# Network: Section 4

HTTP/2 & QUIC

## Background & Motivation

Reducing web latency - PLT

- User experience
- Scaling of web platform

Insecure -> Secure

• TLS/TCP

"In practice, once the user has more than 5 Mbps of bandwidth, further improvements deliver minimal increase in the loading speed of the average Web application....."

- streaming HD video from the Web -> bandwidthbound
- loading the page hosting the HD video, with all of its assets -> latency-bound

Source:

https://queue.acm.org/detail.cfm?id=2555617

# Recap: HTTP/1.X -> HTTP/2.0

Originally developed by Google

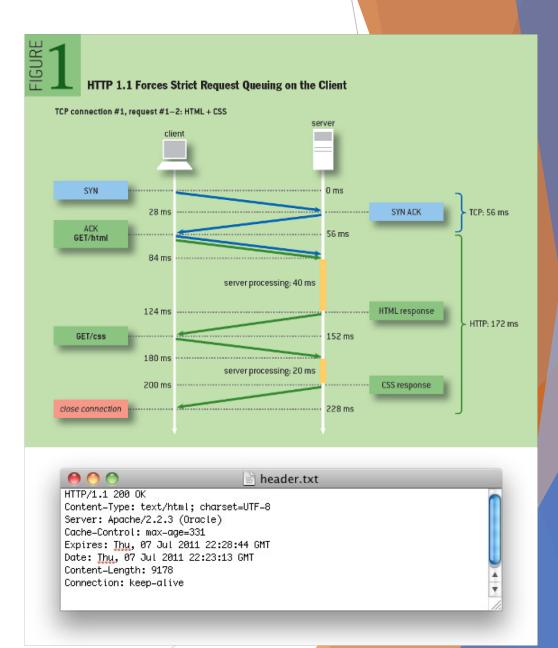
#### Reference:

https://queue.acm.org/detail.cfm?id=2555617

https://developers.google.com/web/fundamentals/performance/http2/

#### HTTP/1.1

- Reuse TCP connection
  - ► HTTP/1.0 one TCP per resource -> overhead
  - ► HTTP/1.1 up to six TCP per origin
- Request pipeline
  - Theoretically, yes; but failed.
  - Community homebrew "optimizations"
    - Multiple origins -> more parallelism
    - Bundle files -> less requests
      - One giant CSS file
      - One giant JS file
      - Code everything directly into HTML
    - Leads to network congestion + poor modularity (caching & page loading)
- ► HTTP Header
  - plain text each char is 1 byte
  - newline-delimited

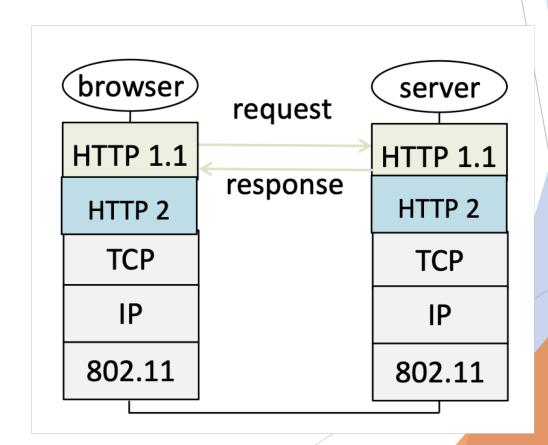


"All problems in computer science can be solved by another level of indirection...

--- David Wheeler

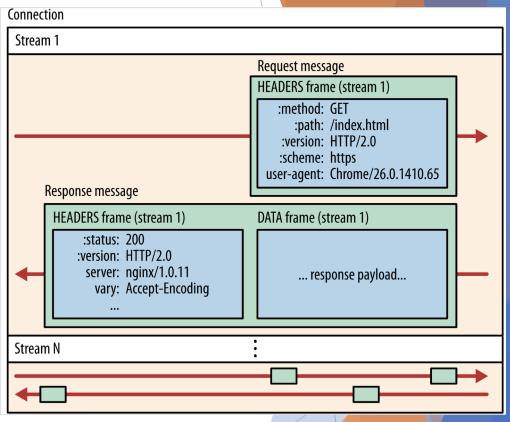
## HTTP/2.0

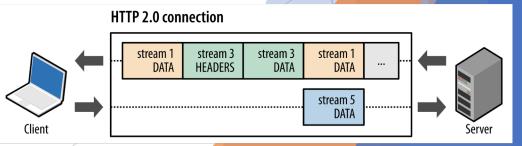
- ► Another layer of abstraction over HTTP/1.1
- Multiplexing
- Prioritization
- ► Header compression
- Server push



## HTTP/2.0: Multiplexing, Prioritization

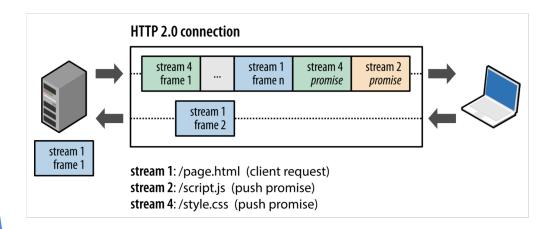
- Streams inside one TCP connection
  - HTTP message -> HEADER frame + DATA frame
  - One stream for one HTTP request + response
  - Multiple streams inside one TCP connection
    - streams can have different priorities
  - Frames from different streams may be interleaved and then reassembled via the embedded stream identifier in the header of each frame.





