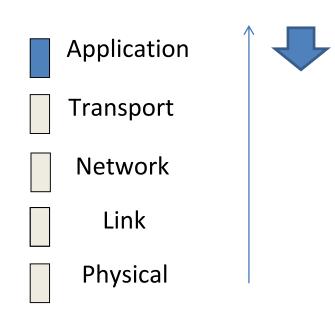
2- Application Level Protocols

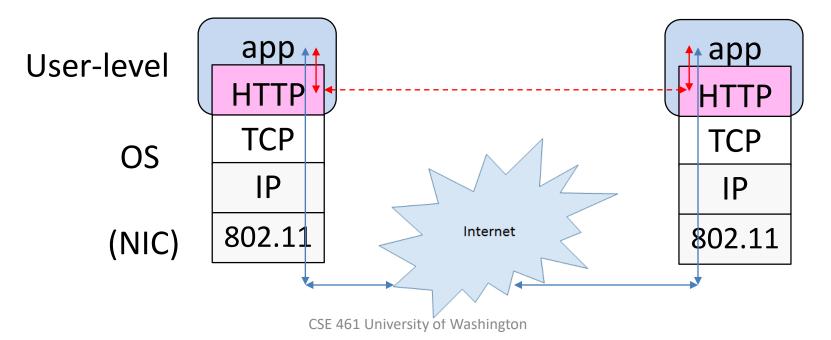
Where we are in the Course



Implementation

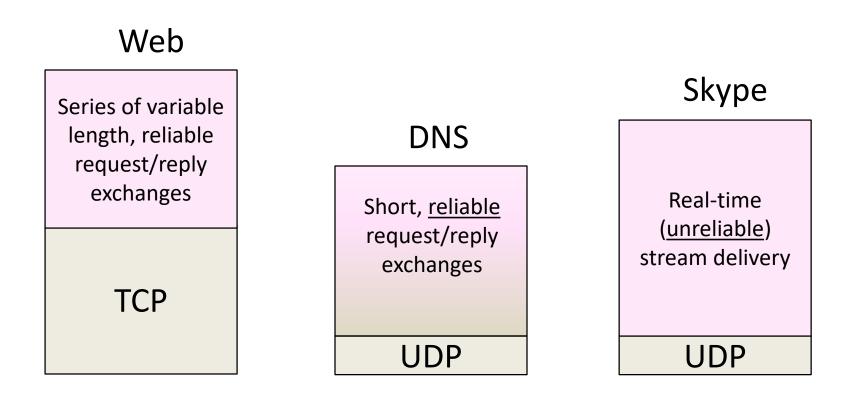
 Application layer protocols are often part of "the app"

- Libraries running in untrusted space



Application Communication Needs

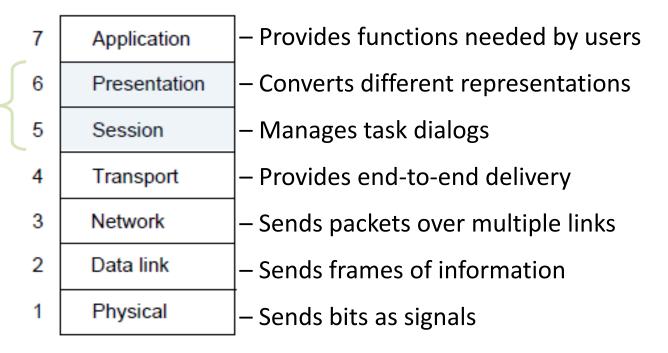
• Vary widely; must build on Transport services



OSI Session/Presentation Layers

• Two relevant concepts ...

Considered part of the application, not strictly layered!



Session Concept

- A session is a series of related network interactions in support of an application task
 - Often informal, not explicit
 Often related to an individual user
- Examples:
 - -Web page fetches multiple resources
 - Skype call involves audio, video, chat

