TCP contd ….

Last class

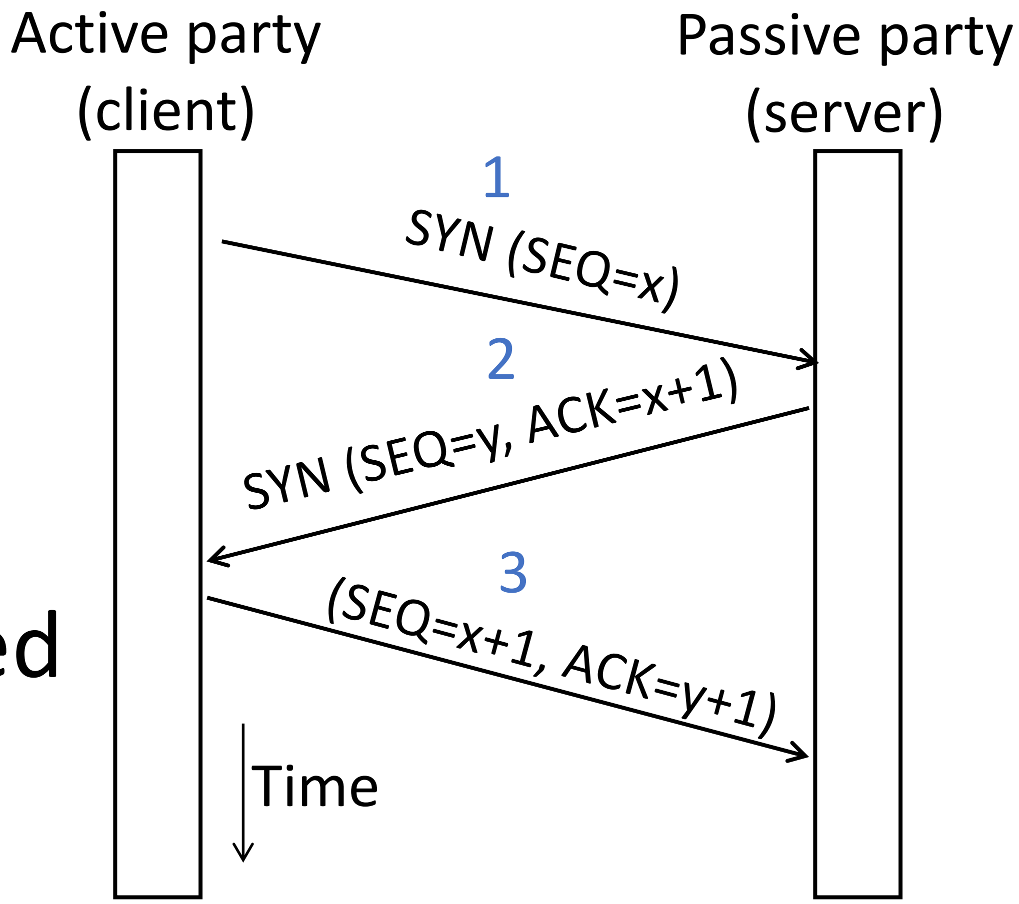
- Connection setup

This class

- Connection release

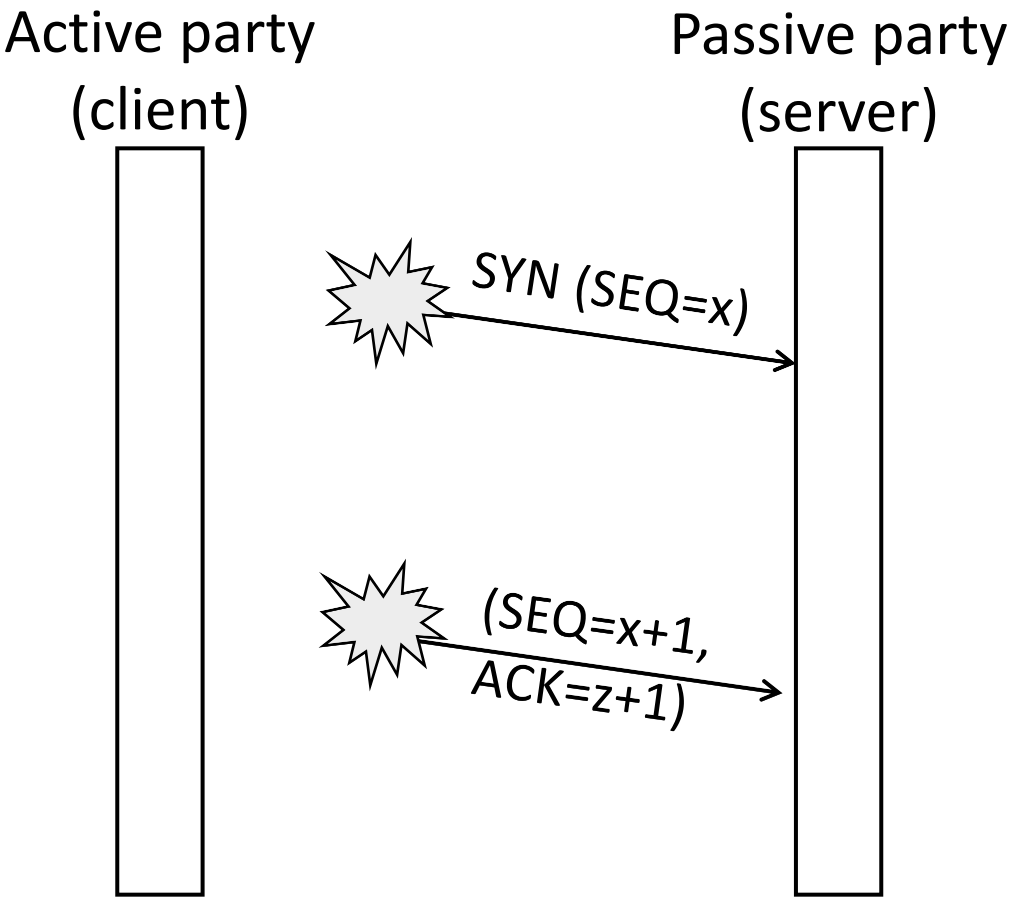
Recap: Connection setup

- Three-way handshake:
  - Client sends SYN(x)
  - Server replies with SYN(y)ACK(x+1)
  - Client replies with ACK(y+1)
  - SYNs are retransmitted if lost
- Sequence and ack numbers carried on further segments



