### Security and Project 3

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#### Administrivia

- HW 4 due 05/23
- Project 1 Updated Submission due 11 PM Friday 05/12
  - In-person/Zoom OH on Friday:
    - 11:30 AM 12:30 PM (Location TBD)
    - 2:30 PM 3:30 PM (Location TBD)
    - Will confirm on ed today

#### Symmetric (Secret Key) Encryption

• Alice and Bob have the same secret key, K<sub>AB</sub> • Anyone with the secret key can encrypt/decrypt



#### Public Key (Asymmetric) Encryption

• Alice and Bob have public/private key pairs ( $K_{B}/K_{B_{-1}}$ )

Public keys are well-known, private keys are secret



#### Public Key Encryption (2)

- Alice encrypts w/ Bob's pubkey K<sub>B</sub>; anyone can send
- Bob decrypts w/ his private key  $K_B^{-1}$ ; only he can



#### Man-in-the-Middle Attacks



#### How can we trust a public key?

#### Certificates

A certificate binds pubkey to identity, e.g., domain
Distributes public keys when signed by a party you trust
Commonly in a format called X.509



#### PKI (Public Key Infrastructure)

Adds hierarchy to certificates to let parties issue
Issuing parties are called CAs (Certificate Authorities)



#### PKI (2)

- •Need public key of PKI root and trust in servers on path to verify a public key of website ABC
  - •Browser has Root's public key
  - •{RA1's key is X} signed Root •{CA1's key is Y} signed RA1
  - •{ABC's key Z} signed CA1



### PKI (3)

# •Browser/OS has public keys of the trusted roots of PKI

#### •>100 root certificates!

- •That's a problem ...
- Inspect your web browser

Certificate for wikipedia.org issued by DigiCert

	Certificate viewer				
eneral <u>D</u> etails					
This certificate has been verified for the following uses:					
SSL Server Certificate					
Issued To					
Common Name (CN)	*.wikipedia.org				
Organization (O)	Wikimedia Foundation, Inc.				
Organizational Unit (OU)	<not certificate="" of="" part=""></not>				
Serial Number	05:DF:E8:FF:15:B8:63:CC:C6:89:C7:8E:64:0C:FE:8B				
Issued By					
Common Name (CN)	DigiCert High Assurance CA-3				
Organization (O)	DigiCert Inc				
Organizational Unit (OU)	www.digicert.com				
Validity					
Issued On	12/08/2011				
Expires On	12/12/2012				
Fingerprints					
SHA1 Fingerprint	03:47:7F:F5:F6:3B:F5:B6:10:C0:7D:65:9A:7B:A9:12:D3:20:83:68				
MD5 Fingerprint	C0:C8:F7:A0:33:20:A2:D4:2E:27:65:73:42:4C:A0:24				
	Close				

#### PKI (4)

Real-world complication:
Public keys may be compromised
Certificates must then be revoked

## •PKI includes a CRL (Certificate Revocation List) •Browsers use to weed out bad keys

$\sim$			•
$(` \alpha r$	t1†1	00	$t \cap$
	ιII	Ua.	ιC

www.google.com	GTS CA 1C3	GTS Root R1	GlobalSign Root CA
Subject Name			
Country	BE		
Organization	GlobalSign nv-sa		
Organizational Unit	Root CA		
Common Name	GlobalSign Root CA		
Issuer Name			
Country	BE		
Organization	GlobalSign nv-sa		
Organizational Unit	Root CA		
Common Name	GlobalSign Root CA		
Validity			
Not Deferre	Tue 01 Cap 1000 10:0	0:00 CMT	

 Not Before
 Tue, 01 Sep 1998 12:00:00 GMT

 Not After
 Fri, 28 Jan 2028 12:00:00 GMT

### Bufferbloat

"Bufferbloat is a cause of high latency in packet-switched networks caused by excess buffering of packets" – Wikipedia

#### Let's revisit network delays

- Previously we studied how to calculate delays for a network
  - RTT: ~typically defines the latency for a small packet round-trip
  - Assumption: Network is unloaded (nothing else going on)
  - This is not often true in practice
  - What changes when something else is going on?
    - <u>https://www.waveform.com/tools/bufferbloat</u>
    - Observe unloaded vs loaded latency
  - What's going on here?