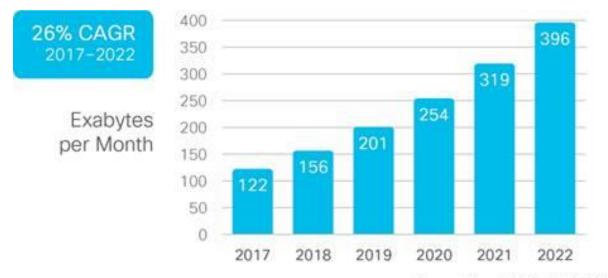
Presentation Concept

 Apps need to identify the type of content, and encode it for transfer

- These are Presentation functions

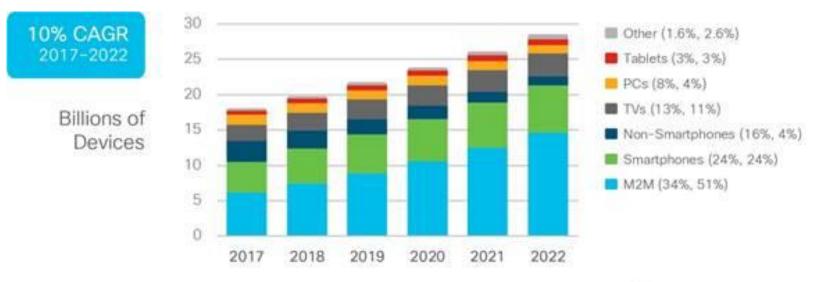
- Examples:
 - Media (MIME) types, e.g., image/jpeg, identify content type
 - Transfer encodings, e.g., gzip, identify the encoding of content
 - Application headers are simple/readable versus packed for efficiency

Which Apps Matter?



Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

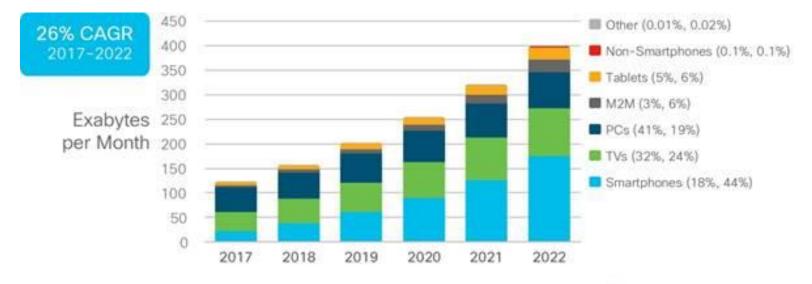
Which Devices Matter?



* Figures (n) refer to 2017, 2022 device share Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

Graph of connected devices

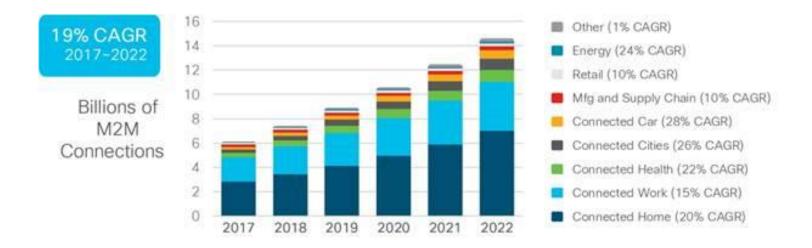
Which Devices Matter?



* Figures (n) refer to 2017, 2022 traffic share Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

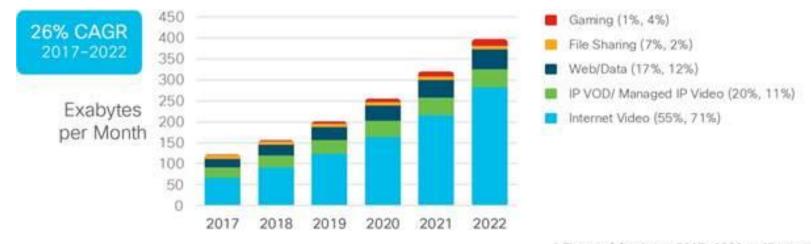
Graph of traffic generated by devices

M2M Connections By Industry



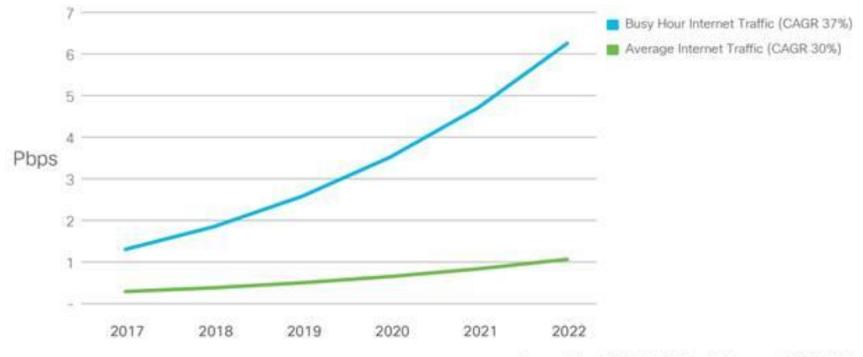
Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