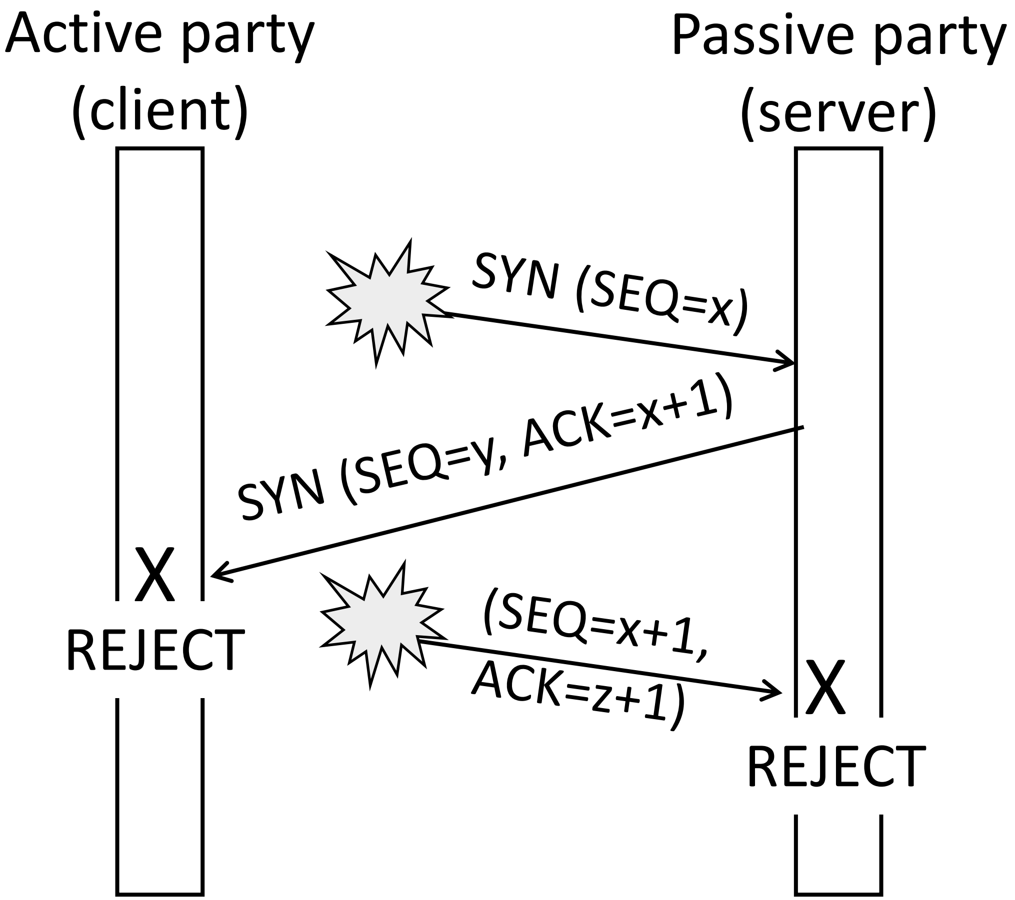
Three-Way Handshake

- Suppose delayed, duplicate copies of the SYN and ACK arrive at the server!
  - Improbable, but anyhow …



Three-Way Handshake

- Suppose delayed, duplicate copies of the SYN and ACK arrive at the server!
  - Improbable, but anyhow …
- Connection will be cleanly rejected on both sides ☺



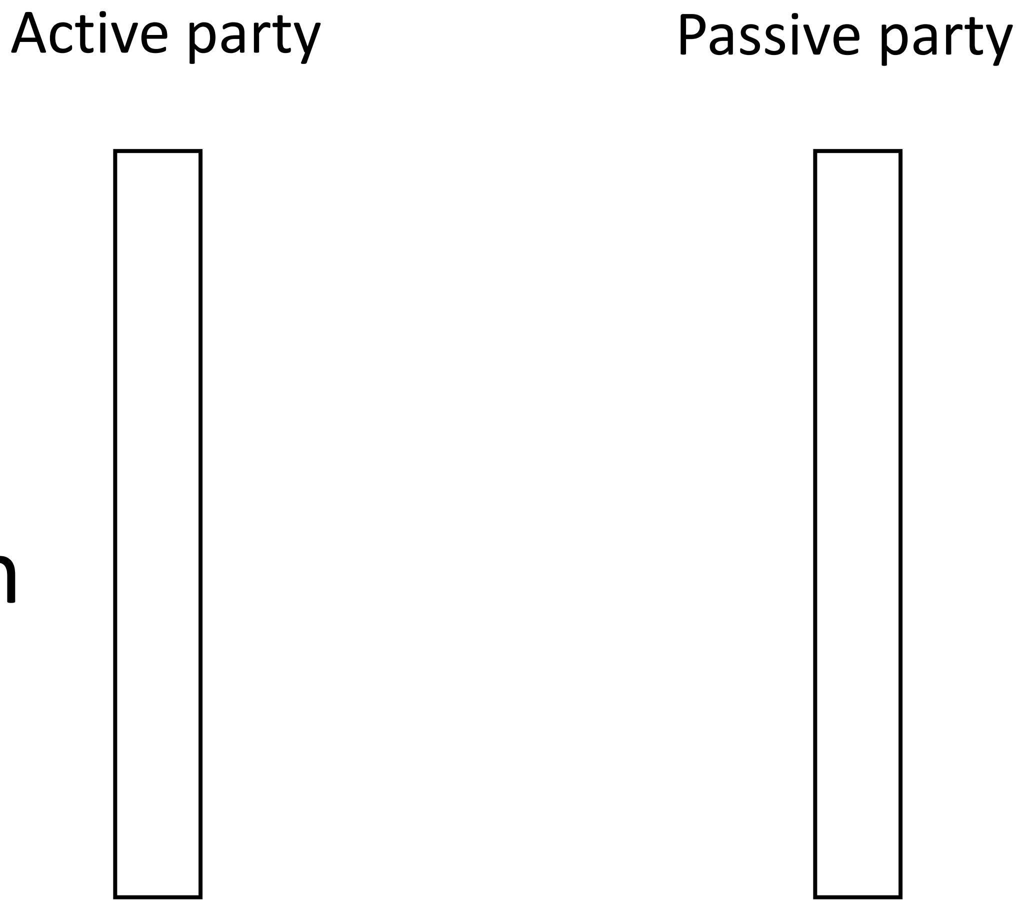
Connection Release

- Orderly release by both parties when done
  - Delivers all pending data and “hangs up”
  - Cleans up state in sender and receiver
- Key problem is to provide reliability while releasing
  - TCP uses a “symmetric” close in which both sides shutdown independently

TCP Connection Release

- Two steps:
  - Active sends FIN(x), passive ACKs
  - Passive sends FIN(y), active ACKs
  - FINs are retransmitted if lost

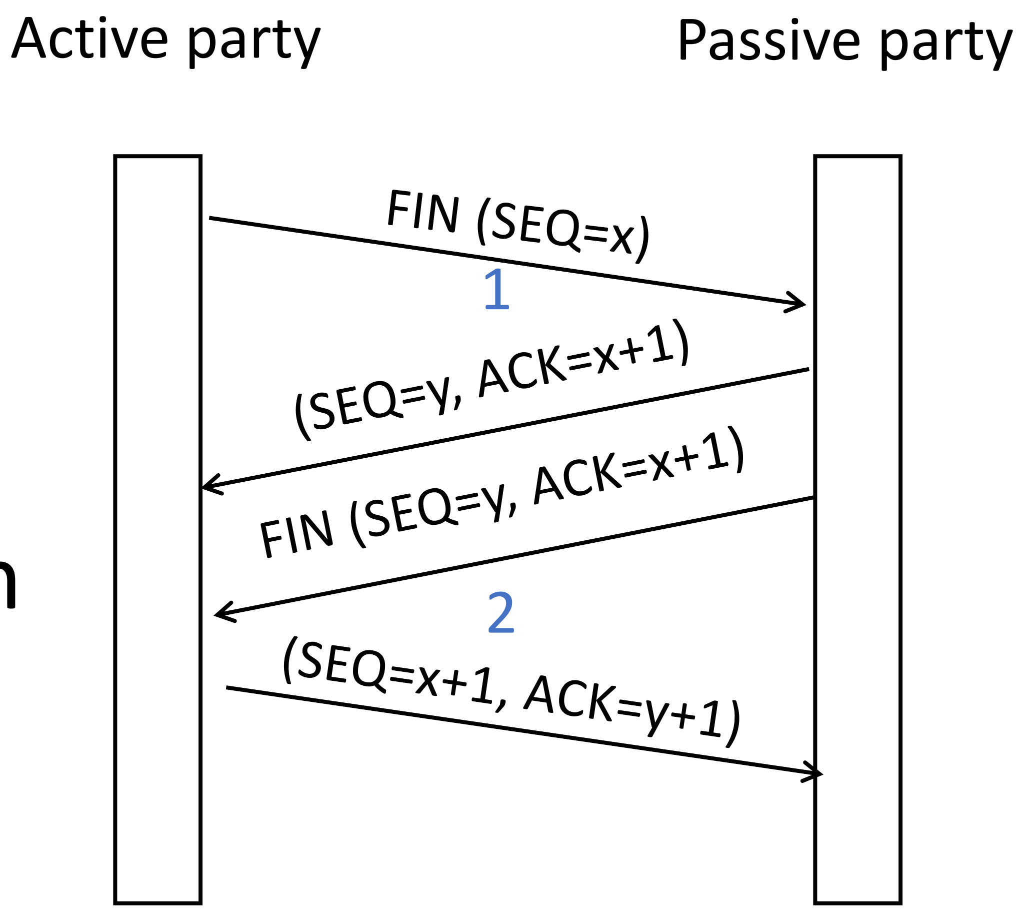
- Each FIN/ACK closes one direction of data transfer



