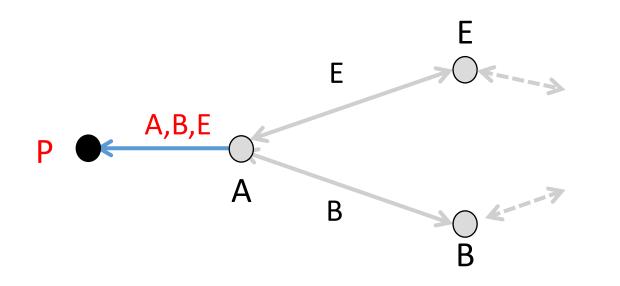
# Rules of Fully Distributed Routing

- All nodes are alike; no controller
- Nodes learn by exchanging messages with neighbors
- Nodes operate concurrently
- There may be node/link/message failures



Simple routing that obeys the rules

Send out routes for hosts you have paths to
And the routes they've sent you



- This works
  - All routers find a path to all hosts
- But scales poorly!

#### Recall: Internet Size

- Over 4 billion people
- 50B devices connect

#### Impact of Network Growth

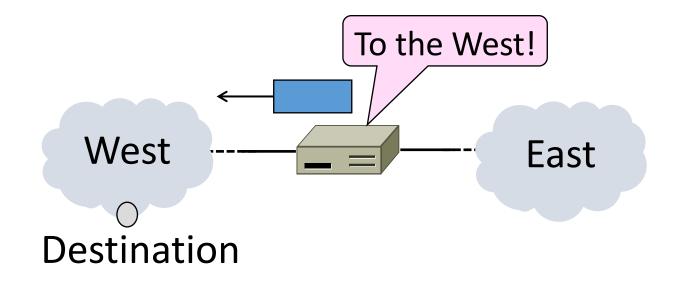
- 1. Forwarding tables grow
  - Larger router memories, may increase lookup time
- 2. Routing messages grow
  - Need to keeps all nodes informed of larger topology
- 3. Routing computation grows
  - Shortest path calculations grow faster than the network

# Techniques to Scale Routing

- First: Network hierarchy
  Route to network regions
- Next: IP prefix aggregation
  Combine, and split, prefixes

# Scaling Idea: Hierarchical Routing

- Scale routing using hierarchy with regions
  - Route to regions, not individual nodes

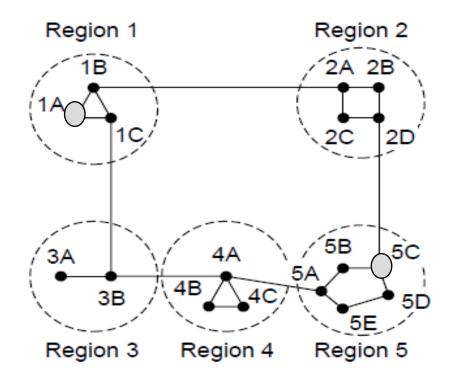


### **Hierarchical Routing**

- Introduce a larger routing unit

  - Region, e.g., ISP network
- Route first to the region, then to the IP prefix within the region
  - Hide details within a region from outside of the region

## Hierarchical Routing (2)



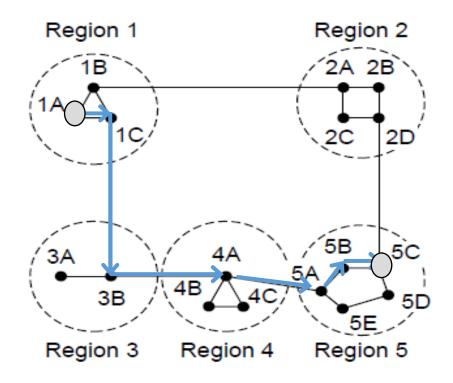
Full table for 1A

Dest.	Line	Hops
1A	-	-
1B	1B	1
1C	1C	1
2A	1B	2
2B	1B	3
2C	1B	3
2D	1B	4
ЗA	1C	3
3B	1C	2
4A	1C	3
4B	1C	4
4C	1C	4
5A	1C	4
5B	1C	5
5C	1B	5
5D	1C	6
5E	1C	5

Hierarchical table for 1A

Dest.	Line	Hops
1A	_	-
1B	1B	1
1C	1C	1
2	1B	2
3	1C	2
4	1C	3
5	1C	4

## Hierarchical Routing (3)



Full	tab	le for	1A
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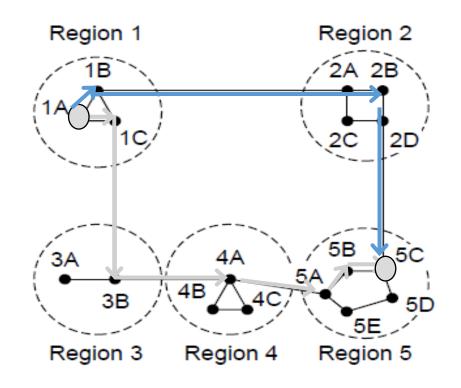
Dest.	Line	Hops
1A	-	-
1B	1B	1
1C	1C	1
2A	1B	2
2B	1B	3
2C	1B	3
2D	1B	4
ЗA	1C	3
3B	1C	2
4A	1C	3
4B	1C	4
4C	1C	4
5A	1C	4
5B	1C	5
5C	1B	5
5D	1C	6
5E	1C	5