Application Data



* Figures (n) refer to 2017, 2022 traffic share Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

Peak to Average Load

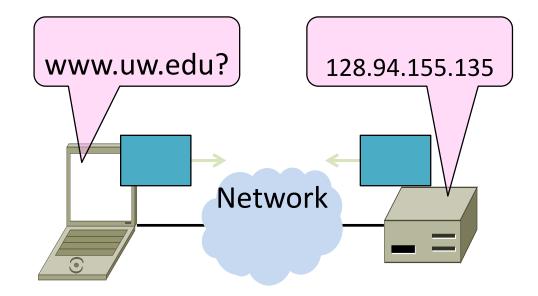


Source: Cisco VNI Global IP Traffic Forecast, 2017-2022

DOMAIN NAME SYSTEM

DNS

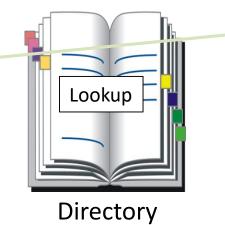
• Human-readable host names, and more



Names and Addresses

- <u>Names</u> are higher-level identifiers for resources
- <u>Addresses</u> are lower-level locators for resources
 - Multiple levels, e.g. full name → email → IP address →
 Ethernet addr
- <u>Resolution</u> (or lookup) is mapping a name to an address

Name, e.g. "Andy Tanenbaum," or "flits.cs.vu.nl"



Address, e.g. "Vrijie Universiteit, Amsterdam" or IPv4 "130.30.27.38"

Before the DNS – HOSTS.TXT

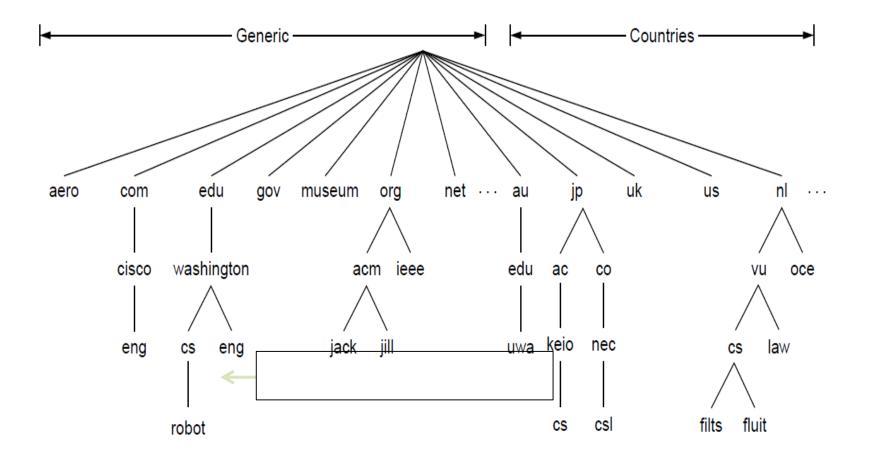
- Directory was a file HOSTS.TXT regularly retrieved for all hosts from a central machine at the NIC (Network Information Center)
- Names were initially flat, became hierarchical (e.g., lcs.mit.edu) ~1985
- Not manageable or efficient as the ARPANET grew ...

DNS

- A naming service to map between host names and their IP addresses (and more)
 - www.uwa.edu.au \rightarrow 130.95.128.140
- Goals:
 - Easy to manage (esp. with multiple parties)
 - Efficient (good performance, few resources)
- Approach:
 - Distributed directory based on a hierarchical namespace
 - Automated protocol to tie pieces together

DNS Namespace

• Hierarchical, starting from "." (dot, typically omitted)

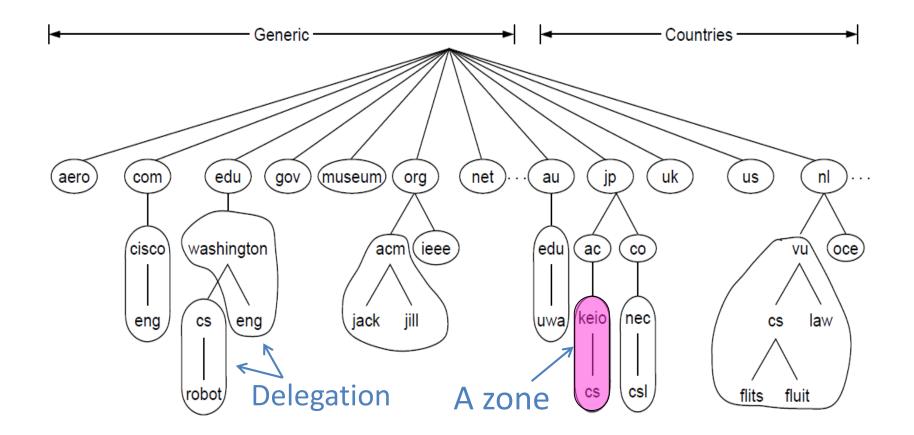


TLDs (Top-Level Domains)