## HTTP/2.0: Server push

- ▶ Push the resource to the client instead of waiting for the client to request it
  - Request html
  - ► Push CSS, JS, IMG...
- "PUSH\_PROMISE"
  - Contains HTTP request header of the pushed elements
  - Server initiates a new stream and start pushing
  - Client can reject by sending RST\_STREAM



## HTTP/2.0: HTTP Header Compression

- HTTP headers are plain texts with a lot of repetitions ("HTTP/1.1", "GET", ...)
- ► HPACK compression algorithm: static table + dynamic table + static Huffman code

#### Request headers Static table Encoded headers :method GET :authority :scheme :method https **GET** :host example.com 51 Huffmann("/resource") :path /resource referer Mozilla/5.0 ... user-agent . . . . . . Mozilla/5.0 Huffmann("custom-hdr") user-agent custom-hdr some-value Huffmann("some-value") 63 :host example.com . . .

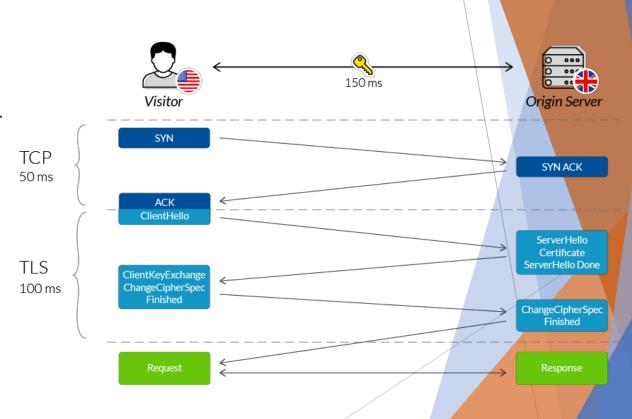
Dynamic table

# QUIC: Quick UDP Internet Connections

(Again) By Google

### Problem with TLS/TCP

- TCP headers unencrypted
  - Middleboxes Firewall, NAT
- ► TCP commonly implemented in OS kernel
  - Update really slow
  - Sizeable user populations lag behind
- Handshake Delay
  - ▶ 1 RTT for TCP + 2 RTT for TLS
  - c is constant
- ► Head-of-line blocking delay
  - ▶ One TCP stream in HTTP/2
  - ► Packet lost? Wait...

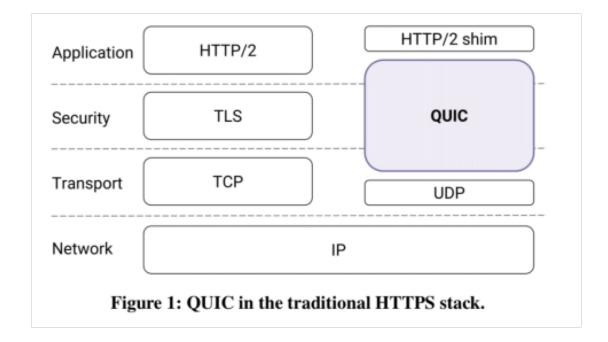


"All problems in computer science can be solved by another level of indirection... Except for the problem of too many layers of indirection."

--- David Wheeler

## QUIC Introduction

- ► TCP+TLS+HTTP2
- Application Layer
- HTTPS Performance+
- ► E2E Encrypted
- Secure
- Rapid Deployment



## QUIC Key Advantages

- Connection establishment latency
- Improved congestion control
- Multiplexing without head-of-line blocking
- Forward error correction
- Connection migration

### QUIC Advantages #1 Connection Establishment Latency

- Combined handshake
  - ► Inchoate and complete CHLO
  - ► SHLO
  - ▶ 1-RTT
- Client cache long-term Diffie-Hellman public key
  - ► 0-RTT

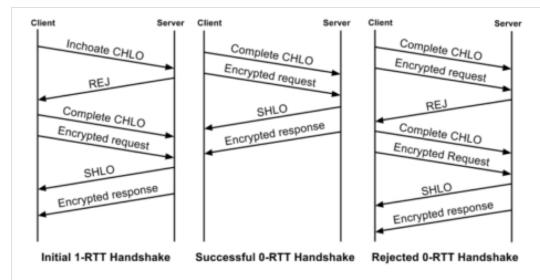


Figure 4: Timeline of QUIC's initial 1-RTT handshake, a subsequent successful 0-RTT handshake, and a failed 0-RTT handshake.