Hierarchical table for 1A

Dest.	Line	Hops
1A	-	-
1B	1B	1
1C	1C	1
2 3	1B	2
	1C	2
4	1C	3
5	1C	4

# Hierarchical Routing (4)

• Penalty is longer paths



Full table for 1A			
Dest.	Line	Hops	
1A	_	-	
1B	1B	1	
1C	1C	1	
2A	1B	2 3	
2B	1B	3	
2C	1B	3	
2D	1B	4	
ЗA	1C	3	
3B	1C	2	
4A	1C	3	
4B	1C	4	
4C	1C	4	
5A	1C	4	
5B	1C	5	
5C	1B	∛5	
5D	1C	6	
5E	1C	5	

Hierarchical table for 1A

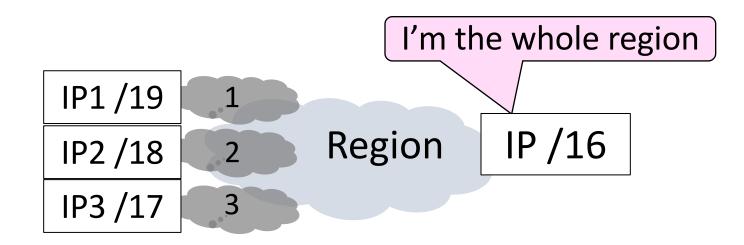
Dest.	Line	Hops
1A	-	-
1B	1B	1
1C	1C	1
2 3	1B	2
	1C	2
4	1C	3
5	1C	4
-		

1C is best route to region 5, except for destination 5C

### Observations

- Outside a region, nodes have <u>one route</u> to all hosts within the region
  - This gives savings in table size, messages and computation
- However, each node may have a <u>different route</u> to an outside region
  - Routing decisions are still made by individual nodes; there is no single decision made by a region

Enabling hierarchical routing: IP Prefix Aggregation and Subnets Scale routing by adjusting the size of IP prefixes
Split (subnets) and join (aggregation)



#### Recall

• IP addresses are allocated in blocks called IP prefixes, e.g., 18.31.0.0/16

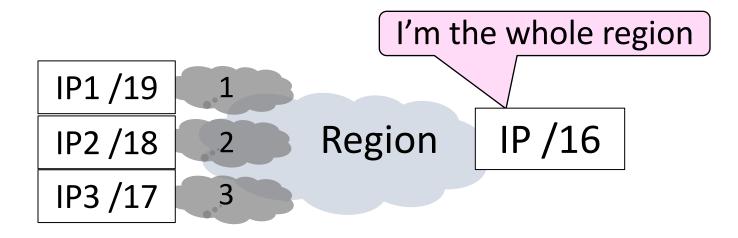
• Hosts on one network in same prefix

- "/N" prefix has the first N bits fixed and contains 2<sup>32-N</sup> addresses
  - E.g., a "/24" has 256 addresses
- Routers keep track of prefix lengths
  - Use it as part of longest prefix matching

Routers can change prefix lengths without affecting hosts

#### Prefixes and Hierarchy

IP prefixes help to scale routing, but can go further
Use a less specific (larger) IP prefix as a name for a region

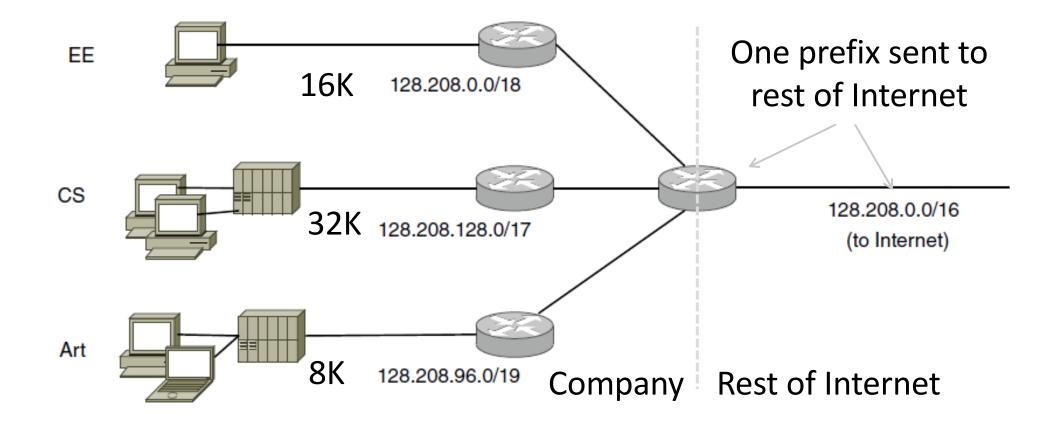


# Subnets and Aggregation

- Two use cases for adjusting the size of IP prefixes; both reduce routing table
- 1. Subnets
  - Internally split one large prefix into multiple smaller ones
- 2. Aggregation
  - Join multiple smaller prefixes into one large prefix

#### Subnets

• Internally split up one IP prefix



#### Aggregation

#### • Externally join multiple separate IP prefixes

