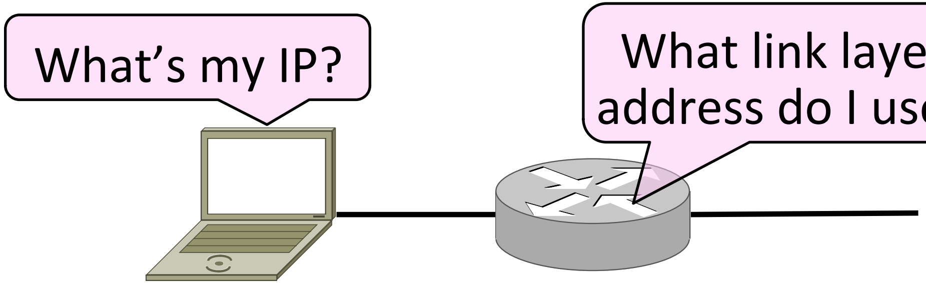
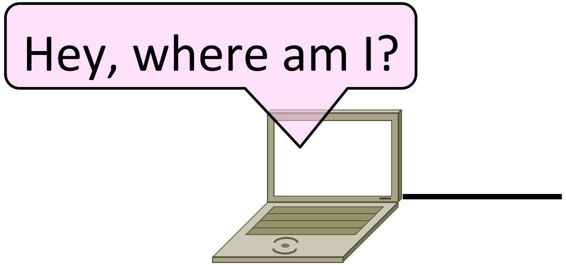
Topic

- Filling in the gaps we need to make for IP forwarding work in practice
  - Getting IP addresses (DHCP) »
  - Mapping IP to link addresses (ARP) »



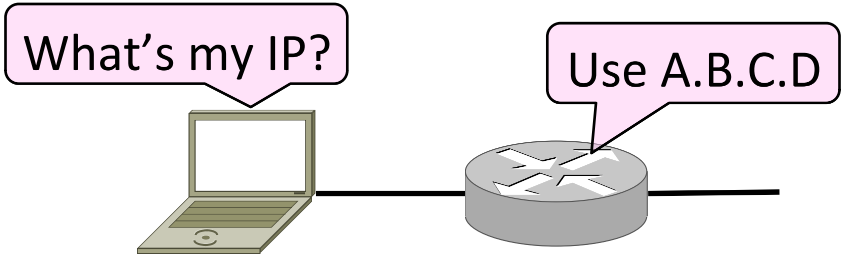
Getting IP Addresses

- Problem:
  - A node wakes up for the first time …
  - What is its IP address? What's the IP address of its router? Etc.
  - At least Ethernet address is on NIC



Getting IP Addresses (2)

1. Manual configuration (old days)
   – Can’t be factory set, depends on use
2. A protocol for automatically configuring addresses (DHCP) »
   – Shifts burden from users to IT folk

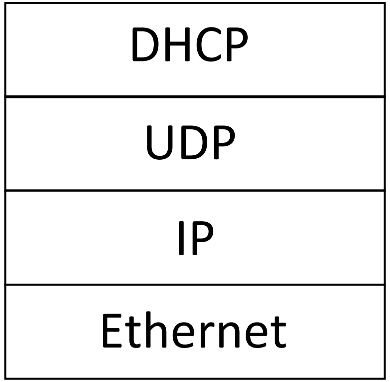


DHCP

- DHCP (Dynamic Host Configuration Protocol), from 1993, widely used
- It leases IP address to nodes
- Provides other parameters too
  - Network prefix
  - Address of local router
  - DNS server, time server, etc.

DHCP Protocol Stack

- DHCP is a client-server application
  - Uses UDP ports 67, 68

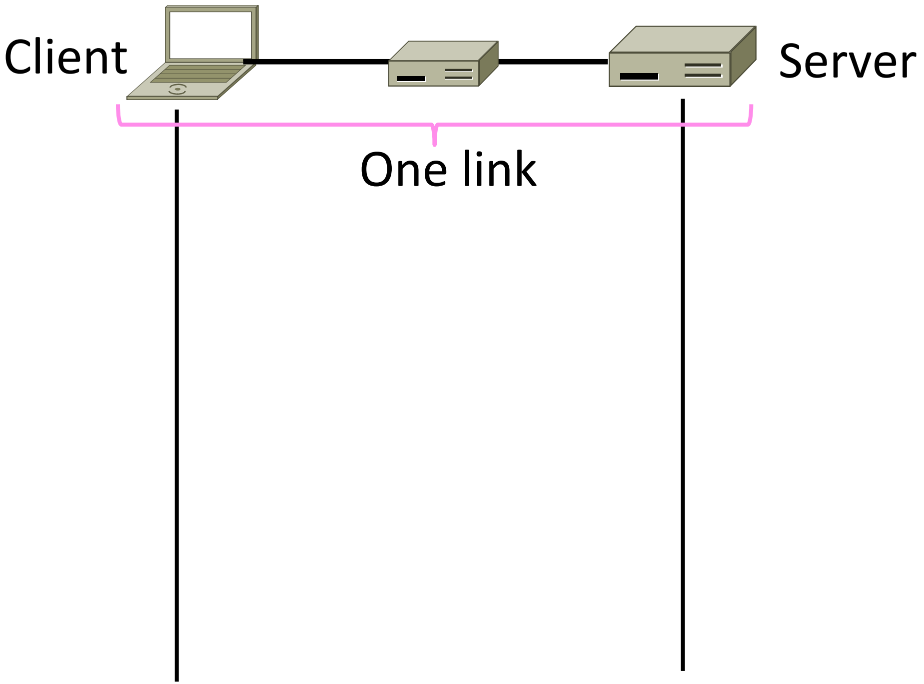


DHCP Addressing

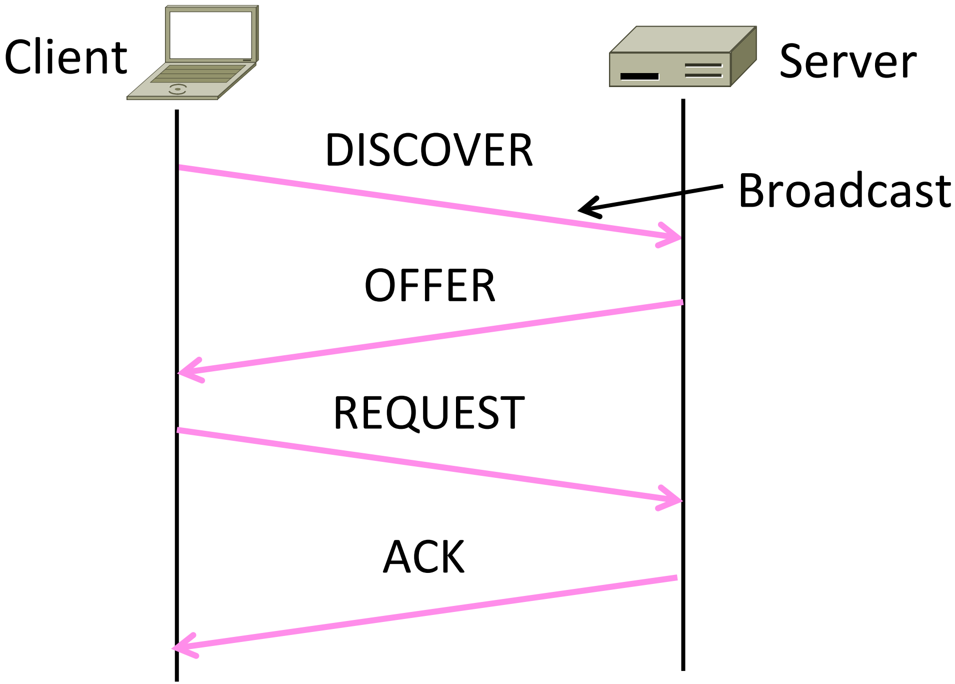
- Bootstrap issue:
  - How does node send a message to DHCP server before it is configured?