TCP Connection Release (2)

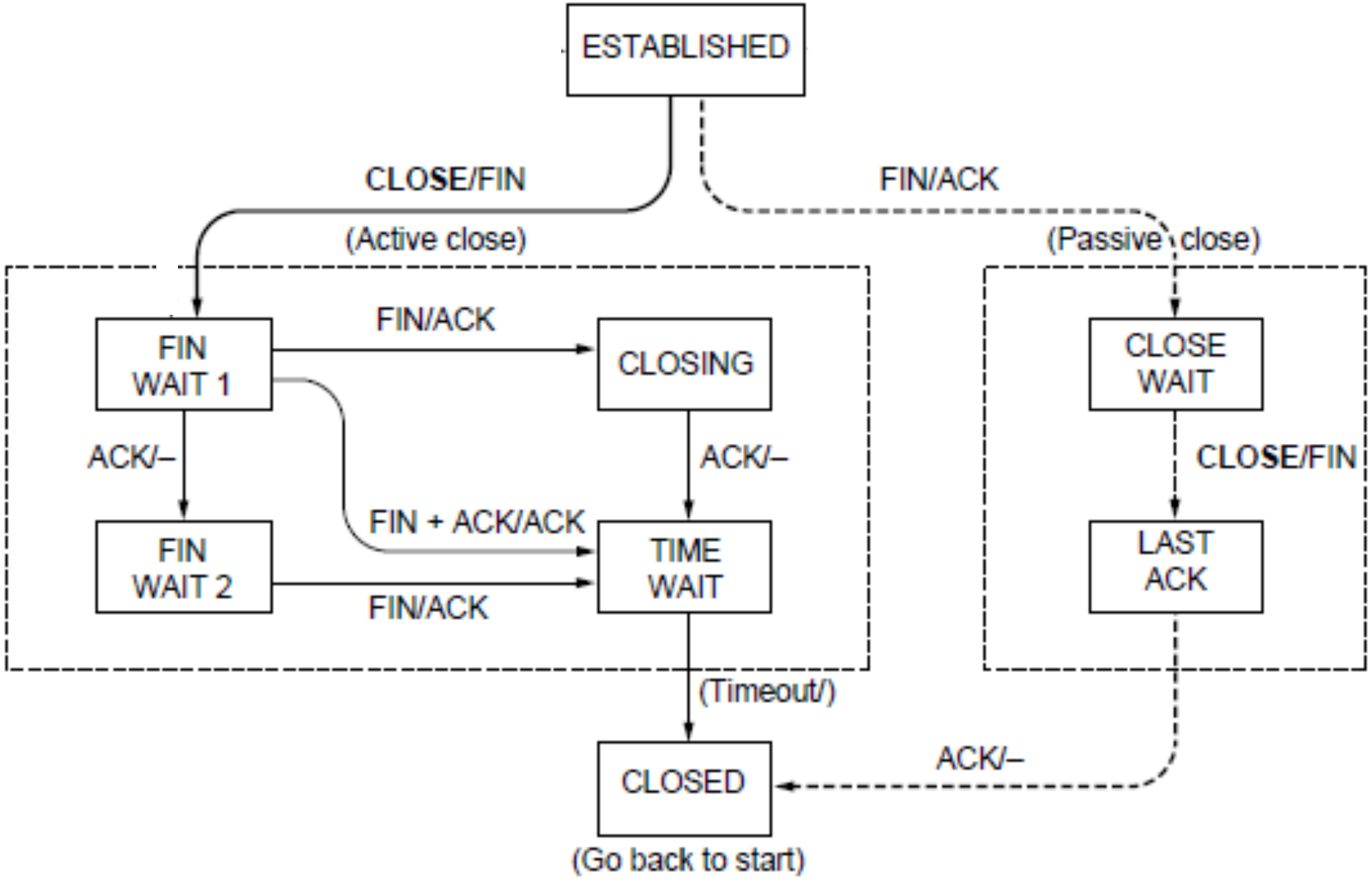
- Two steps:
  - Active sends FIN(x), passive ACKs
  - Passive sends FIN(y), active ACKs
  - FINs are retransmitted if lost