- Run by ICANN (Internet Corp. for Assigned Names and Numbers)
 - Starting in '98; naming is financial, political, and international $\, \odot \,$
- 700+ generic TLDs
 - Initially .com, .edu , .gov., .mil, .org, .net
 - Unrestricted (.com) vs Restricted (.edu)
 - Added regions (.asia, .kiwi), Brands (.apple), Sponsored (.aero) in 2012
- ~250 country code TLDs
 - Two letters, e.g., ".au", plus international characters since 2010
 - Widely commercialized, e.g., .tv (Tuvalu)
 - Many domain hacks, e.g., instagr.am (Armenia), kurti.sh (St. Helena)

DNS Zones

• A <u>zone</u> is a contiguous portion of the namespace



DNS Zones (2)

- Zones are the basis for distribution
 - EDU Registrar administers .edu
 - UW administers washington.edu
 - CSE administers cs.washington.edu
- Each zone has a <u>nameserver</u> to contact for information about it
 - Zone must include contacts for delegations, e.g., .edu knows nameserver for washington.edu

DNS Resource Records

• A zone is comprised of DNS resource records that give information for its domain names

Туре	Meaning
SOA	Start of authority, has key zone parameters
A	IPv4 address of a host
AAAA ("quad A")	IPv6 address of a host
CNAME	Canonical name for an alias
MX	Mail exchanger for the domain
NS	Nameserver of domain or delegated subdomain

DNS Resource Records (2)

; Authoritative data for cs.vu.nl							
cs.vu.nl.	86400	IN	SOA	star boss (9527,7200	,7200,241920,86400)		
cs.vu.nl.	86400	IN	MX	1 zephyr			
cs.vu.nl.	86400	IN	MX	2 top	N I		
cs.vu.nl.	86400	IN	NS	star 🔶	Name		
					server		
star	86400	IN	Α	130.37.56.205			
zephyr	86400	IN	Α	130.37.20.10	IP		
top	86400	IN	Α	130.37.20.11 <			
WWW	86400	IN	CNAME	star.cs.vu.nl	addresses		
ftp	86400	IN	CNAME	zephyr.cs.vu.nl	of		
					01		
flits	86400	IN	Α	130.37.16.112	computers		
flits	86400	IN	Α	192.31.231.165	-		
flits	86400	IN	MX	1 flits			
flits	86400	IN	MX	2 zephyr			
flits	86400	IN	MX	3 top			
rowboat		IN	A	130.37.56.201			
		IN	MX	1 rowboat	-Mail		
		IN	MX	2 zephyr			
little-sister		IN	А	130.37.62.23	gateways		
laserjet		IN	А	192.31.231.216			

dig

\$ dig @june.cs.washington.edu attu.cs.washington.edu ANY

;; Got answer:

;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 34244 ;; flags: qr aa rd; QUERY: 1, ANSWER: 4, AUTHORITY: 5, ADDITIONAL: 5 ;; WARNING: recursion requested but not available

;; OPT PSEUDOSECTION:

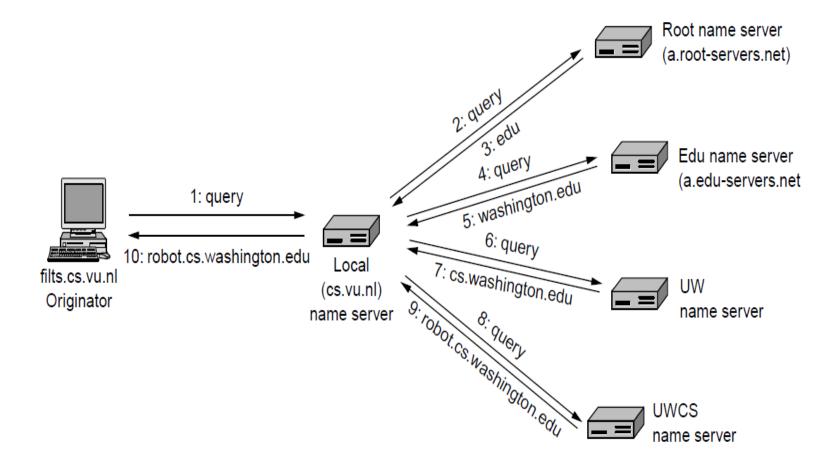
; EDNS: version: 0, flag:	s:; udp: 4096			
;; QUESTION SECTION:				
;attu.cs.washington.edu.		IN	ANY	
;; ANSWER SECTION:				
attu.cs.washington.edu.	60	IN	A	128.208.1.139
attu.cs.washington.edu.	60	IN	A	128.208.1.138
attu.cs.washington.edu.	60	IN	A	128.208.1.137
attu.cs.washington.edu.	60	IN	A	128.208.1.140
;; AUTHORITY SECTION:				
cs.washington.edu.	86400	IN	NS	lumpy.cs.washington.edu.
cs.washington.edu.	86400	IN	NS	<pre>marge.cac.washington.edu.</pre>
cs.washington.edu.	86400	IN	NS	hanna.cac.washington.edu.
cs.washington.edu.	86400	IN	NS	holly.s.uw.edu.
cs.washington.edu.	86400	IN	NS	june.cs.washington.edu.
;; ADDITIONAL SECTION:				
june.cs.washington.edu.	1	IN	AAAA	2607:4000:200:17::104
lumpy.cs.washington.edu.	86400	IN	AAAA	2607:4000:200:17::102
june.cs.washington.edu.	86400	IN	A	128.95.1.4
lumpy.cs.washington.edu.	86400	IN	А	128.95.1.2

DNS Resolution

- DNS protocol lets a host resolve any host name (domain) to IP address
- If unknown, can start with the root nameserver and work down zones
 - But can use cache to do as much resolution as possible
- Let's see an example first ...

DNS Resolution (2)

• flits.cs.vu.nl resolves robot.cs.washington.edu



Iterative vs. Recursive Queries

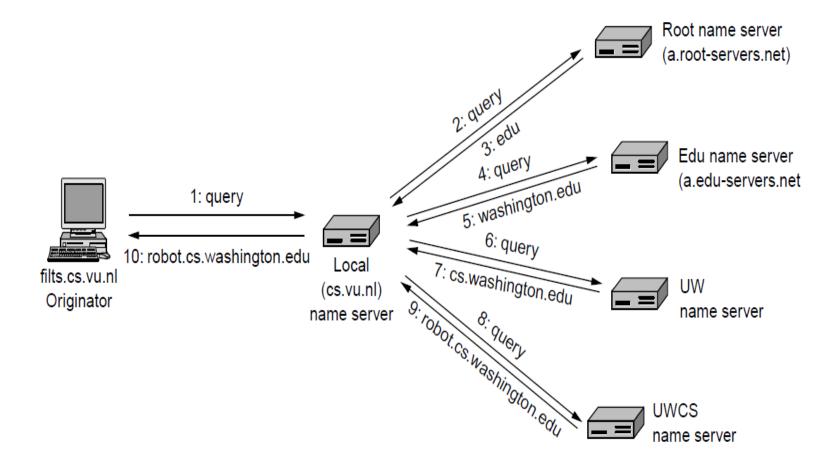
- Recursive query
 - Nameserver resolves and returns final answer
 - E.g., flits \rightarrow local nameserver
- Iterative (Authoritative) query
 - Nameserver returns answer or who to contact for answer
 - E.g., local nameserver \rightarrow all others

Iterative vs. Recursive Queries (2)

- Recursive query
 - Lets server offload client burden (simple resolver) for manageability
 - Lets server cache over a pool of clients for better performance
- Iterative query
 - Lets server "file and forget"
 - Easy to build high load servers

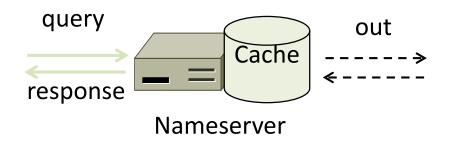
DNS Resolution (2)