## QUIC Advantages #2 Congestion Control

- Pluggable Interface Design
  - ► Easy to switch
  - ► Easy to experiment
  - ► Easy to update
- ► ACK carries more messages
  - ▶ More information
- ► Better use of packet number

## QUIC Advantages #3 Multiplexing

- Solve Head-of-line blocking delay
  - ► Stream lightweight TCP connection without handshakes
  - ► Multiple streams in one connection
  - ▶ One QUIC packets can carries multiple stream frames

#### One QUIC Connection

Stream Frame #1

Stream Frame #1

QUIC Packet #1 (UDP Packet)

Stream Frame #2

Stream Frame #2

QUIC Packet #2 (UDP Packet)

# QUIC Advantages #4 Forward error correction

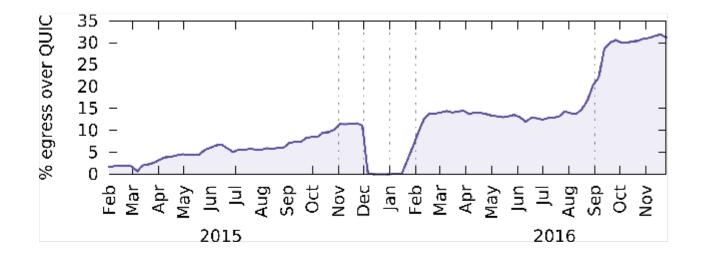
- Skip
- Describe in Sec 7.3 in paper
- Benefits not compelling
- Removed from QUIC in early 2016

## QUIC Advantages #5 Connection Migration

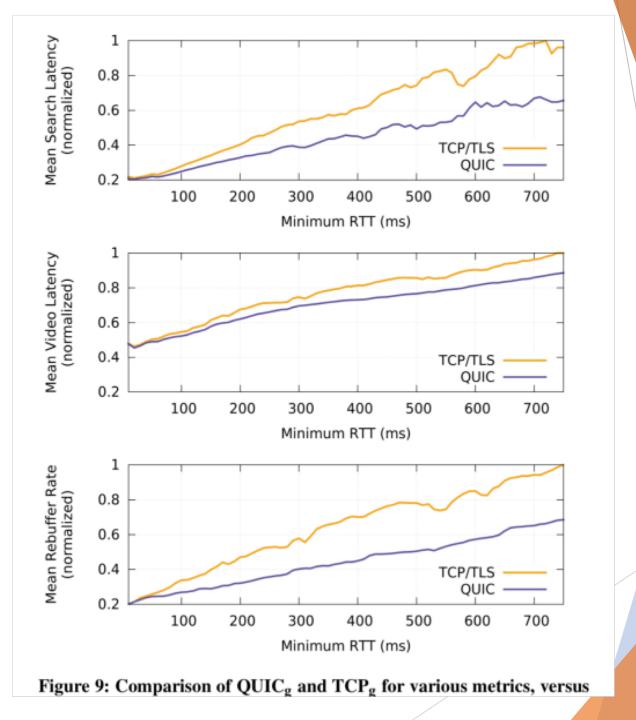
- ► TCP
  - ► Src IP:Port + Dest IP:Port + protocol 5-tuple Identification
  - ► Client change IP, NAT change port... -> connection break
- QUIC
  - Connection ID
  - ► Connection remains even network environment changes

## QUIC Results #1

- Well tested and deployed widely
  - ► Chrome, YouTube, Google Search App, ...
  - ▶ 7% of the internet traffic



## QUIC Results #2



## **QUIC Future Work**

- Alternative congestion control algorithm
- Reduce CPU cost
  - ▶ Twice of the TCP
- Improve performance on mobile devices
  - ► Mobile app -> invisible handshake, compressed content...
  - ► CPU bottleneck
- MTU discover for QUIC
  - ► MTU: maximum packet size
  - ▶ MTU now sets to a fixed tested value: 1450 bytes

## **QUIC** References

- https://tools.ietf.org/html/draft-ietf-quic-transport-09
- https://dl.acm.org/citation.cfm?id=3098842
- https://www.chromium.org/quic
- https://docs.google.com/document/d/1gY9-YNDNAB1eip-RTPbqphgySwSNSDHLq9D5Bty4FSU/edit
- https://static.googleusercontent.com/media/research.google.com/en//pubs/archive/8b935debf13bd176a08326738f5f88ad115a071e.pdf