- Each FIN/ACK closes one direction of data transfer



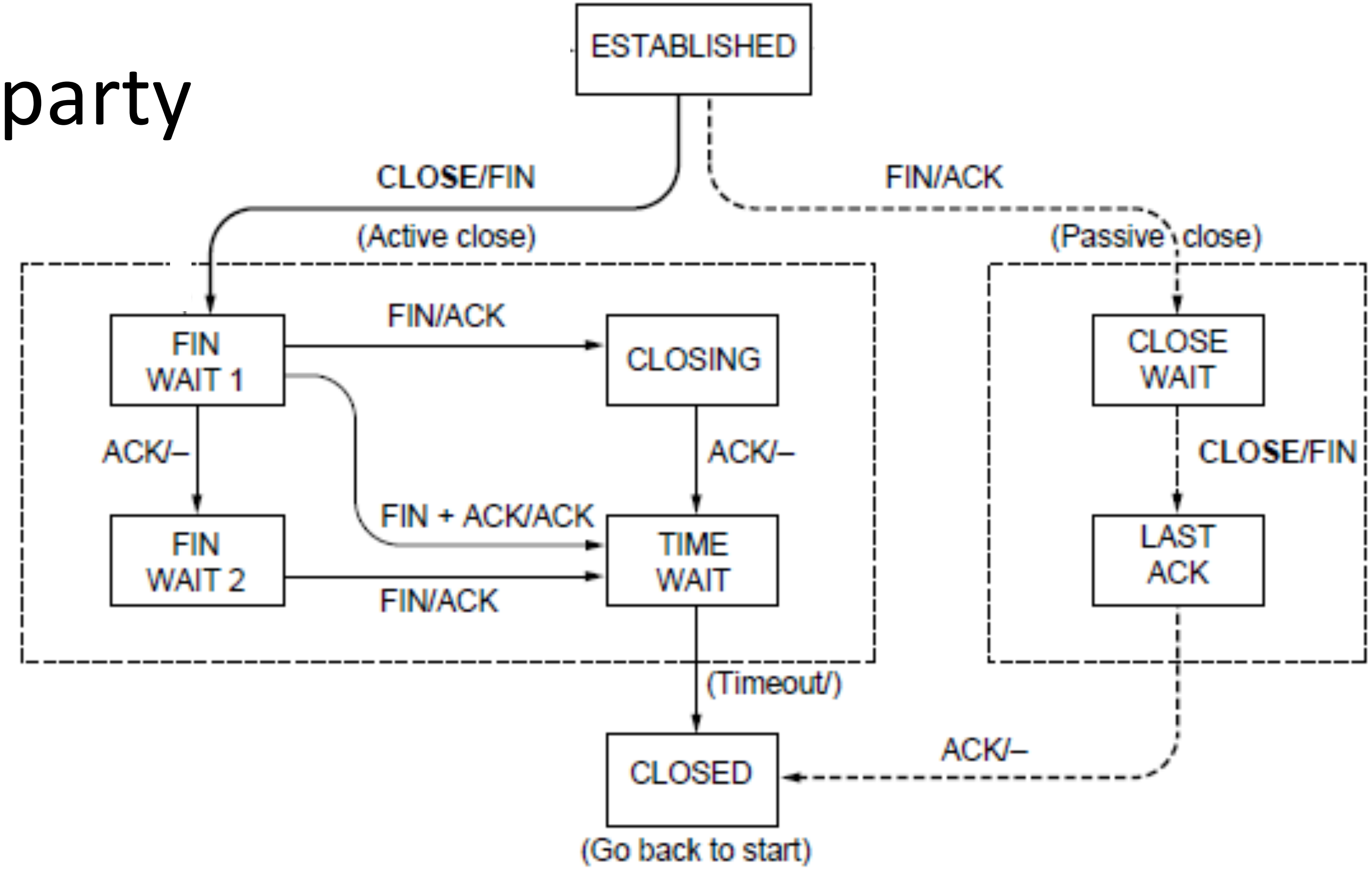
TCP Connection State Machine

Both parties run instances of this state machine



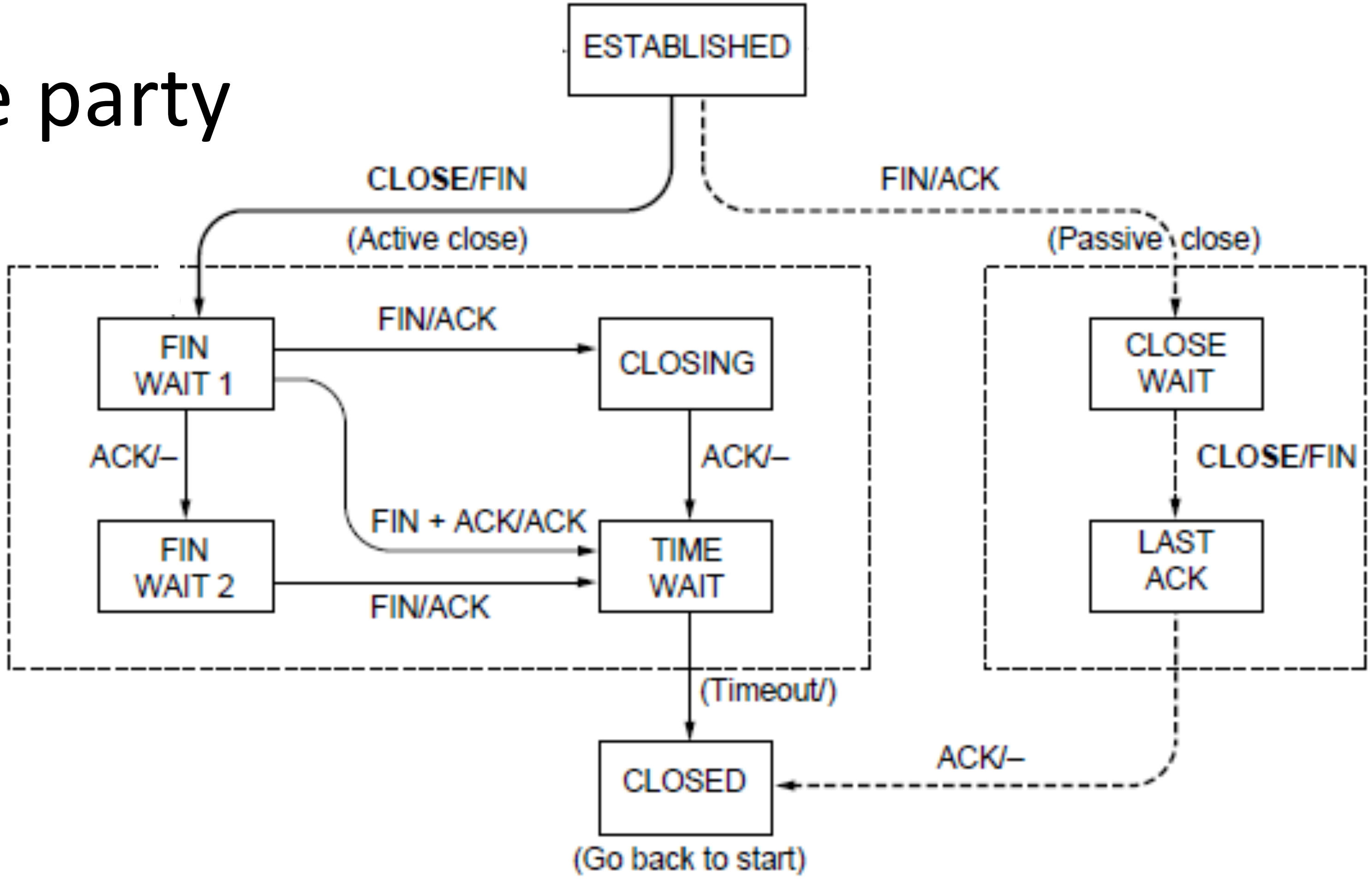
TCP Release

- Follow the active party



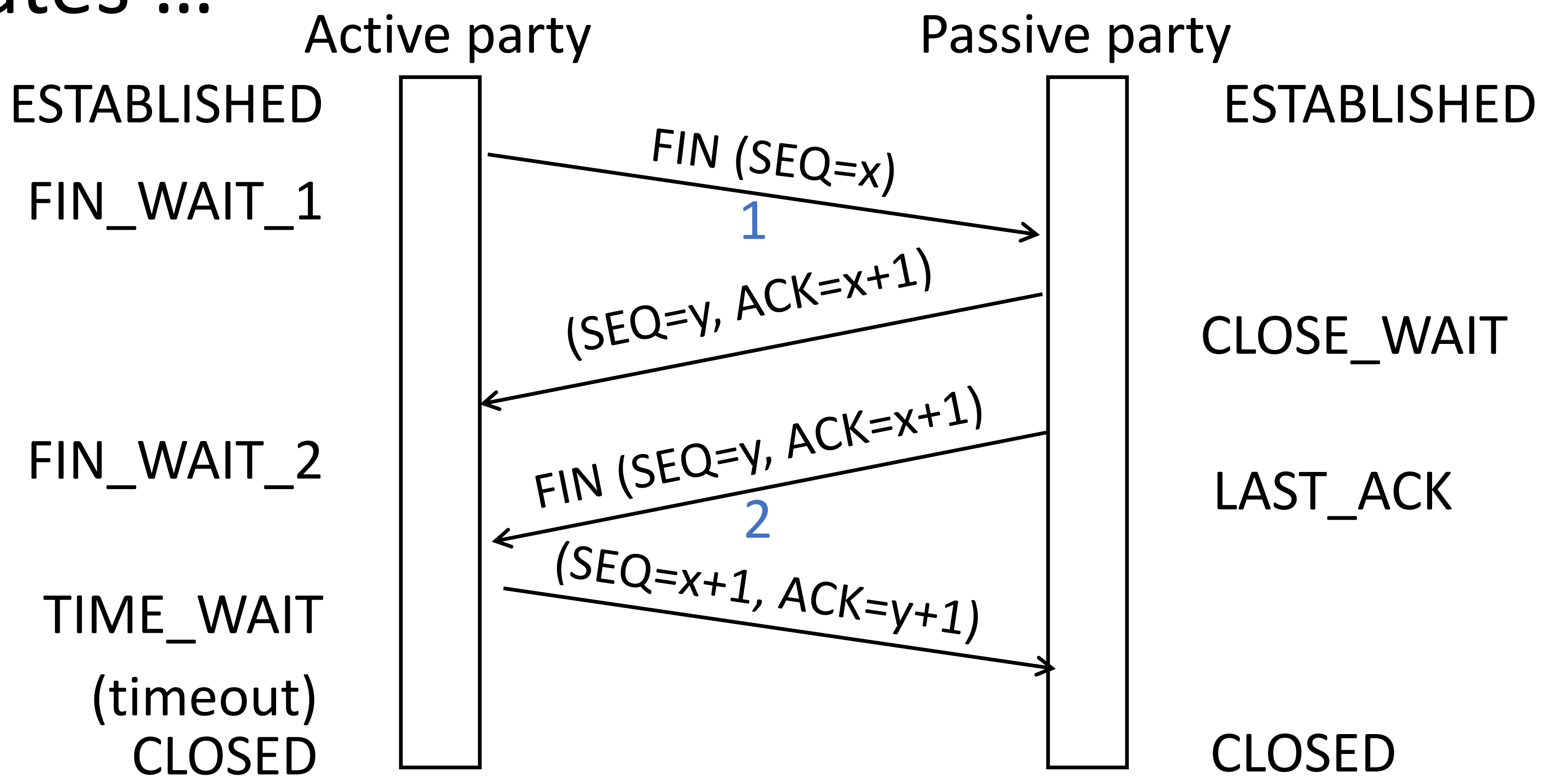
TCP Release (2)

- Follow the passive party



TCP Release (3)

- Again, with states …



TIME_WAIT State

- Wait a long time after sending all segments and before completing the close
  - Two times the maximum segment lifetime of 60 seconds
- Why?

TIME_WAIT State

- Wait a long time after sending all segments and before completing the close
  - Two times the maximum segment lifetime of 60 seconds
- Why?
  - ACK might have been lost, in which case FIN will be resent for an orderly close
  - Could otherwise interfere with a subsequent connection

Flow Control

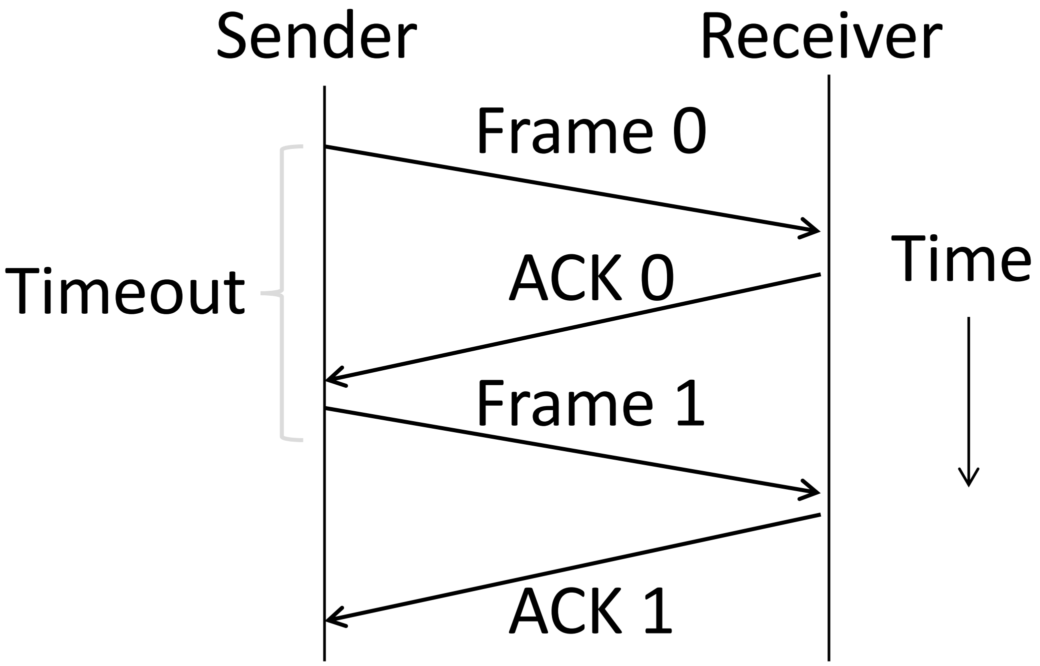
Flow control goal

Match transmission speed to reception capacity