#### Let's investigate this further...

• Recall the equations to calculate latency

```
Latency = Propagation + Transmit + Queue
Propagation = Distance/SpeedOfLight
Transmit = Size/Bandwidth
```

- Queueing delay: How long packet spends in switch buffers
  - Queueing delay  $\sim$  # packets queued
  - For unloaded network: 0
  - For loaded network: depends how packets are sent

#### What happens with TCP Reno

- Host doesn't know the bandwidth of the bottleneck link.
- TCP Reno relies solely on packet losses to guide how fast to send.
  - It keeps sending faster and faster until a packet drops.
- With a queue, this can fill up the queue pretty quickly.
- Larger the switch queue capacity, more the delay
  - Called Bufferbloat



#### **Bufferbloat – Problem**

- Suppose h1 knows to send at 1.5 Mb/s, what's the RTT when the queue is full?
  - ...when it's not full?
- TCP at the end of the day will operate at the bottleneck bandwidth, but is it necessary to fill up the queue?



#### FIGURE 1: DELIVERY RATE AND ROUND-TRIP TIME VS. INFLIGHT



#### How do solve this this?

- Problem: Packet loss is a bad signal of congestion
  - Aside: packet losses because of other reasons can also mess up performance
- What else can we do?
- Proposal #1: Active Queue Management
  - Switch tells sender that it's running close to congestion
  - Explicit Congestion Notification (ECN): Mark a bit in the packet IP header
  - Challenges: Requires switch hardware support; If mixed with regular TCP flows, it takes over all the bandwidth because it works too well
  - Used in datacenters (DCTCP), not on the internet
- Proposal #2: Try to find the optimal point using measurements
  - BBR from Google is one deployed example

#### Bufferbloat aware transport: BBR

- Optimum operating point:
  - (Max BW, Min RTT)
  - Product is called Bandwidth Delay Product (BDP)
- BBR
  - Estimate min RTT, bottleneck bandwidth via probing and measuring performance.
  - Keep inflight bytes equal to BDP



FIGURE 1: DELIVERY RATE AND ROUND-TRIP TIME VS. INFLIGHT



#### Project 3 – Goal

- Simulate bufferbloat problem.
- See the worse performance when queue size is larger
- See the difference between TCP Reno and TCP BBR.

#### **Experiment Setup**

- Long-lived TCP flow from h1 to h2
  - Simulate background traffic
- Back-to-back ping from h1 to h2
  - Measure RTT
- Spawn a webserver on h1 and periodically fetch a page
  - Simulate more important load
  - Measure time
- Plot time series of RTT and number of queued packets.



- Run the experiment with
  - Q=20 and Q=100
  - Reno and BBR
  - 4 experiments total

#### Detour – Hypothesize



In groups of 3-4,...

- In your own words, what is bufferbloat problem?
- For each of the 4 experiments (Q=20 or 100; and with Reno or BBR),
  - How do the webpage fetch time compare?
- How would plot between queue size and time look like for TCP Reno?

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	Q=20	Q=100
Reno	<; <;=;>	=;> <;=;>
BBR	<;	=;>

#### Setup

- Use Mininet VM (same as Project 2)
- Get the starter code and install dependencies

cd ~

wget

https://courses.cs.washington.edu/courses/cse461/22wi/projects/project3/resources/project3.
zip

unzip project3.zip

sudo apt-get

update

```
sudo apt install python3-pip
```

```
sudo python3 -m pip install mininet matplotlib
```

#### Starter Code

- run.sh
  - Run the entire experiment
    - Run bufferbloat.py on q=20 and q=100
    - Generate latency and queue length graphs
- bufferbloat.py
  - Complete the TODOs
    - Setup the mininet topology and the experiment
    - Write shell commands to do the measurements

#### Long-lived TCP Flow

- Starter code sets up iperf server on h2
- Goal: start iperf client on h1, connect to h2
  - Should be "long-lasting", i.e. for time specified by --time parameter
- How do I connect to a certain IP or make the connection long-lasting?
  - man pages are your friend!
  - type `man iperf` in a Linux terminal



#### **Ping Train**

- Goal: Start "ping train" between h1 and h2
  - Pings should occur at 10 per second interval
  - Should run for entire experiment
- How do I specify the ping interval and how long the ping train runs?
  - man pages are your friend!
  - type `man ping` in a Linux terminal
- Write the RTTs recorded from `ping` to {args.dir}/ping.txt
  - See starter code comments for more detail



#### Download Webpage with curl

- Starter code spawns webserver on h1
- Goal: Use `curl` to measure fetch time to download webpage from h1
  - Starter code has hint on formatting curl command
  - Make sure `curl` doesn't output an error
    - Errors report very small latency
- No need to plot fetch times; just need to report average fetch time for each experiment.



#### Plotting

- Starter code contains scripts for plotting, `plot\_queue.py`, `plot\_ping.py`
  - Expects queue occupancy in \$dir/q.txt, ping latency in \$dir/ping.txt
  - Plots are useful for debugging!
- Part 3, run same experiments with TCP BBR instead of TCP Reno
  - How do you expect the graph outputs to differ?





Q = 100