• flits.cs.vu.nl resolves robot.cs.washington.edu



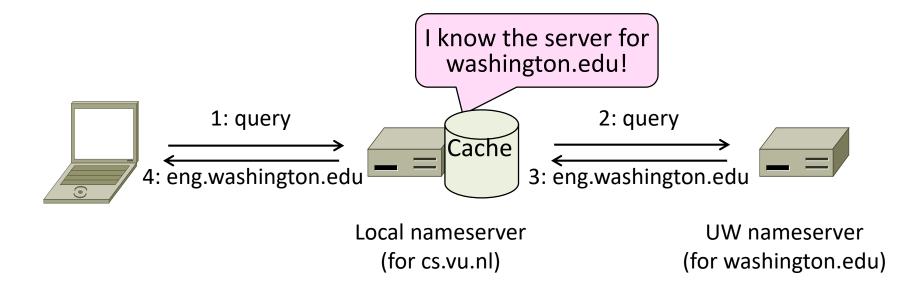
Caching

- Resolution latency should be low
 Adds delay to web browsing
- Cache query/responses to answer future queries immediately
 - Including partial (iterative) answers
 - Responses carry a TTL for caching



Caching (2)

- flits.cs.vu.nl now resolves eng.washington.edu
 - And previous resolutions cut out most of the process



Local Nameservers

 Local nameservers often run by IT (enterprise, ISP)

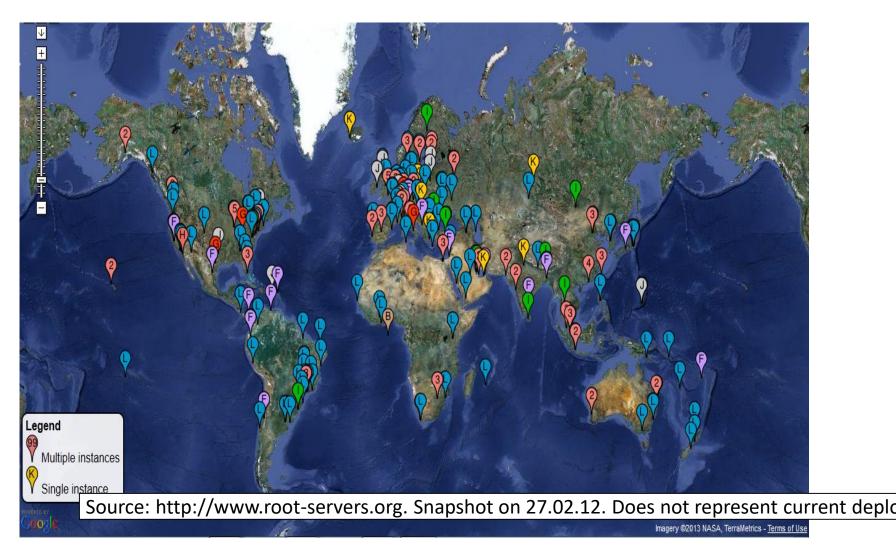
But may be your host or AP

- Or alternatives e.g., Google public DNS
- Clients need to be able to contact local nameservers
 - Typically configured via DHCP

Root Nameservers

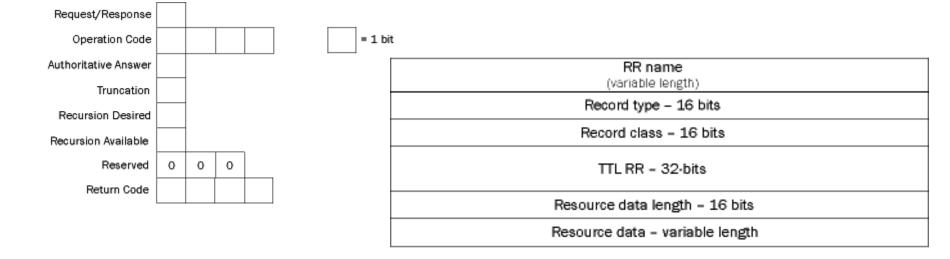
- Root (dot) is served by 13 server names
 - a.root-servers.net to m.root-servers.net
 - All nameservers need root IP addresses
 - Handled via configuration file (named.ca)
- There are >250 distributed server instances
 - Highly reachable, reliable service
 - Most servers are reached by <u>IP anycast</u> (Multiple locations advertise same IP! Routes take client to the closest one.)
 - Servers are IPv4 and IPv6 reachable

Root Server Deployment



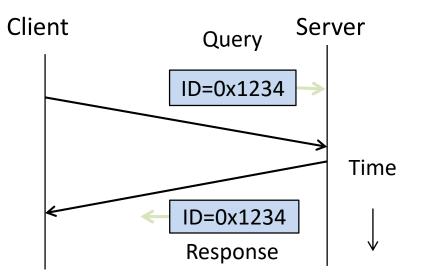
DNS Query Format

Transaction ID	Flags	1		
Question count	Answer RR count	12 bytes		
Authority RR count	Additional RR count	.		
Question entrie	1			
Answer RRs (Variable			
Authority RRs	length			
Additional RRs	↓			