- Otherwise data will be lost

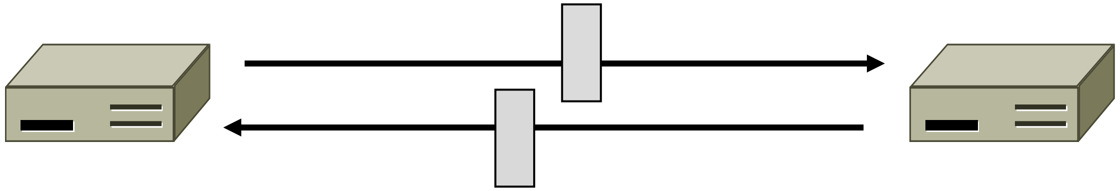
ARQ: Automatic repeat query

- ARQ with one message at a time is Stop-and-Wait



Limitation of Stop-and-Wait

- It allows only a single message to be outstanding from the sender:
  - Fine for LAN (only one frame fits in network anyhow)
  - Not efficient for network paths with longer delays

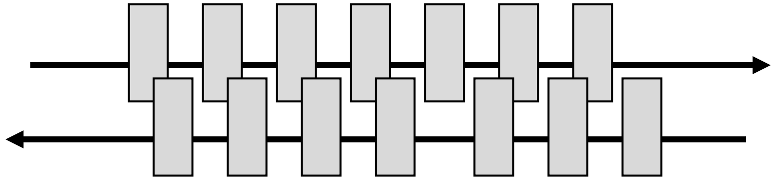


Limitation of Stop-and-Wait (2)

- Example: B=1 Mbps, D = 50 ms
  - RTT (Round Trip Time) = 2D = 100 ms
  - How many packets/sec?
    - 10
  - Usage efficiency if packets are 10kb?
    - $(10,000 \times 10) / (1 \times 10^6) = 10\%$
  - What is the efficiency if B=10 Mbps?
    - 1%

Sliding Window

- Generalization of stop-and-wait
  - Allows W packets to be outstanding
  - Can send W packets per RTT (=2D)



  - Pipelining improves performance
  - Need W=2BD to fill network path

Sliding Window (2)

What W will use the network capacity with 10kb packets?