- Answer:
  - Node sends broadcast messages that delivered to all nodes on the network
  - Broadcast address is all 1s
  - IP (32 bit): 255.255.255.255
  - Ethernet (48 bit): ff:ff:ff:ff:ff:ff

DHCP Messages



DHCP Messages (2)

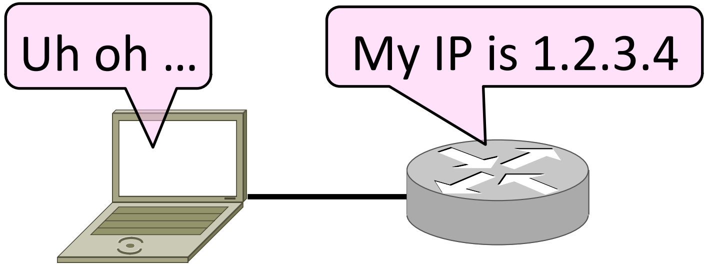


DHCP Messages (3)

- To renew an existing lease, an abbreviated sequence is used:
  - REQUEST, followed by ACK
- Protocol also supports replicated servers for reliability

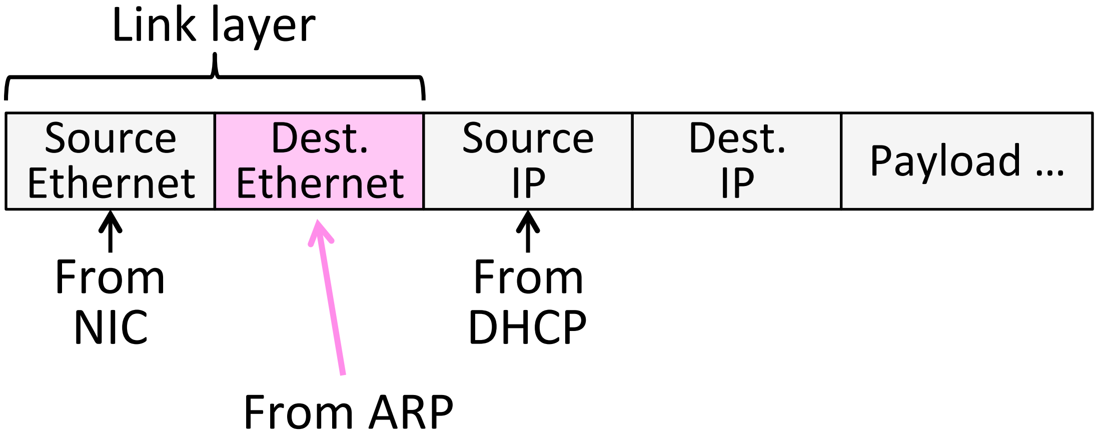
Sending an IP Packet

- Problem:
  - A node needs Link layer addresses to send a frame over the local link
  - How does it get the destination link address from a destination IP address?



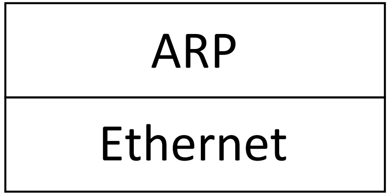
ARP (Address Resolution Protocol)

- Node uses to map a local IP address to its Link layer addresses

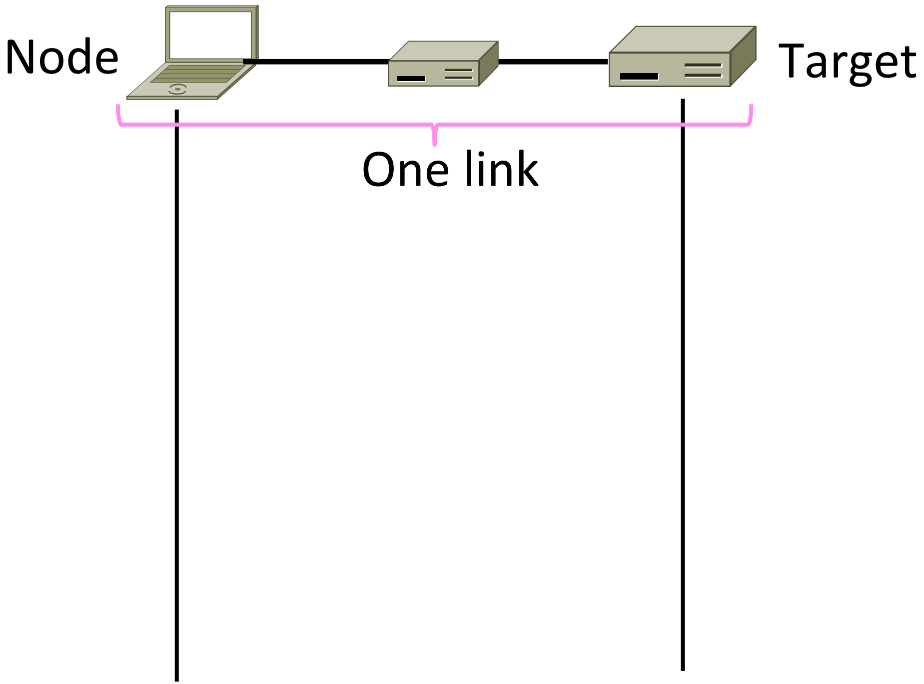


ARP Protocol Stack

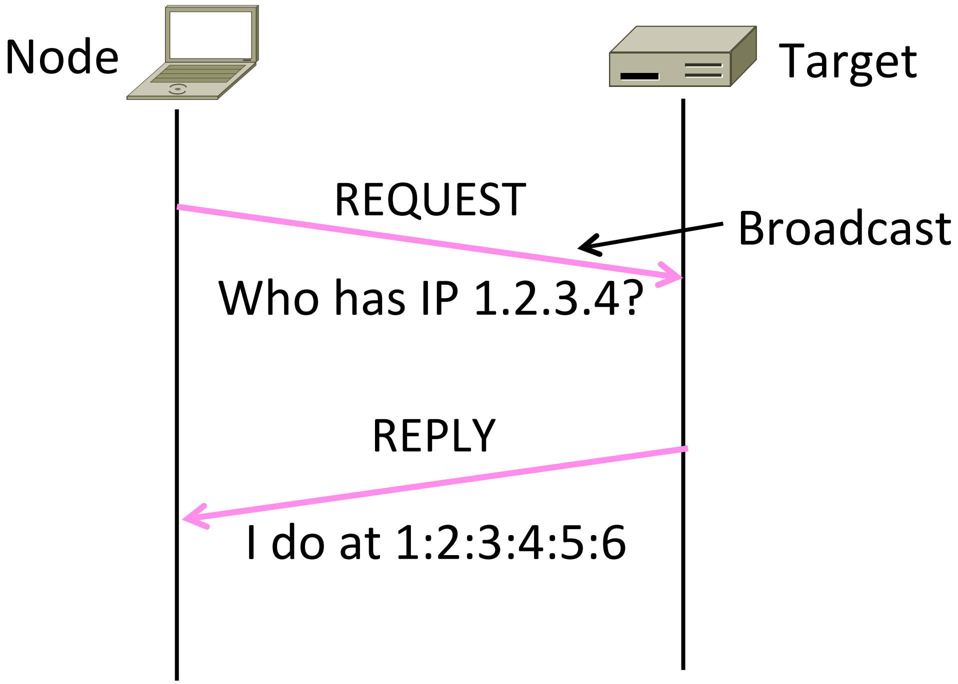
- ARP sits right on top of link layer
  - No servers, just asks node with target IP to identify itself
  - Uses broadcast to reach all nodes



ARP Messages



ARP Messages (2)

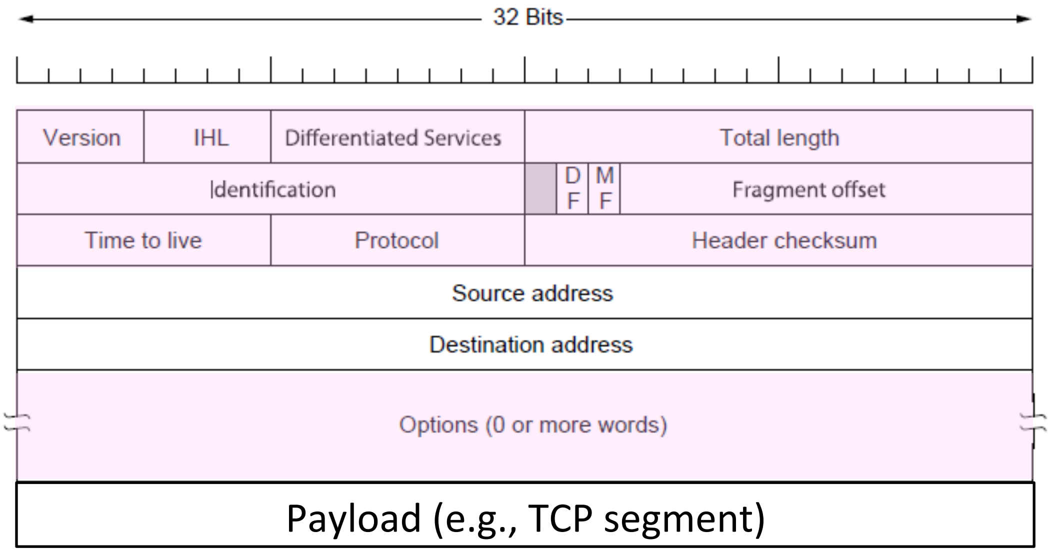


Discovery Protocols

- Help nodes find each other
  - There are more of them!
    - E.g., zeroconf, Bonjour
- Often involve broadcast
  - Since nodes aren't introduced
  - Very handy glue