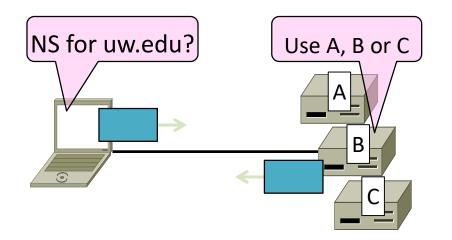
DNS Protocol

- Query and response messages
 - Built on UDP messages, port 53
 - UDP is unreliable
 - time out and repeat request
 - server is stateless (and requests are idem potent)
 - Query/response linked by a 16-bit ID field



DNS Protocol (2)

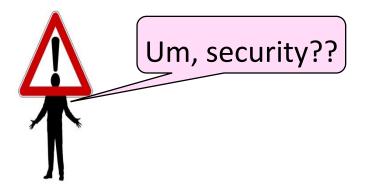
- Service reliability via replicas
 - Run multiple nameservers for domain
 - Return the list; clients use one answer
 - Helps distribute load too



DNS Protocol (3)

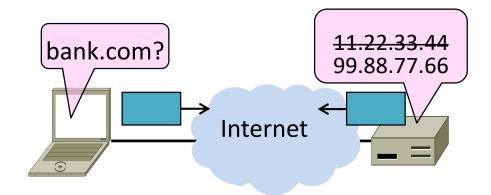
- Security is a major issue
 - Compromise redirects to wrong site!
 - Not part of initial protocols ..
- DNSSEC (DNS Security Extensions)

Mostly deployed



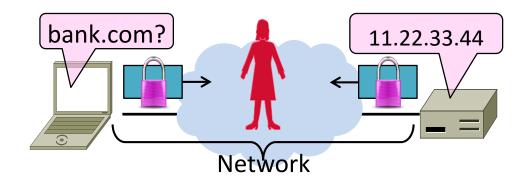
Goal and Threat Model

- Naming is a crucial Internet service
 - Binds host name to IP address
 - Wrong binding can be disastrous ...



Goal and Threat Model (2)

- Goal is to secure the DNS so that the returned binding is correct
 - Integrity/authenticity vs confidentiality
- Attacker can tamper with messages on the network



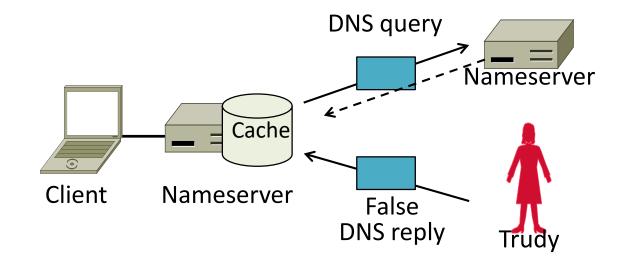
DNS Spoofing

- How can attacker corrupt the DNS?
- Can trick nameserver into caching the wrong binding by using the DNS protocol itself!

This is called <u>DNS spoofing</u>

DNS Spoofing (2)

- To spoof, Trudy returns a fake DNS response that appears to be true
 - Fake response contains bad binding



DNS Spoofing (3)

- Lots of questions!
 - 1. How does Trudy know when the DNS query is sent and what it is for?
 - 2. How can Trudy supply a fake DNS reply that appears to be real?
 - 3. What happens when the real DNS reply shows up?
- There are solutions to each issue ...

DNS Spoofing (4)

- 1. How does Trudy know when the query is sent and what it is for?
- Trudy can make the query herself!
 - Nameserver works for many clients
 - Trudy is just another client

DNS Spoofing (5)

- 2. How can Trudy supply a fake DNS reply that appears to be real?
- A bit more difficult. DNS checks:
 - Reply is from authoritative nameserver (e.g., .com)
 - Reply ID that matches the request
 - Reply is for outstanding query
- (Nothing about content though ...)

DNS Spoofing (6)

- 2. How can Trudy supply a fake DNS reply that appears to be real?
- Techniques:
 - Put IP of authoritative nameserver as the source IP address
 - ID is 16 bits (64K). Send many guesses! (Or if a counter, sample to predict.)
 - Send reply right after query
- Good chance of succeeding!