- Ex: B=1 Mbps, D = 50 ms
  - 2BD = 2 x 10^6 x 50/1000 = 100 Kb
  - W = 100 kb/10 = 10 packets

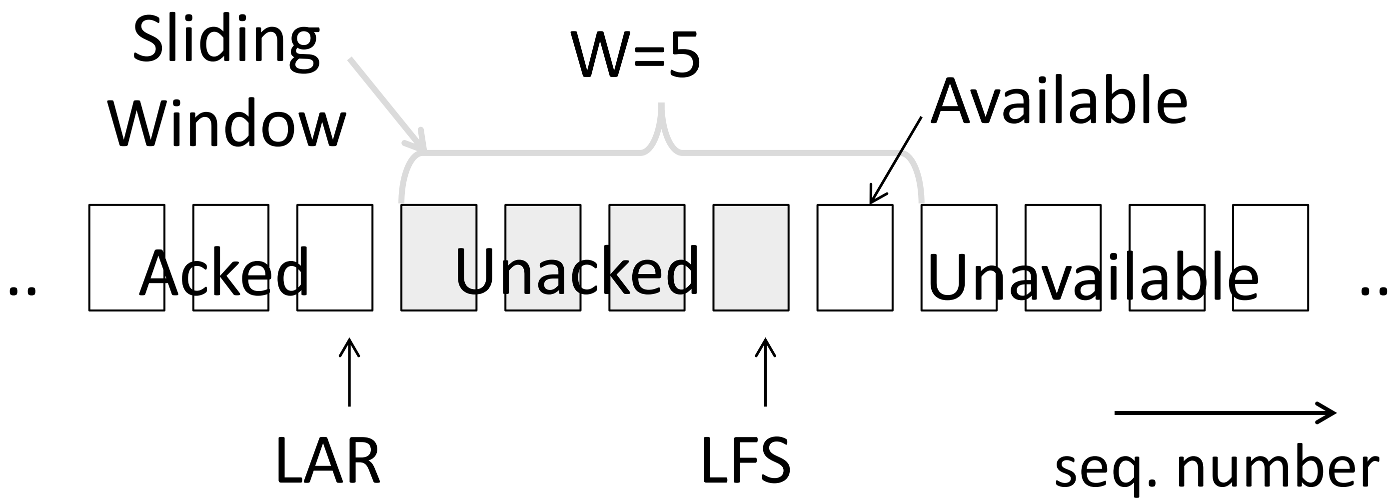
- Ex: What if B=10 Mbps?
  - W = 100 packets

Sliding Window Protocol

- Many variations, depending on how buffers, acknowledgements, and retransmissions are handled
- <u>Go-Back-N</u>
  - Simplest version, can be inefficient
- <u>Selective Repeat</u>
  - More complex, better performance

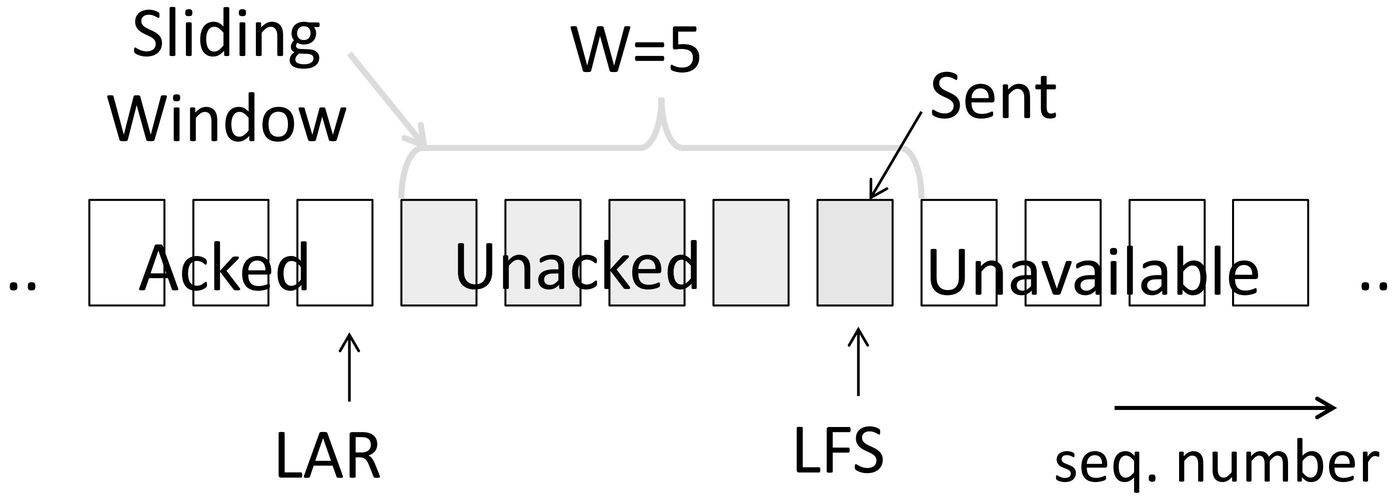
Sender Sliding Window

- Sender buffers up to W segments until they are acknowledged
  - LFS=LAST FRAME SENT, LAR=LAST ACK REC'D
  - Sends while LFS – LAR ≤ W



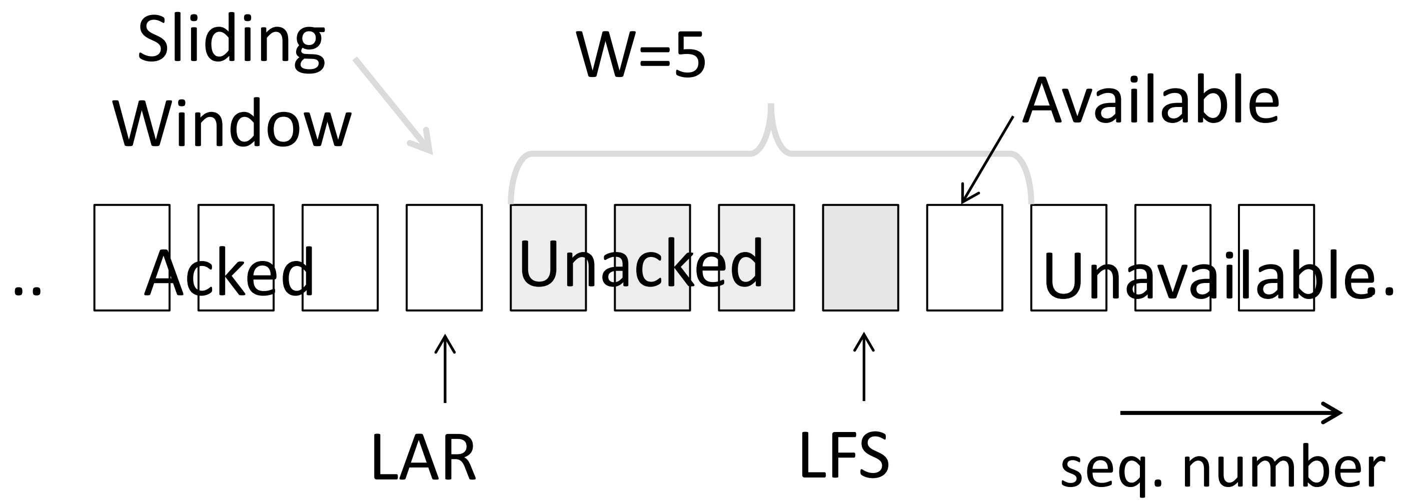
Sender Sliding Window (2)

- Transport accepts another segment of data from the Application ...
  - Transport sends it (LFS–LAR → 5)



Sender Sliding Window (3)

- Next higher ACK arrives from peer…
  - Window advances, buffer is freed
  - LFS–LAR → 4 (can send one more)



Receiver Sliding Window – Go-Back-N

- Receiver keeps only a single packet buffer for the next segment
  - State variable, LAS = LAST ACK SENT
- On receive:
  - If seq. number is LAS+1, accept and pass it to app, update LAS, send ACK
  - Otherwise discard (as out of order)