Other Aspects of Forwarding

- It's not all about addresses …



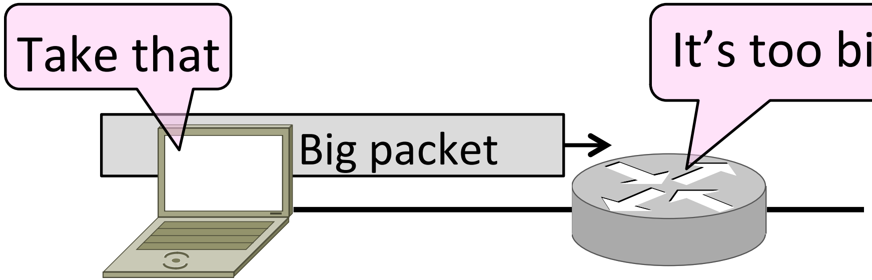
Other Aspects (2)

- Decrement TTL value
  - Protects against loops
- Checks header checksum
  - To add reliability
- Fragment large packets
  - Split to fit it on next link
- Send congestion signals
  - Warns hosts of congestion
- Generates error messages
  - To help mange network
- Handle various options

} Coming later

Topic

- How do we connect networks with different maximum packet sizes?
  - Need to split up packets, or discover the largest size to use



Packet Size Problem

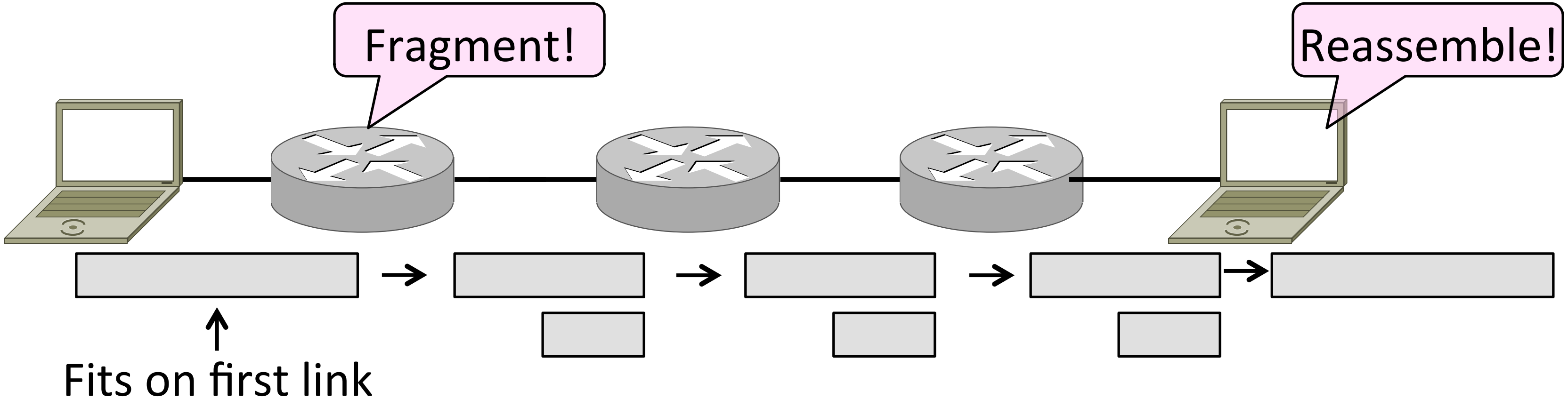
- Different networks have different maximum packet sizes
  - Or MTU (Maximum Transmission Unit)
  - E.g., Ethernet 1.5K, WiFi 2.3K
- Prefer large packets for efficiency
  - But what size is too large?
  - Difficult because node does not know complete network path

Packet Size Solutions

- Fragmentation (now)
  - Split up large packets in the network if they are too big to send
  - Classic method, dated
- Discovery (next)
  - Find the largest packet that fits on the network path and use it
  - IP uses today instead of fragmentation

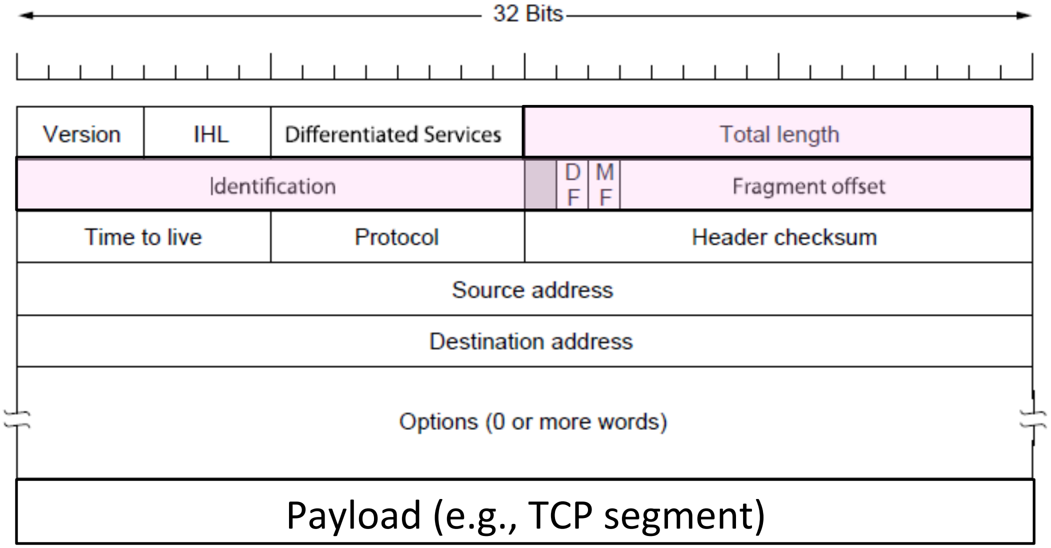
IPv4 Fragmentation

- Routers fragment packets that are too large to forward
- Receiving host reassembles to reduce load on routers



IPv4 Fragmentation Fields

- Header fields used to handle packet size differences
  - Identification, Fragment offset, MF/DF control bits



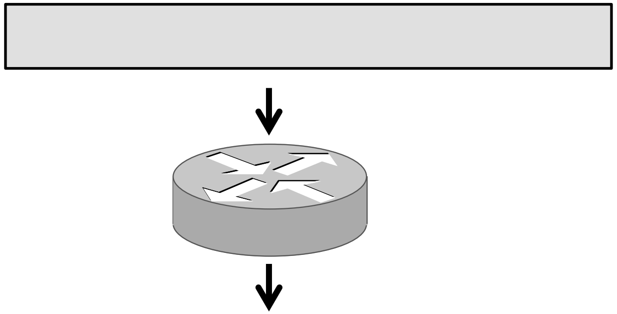
IPv4 Fragmentation Procedure

- Routers split a packet that is too large:
  - Typically break into large pieces
  - Copy IP header to pieces
  - Adjust length on pieces
  - Set offset to indicate position
  - Set MF (More Fragments) on all pieces except last
- Receiving hosts reassembles the pieces:
  - Identification field links pieces together, MF tells receiver when it has all pieces

IPv4 Fragmentation (2)

Before
MTU = 2300

ID = 0x12ef
Data Len = 2300
Offset = 0
MF = 0



(Ignore length of headers)

After
MTU = 1500

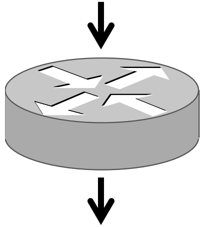
ID =
Data Len =
Offset =
MF =

ID =
Data Len =
Offset =
MF =

IPv4 Fragmentation (3)

Before
MTU = 2300

ID = 0x12ef
Data Len = 2300
Offset = 0
MF = 0



After
MTU = 1500

ID = 0x12ef
Data Len = 1500
Offset = 0
MF = 1

ID = 0x12ef
Data Len = 800
Offset = 1500
MF = 0

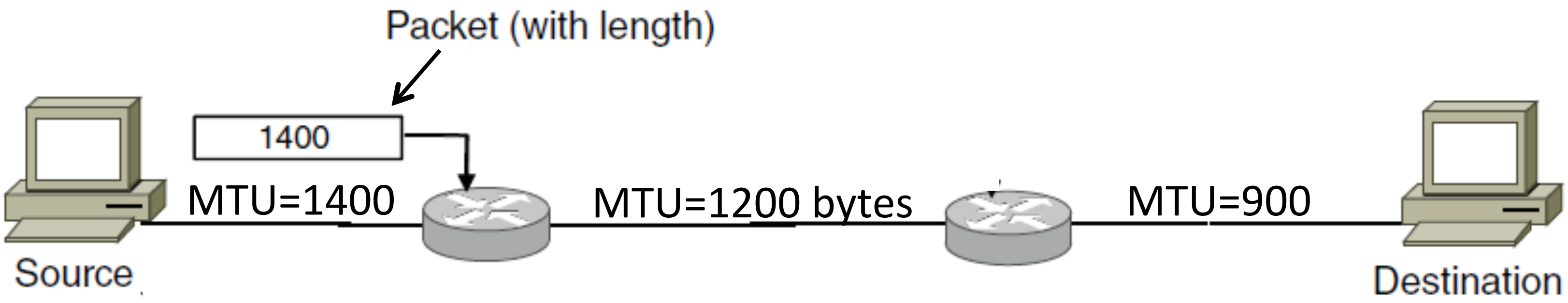
IPv4 Fragmentation (4)