DNS Spoofing (7)

3. What happens when real DNS reply shows up?

- Not likely a problem
 - There is no outstanding query after fake reply is accepted
 - So real reply will be discarded

DNSSEC (DNS Security Extensions)

- Extends DNS with new record types
 - RRSIG for digital signatures of records
 - DNSKEY for public keys for validation
 - DS for public keys for delegation
 - First version in '97, revised by '05
- Deployment requires software upgrade at both client and server
 - Root servers upgraded in 2010
 - Followed by uptick in deployment

\$ dig edu ANY

;; Got answer:

;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 32320

;; flags: qr rd ra ad; QUERY: 1, ANSWER: 12, AUTHORITY: 6, ADDITIONAL: 8

;; OPT PSEUDOSECT ; EDNS: version: 0, f ;; QUESTION SECTIO	lags:; udp: 4096					
;edu.				IN	ANY	
;; ANSWER SECTION edu.			92995		RRSIG	NS 8 1 172800 20171004051028 20170927040028 10478 edu.
edu. LJ3AKIEw1mx6yHSF	G4el9iBgKYZD3Z+aF	lb7fTg/Cvc9Nk02GpL2PxXOs	6595 7tcZ76dV85KDPRy7Ev2jF2	IN ZIIHODfPAG+FUa4b0f	RRSIG paR2fk0f8tDfYdPi ON8cEIKzl	HxVg0ZsyyZSKEIJXbWLyWGpBlo3M8wQXiJfV055JmMOEQhZKzf /EA= DNSKEY 8 1 86400 20171008150000 20170923145500 28065 edu. LbZSCMr6pYowHSp8JoM5YKkzUTigg23S3FHMIZM6ylUM/iO6 s6u6+w3mEx8T/Fa24GEVuwPjRZDkcg/eBSAmq1LGa23bM7G9Ypet5h wle/OA==
edu. AQOurUqXhGxsfZjYo edu.	o0s/I984A/wpP0WH	lqqmFsD9RHm/7nFhKilLG6VR	6595 D a7qTD/VDzNUL/jiaBEh/ 6595	IN mM2QA4y3rUfsxqnJS IN	DNSKEY VweosIPzw0scTahXGSF K9w DNSKEY	256 3 8 v2AvYqYj96dSQ0TJt9+AUkCF8Z+DTa0DvmXYUHyvMqhSLq2hfM7sAz 9mQ2AQ== 257 3 8
			3T+bSdi2HepB3ZuBWyOe		€VBNI7iFZByRKY9yA J8TshA	oBIRq++8Mxvr+2YCSRaCNyrqaU4TnHaPiHLuKe7sGyHQ8K8 \tiUHHwAF7JJtWLW9oO10BqygYG+zS4weKGtkcNpzOUDWs41FW_M8aFs4L1
						DS 8 1 86400 20171011170000 20170928160000 15768 . 88LDENGE7AzyE/iAwZCrukYAwvck9Y88wRnwn6BVM89NRgfInPQ5G 6D+VOqHd/sbtpt4SOmQdgCmXEKPg2f8vfUg4gLMWYp0G06IsG J2H6Pg==
edu. 6595 IN DS 28065 8 2 4172496CDE85534E51129040355BD04B1FCFEBAE996DFDDE652006F6 F8B2CE76						
edu.			92995	IN	NS	a.edu-servers.net.
edu.			92995	IN	NS	l.edu-servers.net.
edu.			92995	IN	NS	g.edu-servers.net.
edu.			92995	IN	NS	c.edu-servers.net.
edu.			92995	IN	NS	d.edu-servers.net.
edu.			92995	IN	NS	f.edu-servers.net.
;; AUTHORITY SECTION	ON:					
edu.			92995	IN	NS	l.edu-servers.net.
edu.			92995	IN	NS	g.edu-servers.net.
edu.			92995	IN	NS	a.edu-servers.net.
edu.			92995	IN	NS	f.edu-servers.net.
edu.			92995	IN	NS	d.edu-servers.net.
edu.			92995	IN	NS	c.edu-servers.net.
;; ADDITIONAL SECTION:						
a.edu-servers.net.	92995	IN	A	192.5.6.30		
c.edu-servers.net.	92995	IN	A	192.26.92.30		
d.edu-servers.net.	92995	IN	A	192.31.80.30		
f.edu-servers.net.	92995	IN	A	192.35.51.30		
g.edu-servers.net.	92995	IN	A	192.42.93.30		
l.edu-servers.net.	92995	IN	A	192.41.162.30		
g.edu-servers.net.	92995	IN	AAAA	2001:503:cc2c	::2:36	