Receiver Sliding Window – Selective Repeat

- Receiver passes data to app in order, and buffers out-of-order segments to reduce retransmissions
- ACK conveys highest in-order segment, plus hints about out-of-order segments
  - Ex: I got everything up to 42 (LAS), and got 44, 45
- TCP uses a selective repeat design; we'll see the details later

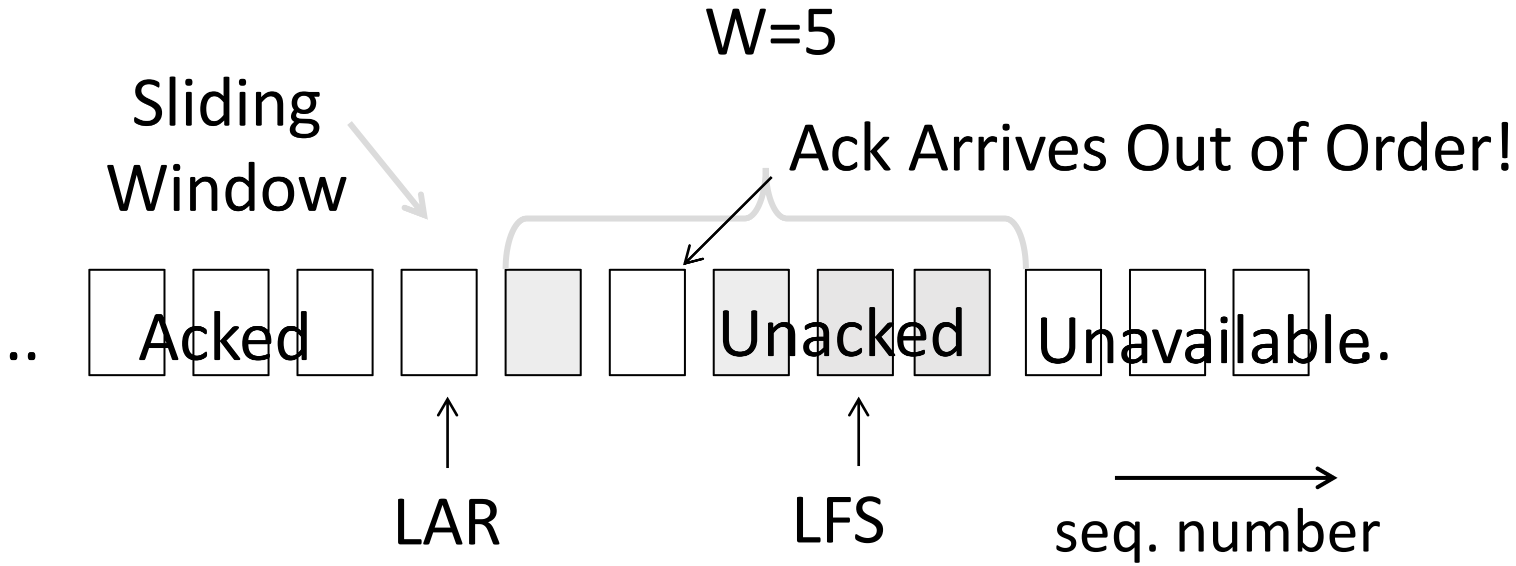
Receiver Sliding Window – Selective Repeat (2)

- Buffers W segments, keeps state variable LAS = LAST ACK SENT

- On receive:
  - Buffer segments [LAS+1, LAS+W]
  - Send app in-order segments from LAS+1, and update LAS
  - Send ACK for LAS regardless

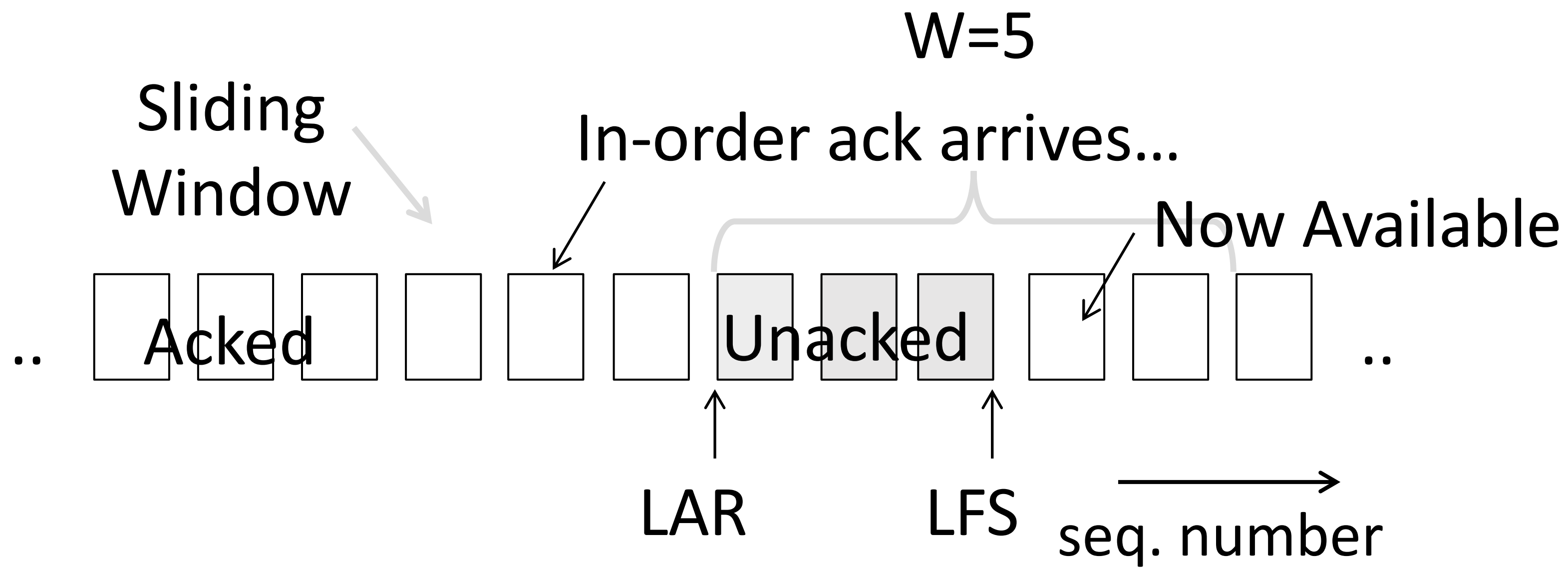
Sender Sliding Window – Selective Repeat

- Keep normal sliding window
- If out-of-order ACK arrives
  - Send LAR+1 again!