- It works!
  - Allows repeated fragmentation
- But fragmentation is undesirable
  - More work for routers, hosts
  - Tends to magnify loss rate
  - Security vulnerabilities too

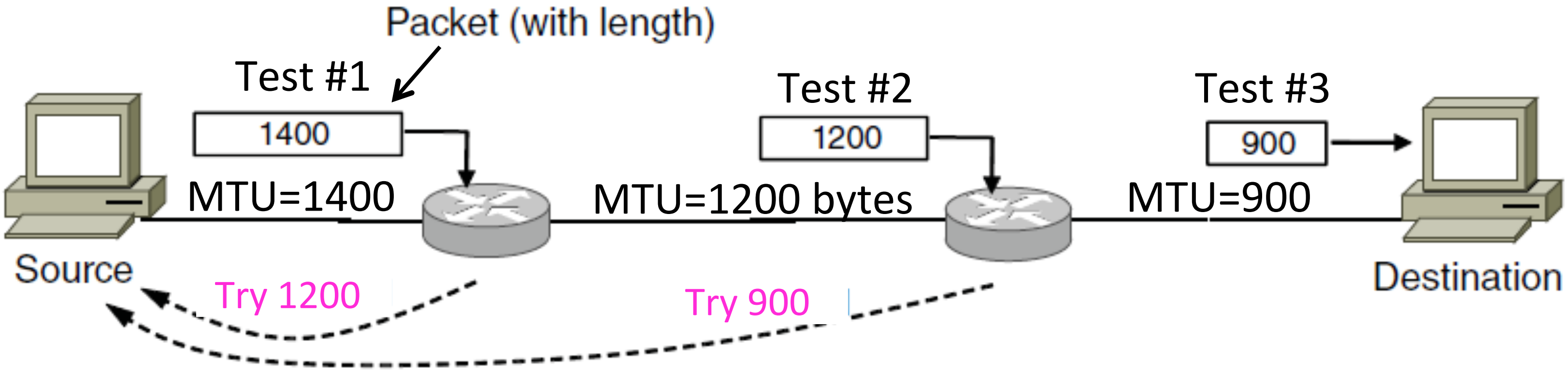
Path MTU Discovery

- Discover the MTU that will fit
  - So we can avoid fragmentation
  - The method in use today
- Host tests path with large packet
  - Routers provide feedback if too large; they tell host what size would have fit

Path MTU Discovery (2)



Path MTU Discovery (3)

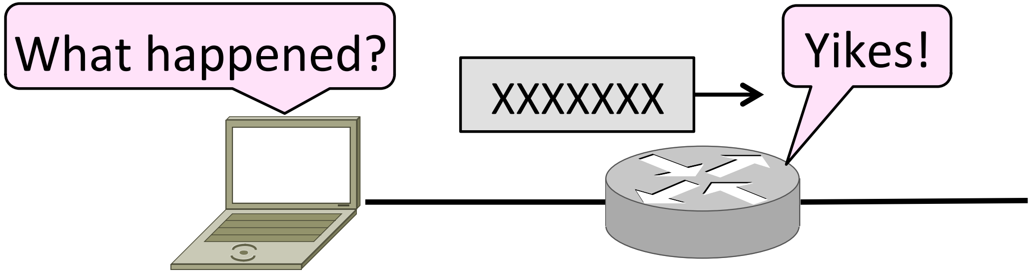


Path MTU Discovery (4)

- Process may seem involved
  - But usually quick to find right size
- Path MTU depends on the path and so can change over time
  - Search is ongoing
- Implemented with ICMP (next)
  - Set DF (Don't Fragment) bit in IP header to get feedback messages

Topic

- What happens when something goes wrong during forwarding?
  - Need to be able to find the problem

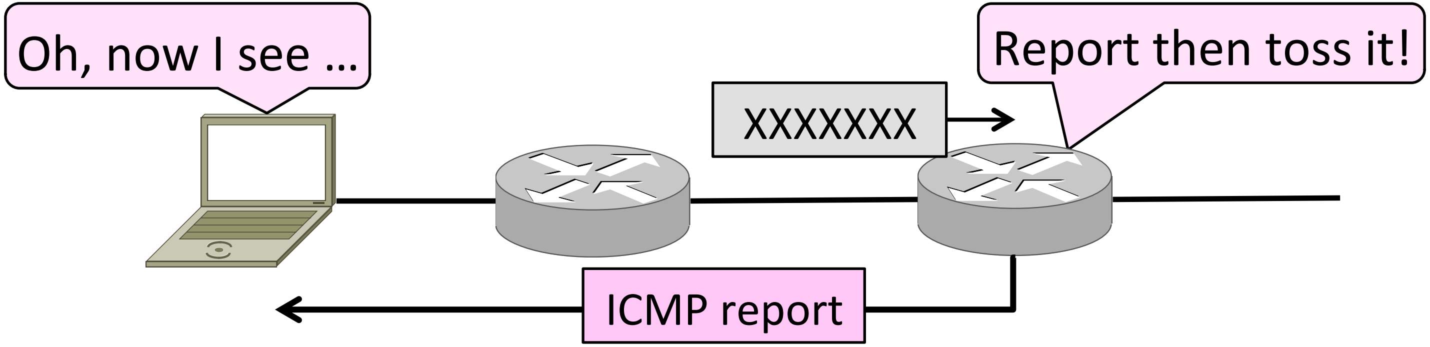


Internet Control Message Protocol

- ICMP is a companion protocol to IP
  - They are implemented together
  - Sits on top of IP (IP Protocol=1)
- Provides error report and testing
  - Error is at router while forwarding
  - Also testing that hosts can use

ICMP Errors

- When router encounters an error while forwarding:
  - It sends an ICMP error report back to the IP source address
  - It discards the problematic packet; host needs to rectify



ICMP Message Format

- Each ICMP message has a Type, Code, and Checksum
- Often carry the start of the offending packet as payload
- Each message is carried in an IP packet

ICMP Message Format (2)

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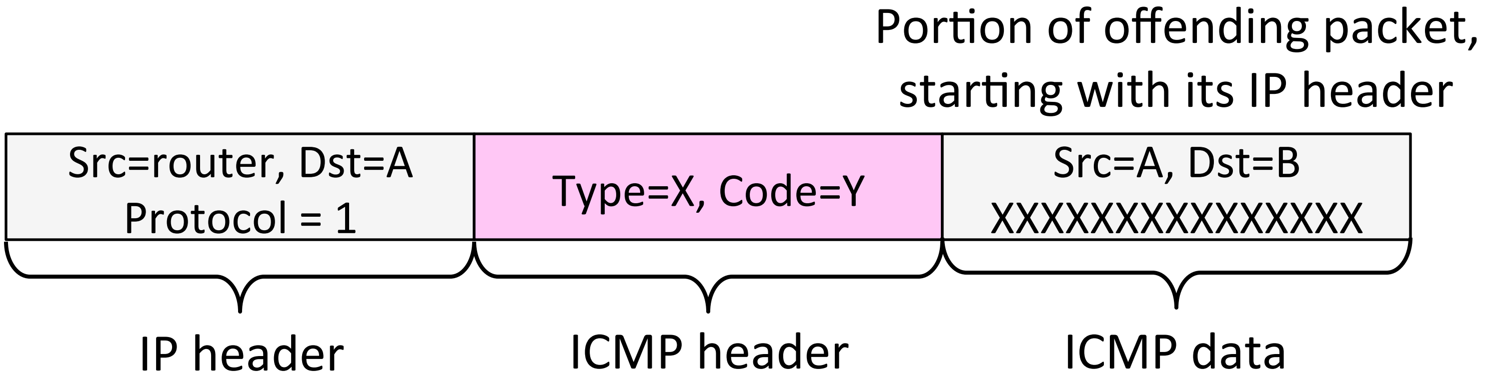


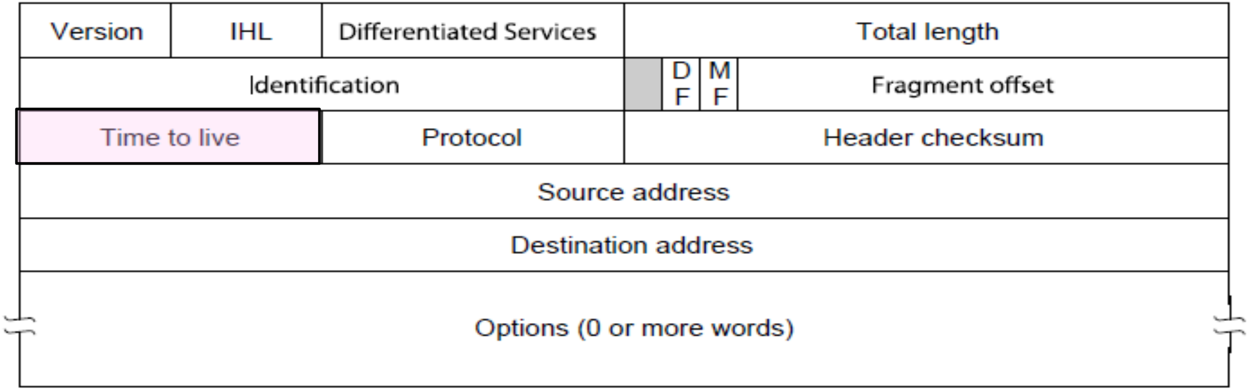
Example ICMP Messages

Name	Type / Code	Usage
Dest. Unreachable (Net or Host)	3 / 0 or 1	Lack of connectivity
Dest. Unreachable (Fragment)	3 / 4	Path MTU Discovery
Time Exceeded (Transit)	11 / 0	Traceroute
Echo Request or Reply	8 or 0 / 0	Ping

Testing, not a forwarding error: Host sends Echo Request, and destination responds with an Echo Reply

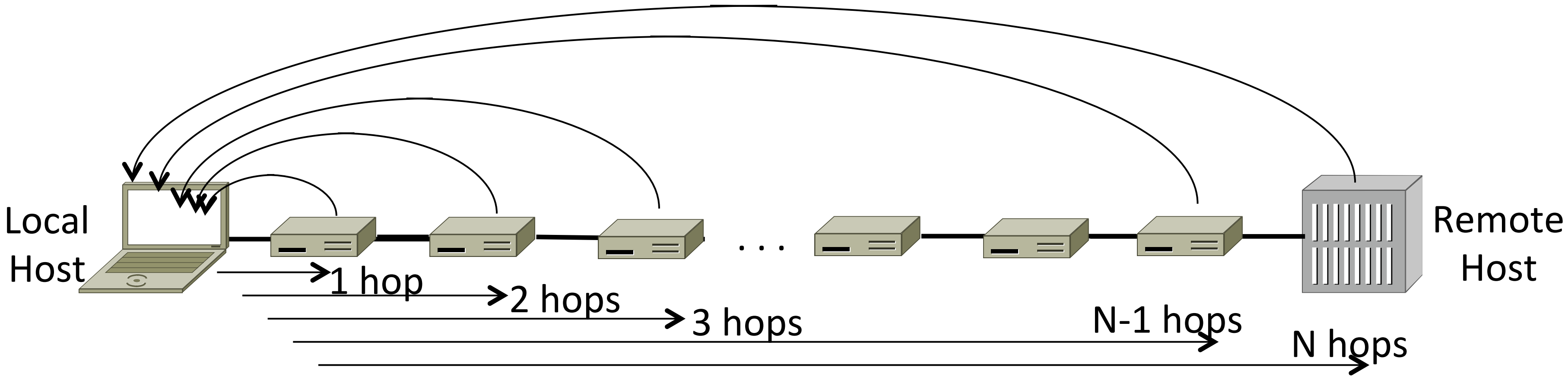
Traceroute

- IP header contains TTL (Time to live) field
  - Decremented every router hop, with ICMP error if it hits zero
  - Protects against forwarding loops



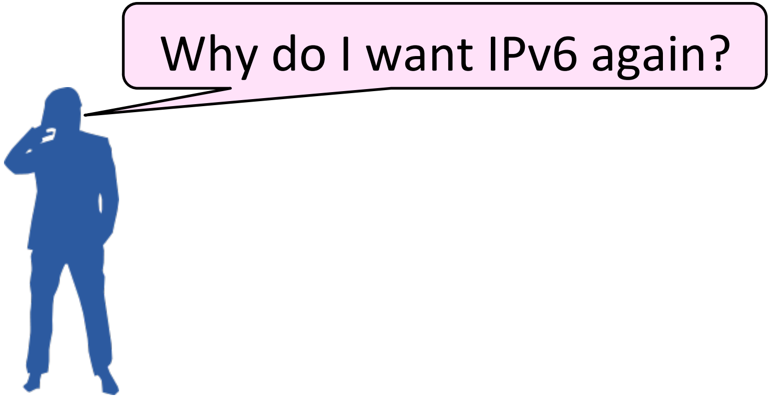
Traceroute (2)

- Traceroute repurposes TTL and ICMP functionality
  - Sends probe packets increasing TTL starting from 1
  - ICMP errors identify routers on the path



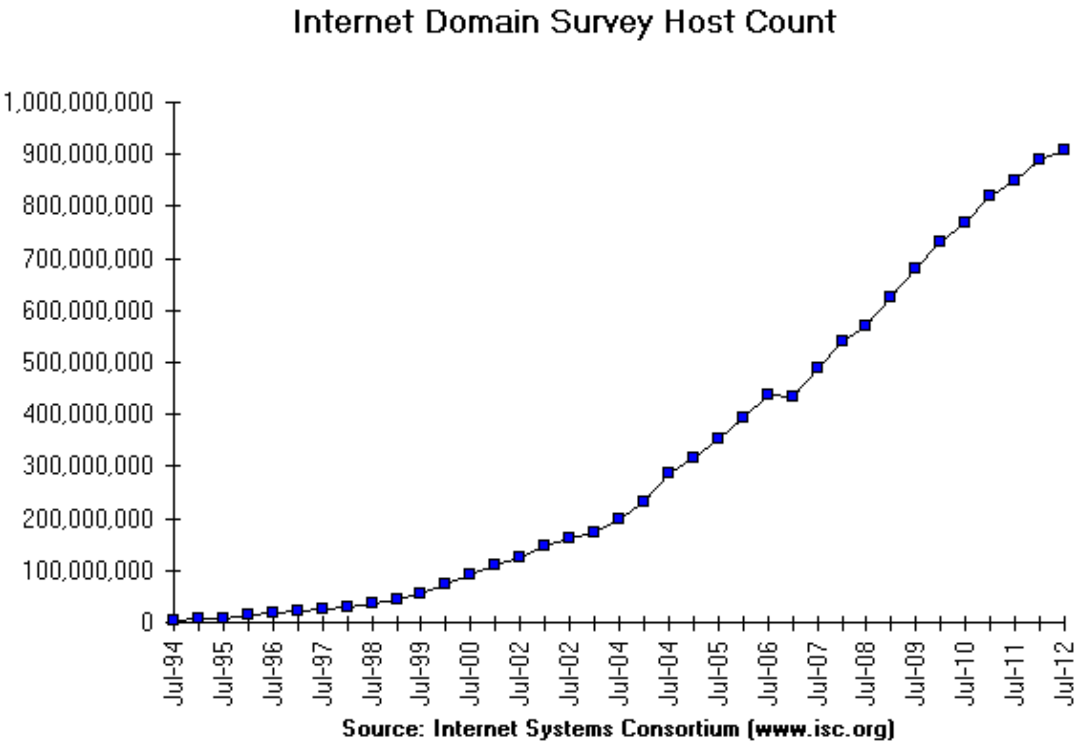
Topic

- IP version 6, the future of IPv4 that is now (still) being deployed



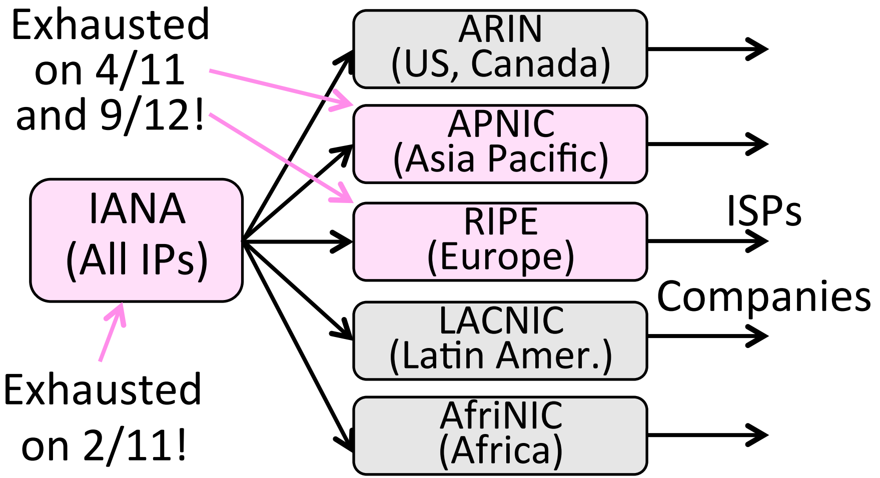
Internet Growth

- At least a billion Internet hosts and growing …
- And we're using 32-bit addresses!



The End of New IPv4 Addresses

- Now running on leftover blocks held by the regional registries; much tighter allocation policies



End of the world ? 12/21/12?