Sender Sliding Window – Selective Repeat (2)

- Keep normal sliding window
- If in-order ACK arrives
  - Move window and LAR, send more messages



Sliding Window – Retransmissions

- Go-Back-N uses a single timer to detect losses
  - On timeout, resends buffered packets starting at LAR+1
- Selective Repeat uses a timer per unacked segment to detect losses
  - On timeout for segment, resend it
  - Hope to resend fewer segments

Sequence Numbers

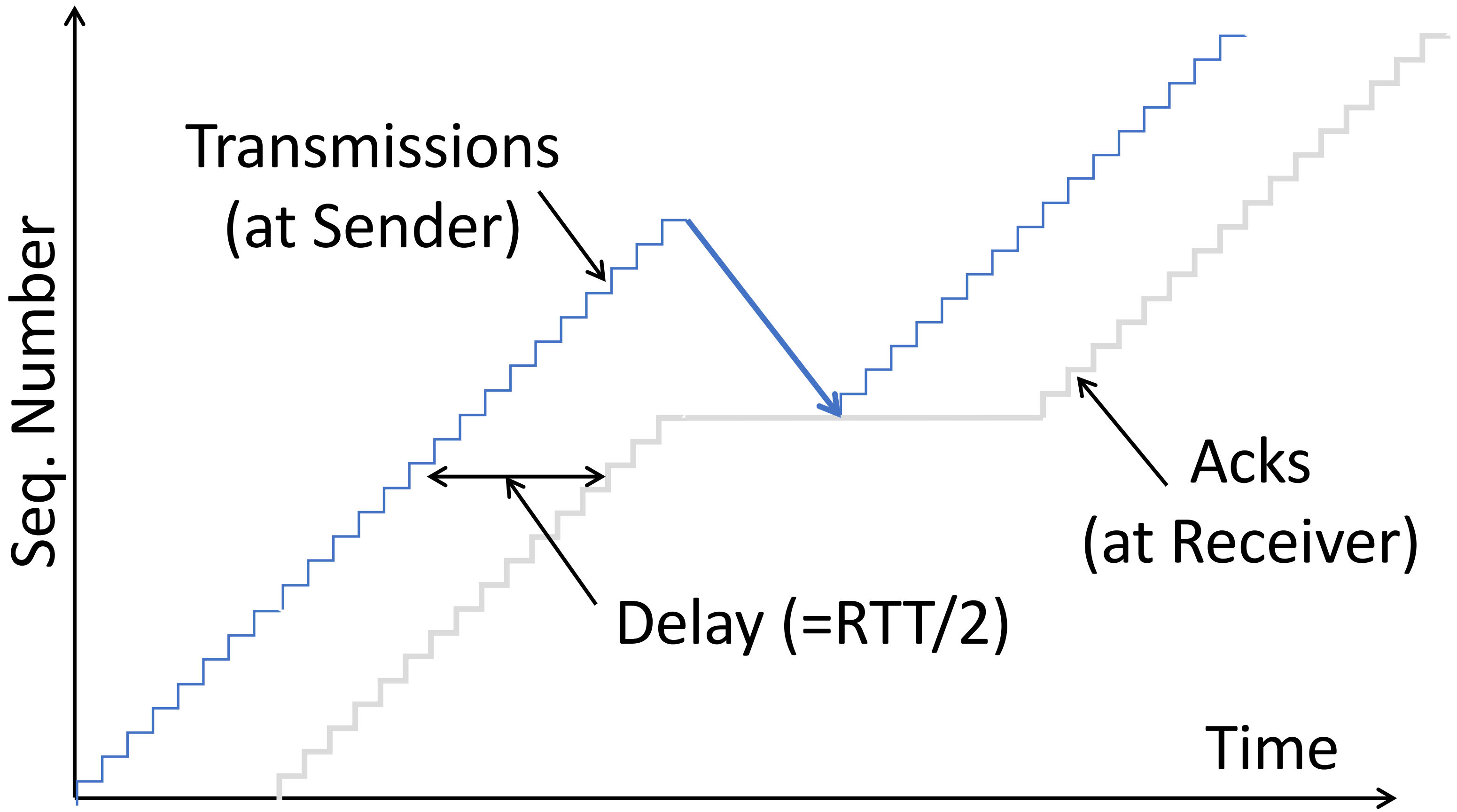
Need more than 0/1 for Stop-and-Wait … but how many?

- For Selective Repeat: 2W seq numbers
  - W for packets, plus W for earlier acks
- For Go-Back-N: W+1 sequence numbers

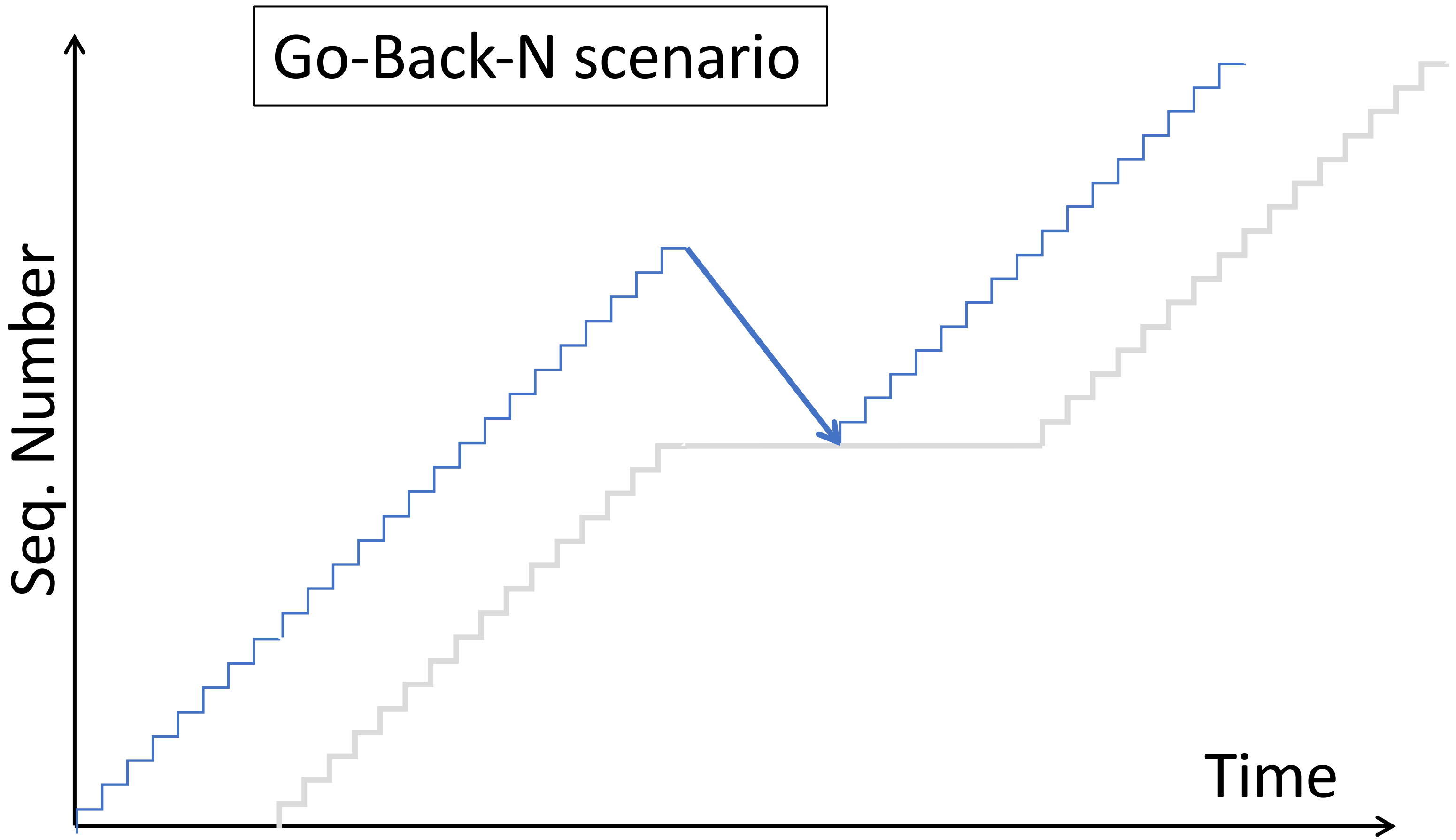
Typically implement seq. number with an N-bit counter that wraps around at 2^N—1

- E.g., N=8: …, 253, 254, 255, 0, 1, 2, 3, …

Sequence Time Plot



Sequence Time Plot (2)



Sequence Time Plot (3)