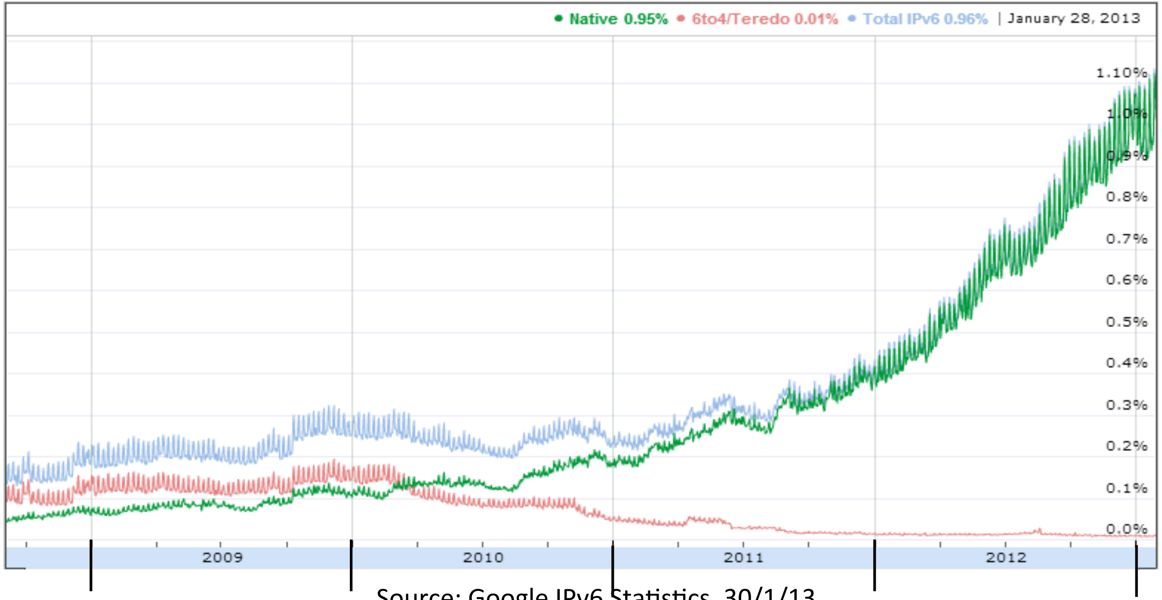
IP Version 6 to the Rescue

- Effort started by the IETF in 1994
  - Much larger addresses (128 bits)
  - Many sundry improvements

- Became an IETF standard in 1998
  - Nothing much happened for a decade
  - Hampered by deployment issues, and a lack of adoption incentives
  - Big push ~2011 as exhaustion looms

IPv6 Deployment

Percentage of users accessing Google via IPv6

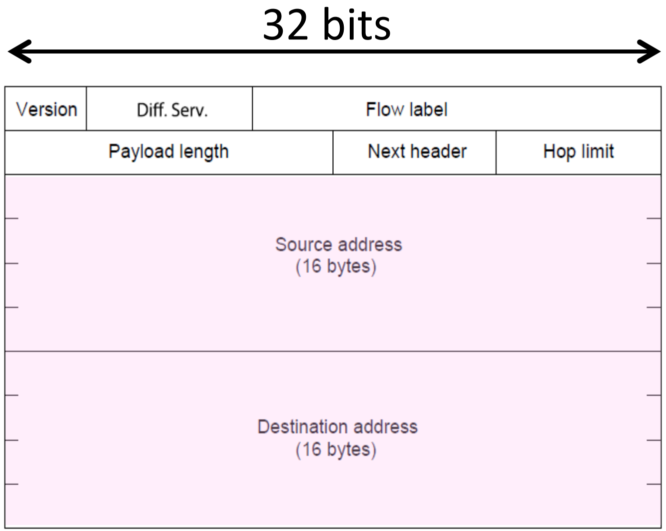


Source: Google IPv6 Statistics, 30/1/13

Time for growth!

IPv6

- Features large addresses
  - 128 bits, most of header
- New notation
  - 8 groups of 4 hex digits (16 bits)
  - Omit leading zeros, groups of zeros

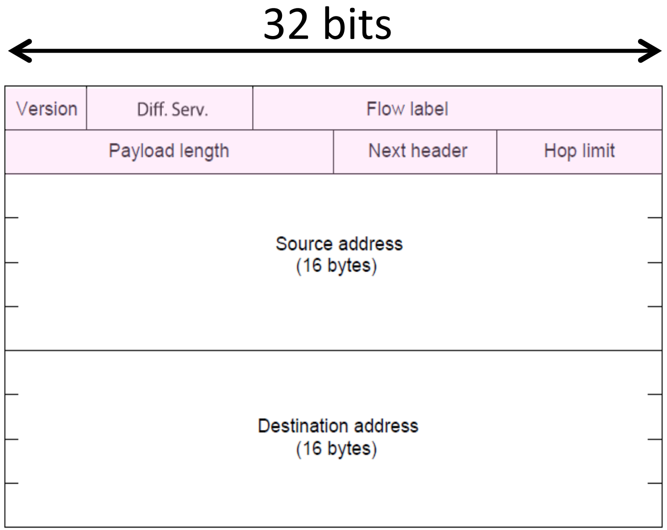


Ex: 2001:0db8:0000:0000:0000:ff00:0042:8329

→

IPv6 (2)

- Lots of other, smaller changes
  - Streamlined header processing
  - Flow label to group of packets
  - Better fit with “advanced” features (mobility, multicasting, security)

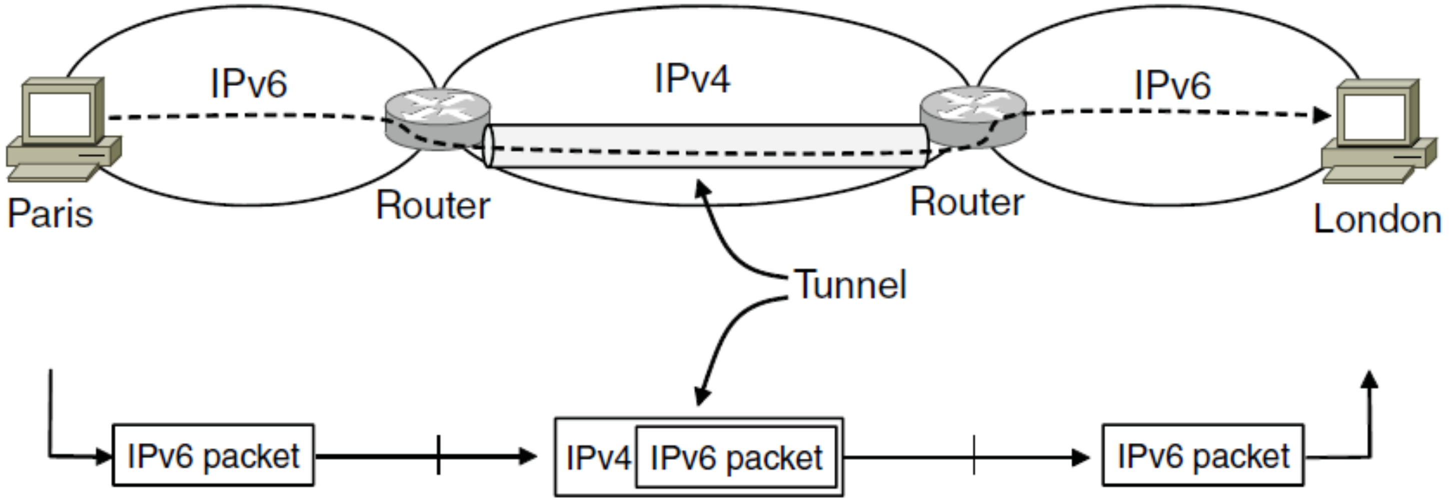


IPv6 Transition

- The Big Problem:
  - How to deploy IPv6?
  - Fundamentally incompatible with IPv4
- Dozens of approaches proposed
  - Dual stack (speak IPv4 and IPv6)
  - Translators (convert packets)
  - Tunnels (carry IPv6 over IPv4) »

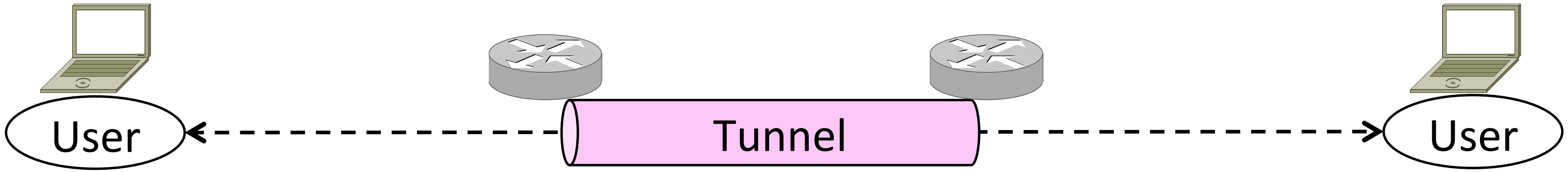
Tunneling

- Native IPv6 islands connected via IPv4
  - Tunnel carries IPv6 packets across IPv4 network



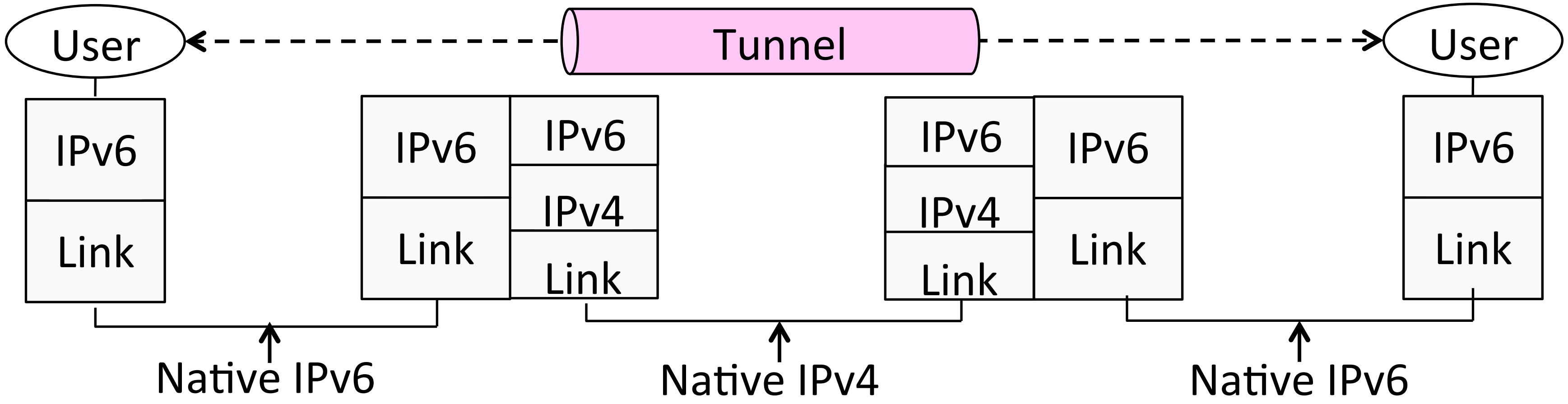
Tunneling (2)

- Tunnel acts as a single link across IPv4 network



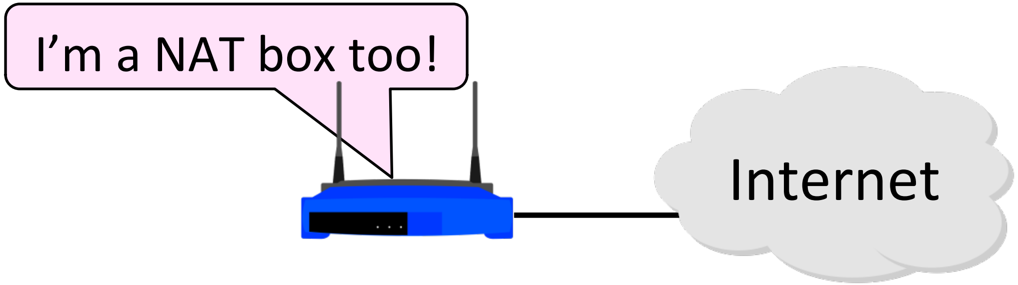
Tunneling (3)

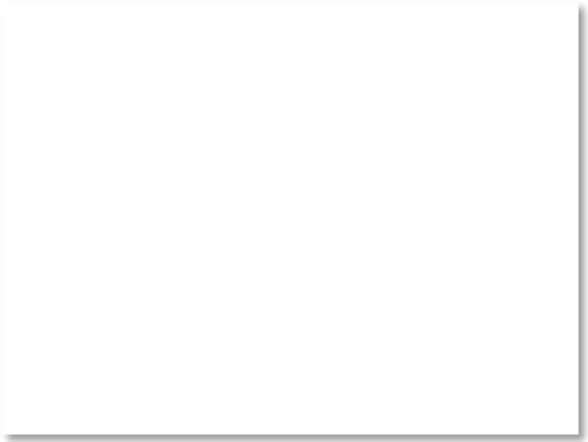
- Tunnel acts as a single link across IPv4 network
  - Difficulty is to set up tunnel endpoints and routing



Topic

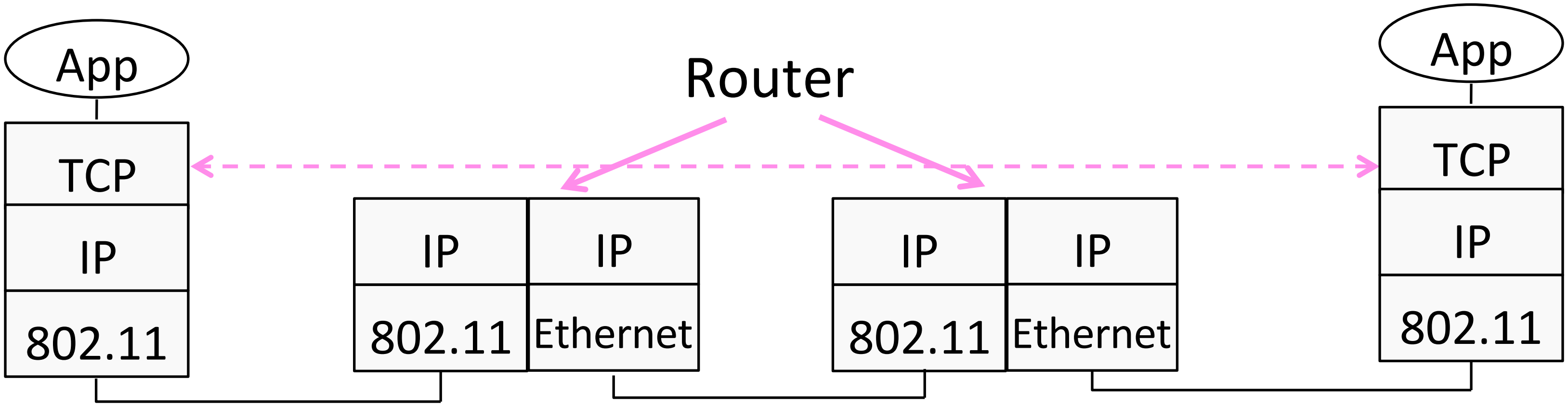
- What is NAT (Network Address Translation)? How does it work?
  - NAT is widely used at the edges of the network, e.g., homes





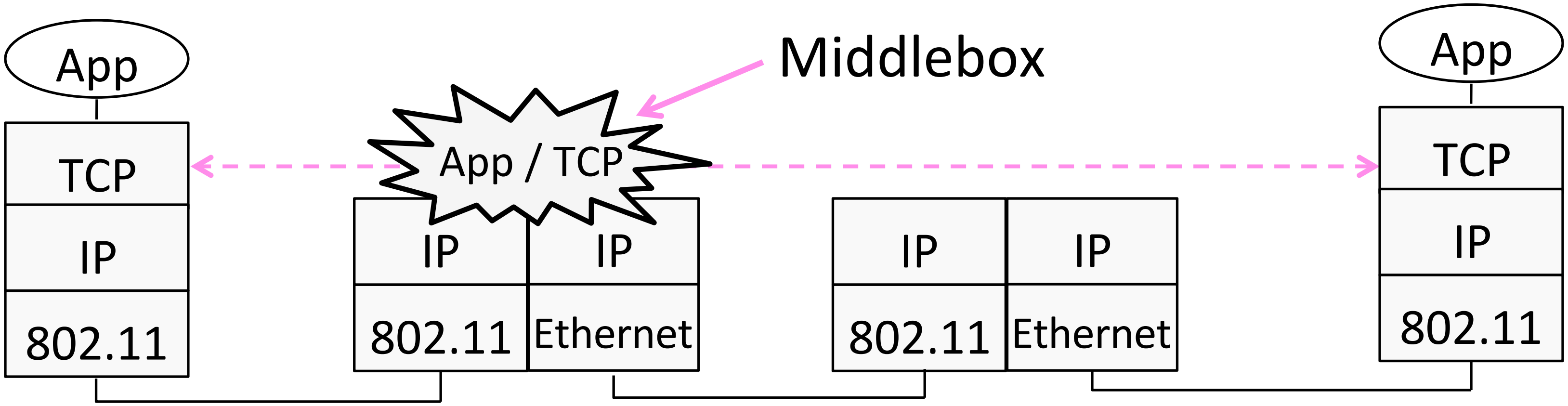
Layering Review

- Remember how layering is meant to work?
  - "Routers don't look beyond the IP header." Well …



Middleboxes

- Sit "inside the network" but perform "more than IP" processing on packets to add new functionality
  - NAT box, Firewall / Intrusion Detection System



Middleboxes (2)

- Advantages
  - A possible rapid deployment path when there is no other option
  - Control over many hosts (IT)

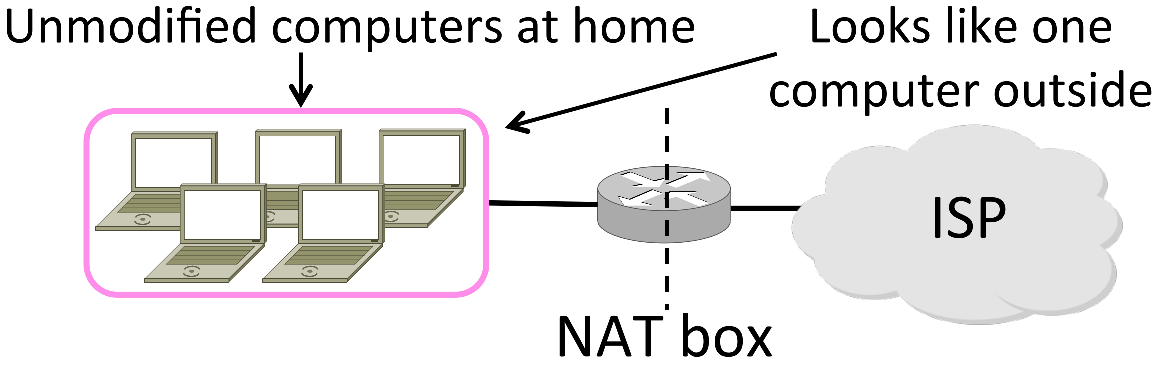
- Disadvantages
  - Breaking layering interferes with connectivity; strange side effects
  - Poor vantage point for many tasks

NAT (Network Address Translation) Box

- NAT box connects an internal network to an external network
  - Many internal hosts are connected using few external addresses
  - Middlebox that "translates addresses"
- Motivated by IP address scarcity
  - Controversial at first, now accepted

NAT (2)

- Common scenario:
  - Home computers use "private" IP addresses
  - NAT (in AP/firewall) connects home to ISP using a single external IP address



How NAT Works

- Keeps an internal/external table
  - Typically uses IP address + TCP port
  - This is address and port translation

What host thinks / What ISP thinks

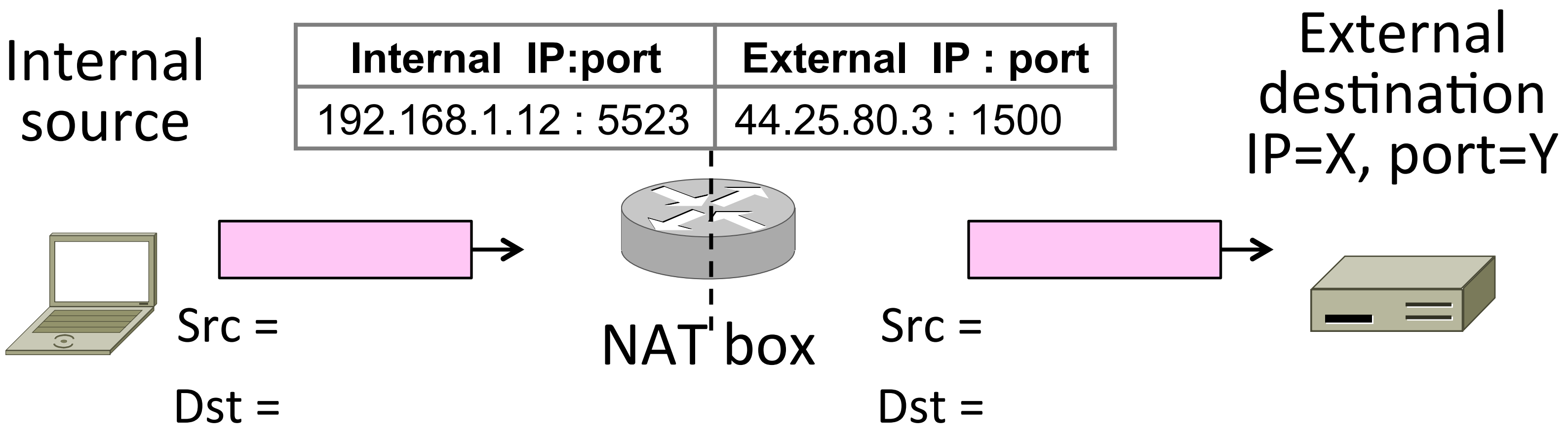
Internal IP:port	External IP : port
192.168.1.12 : 5523	44.25.80.3 : 1500
192.168.1.13 : 1234	44.25.80.3 : 1501
192.168.2.20 : 1234	44.25.80.3 : 1502

- Need ports to make mapping 1-1 since there are fewer external IPs

How NAT Works (2)

- Internal → External:
  - Look up and rewrite Source IP/port

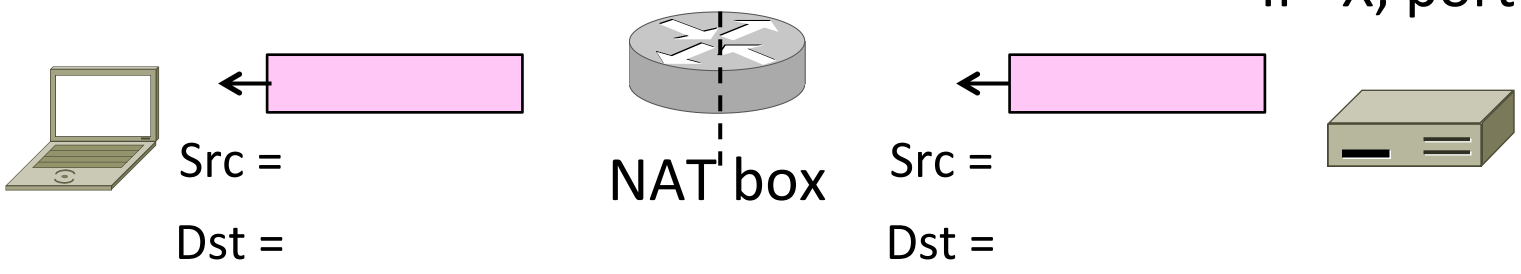
Internal IP:port	External IP : port
192.168.1.12 : 5523	44.25.80.3 : 1500



How NAT Works (3)

- External → Internal
  - Look up and rewrite Destination IP/port

Internal IP:port	External IP : port
192.168.1.12 : 5523	44.25.80.3 : 1500



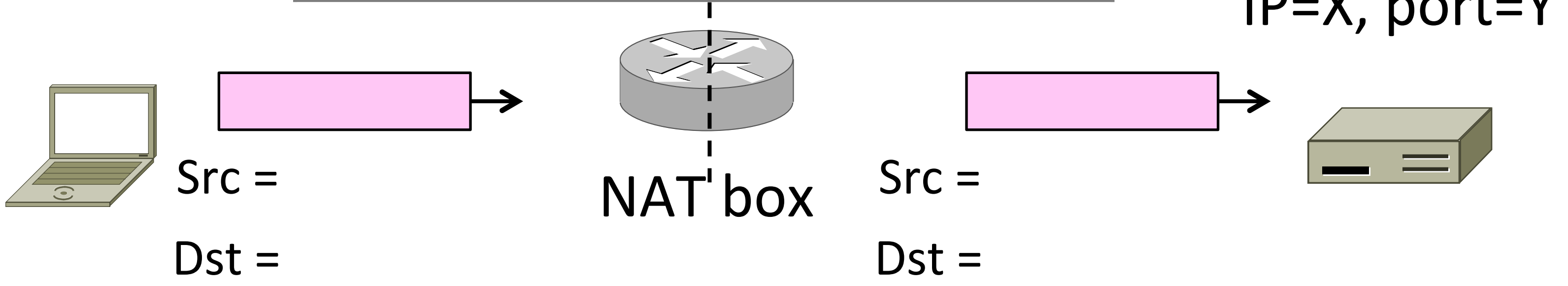
How NAT Works (4)

- Need to enter translations in the table for it to work
  - Create external name when host makes a TCP connection

Internal source

Internal IP:port	External IP : port
192.168.1.12 : 5523	

External destination IP=X, port=Y



NAT Downsides

- Connectivity has been broken!
  - Can only send incoming packets after an outgoing connection is set up
  - Difficult to run servers or peer-to-peer apps (Skype) at home
- Doesn't work so well when there are no connections (UDP apps)
- Breaks apps that unwisely expose their IP addresses (FTP)

NAT Upsides

- Relieves much IP address pressure
  - Many home hosts behind NATs
- Easy to deploy
  - Rapidly, and by you alone
- Useful functionality
  - Firewall, helps with privacy

- Kinks will get worked out eventually
  - "NAT Traversal" for incoming traffic